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DESCRIPTION

FIELD OF THE INVENTION

[0001] The present invention is related to the field of wireless communications. More specifically, the present invention relates to Wireless Local Area Network (WLAN) systems that use a Carrier Sense Multiple Access/Collision Avoidance (CSMA/CA) mechanism and provides means for determining and managing congestion and further enhances network management by providing novel medium access control (MAC) measurements in wireless communications.

BACKGROUND OF THE INVENTION

[0002] Wireless communication systems are well known in the art. Generally, such systems comprise communication stations, which transmit and receive wireless communication signals between each other. Depending upon the type of system, communication stations typically are one of two types: base stations or wireless transmit/receive units (WTRUs), which include mobile units.

[0003] The term base station as used herein includes, but is not limited to, a base station, Node B, site controller, access point or other interfacing device in a wireless environment that provides WTRUs with wireless access to a network with which the base station is associated.

[0004] The term WTRU as used herein includes, but is not limited to, a user equipment, mobile station, fixed or mobile subscriber unit, pager, or any other type of device capable of operating in a wireless environment. WTRUs include personal communication devices, such as phones, video phones, and Internet ready phones that have network connections. In addition, WTRUs include portable personal computing devices, such as PDAs and notebook computers with wireless modems that have similar network capabilities. WTRUs that are portable or can otherwise change location are referred to as mobile units. Generically, base stations are also WTRUs.

[0005] Typically, a network of base stations is provided where each base station is capable of conducting concurrent wireless communications with appropriately configured WTRUs. Some WTRUs are configured to conduct wireless communications directly between each other, i.e., without being relayed through a network via a base station. This is commonly called peer-to-peer wireless communications. Where a WTRU is configured to communicate with other WTRUs it may itself be configured as and function as a base station. WTRUs can be configured for use in multiple networks with both network and peer-to-peer communications capabilities.

[0006] One type of wireless system, called a wireless local area network (WLAN), can be configured to conduct wireless communications with WTRUs equipped with WLAN modems that are also able to conduct peer-to-peer communications with similarly equipped WTRUs. Currently, WLAN modems are being integrated into many traditional communicating and computing devices by manufacturers. For example, cellular phones, personal digital assistants, and laptop computers are being built with one or more WLAN modems.

[0007] A popular local area network environment with one or more WLAN base stations, typically called access points (APs), is built according to the IEEE 802.11 family of standards. An example 802.11 Local Area Network (LAN), as shown in Fig. 1, is based on an architecture, wherein the system is subdivided into cells. Each cell comprises a Basic Service Set (BSS), which comprises at least one AP for communicating with one or more WTRUs which are generally referred to as stations (STAs) in the context of 802.11 systems. Communication between an AP and STAs is conducted in accordance with the IEEE 802.11 standard that defines the air interface between a wireless STA and a wired network.

[0008] A wireless LAN (WLAN) may be formed by a single BSS, with a single AP, having a portal to a distribution system (DS). However, installations are typically composed of several cells, and APs are connected through a backbone, referred to as a DS.

[0009] A mobile ad-hoc network (MANET) is also shown in Figure 1. A MANET is a self-configuring network of mobile routers (and associated hosts) connected by wireless links-the union of which form an arbitrary topology. The routers are free to move randomly and organize themselves arbitrarily; thus, the network's wireless topology may change rapidly and unpredictably. Such a network may operate in a standalone fashion, or may be connected to the larger Internet.

[0010] An interconnected WLAN, including the different cells, their respective APs and the DS, is seen as a single IEEE 802.11 network and is referred to as an Extended Service Set (ESS). IEEE 802.11 networks typically use a Carrier-Sense Multiple Access / Collision Avoidance (CSMA/CA) protocol to exchange information wirelessly between nodes (or STAs) of the WLAN network. In this framework, STAs desiring to transmit must contend for access to the wireless medium. The contention mechanism involves waiting for the medium to remain idle for a certain period of time (according to a set of rules prescribed by the standard) before transmitting a data packet. The time it takes a node to access the channel and transmit its packet increases as the number of stations and data traffic increases. Congestion in such a system can occur when the time to gain access to the medium becomes intolerable due to too many stations competing for the same medium.

[0011] Due to the nature of the CSMA/CA protocol, and considering that most transmissions are best effort, it is quite difficult to determine when a system is classified as experiencing congestion. Determining congestion in such a complex system is not a simple task, as one choice of metrics could indicate congestion while another metric will not.

[0012] Several metrics that can be used to indicate congestion include: collision rate, channel utilization, i.e., the time that the medium is busy, etc. However, these metrics, taken individually do not necessarily give a true picture of the congestion. For example, the channel utilization metric does not give an accurate picture of the congestion situation. One station can be alone on a channel and transmitting all the time. In this case the channel utilization metric would be high. It may seem like the system would not be capable of supporting any more traffic from other stations. However, if a new station were to access the channel, it could still experience good throughput by virtue of the CSMA/CA mechanism, as the channel would then be equally shared between the two stations. A system is in fact congested when there are a number of stations contending for the same channel at a given time and experiencing severe delays due to the longer time each station has to wait for access to the medium, as well as the higher number of collisions.

[0013] In another aspect, there is currently limited network management functionality, particularly in systems compliant with the IEEE 802.11 and IEEE 802.11k standards. The inventors have recognized that there are certain limitations to the usefulness of channel loading information presently employed in the context of network management. There is also a need for an improved method of achieving better network management after considering the limitations of using channel-loading measurements. This present invention provides enhanced network management associated with the IEEE 802.11 and IEEE 802.11k standards in the context of channel loading information.

[0014] EP 1 156 623 discloses a communication system where a traffic load for a number of access points is monitored and a network station selects communication with an access point based on the traffic load.

[0015] Another example is known from Annese. A et al. "Providing delay guarantees in IEEE 802.11e networks", Proceedings /2004 59th Vehicular Technology Conference, VTC 2004-Spring: Towards a global wireless world;17-19 May 2004, Milan, Italy IEEE Operations Center, Piscataway, NJ, vol.4, 17 may 2004 (2004-05-17), pages 22345-2238.

SUMMARY

[0016] The present invention provides a method for determining and advertising congestion in a wireless local area network (WLAN) system. The present invention also provides a method for managing congestion when congestion is detected. One aspect of the present invention applies to wireless systems that use CSMA/CA. Preferably, several metrics are used to determine congestion including: average duration of backoff procedure, in-Basic Service Set (in-BSS) deferral rate, out-of-BSS deferral rate, number of associated stations, mean WTRU channel utilization, and average buffer Medium Access Control (MAC) occupancy. Actions taken to relieve congestion preferably include; sorting the set of WTRUs in order of most wasted time spent trying to transmit acknowledged/unacknowledged packets, and disassociating each WTRU one at a time until the congestion is relieved

[0017] The present invention also provides an improved method of network management, particularly in the context of standards IEEE 802.11 and IEEE 802.11k, preferably through the use of two (2) new MAC measurements. More specifically, the two (2) new measurements include STA uplink traffic loading measurement, and an Access Point (AP) service loading measurement.

[0018] The invention includes considerations of management information base (MIB) representation of the transmit queue size that provides a new measure of the STA transmit load in terms of unserved, queued traffic demand. The invention further includes considerations of MIB representation of the AP service load that provides a new measure of the AP service load to be used to assist STAs with handoff decisions. Implementation of these features can be as software or in any other convenient form. This aspect of the invention is generally applicable, for example, to layers 1 and 2 as applied to an IEEE 802.11k compliant

system in the context of orthogonal frequency division multiplexing (OFDM) and code division multiple access 2000 (CDMA 2000) systems. However, the invention has general applicability to other scenarios as well.

[0019] The methods are advantageously implemented in selectively configured WTRUs of various forms.

[0020] A more detailed understanding of the invention may be had from the following description of the preferred embodiments, given by way of example and to be understood in conjunction with the accompanying drawings.

BRIEF DESCRIPTION OF THE FIGURES

[0021]

Figure 1 is an overview diagram of a conventional IEEE802.11 WLANs with their corresponding components.

Figures 2-9 are flow diagrams illustrating the techniques of the present invention for determining and managing congestion in wireless communications systems. More particularly:

Figures 2 and 2A together present a method for determining congestion using deferral rate (DR) and packet error rate (PER) metrics and disassociating WTRUs based on determining wasted time trying to transmit/retransmit unacknowledged packets.

Figure 3 presents a method for managing load shedding by comparing the load of a node with advertised loads of neighboring nodes.

Figure 4 presents a method for providing an advertised load to WTRUs based on average delay between a packet reaching the head of a queue and transmission of the packet.

Figures 5, 6 and 7 present a method for respectively providing a transmit queue size (TQS), contention-free transmit queue size (CFTQS) and contention transmit queue size (CTQS) to neighboring nodes.

Figure 8 presents a method employed by a node for managing a channel based on evaluation of served and unserved traffic load from WTRUs and for providing a service load scalar for advertisement to WTRUs.

Figure 9 presents a method employed by WTRUs for selecting a node based on load scalars provided by neighboring nodes.

Figure 10 is a diagram of a BSS load element format in accordance with the present invention.

Figure 11 is a diagram of an access category service load element format in accordance with the present invention.

Figure 12 is a communication station configured in accordance with the present invention.

DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

[0022] Although the features and elements of the present invention are described in the preferred embodiments in particular combinations, each feature or element can be used alone (without the other features and elements of the preferred embodiments) or in various combinations with or without other features and elements of the present invention.

[0023] One aspect of the present invention introduces two different approaches to determine the loading metric of channel congestion; first, a Basic Service Set (BSS)-based load metric, which is based primarily on the load of individual APs. Second, a channel-based load metric, which is a metric indicating the load shared amongst different APs.

[0024] BSS-based load metrics are metrics that determine high load condition and channel congestion. The two preferred BSS-based load metrics are: in-BSS deferral rate metric, and packet error rate metric.

[0025] The Deferral Rate (DR) is a measurement that represents the percentage of time that the receiver of the AP is carrier locked (i.e. Clear Channel Assessment (CCA) indicates a busy condition) while the AP has one or more packets to transmit (i.e. it's queue is not empty). In other words, DR represents the amount of time that the AP spends deferring transmission to other WLAN nodes.

[0026] The in-BSS Deferral Rate represents the percentage of time that the receiver of the AP is carrier locked onto an in-BSS packet (i.e. a packet originating from one of its associated WTRUs) while the AP has one or more packets to transmit. In other words, the in-BSS DR represents the amount of time that the AP spends deferring its own transmissions because one of its associated WTRUs has taken control of the medium (i.e. is transmitting a packet).

[0027] The in-BSS deferral rate is indicative of the level of the current load placed in a system, and when there is a need to transmit to another node in the same BSS, measuring the time spent deferring a transmission. A low in-BSS deferral metric indicates that the load for the BSS is low. A high in-BSS deferral rate indicates that there are many nodes transmitting at the same time and that there is thus a significant load.

[0028] In a case where there are only two nodes in the system with a significant amount of data to transmit, the deferral rate could be high and if used alone will indicate congestion. However, since there are only two nodes in the system this is not considered a congestion situation. To address this situation, the present invention uses the packet error rate (PER) in addition to the deferral rate metric.

[0029] The Packet Error Rate (PER) is the ratio of the number of failed transmissions (i.e. packet transmissions for which an ACK was not received) to the total number of transmitted packets. The PER metric is a good indication of the collision rate in the system when conservative data transmission rates are used. The larger the number of nodes in a system, the higher the probability of collision. The use of both the in-BSS deferral rate metric and the PER metric together provide a better indication of the load of an AP than either metric used individually.

[0030] In the present invention, as shown in Figure 2, in-BSS deferral rate metric and PER metric are respectively determined, at steps S1 and S3 and are then averaged over a predefined period of time (e.g. 30 seconds), at steps S2 and S4, respectively. The averages of both metrics are used to signal the occurrence of congestion at steps S5 and S6. More specifically, when in-BSS deferral rate (DR) metric exceeds a first predefined threshold, determined at step S5, and the PER metric exceeds a second predefined threshold, determined at step S6, over a given period (e.g., 30 seconds), then this is an indication of congestion.

[0031] Whether or not congestion is detected based on the criteria as set forth above, or employing other techniques for determining congestion, the present invention provides the following actions; first, the AP at step S7, sorts all WTRUs in the Basic Service Set (BSS) in order of the amount of time spent trying to retransmit. Wasted time is preferably determined in accordance with the wasted time algorithm ALG_{WT} set forth below. More specifically, a set or list of WTRUs with unacknowledged packets is created. For each unacknowledged packet to a WTRUs, the sum of all the wasted time spent trying to transmit and re-transmit the packet (i.e. packet size / packet transmission rate plus a penalty for each retransmitted packet) is recorded. The penalty reflects the increasing delay associated with retransmissions, i.e. the backoff time due to the doubling of the congestion window (CW). The penalty represents the added delay incurred from the time the packet is ready for transmission to the time the packet is actually transmitted over the medium. This retransmit time metric is therefore much greater for stations wasting time retransmitting packets following collisions. The retransmit time metric is normalized over a selected time period.

[0032] An example formula for determining wasted time for a WTRU is given by:

$$wasted_txime_{WTRU} = \sum_{unackPkts} \sum_{i=1}^{\#_pkts_j} \left(\frac{Pkt_size_{ij}}{Pkt_tx_rate_{ij}} + RTx_{i>1} * Penalty \right)$$

where:

$wasted_time_{WTRU}$ = sum of wasted time spent trying to transmit and retransmit unacknowledged packets to a WTRU

$j = j^{th}$ packet

$i = i^{th}$ transmission of j^{th} packet

$\#_pkts_j$ = # of transmissions of j^{th} packet, e.g. 1, 2, 3,... Pkt_size_{ij} = size in bits of i^{th} transmission of j^{th} packet

$Pkt_tx_rate_{ij}$ = transmission rate in bps of i^{th} transmission of j^{th} packet

$RTx_{i>1} = 2^{i-2}$, for $i > 1$, otherwise 0

$Penalty = CW_{min} * slot\ time$, e.g. $CW_{min} = 32$ & slot time = 20 μ s Note: CW will be 2 x CW_{min} after first transmission.

Note that $\#_pkts_j$ corresponds to the number of unacknowledged transmissions of a given packet. If the packet is eventually successfully transmitted, $\#_pkts_j$ corresponds exactly to the number of retransmissions. If the packet is dropped (i.e. never successfully transmitted), $\#_pkts_j$ corresponds to (number of retransmissions + 1).

[0033] An example of the $wasted_tx_time_{STA}$ calculation is given below:

Assume that an AP has 20 packets to send to a particular STA. During the course of the transmissions, the AP monitors and records whether the packet has been successfully acknowledged or not and the number packet retransmissions as, for example, follows:

GGGGG**BBB**↓**BBB**↓GGGGG↑GGGGG↑**BBB**↓GGGG

where:

↑ = rate increase,

↓ = rate decrease,

G = acknowledged or "good" frame,

B = unacknowledged or "bad" frame

The 1st **B** is the sixth packet and there were six transmissions of this sixth (6th) packet, i.e. **BBB**↓**BBB**.

$\#_pkts_6 = 6$

$Pkt_size_{i6} = 12000$ bits

$Pkt_tx_rate_{i6} = \{11.0, 11.0, 11.0, 5.5, 5.5, 5.5\}$ Mbps $RTx_{i>1} * Penalty = \{0.0, 640.0, 1280.0, 2560.0, 5120.0, 10240.0\}$ us

The 7th **B** is the 17th packet and there were three transmissions of this 17th packet, i.e. ↑**BBB**↓.

$\#_pkts_{17} = 3$

$Pkt_size_{i17} = 8000$ bits

$Pkt_tx_rate_{i17} = \{11.0, 11.0, 11.0\}$ Mbps

$RTx_{i>1} * Penalty = \{0.0, 640.0, 1280.0\}$ us

Therefore:

$wasted_tx_time_{STA} = (12000/11e6) + (12000/11e6 + 640.0) + (12000/11e6 + 1280.0) + (12000/5.5e6 + 2560.0) + (12000/5.5e6 + 5120.0) + (12000/5.5e6 + 10240.0) + (8000/11e6) + (8000/11e6 + 640.0) + (8000/11e6 + 1280.0) = \mathbf{33.76\ ms}$

[0034] Preferably, the WTRUs are sorted from greatest to smallest times at step S7-4. The program then advances to step S8. At step S8 (Figure 2), each STA from the sorted list is disassociated greatest time first, until the congestion is relieved.

[0035] The present invention also provides for the use of other metrics including: BSS-based load metrics; the number of associated WTRUs, the time that the Access Point (AP) receives all acknowledgements (ACKS) (e.g. fragmentation) related to that packet at the medium access control (MAC), and the average buffer MAC occupancy (based on the size of the buffer).

[0036] The present invention further provides a method that takes into account the load of the neighboring APs in assessing the system's need to perform any load shedding (i.e. disassociation) or load balancing. For example, as shown in Figure 3, if the load of each of the neighboring APs is also high, as collected at steps S9 and S10, and compared with neighboring APs at steps S11 and S12, load shedding is delayed (step S14) since the user would have a low probability of being served elsewhere, i.e., L1, L2 and L3 are all high (step S13). Load shedding is conducted, at step S16 if L1 or L2 have lower advertised loads (step S15B). If the L3 load is less than L1 and L2, the AP can accept a WRTU, as shown at steps S15A and S17.

[0037] For advertising loading to its stations (WTRUs), an Access Point (AP) can compare its load relative to neighboring APs, i.e. AP(x) and AP(y), for example. When an AP load is high compared to the estimated load of its neighboring APs, then the AP advertises a high load responsive to a determination at step S15A (Figure 3). When the AP load is low compared to the estimated load of its neighbors, the AP advertises a low load responsive to a determination at step S15B.

[0038] Another method of the present invention is to use metrics that determine medium (i.e., channel) load. This metric enables the WTRU to choose the least loaded AP. Medium load metrics are used in cases when the In-BSS channel load is not effective, such as the case when a BSS with an In-BSS channel load could simply be deferring to a neighboring BSS, and therefore, although the load of the AP is low, the medium load is high. In this case, the advertised load should be representative of the medium load. In this case, an AP only advertises a low load when it is able to support the new WTRU.

[0039] A metric that gives an indication of the medium load is the average duration (Avg D) required to execute the backoff procedure that is determined in the manner shown in Figure 4 for downlink transmissions at an AP. More specifically, this metric represents the medium access delay incurred from the time a packet is ready for transmission (i.e. begins CSMA/CA access contention) to the time the packet starts transmission over the medium as determined at steps S18-S23, and advertising AvgD to WTRUs, at step S24.

[0040] The size of the contention window influences the duration needed to execute the backoff procedure. The contention window size is increased whenever an acknowledgement is not received from the receiving node. This aspect covers cases where collisions occur either between nodes of the same BSS or different BSSs. During the countdown of a backoff procedure, the countdown is suspended whenever the medium is sensed to be busy, which increases the duration of the backoff procedure. This additional aspect covers the cases when the medium is highly loaded due to WTRUs of the own BSS and/or neighboring BSSs. This metric taken alone provides a good indication of the congestion as perceived by this node in the BSS. One could consider simply using the time that the medium is busy (channel utilization) as a metric. However, in an example where only one WTRU is associated with the Access Point (AP) and is transmitting or receiving large amounts of data, the channel utilization metric will not give a good indication of the congestion. Channel utilization will indicate a high congestion when in fact the system is only supporting one user. A second user (WTRU) added to this AP could easily be supported. In the single user example, the new proposed Avg. D metric (i.e. the average duration to execute the backoff procedure) would correctly indicate low congestion.

[0041] The AvgD metric is a preferred measure since a short duration required for the backoff procedure indicates a lightly loaded medium, where a long duration indicates a heavily loaded medium. As an example, consider the current IEEE 802.11b standard. The minimum value for a contention window (CW) is $32 \times 20 \mu\text{sec} = 640 \mu\text{sec}$, and the maximum value is $1023 \times 20 \mu\text{sec} = 20.5 \text{msec}$. However, the duration required to execute the backoff may be greater than the maximum size of the CW, caused by the suspension of the countdown due to sensing a busy medium. This increase in duration will give an indication in load due to the activity in the medium.

[0042] The reasons for the use of MAC loading measurements in the context of the present invention include:

- The MAC layer has much information, which is not currently available via the management information base (MIB) or via measurements in the standard IEEE 802.11 and IEEE 802.11k.
- New information items provided by the present invention, which are useful to upper layers, are not presently available although they can be provided within the scope of 802.11k.
- IEEE 802.11e has identified channel utilization (CU) as a useful loading information item.

[0043] The present invention also recognizes that there is need for WTRU uplink loading information and AP service loading information. Some of the limitations of CU information include:

- Loading information is useful for handoff decisions in the WTRU and AP.
- CU information of a potential target AP is useful to WTRU when assessing handoff options.
- CU is the sum of uplink served load (all WTRUs to AP) and downlink served load (AP to all WTRUs), also known as channel utilization.
- Traffic load, however, consists of two parts: served traffic load and unserved (queued) traffic load.
- CU presently does not provide dynamic, unserved, queued traffic load information.

[0044] The network has no current way to access unserved uplink traffic demand (queued traffic load).

[0045] The merits of WTRU uplink traffic loading measurements (UTLM) in network management include:

- A high channel load indicates served traffic close to maximum.
- If unserved traffic demand is low, this is optimal channel management.
- If unserved traffic demand is high, this is sub-optimal.
- Unserved uplink traffic demand is extremely useful to enable an AP to better partition uplink and downlink segments of frame time.
- APs need to manage the channel for maximum traffic utilization and minimal traffic blocking.
- Queued uplink traffic at WTRUs indicates transmission delays and potential channel blockage.
- The volume of data queued in the MAC transmission buffers provide a good measure of queued uplink load.

[0046] The present invention provides a new MAC management information base (MAC MIB) element for transmit traffic load, namely, Transmit Queue Size (TQS). Transmit Queue Size (TQS) is defined as follows: New MIB Information contains three (3) items: Total transmit queue size (TQS) consisting of the sum of Contention-free TQS (CFTQS) and Contention TQS (CTQS).

[0047] TQS contains the current MAC queue size in bytes. TQS can be included in a MAC MIB 802.11 Counters Table. Dot11Counters Table is a defined data structure in the standard. TQS information may be implemented by a counter as shown in Fig. 5, the WTRU, at step S25, initializes the TQS counter to zero upon system start up. The WTRU, at step S26, receives a frame and, at step S27, queues the frame in the MAC layer. At step S28, the WTRU increments the TQS counter by the number of bytes in the queued frame. Alternatively, accumulation may use a software technique wherein a count may be stored in a memory and incremented by replacing a present count (PC) with PC+1, for example, as each byte of the frame is queued.

[0048] The WTRU, at step S29, transmits a frame employing the physical (PHY) layer when a session is initiated and, at step S30, decrements the TQS counter by the number of bytes transmitted, either when operating in the unacknowledged mode or when a frame is acknowledged by an AP after the PHY transmission. The WTRU, at step S31, communicates the TQS count to neighboring APs. TQS is a new MIB element. All MIB elements are transmitted to neighbors as needed via an MIB query performed to retrieve an element from a neighbor's MIB.

[0049] The contention transmit queue size (CTQS) is implemented as shown, for example, in Figure 6, wherein the WTRU, at step S32, initializes the CTQS counter to zero at system startup. The MAC layer of the WTRU, at step S33, receives a contention frame and, at step S34, queues it in the contention queue of the MAC layer. At step S35, the CTQS counter is incremented by the number of bytes in the received frame.

[0050] The WTRU, at step S36, transmits the frame (to an AP, for example) employing the PHY layer when operating either in the unacknowledged mode or when the frame has been acknowledged after PHY transmission and, at step S37, decrements the CTQS counter by the number of bytes transmitted either in unacknowledged mode or when the frame is acknowledged after a PHY layer transmission. At step S38 the WTRU communicates the CTQS count to neighboring APs.

[0051] The contention free transmit queue size (CFTQS) is implemented, as shown in Figure 7, by providing a CFTQS counter wherein the WTRU, at step S39, initializes the CFTQS counter to zero at system startup.

[0052] At step S40, the WTRU MAC layer receives a contention-free frame and, at step S41, queues the frame in the contention free queue (CFQ). At step S42, the WTRU increments the CFTQS counter by the number of bytes in the queued frame.

[0053] At step S43, the WTRU transmits a contention-free frame using the PHY layer and, at step S44, decrements the CFTQS counter by the number of bytes transmitted in the frame in the unacknowledged mode or when the frame is acknowledged after the PHY layer transmission. At step S45 the WTRU communicates the count to neighboring APs.

[0054] Fig. 8 shows one manner in which an AP utilizes the MAC MIB information, wherein the AP, at steps S46, S47 and S48, for example, respectively, receive MAC MIB information including one or more of the TSQ, CTQS and CFTQS counts, from WTRU(x), WTRU(y) and WTRU(z), for example. This data, which represents unserved traffic, is combined with served traffic data such as channel loading which includes both the uplink and downlink load, and is evaluated by the AP, at step S49 and, at step S50, utilizes the served and unserved load data to manage the channel, for example, by adjusting the traffic to maximize traffic utilization and minimize traffic blocking. The AP may adjust the uplink and downlink segments of frame, based upon unserved uplink traffic data, in order to optimize channel utilization.

[0055] The considerations for providing AP service loading measurements in the context of the invention include the following:

[0056] WTRUs may consider multiple APs as target APs for handoff. If two APs have similar channel loading and acceptable signal quality, the WTRU needs a capability of being able to determine which is the better AP. By enabling APs to post information concerning their ability to serve their existing set of WTRUs and their ability to serve additional WTRUs, channel usage can be optimized. This information is similar to a downlink traffic queue measurement for the AP modified by any AP specific information concerning its anticipated capacity.

[0057] The following addresses AP Service Load:

A new MAC MIB information item is provided to assist WTRUs in their handoff decisions.

[0058] A quantitative indication on a 255-value scale (represented by 8 binary bits, for example), from "not currently serving any WTRU", to "can't handle any new services" with a defined middle point indicating that the served load is optimal. For example:

0 == Not serving any WTRU (idle AP or WTRU is not an AP)

1 through 254 == scalar indication of AP Service Load.

255 == unable to accept any new services

[0059] Exact specification of this MIB item is implementation-dependant and need not be specified with exactitude; a detailed definition to obtain maximum utility may be tailored to the characteristics of the particular network.

[0060] The new AP Service Load can be included in MAC dot11Counters Table or elsewhere in the MIB.

[0061] A WTRU having multiple APs that can be chosen as a target AP, in addition to a consideration of channel loading and acceptable signal quality, as shown in Figure 9, can receive load advertisements from AP(x), AP(y) and AP(z), respectively shown at steps S51, S52 and S53, and, at step S54 evaluates the received AP advertised loads (SL scalars) and thus is able to make a decision based upon comparisons of the AP advertised loads received and, at step S55 selects an AP.

[0062] The AP service load (SL) is a scalar value and may, for example, be based upon served and unserved traffic, as well as other data such as signal quality, and anticipated capacity, based on statistical data, for example. The AP SL scalar may be created, as shown in step S50A of Figure 8 and advertised to the neighboring WTRUs, as shown at step S50B.

[0063] The above methods are preferably implemented in selectively configured WTRUs. For example, a WTRU can be configured to assist in channel management in a wireless network by providing a memory device, a processor and a transmitter. The memory device is preferably configured to provide a queue of data frames for a medium access control (MAC) layer of the WTRU. The processor is preferably configured to determine queue size data representing unserved, queued traffic demand at the respective WTRU. The transmitter is preferably configured to communicate the queue size data to access points (APs) of the wireless network whereby a receiving AP utilizes the queue size data to assist in channel management. In particular, the processor is configured to initialize at zero a count representing queued data size at system startup and to increment the count by a number of bytes in a frame when the frame is queued by the medium access control (MAC) layer of the WTRU. Preferably the processor is configured to decrement the count by a number of bytes in a frame when a frame is transmitted by a physical (PHY) layer of the WTRU in an unacknowledged mode. As an alternative, the processor can be configured to decrement the count by a number of bytes in a frame when a frame is transmitted by a physical (PHY) layer of the WTRU when the frame has been acknowledged after a PHY transmission.

[0064] In such a WTRU, the memory is preferably configured with contention and contention free queues of the medium access control (MAC) layer and the processor is configured to determine contention transmit queue-size (CFTQS) data representing unserved, queued traffic demand for the contention queue, contention free transmit queue-size (CTQS) data representing unserved, queued traffic demand for the contention free queue and total transmit queue-size (TQS) data representing unserved, queued traffic demand for all transmit data queues of a medium access control (MAC) layer.

[0065] Such a WTRU preferably also includes a receiver configured to receive from APs service load indicators formulated

based on queue size data received from WTRUs by the APs and a controller configured to select an AP for wireless communication based on the received load indicators.

[0066] An access point (AP) can be provided configured to provide channel management in a wireless network for both access points (APs) and wireless transmit receive units (WTRUs) capable of wireless communications with the APs over wireless channels. A receiver is configured to receive unserved traffic demand data received from WTRUs located within a wireless service range of the AP. The AP preferably has a processor configured to calculate a service load indicator based on unserved traffic demand data received from WTRUs. A transmitter is included that is configured to advertise the service load indicator to WTRUs within the AP wireless service range whereby WTRUs located within the AP wireless service range of the AP can use the advertised service load indicator to assist in selection of an AP with which to conduct a wireless communication. In such an AP, the receiver is preferably configured to receive advertised service load indicators from other APs and the processor is preferably configured to use the advertised service load indicators received from other APs to assist in decisions regarding disassociating operatively associated WTRUs from communications with the AP.

[0067] In another embodiment, a wireless transmit receive unit (WTRU) is configured to manage congestion in a wireless communication system defined by a base service set (BSS). The WTRU has a processor configured to determine an in-base service set (in-BSS) deferral rate (DR) and average said DR over a given time interval. Preferably, the processor is configured to also determine packet error rate (PER) and average said PER over said time interval. A memory is configured to store comparative values reflecting wasted time spent trying to transmit data for each of the WTRUs operatively associated with the WTRU in the BSS. A transceiver is included that is configured to disassociate operatively associated WTRUs from the WTRU commencing with a WTRU having a stored comparative value reflective of the greatest time spent trying to transmit data when said average DR and said average PER are greater than given thresholds.

[0068] In such a WTRU, the processor is preferably configured to average the DR and the PER over a time interval of the order of thirty seconds and the transceiver is configured to periodically receive and update the memory with comparative values reflecting wasted time spent trying to transmit data for each WTRU operatively associated with the WTRU.

[0069] In such a WTRU, the processor may also be configured to determine a comparative wasted time value by measuring the time it takes the WTRU to receive either a successful acknowledge (ACK) or negative acknowledgment (NACK) responsive to a transmitted data packet, summing the measured times during a beacon period and normalizing the sum by the beacon period. The transceiver is then preferably configured to periodically transmit current comparative values reflecting wasted time spent trying to transmit data to other WTRUs.

[0070] An access point AP may also be configured to assist wireless transmit receive stations (WTRUs) in selecting an access point AP with which to conduct wireless communication in a wireless communication system by providing it with selectively configured components. Preferably, a receiver is configured to receiving advertised load indicators of other APs. A processor is included that is configured to compare a communication load of the AP with received advertised load indicators from other APs and to determine an adjusted load of the AP based on said comparison. A transmitter is configured to advertise the adjusted AP load to WTRUs. Preferably, the processor is configured to periodically perform said comparing and determining operations in order to update the load that transmitter advertises to WTRUs.

[0071] In such an AP, the transmitter may be configured to advertise a low load when the processor determines that the communication load of the AP is low compared to the advertised load of other APs and to advertise a high load when the processor determines that the communication load of the AP is high compared to the advertised load of other APs. Also, the processor can be configured to determine a communication load of the AP by measuring delay between a time when a data packet is ready for transmission and a time when the packet is actually transmitted to a WTRU, averaging said delay over a given period, and utilizing the average delay to indicate load.

[0072] In another embodiment, a base station is configured to disassociate WTRUs from operative association therewith when a congestion condition is detected in a wireless network. The base station has a processor configured to determine wasted time (Tw) spent attempting to transmit/retransmit unacknowledged packets for each associated WTRU and to normalize wasted time Tw for each associated WTRU over a given time period. A memory is provided that is configured to store a list of associated WTRUs and their respective normalized wasted times. A transceiver is configured to disassociate WTRUs to relieve said congestion based on their respective normalized wasted times whereby a WTRU having a greatest Tw is disassociated first. Preferably, the processor is configured to add a penalty to said Tw representing increasing delay associated with retransmissions such as by being configured to calculate wasted transmission time (Tw) of WTRUs according to the formula set forth above.

[0073] IEEE 802.11e supports several access categories such as, for example, voice, video, best effort, and background traffic. In one embodiment, the present invention preferably utilizes the AP service load per access category. The BSS Load element contains information on the current station population, traffic level, and service level in the BSS. Figure 10 shows an example of the element information fields in accordance with the present invention.

[0074] The Length field shall be set to the number of octets in the following fields. The Station Count field is interpreted as an unsigned integer that indicates the total number of STAs currently associated with this BSS. The Station Count field shall not be present in beacon or probe response frames if, purely by way of example, dot11QoSOptionImplemented, dot11QBSSLoadImplemented, and dot11RadioMeasurementEnabled are all true.

[0075] The Channel Utilization field is defined as the percentage of time the AP sensed the medium busy, as indicated by either the physical or virtual carrier sense mechanism. This percentage is represented as a moving average of $((\text{channel busy time} / (\text{dot11ChannelUtilizationBeaconIntervals} * \text{dot11BeaconPeriod} * 1024)) * 255)$, where channel busy time is defined to be the number of microseconds during which the carrier sense mechanism has indicated a channel busy indication, and dot11ChannelUtilizationBeaconIntervals represents the number of consecutive beacon intervals during which the average should be calculated. The Channel Utilization field shall not be present in beacon or probe response frames if dot11QoSOptionImplemented, dot11QBSSLoadImplemented, and dot11RadioMeasurementEnabled are all true.

[0076] The AP Service Load shall be a scalar indication of the relative level of service loading at an AP. A low value shall indicate more available service capacity than a higher value. The value 0 shall indicate that this AP is not currently serving any STA. The values between 0 and 254 shall be a logarithmically scaled representation of the average medium access delay for DCF transmitted packets measured from the time the DCF packet is ready for transmission (i.e. begins CSMA/CA access) until the actual packet transmission start time. A value of 1 shall represent a 50 μ s delay while a value of 253 shall represent a 5.5 ms delay or any delay greater than 5.5 ms. The value 254 shall indicate no additional AP service capacity is available. The value 255 shall indicate that the AP Service Load is not available. The AP shall measure and average the medium access delay for all transmit packets using DCF access mechanism over a predetermined time window, such as a thirty second measurement window. The accuracy for the average medium access delay shall be $\pm 200 \mu$ s or better when averaged over at least 200 packets.

[0077] The Access Category (AC) Service Load elements may be provided in the BSS Load only at QoS enhanced APs (QAPs). The AC Service Load shall be a scalar indication of the Average Access Delay (AAD) at a QAP for services of the indicated Access Category. A low value shall indicate shorter access delay than a higher value. The value 0 shall indicate that this QAP is not currently providing services of the indicated AC. The values between 0 and 254 shall be a logarithmically scaled representation of the average medium access delay for transmitted packets in the indicated AC measured from the time the EDCF packet is ready for transmission (i.e. begins CSMA/CA access) until the actual packet transmission start time. A value of 1 shall represent a 50 μ s delay while a value of 253 shall represent a 5.5 ms delay or any delay greater than 5.5 ms. The value 254 shall indicate that services at the indicated AC are currently blocked or suspended. The value 255 shall indicate that the AC Service Load is not available.

[0078] The QAP shall measure and average the medium access delay for all transmit packets of the indicated AC using EDCF access mechanism over a predetermined time window, such as a continuous thirty second measurement window. The accuracy for the average medium access delay shall be $\pm 200 \mu$ s or better when averaged over at least 200 packets. The AC Service load is preferably formatted as shown in Figure 11, as two octet sub elements with the first octet containing the AC Indication (ACI) and the second octet containing the measured value of the AAD for the indicated AC. It should be noted that the octets shown in Figures 10 and 11 are provided just as an example and any other octet may be utilized. Table 1 shows an example of ACI encoding.

Table 1

Access Category (AC)	ACI
Best Effort	0
Background	1
Video	2
Voice	3
Reserved	4-255

[0079] Referring now to Figure 12, there is shown a communication station 100 configured in accordance with the present invention. It is noted that the communication station 100 may be an access point (AP), WTRU, or any other type of device capable

of operating in a wireless environment. The communication station 100 preferably includes a receiver 102 configured to receive unserved traffic demand data from WTRUs located within a wireless service range 108 of the communication station 100. The communication station 100 also includes a processor 104. The processor 104 is preferably coupled to the receiver 102 and is configured to calculate a BSS load element for each of plurality of access categories. The communication station 100 also includes a transmitter 106. The transmitter 106 is preferably configured to advertise the BSS load element within a service range 108 of the communication station 100. The BSS load element may then be received by other communication stations (e.g. access points and/or WTRUs) within the service range 108 of the communication station 100 thereby providing them with information regarding the BSS.

Embodiments**[0080]**

1. A method for providing channel management in a wireless network for optimizing network utilization by both access points (APs) and wireless transmit receive units (WTRUs) capable of wireless communications with each other on wireless channels, comprising creating a service load indicator by a first AP for each access category.
2. The method of embodiment 1 further comprising advertising the service load indicator to WTRUs within a service range of the first AP.
3. The method of any preceding embodiment further comprising selecting an AP by the WTRU based on the service load indicator.
4. The method of any preceding embodiment wherein the service load indicator is an indication of average access delay at the first AP.
5. The method of embodiment 4, wherein the average access delay is measured in a predetermined time period.
6. The method of embodiment 5, wherein the time period is thirty (30) seconds.
7. The method of any preceding embodiment wherein the access categories include voice, video, best effort, and/or background traffic.
8. The method of any preceding embodiment further comprising receiving the advertised service load indicator from by a second AP.
9. The method of embodiment 8, further comprising using the advertised service load indicator in deciding disassociation of WTRUs by the second AP.
10. The method of any of embodiments 8-9, wherein the second AP disassociates WTRUs with the second AP where the service load indicator from the first AP is low compared to a service load indicator determined by the second AP.
11. An access point (AP) configured to provide channel management according to a method of any of the preceding embodiments.
12. The AP of embodiment 11, comprising a processor configured to calculate a service load indicator for each access category.
13. The AP of any preceding embodiment comprising a transmitter configured to advertise the service load indicator to WTRUs within the AP wireless service range.
14. The AP of any preceding embodiment whereby WTRUs located within the AP wireless service range of the AP can use the advertised service load indicator to assist in selection of an AP with which to conduct a wireless communication.
15. The AP of any preceding embodiment comprising a receiver configured to receive advertised service load indicators from other APs.
16. The AP of any preceding embodiment, wherein the processor is configured to use the advertised service load indicators received from other APs to assist in decisions regarding disassociating WTRUs with the AP.
17. A wireless transmit/receive unit (WTRU) configured to provide channel management in a wireless network according to a method of any of the preceding embodiments.

18. The WTRU of embodiment 17 comprising a receiver for receiving a service load indicator for each access category from an AP.
19. The WTRU of any of embodiments 17-18 comprising a processor configured to utilize the service load indicator in selection of an AP with which to conduct a wireless communication.
20. A method for providing channel management in a wireless network to optimize network utilization by communication stations capable of wireless communications with each other on wireless channels, comprising a first communication station providing a basic service set (BSS) load element for each of a plurality of access categories.
21. The method of embodiment 20, further comprising advertising the BSS load element to other communication stations within a service range of the first communication station.
22. The method of any of embodiments 20-21, further comprising at least one communication station selecting another communication station with which to communicate based on the BSS load element.
23. The method of any of embodiments 20-22 wherein the BSS load element includes an element identification field.
24. The method of any of embodiments 20-23 wherein the BSS load element includes a communication station, AP, or WTRU service load field, wherein said communication station, AP or WTRU service load field is a scalar indication of a relative level of service loading at the first communication station.
25. The method of any of embodiments 20-24 wherein the BSS load element includes a length field whose value is set to a total number of octets included in all fields of the BSS load element.
26. The method of any of embodiments 20-25 wherein the BSS load element further includes a station count field, wherein said station count field is an unsigned integer that indicates a total number of communication stations associated with a current BSS.
27. The method of any of embodiments 20-26 wherein the first communication station is a quality of service (QoS) enhanced communication station (QCS) or QoS enhanced AP (QAP).
28. The method of embodiment 27 wherein said BSS load element further includes an access category (AC) service load field, said AC service load field being formatted as four sub-fields, one each for providing a scalar indication of an average-access-delay (AAD) at the QCS or QAP for services of one of the access categories.
29. The method of embodiment 28 wherein the AC service load field is included in the BSS load element only if a QoS-Option-Implemented parameter is true.
30. The method of any of embodiments 28-29 wherein the four sub-fields comprise an AAD for best-effort (AADBE) field, an AAD for background (AADBG) field, an AAD for video (AADVI) field, and/or an AAD for voice (AADVO) field.
31. The method of any of embodiments 28-30 wherein a low AAD value indicates a shorter access delay than a higher AAD value.
32. The method of any of embodiments 28-31, further comprising setting an AAD value for a first of the four sub-fields to an AAD value of the sub-field that is adjacent and to the right of said first sub-field when the QCS or QAP is not providing services for an indicated access category.
33. The method of any preceding embodiment, further comprising measuring and/or averaging a medium access delay (MAD) value for all transmit packets of an indicated access category.
34. The method of embodiment 33 wherein said MAD value is measured and/or averaged using an EDCF access mechanism over a continuous window of time, wherein an averaged MAD has a predetermined accuracy range and is based on a minimum number of transmit packet delay measurements.
35. The method of embodiment 34 wherein said window of time is a thirty (30) second measurement window, wherein the predetermined accuracy range is two-hundred (200) μ s, and/or wherein said MAD average is based on at least two-hundred transmit packet delay measurements.
36. The method of any of embodiments 28-35 wherein an AAD value within a predetermined range of values in one of the four sub-fields is a logarithmically scaled representation of an average MAD for transmitted packets in an indicated access category, said average MAD being measured from a time an EDCF packet is ready for transmission until a time the EDCF packet is actually transmitted.
37. The method of embodiment 36 wherein said range of values is between zero (0) and two-hundred and fifty-four (254).

38. The method of any of embodiments 28-37 wherein a predetermined AAD value in any of the four sub-fields indicates that a QCS or QAP is not providing services to an indicated access category or to any higher priority access category.
39. The method of embodiment 39 wherein said predetermined AAD value is zero (0).
40. The method of any of embodiments 28-39 wherein other predetermined AAD values represent various average MAD times.
41. The method of any of embodiments 28-40 wherein an AAD value of one (1) represents an average MAD of fifty (50) μ s.
42. The method of any of embodiments 28-41 wherein an AAD value of two-hundred and fifty-three (253) represents an average MAD of five and one-half (5.5) μ s or greater.
43. The method of any of embodiments 28-42 wherein an AAD value of two-hundred and fifty-four (254) indicates that services at an indicated access category are currently blocked.
44. The method of any of embodiments 28-43 wherein an AAD value of two-hundred and fifty-five (255) indicates that an AC service load is not available.
45. The method of any preceding embodiment wherein the BSS load element further includes a channel utilization field.
46. The method of embodiment 45 wherein said channel utilization field defines a percentage of time the first communication station sensed a transmit medium as being busy, as indicated by a carrier sense mechanism.
47. The method of embodiment 46 wherein the percentage of time is a moving average.
48. The method of embodiment 47 wherein the moving average is defined using at least one parameter selected from the group consisting of a channel-busy-time parameter, a channel-utilization-beacon-interval parameter, and/or a beacon-period parameter.
49. The method of any of embodiments 47-48 wherein said moving average is defined as a product of a channel-busy-time parameter and two-hundred fifty-five (255), divided by a product of a channel-utilization-beacon-interval parameter, a beacon period, and one-thousand twenty-four (1024).
50. The method of any of embodiments 48-49 wherein the channel-busy-time parameter is defined as a number of microseconds during which a carrier-sense mechanism has indicated a channel busy indication.
51. The method of any of embodiments 48-50 wherein the channel-utilization-beacon-interval parameter is defined as a number of consecutive beacon intervals during which an average may be calculated.
52. The method of any of embodiments 48-51 wherein the channel utilization field is included in the BSS load element when at least one of a QoS-Option-Implemented parameter and a PBSS-Load-Implemented parameter is false.
53. A method of determining a medium access delay (MAD) timing for single access to a communication station, the method comprising determining a first time at which a data packet is ready for transmission.
54. The method of embodiment 53 wherein said first time is a time at which a Carrier-Sense Multiple Access / Collision Avoidance (CSMA/CA) protocol is initiated.
55. The method of any of embodiments 53-54 comprising determining a second time at which a transmission request is made to a physical (PHY) layer transmission process.
56. The method of any of embodiments 53-55 comprising determining a third time at which said transmission request is acknowledged.
57. The method of any of embodiments 53-56 comprising calculating a packet transmission and acknowledgement timing as a difference between the second time and the third time.
58. The method of any of embodiments 53-57 comprising calculating a total access timing as a difference between the third time and the first time.
- 58". The method of any of embodiments 53-58 comprising calculating the MAD timing by subtracting the packet transmission and acknowledgement timing from the total access timing.
59. The method of any of embodiments 53-58 and 58" wherein the transmission request is preceded by a Request-to-Send/Clear-to-Send (RTS/CTS) handshake.

60. A method of determining a MAD timing for data packet retransmissions.
61. The method of embodiment 60 comprising determining a first time at which a data packet enters a medium access control (MAC) queue.
62. The method of any of embodiments 60-61 comprising determining a second time at which the data packet is at a head of the MAC queue.
63. The method of any of embodiments 60-62 comprising calculating a MAC queuing delay as a difference between the second time and the first time.
64. The method of any of embodiments 60-63 comprising determining a first retransmission timing as a difference between a first transmission start time and a first transmission end time.
65. The method of embodiment 64 wherein said first transmission start time indicates a commencement of a first transmission of the data packet and said first transmission end time indicates a conclusion to said first transmission without receiving a transmission acknowledgement.
66. The method of any of embodiments 60-64 comprising determining a second retransmission timing as a difference between a second transmission start time and a second transmission end time.
67. The method of embodiment 66 wherein said second transmission start time commences after a deferral and back-off period and indicates a commencement of a second transmission of the data packet and said second transmission end time indicates a conclusion to said second transmission without receiving a transmission acknowledgement.
68. The method of any of embodiments 60-67 comprising determining an Nth retransmission timing as a difference between an Nth transmission start time and a Nth transmission end time.
69. The method of embodiment 68 wherein said Nth transmission start time commences after a deferral and back-off period and indicates a commencement of a Nth transmission of the data packet and said Nth transmission end time indicates a receipt of a transmission acknowledgement.
70. The method of any of embodiments 60-69 comprising calculating a total retransmission timing as a sum of the first, second, and Nth retransmission timings.
71. The method of any of embodiments 60-70 comprising determining a finished time, said finished time indicating a time at which the acknowledgement is received.
72. The method of any of embodiments 60-71 comprising calculating a MAD timing for the data packet as a difference between the finished time and the first time, less the MAC queuing delay, less the total retransmission timing, all divided by N.
73. The method of any of embodiments 20-52 wherein the first communication station is an access point (AP) and wherein the features of the BSS load element are configured for use in and/or by an AP.
74. The method of any of embodiments 20-53 wherein any of the other communication stations is an AP.
75. The method of any of embodiments 20-54 wherein the first communication station is a WTRU and wherein the features of the BSS load element are configured for use by a WTRU.
76. The method of any of embodiments 20-55 wherein any of the other communication stations in and/or by a WTRU.
77. The method of any of embodiments 53-72 wherein the communication station is an AP.
78. The method of any of embodiments 53-72 wherein the communication station is a WTRU.
79. A communication station configured to provide channel management according to any of the methods of embodiments 20-52 and 73-76.
80. The communication station of embodiment 79 comprising a receiver configured to receive unserved traffic demand data from other communication stations located within a wireless service range of said communication station.
81. The communication station of any of embodiments 79-80 comprising a processor configured to calculate a BSS load element for each of a plurality of access categories.
82. The communication station of any of embodiments 79-81 comprising a transmitter configured to advertise the BSS load

element to the other communication stations within a service range of said communication station.

83. The communication station of any of embodiments 79-82 wherein the receiver is configured to receive advertised BSS load elements from other communication stations.

84. The communication station of any of embodiments 79-83 wherein the processor is further configured to utilize the received BSS load elements from other communication stations to assist communication stations in making disassociation decisions.

85. The communication station of any of embodiments 79-84 wherein said communication station is an AP.

86. The communication of any of embodiments 79-84 wherein said communication station is a WTRU.

87. The communication station of any of embodiments 79-86 wherein any of the other communication stations is an AP.

88. The communication station of any of embodiments 79-87 wherein any of the other communication stations is a WTRU.

89. A communication station configured to determine medium access delay according to any of the methods and/or features of embodiments 53-72 and 77-78.

90. The communication station of embodiment 89 wherein said communication station is an AP.

91. The communication station of embodiment 89 wherein said communication station is a WTRU.

92. The communication station of any of embodiments 90-91 comprising a processor configured to determine medium access delay according to any of the methods and/or features of embodiments 53-72 and 77-78.

93. A method of determining an average MAD timing evaluated over a predetermined period of duration comprising defining a period of duration.

94. The method of embodiment 93 comprising determining a total packet transmission duration by summing a packet transmission time and a time spent waiting for and/or receiving an acknowledgement for a quantity of packet transmissions occurring during said period of duration.

95. The method of any of embodiments 93-94 wherein packet transmissions include packet retransmissions.

96. The method of any of embodiments 93-95 comprising determining a total empty-transmit-queue time for a plurality of access categories.

97. The method of any of embodiments 96 wherein total empty-transmit-queue time includes periods of time during which transmit queues of the access categories remain empty.

98. The method of any of embodiments 93-96 comprising subtracting the total packet transmission duration, the total empty-transmit-queue time, and/or the total transmit-queue-deferral time from the period of duration to yield a total difference.

99. The method of any of embodiments 93-97 comprising dividing the total difference by the quantity of packet transmissions to obtain an average MAD timing.

100. The method of any of embodiments 93-99 comprising determining a total transmit-queue-deferral time for the plurality of access categories, wherein said transmit-queue-deferral time includes periods of time during which the access categories deferred their respective transmissions to higher priority queues.

101. The method of embodiment 100 comprising subtracting said total transmit-queue-deferral time from the total difference before said total difference is divided by the quantity of packet transmissions to obtain the average MAD timing.

102. A communication station configured to determine MAD timing according to any of the methods and/or features of embodiments 93-101.

103. The communication station of embodiment 102 comprising a processor.

104. The communication station of any of embodiments 102-103 wherein said communication station is an AP.

105. The communication station of any of embodiments 102-103 wherein said communication station is a WTRU.

106. A communication station configured to perform any of the methods and/or features described in any of the preceding embodiments and/or comprising any of the features described in any of the preceding embodiments.

107. The communication station of embodiment 106 wherein said communication station is an AP.
108. The communication station of embodiment 106 wherein said communication station is a WTRU.
109. The method of any of embodiments 20-22 and 73-76, wherein the access categories include voice, video, best effort and background traffic.
110. The AP of embodiment 87, wherein the access categories include voice, video, best effort and background traffic.
111. The method of any of embodiments 20-52, wherein the access categories include voice, video, best effort and background traffic.
112. The communication station of any of embodiments 79-82, wherein the access categories include voice, video, best effort and background traffic.

[0081] While this invention has been particularly shown and described with reference to preferred embodiments, it will be understood by those skilled in the art that various changes in form and details may be made therein without departing from the scope of the invention as described hereinabove.

REFERENCES CITED IN THE DESCRIPTION

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P A T E N T K R A V

1. Fremgangsmåde til trådløs kommunikation, hvilken fremgangsmåde omfatter:
 at generere en tjenestebelastningsindikator i et adgangspunkt, AP, hvor tjeneste-
 belastningsindikatoren indbefatter:
 et bedste præstations, BE, -forsinkelsesfelt, der indikerer en repræsentation af: en
 gennemsnitlig adgangsforsinkelse for BE-rammer, som er sendt i løbet af et må-
 lingsvindue, at en BE-tjeneste er utilgængelig, eller at den gennemsnitlige ad-
 gangsforsinkelse for BE-rammer ikke er tilgængelig;
 et baggrunds, BK, -forsinkelsesfelt, som indikerer en repræsentation af: en gen-
 gennemsnitlig adgangsforsinkelse for BK-rammer, som er sendt i løbet af målingsvin-
 duet, at en BK-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsin-
 kelse for BK-rammer ikke er tilgængelig;
 et video, VI, -forsinkelsesfelt, der indikerer en repræsentation af: en gennemsnitlig
 adgangsforsinkelse for VI-rammer, der er sendt i løbet af målingsvinduet, at en VI-
 tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse for VI-
 rammer ikke er tilgængelig; og
 et stemme, VO, -forsinkelsesfelt, som indikerer en repræsentation af: en gennem-
 snitlig adgangsforsinkelse for VO-rammer, som er sendt i løbet af målingsvinduet,
 at en VO-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse
 for VO-rammer ikke er tilgængelig.
2. Fremgangsmåde ifølge krav 1, hvor tjenestebelastningsindikatoren er 4 okteter, og
 BE-forsinkelsesfeltet, BK-forsinkelsesfeltet, VO-forsinkelsesfeltet og VI-forsinkelses-
 feltet hver er en forskellig oktet af de 4 okteter.
3. Fremgangsmåde ifølge krav 2, hvor en værdi i en oktet indikerer en skaleret re-
 præsentation af en gennemsnitlig adgangsforsinkelse for en tilsvarende adgangskat-
 tegori.
4. Fremgangsmåde ifølge krav 2, hvor en værdi i en oktet indikerer en gennemsnitlig
 adgangsforsinkelse, der er større end en grænseværdi, for en tilsvarende adgangskat-
 tegori.
5. Fremgangsmåde ifølge krav 2, hvor en værdi i en oktet indikerer, at en tjeneste for
 en tilsvarende adgangskategori er utilgængelig.
6. Fremgangsmåde ifølge krav 2, hvor en værdi i en oktet indikerer, at en gennem-
 snitlig adgangsforsinkelse for en tilsvarende adgangskategori.
7. Fremgangsmåde ifølge krav 1, hvor målingsvinduet er en forudbestemt tidsperiode.
8. Adgangspunkt, AP, til trådløs kommunikation, hvilket adgangspunkt omfatter:
 en processor, der er indrettet til at generere en tjenestebelastningsindikator, som
 indbefatter:

- et bedste præstations, BE, -forsinkelsesfelt, der indikerer en repræsentation af: en gennemsnitlig adgangsforsinkelse for BE-rammer, som er sendt i løbet af et målingsvindue, at en BE-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse for BE-rammer ikke er tilgængelig;
- 5 et baggrunds, BK, -forsinkelsesfelt, som indikerer en repræsentation af: en gennemsnitlig adgangsforsinkelse for BK-rammer, som er sendt i løbet af målingsvinduet, at en BK-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse for BK-rammer ikke er tilgængelig;
- 10 et video, VI, -forsinkelsesfelt, der indikerer en repræsentation af: en gennemsnitlig adgangsforsinkelse for VI-rammer, der er sendt i løbet af målingsvinduet, at en VI-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse for VI-rammer ikke er tilgængelig; og
- 15 et stemme, VO, -forsinkelsesfelt, som indikerer en repræsentation af: en gennemsnitlig adgangsforsinkelse for VO-rammer, som er sendt i løbet af målingsvinduet, at en VO-tjeneste er utilgængelig, eller at den gennemsnitlige adgangsforsinkelse for VO-rammer ikke er tilgængelig; og
- en sendeenhed, der er indrettet til at sende tjenestebelastningsindikatoren som en del af en enkelt besked.
9. AP ifølge krav 8, hvor processoren er indrettet til at skabe en tjenestebelastningsindikator som er 4 okteter, og BE-forsinkelsesfeltet, BK-forsinkelsesfeltet, VO-forsinkelsesfeltet og VI-forsinkelsesfeltet hver er en forskellig oktet af de 4 okteter.
10. AP ifølge krav 9, hvor en værdi i en oktet indikerer en skaleret repræsentation af en gennemsnitlig adgangsforsinkelse for en tilsvarende adgangskategori.
11. AP ifølge krav 9, hvor en værdi i en oktet indikerer en gennemsnitlig adgangsforsinkelse, der er større end en grænseværdi, for en tilsvarende adgangskategori.
12. AP ifølge krav 9, hvor en værdi i en oktet indikerer, at en tjeneste for en tilsvarende adgangskategori er utilgængelig.
13. AP ifølge krav 9, hvor en værdi i en oktet indikerer, at en gennemsnitlig adgangsforsinkelse for en tilsvarende adgangskategori.
14. AP ifølge krav 8, hvor målingsvinduet er en forudbestemt tidsperiode.
15. AP ifølge krav 14, hvor tidsperioden er tredive sekunder.

DRAWINGS

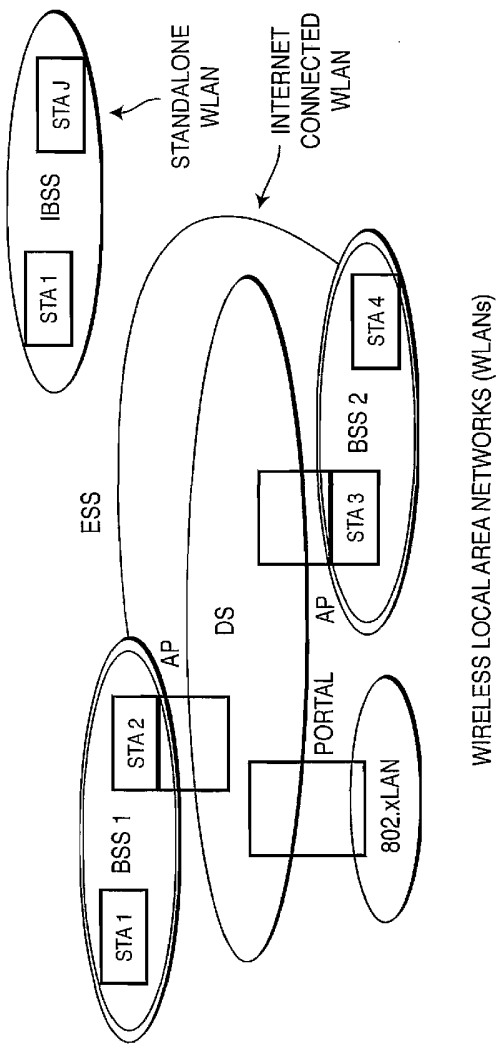
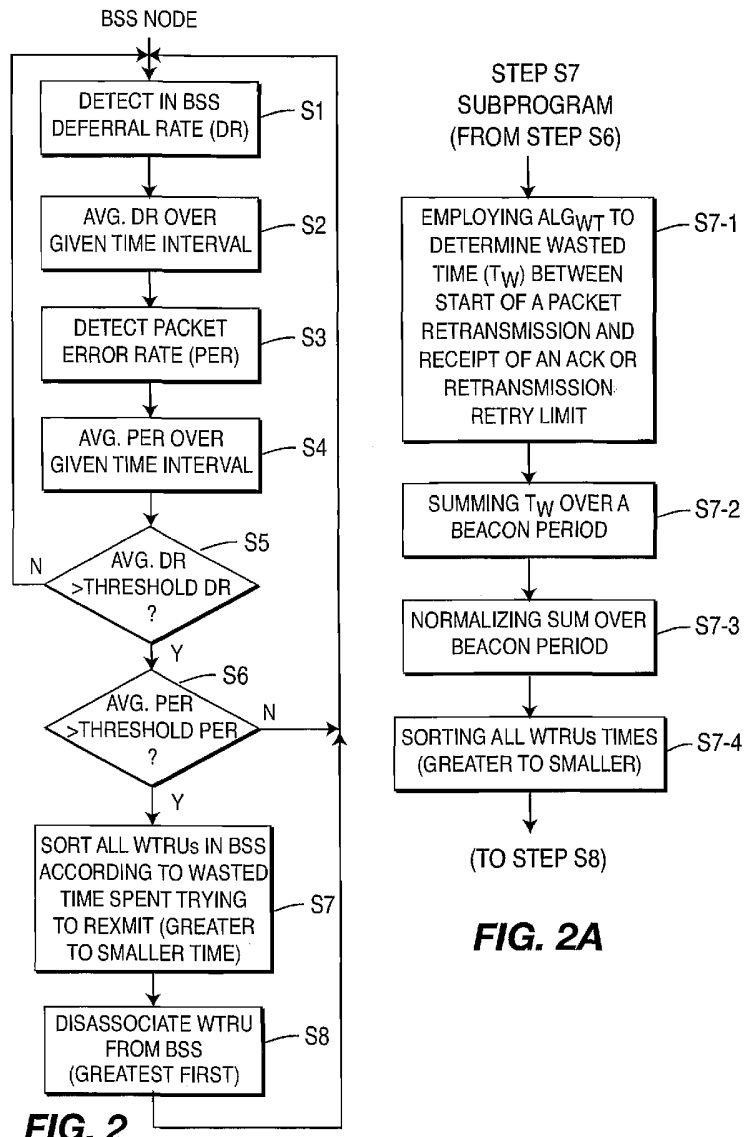
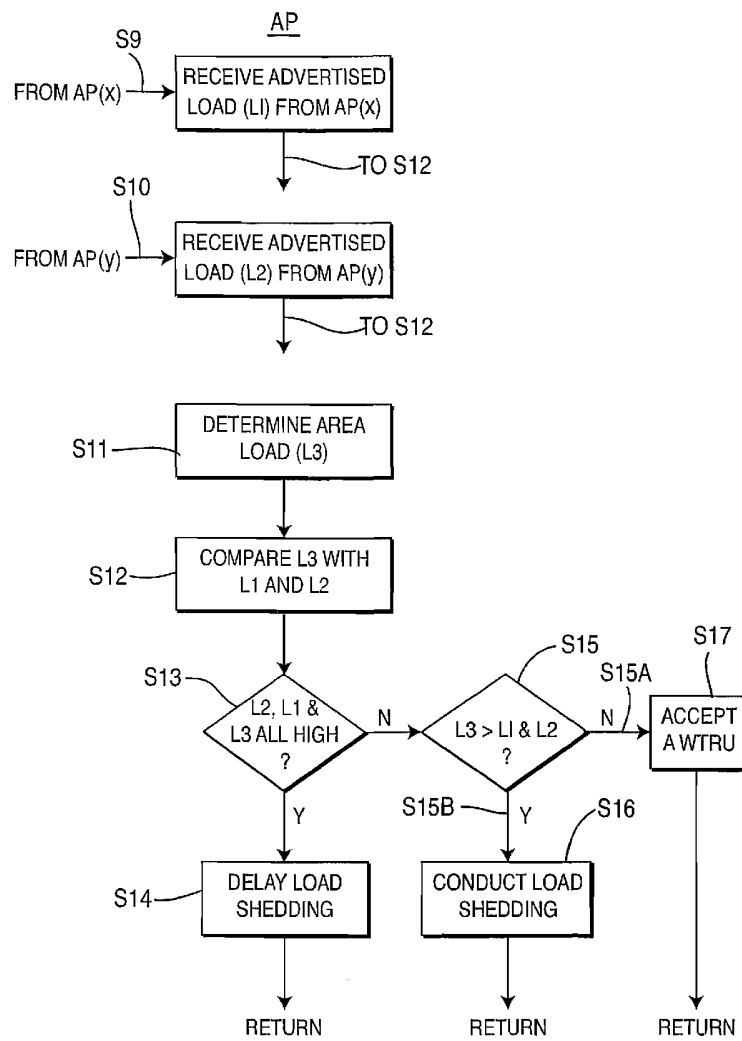
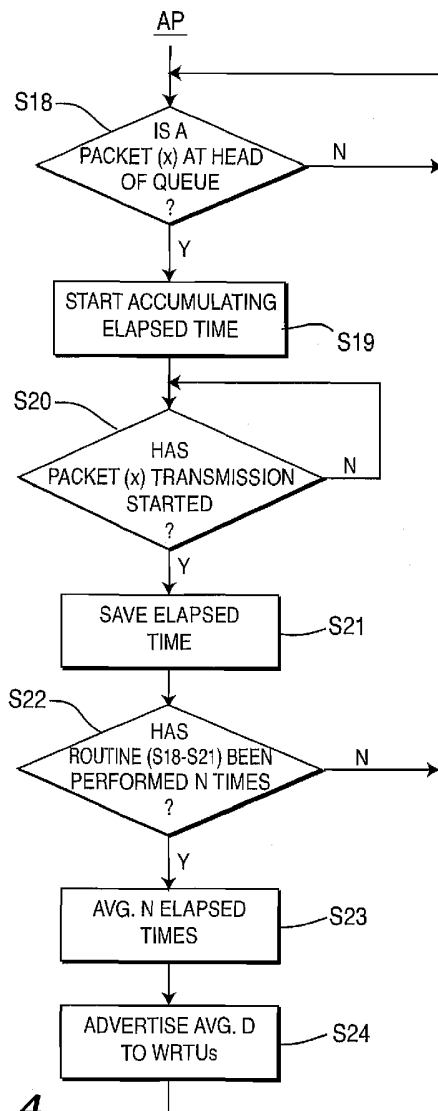
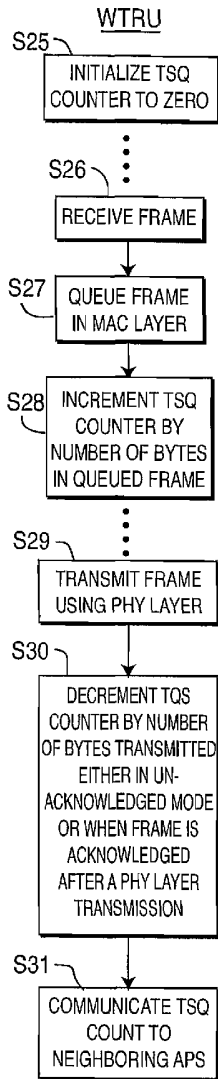
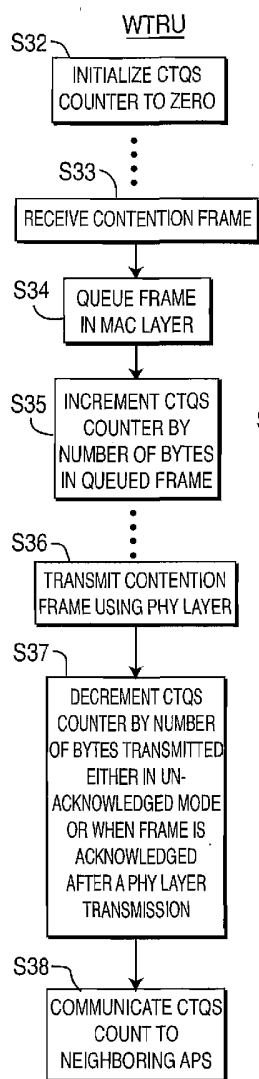
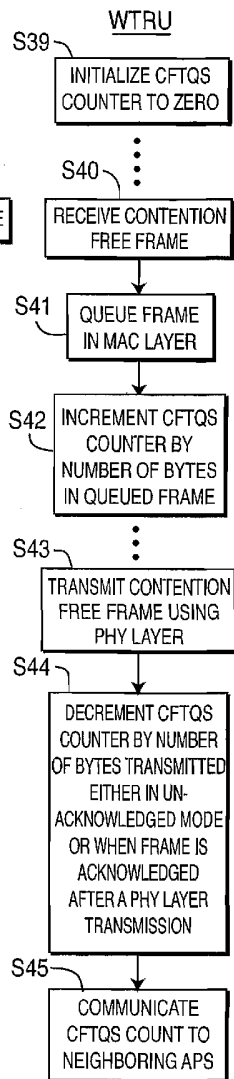


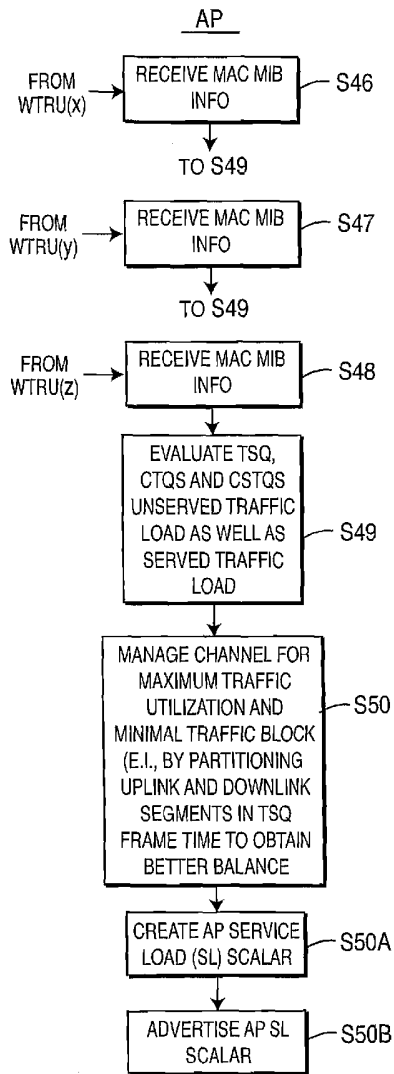
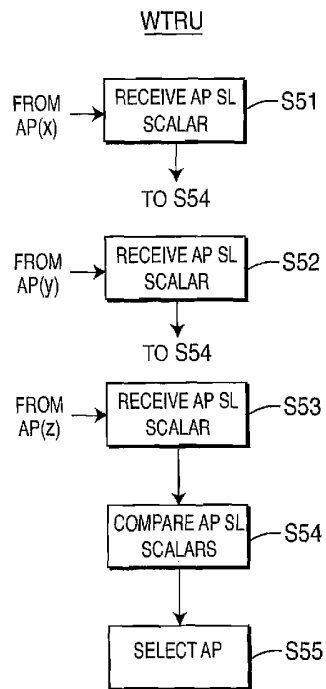
FIG. 1





**FIG. 4**

**FIG. 5****FIG. 6****FIG. 7**

**FIG. 8****FIG. 9**

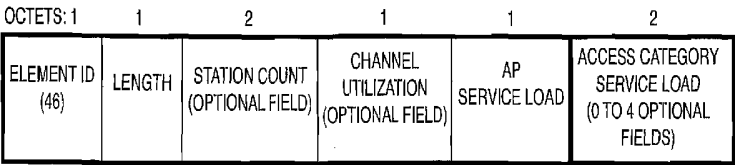


FIG. 10

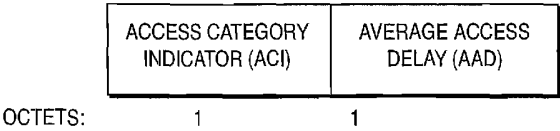


FIG. 11

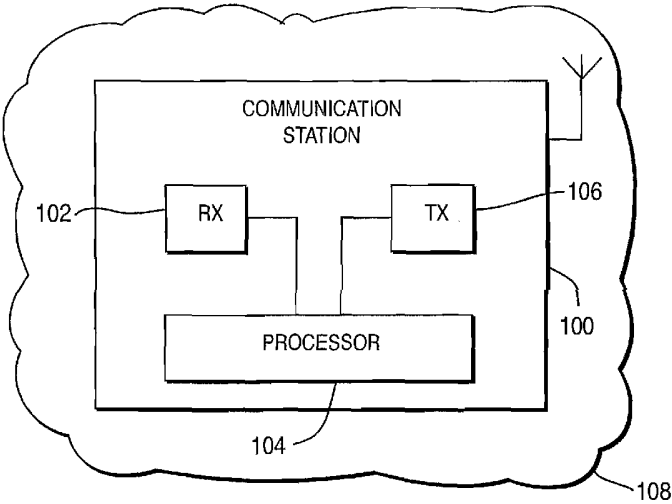


FIG. 12