(12) INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(19) World Intellectual Property Organization

International Bureau





(10) International Publication Number WO 2012/069276 A1

- (43) International Publication Date 31 May 2012 (31.05.2012)
- (51) International Patent Classification: *G06F 9/455* (2006.01)
- (21) International Application Number:

PCT/EP2011/068987

(22) International Filing Date:

28 October 2011 (28.10.2011)

(25) Filing Language:

English

(26) Publication Language:

English

(30) Priority Data:

12/953,483 24 November 2010 (24.11.2010)

US

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- (81) Designated States (unless otherwise indicated, for every kind of national protection available): AE, AG, AL, AM, AO, AT, AU, AZ, BA, BB, BG, BH, BR, BW, BY, BZ, CA, CH, CL, CN, CO, CR, CU, CZ, DE, DK, DM, DO, DZ, EC, EE, EG, ES, FI, GB, GD, GE, GH, GM, GT, HN, HR, HU, ID, IL, IN, IS, JP, KE, KG, KM, KN, KP, KR, KZ, LA, LC, LK, LR, LS, LT, LU, LY, MA, MD, ME, MG, MK, MN, MW, MX, MY, MZ, NA, NG, NI, NO, NZ, OM, PE, PG, PH, PL, PT, QA, RO, RS, RU, RW, SC, SD, SE, SG, SK, SL, SM, ST, SV, SY, TH, TJ, TM, TN, TT, TZ, UA, UG, US, UZ, VC, VN, ZA, ZM, ZW.
- (84) Designated States (unless otherwise indicated, for every kind of regional protection available): ARIPO (BW, GH,

[Continued on next page]

(54) Title: DEVICE EMULATION IN A VIRTUALIZED COMPUTING ENVIRONMENT

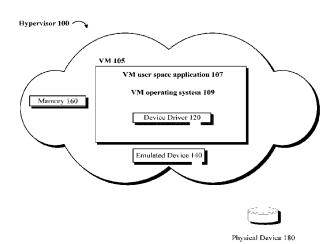


FIG. 1

(57) Abstract: Systems and methods for optimizing operation of a physical device in a virtualized computing environment are provided. The method comprises monitoring instructions issued by an application running on a virtual machine's (VM) operating system, wherein the VM is hosted by a hypervisor providing access to a physical device connected to a virtualized computing environment; wherein a device driver supported by the operating system issues one or more instructions to an emulated device supported by the hypervisor to control the physical device according to the issued instructions, wherein the emulated device's registers are implemented in memory of the hypervisor and the instructions issued by the device driver are stored in said registers, wherein the hypervisor handles the instructions, in response to learning that the instructions are stored in said registers.



GM, KE, LR, LS, MW, MZ, NA, RW, SD, SL, SZ, TZ, UG, ZM, ZW), Eurasian (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European (AL, AT, BE, BG, CH, CY, CZ, DE, DK, EE, ES, FI, FR, GB, GR, HR, HU, IE, IS, IT, LT, LU, LV, MC, MK, MT, NL, NO, PL, PT, RO, RS,

SE, SI, SK, SM, TR), OAPI (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, ML, MR, NE, SN, TD, TG).

Published:

— with international search report (Art. 21(3))

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DEVICE EMULATION IN A VIRTUALIZED COMPUTING ENVIRONMENT

TECHNICAL FIELD

The present invention relates generally to virtual machines and, more particularly, to a system and method for optimizing data operations for a device emulated in a virtualized environment.

BACKGROUND

A virtual machine (VM) is a software implementation of a machine (i.e., a computer) that executes programs like a physical machine would. Generally, access to resources in virtualized environments is associated with a high level of overhead because the VM typically cannot directly communicate with a virtualized resource. A hypervisor that hosts the VM generally has the burden of emulating the needed interface or driver on behalf of the VM. In other words, the hypervisor presents the VM with a virtual platform and monitors the execution of the VM and how the VM accesses the available hardware resources and devices.

For example, for a VM to read or write data to a virtualized storage device, control of the read or write operations will have to be transferred between the VM and the hypervisor. In a purely emulated case, the VM has no knowledge that the emulated device is not a real physical device, because the hypervisor emulates the interface in software. A high overhead is associated with such emulation. The operational overhead is even more costly, if there is a need for transferring control intermittently between the VM and the hypervisor when accessing a device.

SUMMARY

For purposes of summarizing, certain aspects, advantages, and novel features have been described herein. It is to be understood that not all such advantages may be achieved in accordance with any one particular embodiment. Thus, the disclosed subject matter may be embodied or carried out in a manner that achieves or optimizes one advantage or group of advantages without achieving all advantages as may be taught or suggested herein.

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Systems and methods for optimizing operation of a physical device in a virtualized computing environment are provided. The method comprises monitoring instructions issued by an application running on a virtual machine's (VM) operating system, wherein the VM is hosted by a hypervisor providing access to a physical device connected to a virtualized computing environment; wherein a device driver supported by the operating system issues one or more instructions to an emulated device supported by the hypervisor to control the physical device according to the issued instructions, wherein the emulated device's registers are implemented in memory of the hypervisor and the instructions issued by the device driver are stored in said registers, wherein the hypervisor handles the instructions, in response to learning that the instructions are stored in said registers.

In accordance with one or more embodiments, a system comprising one or more logic units is provided. The one or more logic units are configured to perform the functions and operations associated with the above-disclosed methods. In yet another embodiment, a computer program product comprising a computer readable storage medium having a computer readable program is provided. The computer readable program when executed on a computer causes the computer to perform the functions and operations associated with the above-disclosed methods.

One or more of the above-disclosed embodiments in addition to certain alternatives are provided in further detail below with reference to the attached figures. The disclosed subject matter is not, however, limited to any particular embodiment disclosed.

BRIEF DESCRIPTION OF THE DRAWINGS

The disclosed embodiments may be better understood by referring to the figures in the attached drawings, as provided below.

FIGS. 1 and 2 illustrate exemplary computing environments in accordance with one or more embodiments, wherein a VM hosted by a hypervisor is implemented to communicate with a device accessible via the hypervisor.

FIGS. 3A and 3B are exemplary flow diagrams of methods for handling instructions to control the device illustrated FIGS. 1 and 2 in accordance with one embodiment.

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FIGS. 4A and 4B are block diagrams of hardware and software environments in which the disclosed systems and methods may operate, in accordance with one or more embodiments.

Features, elements, and aspects that are referenced by the same numerals in different figures represent the same, equivalent, or similar features, elements, or aspects, in accordance with one or more embodiments.

DETAILED DESCRIPTION OF EXAMPLE EMBODIMENTS

In the following, numerous specific details are set forth to provide a thorough description of various embodiments. Certain embodiments may be practiced without these specific details or with some variations in detail. In some instances, certain features are described in less detail so as not to obscure other aspects. The level of detail associated with each of the elements or features should not be construed to qualify the novelty or importance of one feature over the others.

Referring to FIG. 1, in accordance with one embodiment, a VM 105 may be hosted by a hypervisor 100, in a computing environment. The hypervisor 100 may be implemented to support the emulation of a physical device (e.g., a disk storage device) to allow the VM 105 communicate with a resource (e.g., physical device 180). Emulated device 140 acts as an interface between the device driver 120 and physical device 180. For example, emulated device 140 may support a standard protocol for interfacing with device driver 120, where the protocol defines the manner of use of the device registers and the respective I/O functionality.

Emulated device 140 may be implemented in software or firmware to emulate some physical device so that a VM may communicate with the physical device 180 without any knowledge of the exact structure or nature of the particular interface requirements of physical device 180. Depending on implementation, the overhead associated with the emulation and the need for transfer of control between VM 105 and the hypervisor 100 may be reduced by way

of using a para-virtualized device to act as an interface between the device driver 120 and physical device 180. In one embodiment, para-virtualization of a device may be accomplished by using a virtualized component (e.g., a custom driver) to enhance the system's performance by, for example, batching multiple instructions or I/O requests and using shared memory to avoid multiple data copies.

As shown in FIG. 1, a computer-implemented application such as a software code (e.g., user space application 107) may be running on a VM operating system 109 executed on VM 105. The VM operating system 109 may support a device driver 120 and the hypervisor 100 may support an emulated device 140 capable of communicating with the device driver 120 and physical device 180. A device assignment option may be utilized, in one implementation, to allow the VM 105 access physical device 180. A device assignment grants the VM 105, desirably, full access to physical device 180 so that the VM 105 may send I/O requests directly to physical device 180, where interrupts may be routed through the hypervisor 100. Device assignment option may not be desirable if physical device 180 is intended for sharing among several VMs and the hypervisor and does not support VM migration due to the VM's direct interface assignment to a particular device.

In one implementation, the emulated device 140 has registers that are virtually implemented in the hypervisor 105 memory 160. That is, a portion of hypervisor 105 memory 160 may be shared or reserved for the emulated device 140, wherein any instructions issued by the device driver 120 are written to said shared memory area. In this manner, the device driver 120 and the emulated device 140 may be configured to allow hypervisor 100 to monitor or control corresponding operations on physical device 180.

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In one embodiment, an instruction written to the emulated device 140 registers (i.e., shared hypervisor 100 memory 160) may be controlled by the VM operating system 109, or monitored by the hypervisor 100 or the emulated device 140. It is noteworthy that the emulated device 160 registers may be fully accessible from the hypervisor 100 as the registers are implemented in the hypervisor 105 memory 160. In accordance with one embodiment, device driver 120 is unaware that it is communicating with emulated device 140 and not physical device 180 when issuing the respective instructions.

As provided in further detail below, the emulated device 140 or the hypervisor 100 may monitor the shared memory 160 to determine when an instruction is sent from device driver 120 to emulated device 140. Upon detecting such instruction, the hypervisor 100 is engaged to handle the instruction by emulating the proper behavior for the physical device 180. In this manner, the overall operational performance of the physical device 180 is improved as the number of control switches between the VM 105 and the hypervisor 100 is minimized. In one embodiment, the monitoring of the shared memory 160 takes place by way executing a separate thread of execution from the VM 105 to allow the emulated device 140 or the hypervisor 100 to determine if instructions are written to the memory 160, while the VM 105 continues to run.

As such, in the above exemplary scenario, when the device driver 120 writes to a device register implemented in shared memory 160, the VM 105 continues to run, while in a separate thread of execution, the emulated device 140 or the hypervisor 100 continues to monitor data written to the device registers in shared memory 160. In response to noticing a change in the register value (i.e., resulting from a write operation from the device driver 120) the emulated device 140 or the hypervisor 100 performs the task to handle the instruction without the VM 105 having to exit and explicitly notify the hypervisor of changes in the register value (i.e., cause a switch in CPU control from the VM to the hypervisor).

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Referring to FIGS. 2 and 3A, in one embodiment, VM user space application 107 may issue a system call. The VM's operating system 109 monitors the system calls for input/output (I/O) operations or other issued instructions and determines whether an instruction was sent from the device driver 120 to emulated device 140 (S310). For example, the device driver 120 in VM 105 may send and instruction to the emulated device 140 using a programmed input/output (PIO) or a memory-mapped I/O (MMIO) method.

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PIO and MMIO are methods of transferring data and performing I/O between a CPU and one or more physical devices 180 (e.g., network adapters, storage devices, etc.). For example, a PIO may happen when VM user space application 107 issues CPU instructions that access I/O address space to perform data transfers to or from physical device 180. MMIO may use an address interface (e.g., a bus) to address both memory 160 and physical device 180. The

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CPU instructions used to access memory 160 may be also used for accessing physical device 180.

In one embodiment, areas of the CPU's addressable space may be reserved for I/O. The reservation may be temporary or permanent, for example. Depending on implementation, physical device 180 may monitor the CPU's address bus and respond to the CPU's access of device-assigned address space, connecting the data bus to the physical device 180 hardware register.

In the virtualized system illustrated in FIG. 2, the VM operating system 109 may detect that a CPU instruction is issued from the device driver 120 to the emulated device 140 (S320). If so, the VM operating system 109 generates a signal (e.g., triggers a trap) to notify the emulated device 140 of the arrival of the instruction (S330).

It is noteworthy that since the VM user space application 107 is executed within the context of the VM 105, the VM 105 is responsible for handling the respective CPU instructions. In one embodiment, the hypervisor 100 causes control of the CPU to be switched from the VM 105 to the hypervisor 100 (S340), in response to the above-mentioned trap being triggered. As such, depending on implementation, the emulated device 140 or hypervisor 100 may directly handle the instruction (S350).

Referring to FIGS. 2 and 3B, in an alternative implementation, the emulated device 140 or hypervisor 100 actively monitors (e.g., polls) the emulated device 140 registers (i.e., areas of memory 160 shared with the hypervisor) for instructions issued by the device driver 120 for the emulated device 140 (S410). If the registers are written to (S420), then the emulated device 140 or hypervisor 100 handles the instruction (S430).

In different embodiments, the claimed subject matter may be implemented as a combination of both hardware and software elements, or alternatively either entirely in the form of hardware or entirely in the form of software. Further, computing systems and program software disclosed herein may comprise a controlled computing environment that may be presented in terms of hardware components or logic code executed to perform methods and

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processes that achieve the results contemplated herein. Said methods and processes, when performed by a general purpose computing system or machine, convert the general purpose machine to a specific purpose machine.

Referring to FIGS. 4A and 4B, a computing system environment in accordance with an exemplary embodiment may be composed of a hardware environment 1110 and a software environment 1120. The hardware environment 1110 may comprise logic units, circuits or other machinery and equipments that provide an execution environment for the components of software environment 1120. In turn, the software environment 1120 may provide the execution instructions, including the underlying operational settings and configurations, for the various components of hardware environment 1110.

Referring to FIG. 4A, the application software and logic code disclosed herein may be implemented in the form of computer readable code executed over one or more computing systems represented by the exemplary hardware environment 1110. As illustrated, hardware environment 110 may comprise a processor 1101 coupled to one or more storage elements by way of a system bus 1100. The storage elements, for example, may comprise local memory 1102, storage media 1106, cache memory 1104 or other computer-usable or computer readable media. Within the context of this disclosure, a computer usable or computer readable storage medium may include any recordable article that may be utilized to contain, store, communicate, propagate or transport program code.

A computer readable storage medium may be an electronic, magnetic, optical, electromagnetic, infrared, or semiconductor medium, system, apparatus or device. The computer readable storage medium may also be implemented in a propagation medium, without limitation, to the extent that such implementation is deemed statutory subject matter. Examples of a computer readable storage medium may include a semiconductor or solid-state memory, magnetic tape, a removable computer diskette, a random access memory (RAM), a read-only memory (ROM), a rigid magnetic disk, an optical disk, or a carrier wave, where appropriate. Current examples of optical disks include compact disk, read only memory (CD-ROM), compact disk read/write (CD-R/W), digital video disk (DVD), high definition video disk (HD-DVD) or Blue-rayTM disk.

In one embodiment, processor 1101 loads executable code from storage media 1106 to local memory 1102. Cache memory 1104 optimizes processing time by providing temporary storage that helps reduce the number of times code is loaded for execution. One or more user interface devices 1105 (e.g., keyboard, pointing device, etc.) and a display screen 1107 may be coupled to the other elements in the hardware environment 1110 either directly or through an intervening I/O controller 1103, for example. A communication interface unit 1108, such as a network adapter, may be provided to enable the hardware environment 1110 to communicate with local or remotely located computing systems, printers and storage devices via intervening private or public networks (e.g., the Internet). Wired or wireless modems and Ethernet cards are a few of the exemplary types of network adapters.

It is noteworthy that hardware environment 1110, in certain implementations, may not include some or all the above components, or may comprise additional components to provide supplemental functionality or utility. Depending on the contemplated use and configuration, hardware environment 1110 may be a desktop or a laptop computer, or other computing device optionally embodied in an embedded system such as a set-top box, a personal digital assistant (PDA), a personal media player, a mobile communication unit (e.g., a wireless phone), or other similar hardware platforms that have information processing or data storage capabilities.

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In some embodiments, communication interface 1108 acts as a data communication port to provide means of communication with one or more computing systems by sending and receiving digital, electrical, electromagnetic or optical signals that carry analog or digital data streams representing various types of information, including program code. The communication may be established by way of a local or a remote network, or alternatively by way of transmission over the air or other medium, including without limitation propagation over a carrier wave.

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As provided here, the disclosed software elements that are executed on the illustrated hardware elements are defined according to logical or functional relationships that are exemplary in nature. It should be noted, however, that the respective methods that are implemented by way of said exemplary software elements may be also encoded in said

hardware elements by way of configured and programmed processors, application specific integrated circuits (ASICs), field programmable gate arrays (FPGAs) and digital signal processors (DSPs), for example.

Referring to FIG. 4B, software environment 1120 may be generally divided into two classes comprising system software 1121 and application software 1122 as executed on one or more hardware environments 1110. In one embodiment, the methods and processes disclosed here may be implemented as system software 1121, application software 1122, or a combination thereof. System software 1121 may comprise control programs, such as an operating system (OS) or an information management system, that instruct one or more processors 1101 (e.g., microcontrollers) in the hardware environment 1110 on how to function and process information. Application software 1122 may comprise but is not limited to program code, data structures, firmware, resident software, microcode or any other form of information or routine that may be read, analyzed or executed by a processor 1101.

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In other words, application software 1122 may be implemented as program code embedded in a computer program product in form of a computer-usable or computer readable storage medium that provides program code for use by, or in connection with, a computer or any instruction execution system. Moreover, application software 1122 may comprise one or more computer programs that are executed on top of system software 1121 after being loaded from storage media 1106 into local memory 1102. In a client-server architecture, application software 1122 may comprise client software and server software. For example, in one embodiment, client software may be executed on a client computing system that is distinct and separable from a server computing system on which server software is executed.

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Software environment 1120 may also comprise browser software 1126 for accessing data available over local or remote computing networks. Further, software environment 1120 may comprise a user interface 1124 (e.g., a graphical user interface (GUI)) for receiving user commands and data. It is worthy to repeat that the hardware and software architectures and environments described above are for purposes of example. As such, one or more embodiments may be implemented over any type of system architecture, functional or logical platform or processing environment.

It should also be understood that the logic code, programs, modules, processes, methods and the order in which the respective processes of each method are performed are purely exemplary. Depending on implementation, the processes or any underlying sub-processes and methods may be performed in any order or concurrently, unless indicated otherwise in the present disclosure. Further, unless stated otherwise with specificity, the definition of logic code within the context of this disclosure is not related or limited to any particular programming language, and may comprise one or more modules that may be executed on one or more processors in distributed, non-distributed, single or multiprocessing environments

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As will be appreciated by one skilled in the art, a software embodiment may include firmware, resident software, micro-code, etc. Certain components including software or hardware or combining software and hardware aspects may generally be referred to herein as a "circuit," "module" or "system." Furthermore, the subject matter disclosed may be implemented as a computer program product embodied in one or more computer readable storage medium(s) having computer readable program code embodied thereon. Any combination of one or more computer readable storage medium(s) may be utilized. The computer readable storage medium may be a computer readable signal medium or a computer readable storage medium. A computer readable storage medium may be, for example, but not limited to, an electronic, magnetic, optical, electromagnetic, infrared, or semiconductor system, apparatus, or device, or any suitable combination of the foregoing.

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In the context of this document, a computer readable storage medium may be any tangible medium that can contain, or store a program for use by or in connection with an instruction execution system, apparatus, or device. A computer readable signal medium may include a propagated data signal with computer readable program code embodied therein, for example, in baseband or as part of a carrier wave. Such a propagated signal may take any of a variety of forms, including, but not limited to, electro-magnetic, optical, or any suitable combination thereof. A computer readable signal medium may be any computer readable medium that is not a computer readable storage medium and that can communicate, propagate, or transport a program for use by or in connection with an instruction execution system, apparatus, or device.

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Program code embodied on a computer readable storage medium may be transmitted using any appropriate medium, including but not limited to wireless, wireline, optical fiber cable, RF, etc., or any suitable combination of the foregoing. Computer program code for carrying out the disclosed operations may be written in any combination of one or more programming languages, including an object oriented programming language such as Java, Smalltalk, C++ or the like and conventional procedural programming languages, such as the "C" programming language or similar programming languages.

The program code may execute entirely on the user's computer, partly on the user's computer, as a stand-alone software package, partly on the user's computer and partly on a remote computer or entirely on the remote computer or server. In the latter scenario, the remote computer may be connected to the user's computer through any type of network, including a local area network (LAN) or a wide area network (WAN), or the connection may be made to an external computer (for example, through the Internet using an Internet Service Provider).

Certain embodiments are disclosed with reference to flowchart illustrations and/or block diagrams of methods, apparatus (systems) and computer program products according to embodiments. It will be understood that each block of the flowchart illustrations and/or block diagrams, and combinations of blocks in the flowchart illustrations and/or block diagrams, can be implemented by computer program instructions. These computer program instructions may be provided to a processor of a general purpose computer, special purpose computer, or other programmable data processing apparatus to produce a machine, such that the instructions, which execute via the processor of the computer or other programmable data processing apparatus, create means for implementing the functions/acts specified in the flowchart and/or block diagram block or blocks.

These computer program instructions may also be stored in a computer readable storage medium that can direct a computer, other programmable data processing apparatus, or other devices to function in a particular manner, such that the instructions stored in the computer readable storage medium produce an article of manufacture including instructions which implement the function/act specified in the flowchart and/or block diagram block or blocks.

The computer program instructions may also be loaded onto a computer, other programmable data processing apparatus, or other devices to cause a series of operational steps to be performed on the computer, other programmable apparatus or other devices to produce a computer implemented process such that the instructions which execute on the computer or other programmable apparatus provide processes for implementing the functions/acts specified in the flowchart and/or block diagram block or blocks.

The flowchart and block diagrams in the figures illustrate the architecture, functionality, and operation of possible implementations of systems, methods and computer program products according to various embodiments. In this regard, each block in the flowchart or block diagrams may represent a module, segment, or portion of code, which comprises one or more executable instructions for implementing the specified logical function(s). It should also be noted that, in some alternative implementations, the functions noted in the block may occur out of the order noted in the figures.

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For example, two blocks shown in succession may, in fact, be executed substantially concurrently, or the blocks may sometimes be executed in the reverse order, depending upon the functionality involved. It will also be noted that each block of the block diagrams and/or flowchart illustration, and combinations of blocks in the block diagrams and/or flowchart illustration, can be implemented by special purpose hardware-based systems that perform the specified functions or acts, or combinations of special purpose hardware and computer instructions.

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The claimed subject matter has been provided here with reference to one or more features or embodiments. Those skilled in the art will recognize and appreciate that, despite of the detailed nature of the exemplary embodiments provided here, changes and modifications may be applied to said embodiments without limiting or departing from the generally intended scope. These and various other adaptations and combinations of the embodiments provided here are within the scope of the disclosed subject matter as defined by the claims and their full set of equivalents.

CLAIMS

1. A computer implemented method for optimizing operation of a physical device in a virtualized computing environment, the method comprising:

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monitoring instructions issued by an application running on a virtual machine's (VM) operating system, wherein the VM is hosted by a hypervisor providing access to a physical device connected to a virtualized computing environment;

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wherein a device driver supported by the operating system issues one or more instructions to an emulated device supported by the hypervisor to control the physical device according to the issued instructions,

wherein the emulated device's registers are implemented in memory of the hypervisor and the instructions issued by the device driver are stored in said registers,

wherein the hypervisor handles the instructions, in response to learning that the instructions are stored in said registers.

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2. The method of claim 1, wherein the operating system controls whether the issued instructions are stored in said registers.

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3. The method of claim 1, wherein the emulated device monitors whether the issued instructions are stored in said registers.

4. The method of claim 1, wherein the hypervisor monitors whether the issued instructions are stored in said registers and notifies the emulated device the instructions are stored in said registers.

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5. A computer implemented method of handling instructions for managing a physical device in a virtualized computing environment, the method comprising:

monitoring whether an instruction is issued by a device driver to an emulated device, wherein the device driver is running over a virtual machine (VM) operating system, and the emulated device is running over a hypervisor, which host the VM;

device's register;

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wherein the hypervisor switches control of CPU handling the instruction from the VM to the hypervisor.

- 6. The method of claim 5, wherein the hypervisor handles the instruction.
- 7. The method of claim 5, wherein the emulated device handles the instruction.
- 8. The method of claim 5, wherein the VM operating system monitors whether an instruction is issued by the device driver.
- 9. The method of claim 5, wherein the emulated device's registers are implemented in the hypervisor's memory.
- 10. A computer implemented method of handling instructions for managing a physical device in a virtualized computing environment, the method comprising:

monitoring whether an instruction issued by a device driver to an emulated device is written to the emulated device's registers, wherein the device driver is running over a virtual machine (VM) operating system, and the emulated device is running over a hypervisor, which hosts the VM;

handling the issued instruction.

- 25 11. The method of claim 10, wherein the emulated device's registers are implemented in the hypervisor's memory.
 - 12. The method of claim 10, wherein the hypervisor performs the monitoring.
 - 13. The method of claim 10, wherein the emulated device performs the monitoring.
 - 14. The method of claim 10, wherein the hypervisor handles the issued instruction.

- 15. The method of claim 10, wherein the emulated device handles the issued instruction.
- 16. A computer-implemented system for optimizing operation of a physical device in a virtualized computing environment, the system comprising:

a logic unit for monitoring instructions issued by an application running on a virtual machine's (VM) operating system, wherein the VM is hosted by a hypervisor providing access to a physical device connected to a virtualized computing environment;

a device driver supported by the operating system issuing one or more instructions to an emulated device supported by the hypervisor to control the physical device according to the issued instructions,

wherein the emulated device's registers are implemented in memory of the hypervisor and the instructions issued by the device driver are stored in said registers,

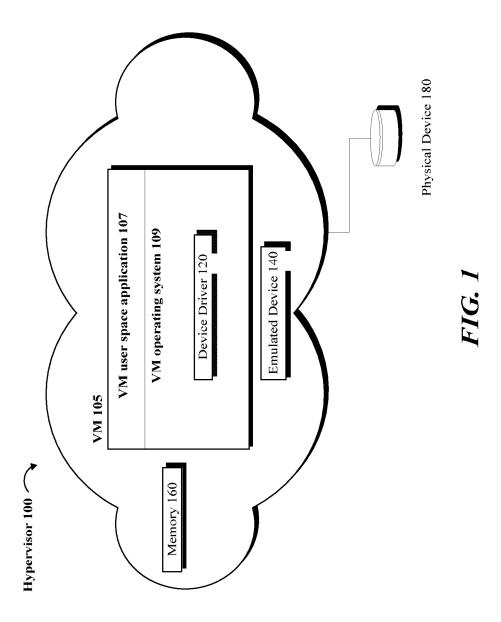
wherein the hypervisor handles the instructions, in response to learning that the instructions are stored in said registers.

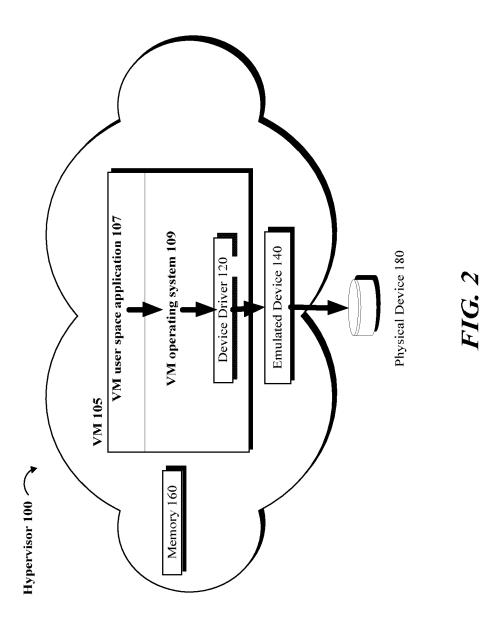
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- 17. The system of claim 16, wherein the operating system controls whether the issued instructions are stored in said registers.
- 18. The system of claim 16, wherein the emulated device monitors whether the issued instructions are stored in said registers.
- 19. The system of claim 16, wherein the hypervisor monitors whether the issued instructions are stored in said registers and notifies the emulated device the instructions are stored in said registers.





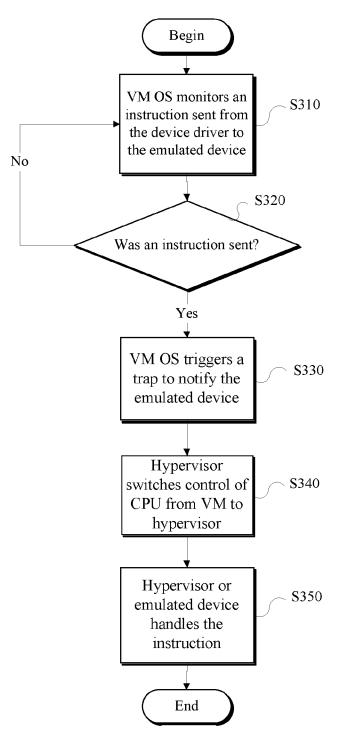


FIG. 3A

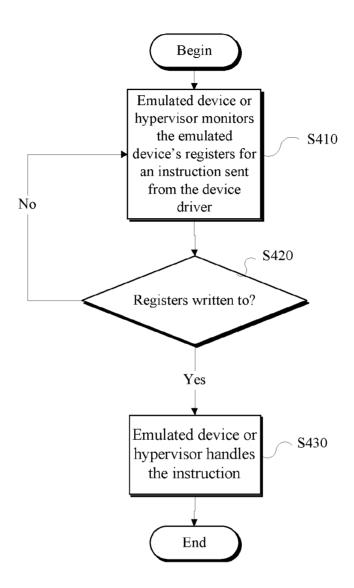


FIG. 3B

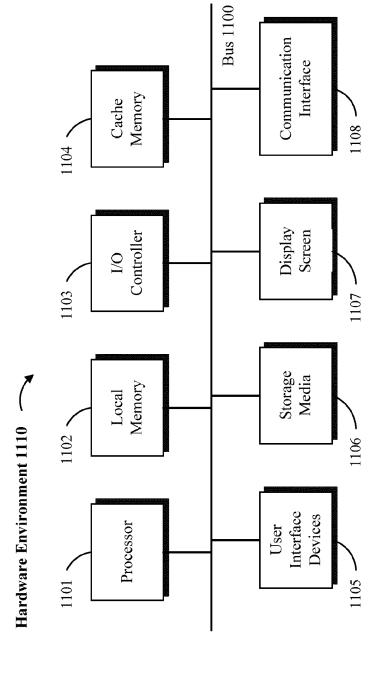


FIG. 4A

Software Environment 1120

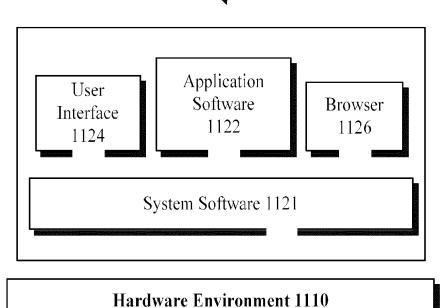


FIG. 4B

INTERNATIONAL SEARCH REPORT

International application No PCT/EP2011/068987

A. CLASSIFICATION OF SUBJECT MATTER INV. G06F9/455 ADD.					
According to International Patent Classification (IPC) or to both national classification and IPC					
	SEARCHED	MION AND IFO			
Minimum documentation searched (classification system followed by classification symbols) $606F$					
Dooumentation searched other than minimum documentation to the extent that such documents are included in the fields searched					
Eleotronio d	ata base consulted during the international search (name of data bas	se and, where practical, search terms used)			
EPO-Internal, WPI Data, COMPENDEX, INSPEC, IBM-TDB					
C. DOCUME	ENTS CONSIDERED TO BE RELEVANT		-		
Category*	Citation of document, with indication, where appropriate, of the rele	evant passages	Relevant to claim No.		
X	US 2009/006074 A1 (GREEN DUSTIN L [US]) 1 January 2009 (2009-01-01) paragraph [0001] - paragraph [0004] paragraph [0011] paragraph [0015] paragraph [0032]		1-19		
А	XINMEI HUANG ET AL: "Accelerating the Emulational Device Model in Virtualization System", EMBEDDED SOFTWARE AND SYSTEMS, 2008. ICESS '08. INTERNATIONAL CONFERENCE ON, IEEE, PISCATAWAY, NJ, USA, 29 July 2008 (2008-07-29), pages 181-186, XP031303449, ISBN: 978-0-7695-3287-5 the whole document		1-19		
Furth	ner documents are listed in the continuation of Box C.	X See patent family annex.			
"A" document defining the general state of the art which is not considered to be of particular relevance "E" earlier document but published on or after the international filing date "L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified) "O" document referring to an oral disclosure, use, exhibition or other means "P" document published prior to the international filing date but later than the priority date claimed		"T" later document published after the international filing date or priority date and not in conflict with the application but oited to understand the principle or theory underlying the invention "X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art. "&" document member of the same patent family			
	actual completion of the international search	Date of mailing of the international sea	rch report		
7 December 2011		15/12/2011			
Name and mailing address of the ISA/ European Patent Office, P.B. 5818 Patentlaan 2 NL - 2280 HV Rijswijk Tel. (+31-70) 340-2040, Fax: (+31-70) 340-3016		Authorized officer Lo Turco, Salvatore			

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INTERNATIONAL SEARCH REPORT

Information on patent family members

International application No
PCT/EP2011/068987

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
US 2009006074 A1	01-01-2009	NONE	