Innovation, Science and **Economic Development Canada** Office de la Propriété Intellectuelle du Canada Canadian Intellectual Property Office

CA 3045694 C 2021/09/07

(11)(21) 3 045 694

# (12) BREVET CANADIEN CANADIAN PATENT

(13) **C** 

(22) Date de dépôt/Filing Date: 2007/11/01

(41) Mise à la disp. pub./Open to Public Insp.: 2008/05/08

(45) Date de délivrance/Issue Date: 2021/09/07

(62) Demande originale/Original Application: 2 668 025

(30) Priorité/Priority: 2006/11/02 (US60/856,212)

(51) Cl.Int./Int.Cl. H04L 12/721 (2013.01), G06F 16/30 (2019.01), H04L 12/66 (2006.01), H04M 15/00 (2006.01)

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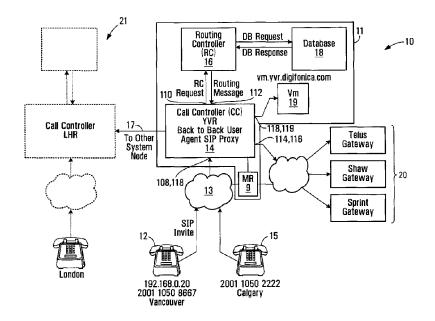
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(54) Titre: PRODUCTION DE MESSAGES DE ROUTAGE POUR DES COMMUNICATIONS PAR VOIX SUR IP

(54) Title: PRODUCING ROUTING MESSAGES FOR VOICE OVER IP COMMUNICATIONS



#### (57) Abrégé/Abstract:

Routing a communication between first and second participants involves causing a controller to receive on an IP network, from a first participant device, a first participant identifier and a second participant identifier to locate first and second participant profiles respectively. At least part of the second participant identifier is compared with at least one attribute contained in the first participant user profile and, based on the comparison, the communication is classified as a system communication or an external network communication whereupon a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant is produced, or an external network routing message identifying an Internet address of a gateway to an external network is produced. The produced routing message thus indicates how the communication should be routed.





#### ABSTRACT

Routing a communication between first and second participants involves causing a controller to receive on an IP network, from a first participant device, a first participant identifier and a second participant identifier to locate first and second participant profiles respectively. At least part of the second participant identifier is compared with at least one attribute contained in the first participant user profile and, based on the comparison, the communication is classified as a system communication or an external network communication whereupon a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant is produced, or an external network routing message identifying an Internet address of a gateway to an external network is produced. The produced routing message thus indicates how the communication should be routed.

#### PRODUCING ROUTING MESSAGES FOR VOICE OVER IP COMMUNICATIONS

#### **BACKGROUND**

## 1. Field

This invention relates to voice over IP communications and methods and apparatus for routing and billing.

## 2. Description of Related Art

Internet protocol (IP) telephones are typically personal computer (PC) based telephones connected within an IP network, such as the public Internet or a private network of a large organization. These IP telephones have installed "voice-over-IP" (VoIP) software enabling them to make and receive voice calls and send and receive information in data and video formats.

15 IP telephony switches installed within the IP network enable voice calls to be made within or between IP networks, and between an IP network and a switched circuit network (SCN), such as the public switched telephone network (PSTN). If the IP switch supports the Signaling System 7 (SS7) protocol, the IP telephone can also access PSTN databases.

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The PSTN network typically includes complex network nodes that contain all information about a local calling service area including user authentication and call routing. The PSTN network typically aggregates all information and traffic into a single location or node, processes it locally and then passes it on to other network nodes, as necessary, by maintaining route tables at the node. PSTN nodes are redundant by design and thus provide reliable service, but if a node should fail due to an earthquake or other natural disaster, significant, if not complete service outages can occur, with no other nodes being able to take up the load.

Existing VoIP systems do not allow for high availability and resiliency in delivering Voice Over IP based Session Initiation Protocol (SIP) Protocol service over a geographically dispersed area such as a city, region or continent. Most resiliency originates from the provision of IP based telephone services to one location or a small number of locations such as a single office or network of branch offices.

#### SUMMARY

In one embodiment, there is provided a method for routing a communication in a communication system between an Internet-connected first participant device associated with a first participant and a second participant device associated with a second participant. The method involves in response to initiation of the communication by the first participant device, receiving, by a controller comprising at least one processor, over an Internet protocol (IP) network a first participant identifier and a second participant identifier. The method further involves causing the at least 15 one processor to access at least one database comprising user profiles using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes. The method further involves comparing at least a portion of the second participant identifier, using the at least 20 one processor, with at least one of the plurality of first participant attributes obtained from the user profile for the first participant. The method further involves causing the at least one processor to access the at least one database to search for a user profile for the second participant, and classifying the communication, based on the comparing, as a system communication or an external network communication, using the at least one processor. The method further involves when the communication is classified as a system communication, producing a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant, using the at least one processor, wherein the system routing message causes the communication to be established to the second participant

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device. The method further involves when the communication is classified as an external network communication, producing an external network routing message identifying an Internet address associated with a gateway to an external network, using the at least one processor, wherein the external network routing message causes the communication to the second participant device to be established using the gateway to the external network.

In another embodiment, there is provided a non-transitory computer readable medium encoded with program code for directing the at least one processor to execute the method above.

In another embodiment, there is provided an apparatus for routing communications in a communication system that includes an Internet-connected first participant device associated with a first participant, wherein the first participant device is operable to initiate a communication to a second participant device associated with a second participant. The apparatus includes a controller comprising at least one processor in communication with at least one memory storing processor readable instructions. The at least one processor is operably configured by the processor readable instructions to, in response to initiation of the communication by the first participant device, receive over an Internet protocol (IP) network a first participant identifier and a second participant identifier, and access at least one database comprising user profiles using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes. The at least one processor is further operably configured by the processor readable instructions to compare at least a portion of the second participant identifier with at least one of the plurality of first participant attributes obtained from the user profile for the first participant to generate a comparison result, access the at least one database to search for a user profile for the second participant, and classify the communication, based on the comparison result, as a system communication or an

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external network communication. The at least one processor is further operably configured by the processor readable instructions to, when the communication is classified as a system communication, produce a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant, wherein the system routing message causes the communication to be established to the second participant device. The at least one processor is operably configured by the processor readable instructions to, when the communication is classified as an external network communication, produce an external network routing message identifying an Internet address associated with a gateway to an external network, wherein the external network routing message causes the communication to the second participant device to be established using the gateway to the external network.

15 In another embodiment, there is provided an apparatus for routing communications in a communication system that includes an Internet-connected first participant device associated with a first participant, the first participant device operable to initiate a communication to a second participant device associated with a second participant. The apparatus includes a controller comprising at least one processor in 20 communication with at least one memory storing processor-readable instruction codes. The at least one processor is operably configured by the processor-readable instruction codes to, in response to initiation of the communication by the first participant device over the Internet, receive a first participant identifier and a second participant identifier, and access at least one database comprising user profiles 25 using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes. The at least one processor is further operably configured by the processor-readable instruction codes to compare at least a portion of the second participant identifier with at least one of the plurality of first participant attributes obtained from the user profile for the first participant to 30

generate a comparison result, and access the at least one database to search for a user profile for the second participant and communication blocking information for the second participant. The at least one processor is also operably configured by the processor-readable instruction codes to classify the communication, based on at least one of the comparison result and the communication blocking information for the second participant, as a system communication, an external network communication or a blocked communication, and when the communication is classified as a system communication, produce a system routing message identifying a first Internet address associated with the second participant device, 10 causing the communication to be established entirely over an Internet protocol (IP) network. The at least one processor is operably configured by the processorreadable instruction codes to, when the communication is classified as an external network communication, produce an external routing message identifying an Internet address associated with a gateway to a network that is external to the 15 communication system, causing a portion of a path taken by the communication to be established over a circuit switched network.

Other aspects and features of the embodiments described herein will become apparent to those ordinarily skilled in the art upon review of the following description in conjunction with the accompanying figures.

#### **BRIEF DESCRIPTION OF THE DRAWINGS**

In drawings which illustrate embodiments,

- 25 Figure 1 is a block diagram of a system according to a first embodiment;
  - Figure 2 is a block diagram of a caller telephone according to the first embodiment;

- Figure 3 is a schematic representation of a SIP invite message transmitted between the caller telephone and a controller shown in Figure 1;
- Figure 4 is a block diagram of a call controller shown in Figure 1;

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- Figure 5 is a flowchart of a process executed by the call controller shown in Figure 1;
- Figure 6 is a schematic representation of a routing, billing and rating (RC) request message produced by the call controller shown in Figure 1;
  - Figure 7 is a block diagram of a processor circuit of a routing, billing, rating element of the system shown in Figure 1;
- 15 Figures 8A-8D is a flowchart of a RC request message handler executed by the RC processor circuit shown in Figure 7;
  - Figure 9 is a tabular representation of a dialing profile stored in a database accessible by the RC shown in Figure 1;

- Figure **10** is a tabular representation of a dialing profile for a caller using the caller telephone shown in Figure **1**;
- Figure **11** is a tabular representation of a callee profile for a callee located in Calgary;
  - Figure 12 is a tabular representation of a callee profile for a callee located in London;

Figure **13** is a tabular representation of a Direct-in-Dial (DID) bank table record stored in the database shown in Figure 1; Figure 14 is a tabular representation of an exemplary DID bank table record for the 5 Calgary callee referenced in Figure 11; Figure 15 is a tabular representation of a routing message transmitted from the RC to the call controller shown in Figure 1; Figure **16** is a schematic representation of a routing message buffer holding a 10 routing message for routing a call to the Calgary callee referenced in Figure 11; Figure 17 is a tabular representation of a prefix to supernode table record stored in 15 the database shown in Figure 1; Figure 18 is a tabular representation of a prefix to supernode table record that would be used for the Calgary callee referenced in Figure 11; 20 Figure 19 is a tabular representation of a master list record stored in a master list table in the database shown in Figure 1; Figure 20 is a tabular representation of a populated master list record; 25 Figure 21 is a tabular representation of a suppliers list record stored in the database shown in Figure 1;

is a tabular representation of a specific supplier list record for a first

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Figure 22

supplier;

is a tabular representation of a specific supplier list record for a second

supplier; Figure 24 is a tabular representation of a specific supplier list record for a third 5 supplier; Figure 25 is a schematic representation of a routing message, held in a routing message buffer, identifying to the controller a plurality of possible suppliers that may carry the call; 10 Figure 26 is a tabular representation of a call block table record; Figure 27 is a tabular representation of a call block table record for the Calgary callee; 15 Figure 28 is a tabular representation of a call forwarding table record; Figure 29 is a tabular representation of a call forwarding table record specific for the Calgary callee; 20 Figure 30 is a tabular representation of a voicemail table record specifying voicemail parameters to enable the caller to leave a voicemail message for the callee; 25 Figure 31 is a tabular representation of a voicemail table record specific to the Calgary callee; Figure **32** is a schematic representation of an exemplary routing message, held in a routing message buffer, indicating call forwarding numbers and a 30 voicemail server identifier;

Figure 23

Figures 33A and 33B	are	respective	portions	of	а	flowchart	of	а	process
executed by	the R	C processor	for deterr	ninii	ng :	a time to liv	e va	alue	э;

- 5 Figure **34** is a tabular representation of a subscriber bundle table record;
  - Figure **35** is a tabular representation of a subscriber bundle record for the Vancouver caller;
- 10 Figure **36** is a tabular representation of a bundle override table record;
  - Figure **37** is a tabular representation of bundle override record for a located master list ID;
- 15 Figure **38** is a tabular representation of a subscriber account table record;
  - Figure **39** is a tabular representation of a subscriber account record for the Vancouver caller;
- 20 Figure **40** is a flowchart of a process for producing a second time value executed by the RC processor circuit shown in Figure **7**;
  - Figure **41** is a flowchart for calculating a call cost per unit time;
- 25 Figure **42** is a tabular representation of a system operator special rates table record;
  - Figure **43** is a tabular representation of a system operator special rates table record for a reseller named Klondike;

	Figure <b>44</b>	is a tabular representation of a system operator mark-up table record;
5	Figure <b>45</b>	is a tabular representation of a system operator mark-up table record for the reseller Klondike;
	Figure <b>46</b>	is a tabular representation of a default system operator mark-up table record;
10	Figure <b>47</b>	is a tabular representation of a reseller special destinations table record;
	Figure <b>48</b>	is a tabular representation of a reseller special destinations table record for the reseller Klondike;
15	Figure <b>49</b>	is a tabular representation of a reseller global mark-up table record;
	Figure <b>50</b>	is a tabular representation of a reseller global mark-up table record for the reseller Klondike;
20	Figure <b>51</b>	is a tabular representation of a SIP bye message transmitted from either of the telephones shown in Figure 1 to the call controller;
	Figure <b>52</b>	is a tabular representation of a SIP bye message sent to the controller from the Calgary callee;
25	Figure <b>53</b>	is a flowchart of a process executed by the call controller for producing a RC stop message in response to receipt of a SIP bye message;
	Figure <b>54</b>	is a tabular representation of an exemplary RC call stop message;

Figure **55** is a tabular representation of an RC call stop message for the Calgary callee;

Figures **56**A and **56**B are respective portions of a flowchart of a RC call stop

message handling routine executed by the RC shown in Figure **1**;

Figure 57 is a tabular representation of a reseller accounts table record;

Figure **58** is a tabular representation of a reseller accounts table record for the reseller Klondike;

Figure **59** is a tabular representation of a system operator accounts table record; and

15 Figure **60** is a tabular representation of a system operator accounts record for the system operator described herein.

## **DETAILED DESCRIPTION**

Referring to Figure 1, a system for making voice over IP telephone/videophone calls is shown generally at 10. The system includes a first super node shown generally at 11 and a second super node shown generally at 21. The first super node 11 is located in geographical area, such as Vancouver, B.C., Canada for example and the second super node 21 is located in London, England, for example. Different super nodes may be located in different geographical regions throughout the world to provide telephone/videophone service to subscribers in respective regions. These super nodes may be in communication with each other by high speed/ high data throughput links including optical fiber, satellite and/or cable links, forming a backbone to the system. These super nodes may alternatively or, in addition, be in communication with each other through conventional internet services.

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In the embodiment shown. the Vancouver supernode 11 provides telephone/videophone service to western Canadian customers from Vancouver Island to Ontario. Another node (not shown) may be located in Eastern Canada to provide services to subscribers in that area.

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Other nodes of the type shown may also be employed within the geographical area serviced by a supernode, to provide for call load sharing, for example within a region of the geographical area serviced by the supernode. However, in general, all nodes are similar and have the properties described below in connection with the 10 Vancouver supernode 11.

In this embodiment, the Vancouver supernode includes a call controller (C) 14, a routing controller (RC) 16, a database 18 and a voicemail server 19 and a media relay 9. Each of these may be implemented as separate modules on a common computer system or by separate computers, for example. The voicemail server 19 15 need not be included in the node and can be provided by an outside service provider.

Subscribers such as a subscriber in Vancouver and a subscriber in Calgary 20 communicate with the Vancouver supernode using their own internet service providers which route internet traffic from these subscribers over the internet shown generally at 13 in Figure 1. To these subscribers the Vancouver supernode is accessible at a pre-determined internet protocol (IP) address or a fully qualified domain name that can be accessed in the usual way through a subscriber's internet service provider. The subscriber in Vancouver uses a telephone 12 that is capable of communicating with the Vancouver supernode 11 using Session Initiation Protocol (SIP) messages and the Calgary subscriber uses a similar telephone 15, in Calgary AB.

It should be noted that throughout this description, the IP/UDP addresses of all elements such as the caller and callee telephones, call controller, media relay, and any others, will be assumed to be valid IP/UDP addresses directly accessible via the Internet or a private IP network, for example, depending on the specific implementation of the system. As such, it will be assumed, for example, that the caller and callee telephones will have IP/UDP addresses directly accessible by the call controllers and the media relays on their respective supernodes, and those addresses will not be obscured by Network Address Translation (NAT) or similar mechanisms. In other words, the IP/UDP information contained in SIP messages (for example the SIP Invite message or the RC Request message which will be described below) will match the IP/UDP addresses of the IP packets carrying these SIP messages.

It will be appreciated that in many situations, the IP addresses assigned to various elements of the system may be in a private IP address space, and thus not directly accessible from other elements. Furthermore, it will also be appreciated that NAT is commonly used to share a "public" IP address between multiple devices, for example between home PCs and IP telephones sharing a single Internet connection. For example, a home PC may be assigned an IP address such as 192.168.0.101 20 and a Voice over IP telephone may be assigned an IP address of 192,168.0.103. These addresses are located in so called "non-routable" (IP) address space and cannot be accessed directly from the Internet. In order for these devices to communicate with other computers located on the Internet, these IP addresses have to be converted into a "public" IP address, for example 24.10.10.123 assigned by the Internet Service Provider to the subscriber, by a device performing NAT, typically a home router. In addition to translating the IP addresses, NAT typically also translates UDP port numbers, for example an audio path originating at a VoIP telephone and using a UDP port 12378 at its private IP address, may have be translated to a UDP port 23465 associated with the public IP address of the NAT device. In other words, when a packet originating from the above VoIP telephone arrives at an Internet-

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based supernode, the source IP/UDP address contained in the IP packet header will be 24.10.10.1:23465, whereas the source IP/UDP address information contained in the SIP message inside this IP packet will be 192.168.0.103:12378. The mismatch in the IP/UDP addresses may cause a problem for SIP-based VoIP systems because, for example, a supernode will attempt to send messages to a private address of a telephone but the messages will never get there.

Referring to Figure 1, in an attempt to make a call by the Vancouver telephone/videophone 12 to the Calgary telephone/videophone 15, the Vancouver telephone/videophone sends a SIP invite message to the Vancouver supernode 11 and in response, the call controller 14 sends an RC request message to the RC 16 which makes various enquiries of the database 18 to produce a routing message which is sent back to the call controller 14. The call controller 14 then communicates with the media relay 9 to cause a communications link including an audio path and a videophone (if a videopath call) to be established through the media relay to the same node, a different node or to a communications supplier gateway as shown generally at 20 to carry audio, and where applicable, video traffic to the call recipient or callee.

Generally, the RC **16** executes a process to facilitate communication between callers and callees. The process involves, in response to initiation of a call by a calling subscriber, receiving a callee identifier from the calling subscriber, using call classification criteria associated with the calling subscriber to classify the call as a public network call or a private network call and producing a routing message identifying an address on the private network, associated with the callee when the call is classified as a private network call and producing a routing message identifying a gateway to the public network when the call is classified as a public network call.

### 30 Subscriber Telephone

In greater detail, referring to Figure 2, in this embodiment, the telephone/videophone 12 includes a processor circuit shown generally at 30 comprising a microprocessor 32, program memory 34, an input/output (I/O) port 36, parameter memory 38 and temporary memory 40. The program memory 34, I/O port 36, parameter memory 38 and temporary memory 40 are all in communication with the microprocessor 32. The I/O port 36 has a dial input 42 for receiving a dialled telephone/videophone number from a keypad, for example, or from a voice recognition unit or from pre-stored telephone/videophone numbers stored in the parameter memory 38, for example. For simplicity, in Figure 2 a box labelled dialing functions 44 represents any device capable of informing the microprocessor 32 of a callee identifier, e.g., a callee telephone/videophone number.

The processor 32 stores the callee identifier in a dialled number buffer 45. In this case, assume the dialled number is 2001 1050 2222 and that it is a number 15 associated with the Calgary subscriber. The I/O port **36** also has a handset interface 46 for receiving and producing signals from and to a handset that the user may place to his ear. This interface 46 may include a BLUETOOTH<sup>TM</sup> wireless interface, a wired interface or speaker phone, for example. The handset acts as a termination point for an audio path (not shown) which will be appreciated later. The I/O port 36 also has an internet connection 48 which is preferably a high speed internet connection and is operable to connect the telephone/videophone to an internet service provider. The internet connection 48 also acts as a part of the voice path, as will be appreciated later. It will be appreciated that where the subscriber device is a videophone, a separate video path is established in the same way an audio path is established. For simplicity, the following description refers to a telephone call, but it is to be understood that a videophone call is handled similarly, with the call controller causing the media relay to facilitate both an audio path and a video path instead of only an audio path.

The parameter memory 38 has a username field 50, a password field 52 an IP address field 53 and a SIP proxy address field 54, for example. The user name field 50 is operable to hold a user name, which in this case is 2001 1050 8667. The user name is assigned upon subscription or registration into the system and, in this embodiment, includes a twelve digit number having a continent code 61, a country code 63, a dealer code 70 and a unique number code 74. The continent code 61 is comprised of the first or left-most digit of the user name in this embodiment. The country code 63 is comprised of the next three digits. The dealer code 70 is comprised of the next four digits and the unique number code 74 is comprised of the last four digits. The password field 52 holds a password of up to 512 characters, in this example. The IP address field 53 stores an IP address of the telephone, which for this explanation is 192.168.0.20. The SIP proxy address field 54 holds an IP protocol compatible proxy address which may be provided to the telephone through the internet connection 48 as part of a registration procedure.

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The program memory 34 stores blocks of codes for directing the processor 32 to carry out the functions of the telephone, one of which includes a firewall block 56 which provides firewall functions to the telephone, to prevent access by unauthorized persons to the microprocessor 32 and memories 34, 38 and 40 through the internet connection 48. The program memory 34 also stores codes 57 for establishing a call ID. The call ID codes 57 direct the processor 32 to produce a call identifier having a format comprising a hexadecimal string at an IP address, the IP address being the IP address of the telephone. Thus, an exemplary call identifier might be FF10@192.168.0.20.

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Generally, in response to picking up the handset interface **46** and activating a dialing function **44**, the microprocessor **32** produces and sends a SIP invite message as shown in Figure **3**, to the routing controller **16** shown in Figure **1**. This SIP invite message is essentially to initiate a call by a calling subscriber.

Referring to Figure 3, the SIP invite message includes a caller ID field 60, a callee identifier field 62, a digest parameters field 64, a call ID field 65 an IP address field 67 and a caller UDP port field 69. In this embodiment, the caller ID field 60 includes the user name 2001 1050 8667 that is the Vancouver user name stored in the user name field 50 of the parameter memory 38 in the telephone 12 shown in Figure 2. In addition, referring back to Figure 3, the callee identifier field 62 includes a callee identifier which in this embodiment is the user name 2001 1050 2222 that is the dialled number of the Calgary subscriber stored in the dialled number buffer 45 shown in Figure 2. The digest parameters field 64 includes digest parameters and the call ID field 65 includes a code comprising a generated prefix code (FF10) and a suffix which is the Internet Protocol (IP) address of the telephone 12 stored in the IP address field 53 of the telephone. The IP address field 67 holds the IP address assigned to the telephone, in this embodiment 192.168.0.20, and the caller UDP port field 69 includes a UDP port identifier identifying a UDP port at which the audio path will be terminated at the caller's telephone.

## Call Controller

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Referring to Figure 4, a call controller circuit of the call controller 14 (Figure 1) is shown in greater detail at 100. The call controller circuit 100 includes a microprocessor 102, program memory 104 and an I/O port 106. The circuit 100 may include a plurality of microprocessors, a plurality of program memories and a plurality of I/O ports to be able to handle a large volume of calls. However, for simplicity, the call controller circuit 100 will be described as having only one microprocessor 102, program memory 104 and I/O port 106, it being understood that 25 there may be more.

Generally, the I/O port 106 includes an input 108 for receiving messages such as the SIP invite message shown in Figure 3, from the telephone shown in Figure 2. The I/O port 106 also has an RC request message output 110 for transmitting an RC request message to the RC 16 of Figure 1, an RC message input 112 for receiving routing messages from the RC 16, a gateway output 114 for transmitting messages to one of the gateways 20 shown in Figure 1 to advise the gateway to establish an audio path, for example, and a gateway input 116 for receiving messages from the gateway. The I/O port 106 further includes a SIP output 118 for transmitting messages to the telephone 12 to advise the telephone of the IP addresses of the gateways which will establish the audio path. The I/O port 106 further includes a voicemail server input and output 117, 119 respectively for communicating with the voicemail server 19 shown in Figure 1.

10 While certain inputs and outputs have been shown as separate, it will be appreciated that some may be a single IP address and IP port. For example, the messages sent to the RC **16** and received from the RC **16** may be transmitted and received on the same single IP port.

The program memory **104** includes blocks of code for directing the microprocessor **102** to carry out various functions of the call controller **14**. For example, these blocks of code include a first block **120** for causing the call controller circuit **100** to execute a SIP invite to RC request process to produce an RC request message in response to a received SIP invite message. In addition, there is a routing message to gateway message block **122** which causes the call controller circuit **100** to produce a gateway query message in response to a received routing message from the RC **16**.

Referring to Figure 5, the SIP invite to RC request process is shown in more detail at 120. On receipt of a SIP invite message of the type shown in Figure 3, block 122 of Figure 5 directs the call controller circuit 100 of Figure 4 to authenticate the user. This may be done, for example, by prompting the user for a password, by sending a message back to the telephone 12 which is interpreted at the telephone as a request for a password entry or the password may automatically be sent to the call controller 14 from the telephone, in response to the message. The call controller 14 may then make enquiries of databases to which it has access, to determine whether or not the

user's password matches a password stored in the database. Various functions may be used to pass encryption keys or hash codes back and forth to ensure that the transmittal of passwords is secure.

Should the authentication process fail, the call controller circuit 100 is directed to an error handling routine 124 which causes messages to be displayed at the telephone 12 to indicate there was an authentication problem. If the authentication procedure is passed, block 121 directs the call controller circuit 100 to determine whether or not the contents of the caller ID field 60 of the SIP invite message received from the telephone is an IP address. If it is an IP address, then block 123 directs the call controller circuit 100 to set the contents of a type field variable maintained by the microprocessor 102 to a code representing that the call type is a third party invite. If at block 121 the caller ID field contents do not identify an IP address, then block 125 directs the microprocessor to set the contents of the type field to a code indicating 15 that the call is being made by a system subscriber. Then, block 126 directs the call controller circuit to read the call identifier 65 provided in the SIP invite message from the telephone 12, and at block 128 the processor is directed to produce an RC request message that includes that call ID. Block 129 then directs the call controller circuit 100 to send the RC request to the RC 16.

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Referring to Figure 6, an RC request message is shown generally at 150 and includes a caller field 152, a callee field 154, a digest field 156, a call ID field 158 and a type field 160. The caller, callee, digest call ID fields 152, 154, 156 and 158 contain copies of the caller, callee, digest parameters and call ID fields 60, 62, 64 and 65 of the SIP invite message shown in Figure 3. The type field 160 contains the type code established at blocks 123 or 125 of Figure 5 to indicate whether the call is from a third party or system subscriber, respectively. The caller identifier field may include a PSTN number or a system subscriber username as shown, for example.

## 30 Routing Controller (RC)

Referring to Figure 7, the RC 16 is shown in greater detail and includes an RC processor circuit shown generally at 200. The RC processor circuit 200 includes a processor 202, program memory 204, a table memory 206, buffer memory 207, and an I/O port 208, all in communication with the processor 202. (As earlier indicated, there may be a plurality of processor circuits (202), memories (204), etc.)

The buffer memory 207 includes a caller id buffer 209 and a callee id buffer 211.

The I/O port 208 includes a database request port 210 through which a request to
10 the database (18 shown in Figure 1) can be made and includes a database
response port 212 for receiving a reply from the database 18. The I/O port 208
further includes an RC request message input 214 for receiving the RC request
message from the call controller (14 shown in Figure 1) and includes a routing
message output 216 for sending a routing message back to the call controller 14.
15 The I/O port 208 thus acts to receive caller identifier and a callee identifier contained
in the RC request message from the call controller, the RC request message being
received in response to initiation of a call by a calling subscriber.

The program memory **204** includes blocks of codes for directing the processor **202** to carry out various functions of the RC (**16**). One of these blocks includes an RC request message handler **250** which directs the RC to produce a routing message in response to a received RC request message. The RC request message handler process is shown in greater detail at **250** in Figures **8**A through **8**D.

## 25 RC Request Message Handler

Referring to Figure 8A, the RC request message handler begins with a first block 252 that directs the RC processor circuit (200) to store the contents of the RC request message (150) in buffers in the buffer memory 207 of Figure 7, one of which includes the caller ID buffer 209 of Figure 7 for separately storing the contents of the callee field 154 of the RC request message. Block 254 then directs the RC

processor circuit to use the contents of the caller field **152** in the RC request message shown in Figure **6**, to locate and retrieve from the database **18** a record associating calling attributes with the calling subscriber. The located record may be referred to as a dialing profile for the caller. The retrieved dialing profile may then be stored in the buffer memory **207**, for example.

Referring to Figure 9, an exemplary data structure for a dialing profile is shown generally at 253 and includes a user name field 258, a domain field 260, and calling attributes comprising a national dialing digits (NDD) field 262, an international dialing digits (IDD) field 264, a country code field 266, a local area codes field 267, a caller minimum local length field 268, a caller maximum local length field 270, a reseller field 273, a maximum number of concurrent calls field 275 and a current number of concurrent calls field 277. Effectively the dialing profile is a record identifying calling attributes of the caller identified by the caller identifier. More generally, dialing profiles represent calling attributes of respective subscribers.

An exemplary caller profile for the Vancouver subscriber is shown generally at 276 in Figure 10 and indicates that the user name field 258 includes the user name (2001 1050 8667) that has been assigned to the subscriber and is stored in the user name field 50 in the telephone as shown in Figure 2.

Referring back to Figure 10, the domain field 260 includes a domain name as shown at 282, including a node type identifier 284, a location code identifier 286, a system provider identifier 288 and a domain portion 290. The domain field 260 effectively identifies a domain or node associated with the user identified by the contents of the user name field 258.

In this embodiment, the node type identifier 284 includes the code "sp" identifying a supernode and the location identifier 286 identifies the supernode as being in

Vancouver (YVR). The system provider identifier **288** identifies the company supplying the service and the domain portion **290** identifies the "com" domain.

The national dialled digit field **262** in this embodiment includes the digit "**1**" and, in general, includes a number specified by the International Telecommunications Union (ITU) Telecommunications Standardization Sector (ITU-T) E. **164** Recommendation which assigns national dialing digits to countries.

The international dialing digit field **264** includes a code also assigned according to the ITU-T according to the country or location of the user.

The country code field **266** also includes the digit "1" and, in general, includes a number assigned according to the ITU-T to represent the country in which the user is located.

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The local area codes field 267 includes a list of area codes that have been assigned by the ITU-T to the geographical area in which the subscriber is located. The caller minimum and maximum local number length fields 268 and 270 hold numbers representing minimum and maximum local number lengths permitted in the area code(s) specified by the contents of the local area codes field 267. The reseller field 273 is optional and holds a code identifying a retailer of the services, in this embodiment "Klondike". The maximum number of concurrent calls field 275 holds a code identifying the maximum number of concurrent calls that the user is entitled to cause to concurrently exist. This permits more than one call to occur concurrently while all calls for the user are billed to the same account. The current number of concurrent calls field 277 is initially 0 and is incremented each time a concurrent call is terminated.

The area codes associated with the user are the area codes associated with the location code identifier **286** of the contents of the domain field **260**.

A dialing profile of the type shown in Figure 9 is produced whenever a user registers with the system or agrees to become a subscriber to the system. Thus, for example, a user wishing to subscribe to the system may contact an office maintained by a system operator and personnel in the office may ask the user certain questions about his location and service preferences, whereupon tables can be used to provide office personnel with appropriate information to be entered into the user name 258, domain 260, NDD 262, IDD 264, country code 266, local area codes 267, caller minimum and maximum local length fields 268 and 270 reseller field 273 and concurrent call fields 275 and 277 to establish a dialing profile for the user.

Referring to Figures **11** and **12**, callee dialing profiles for users in Calgary and London, respectively for example, are shown.

In addition to creating dialing profiles when a user registers with the system, a direct-in-dial (DID) record of the type shown at **278** in Figure **13** is added to a direct-in-dial bank table in the database (**18** in Figure **1**) to associate the username and a host name of the supernode with which the user is associated, with an E.**164** number associated with the user on the PSTN network.

An exemplary DID table record entry for the Calgary callee is shown generally at 300 in Figure 14. The user name field 281 and user domain field 272 are analogous to the user name and user domain fields 258 and 260 of the caller dialing profile shown in Figure 10. The contents of the DID field 274 include a E.164 public telephone number including a country code 283, an area code 285, an exchange code 287 and a number 289. If the user has multiple telephone numbers, then multiple records of the type shown at 300 would be included in the DID bank table, each having the

same user name and user domain, but different DID field **274** contents reflecting the different telephone numbers associated with that user.

In addition to creating dialing profiles as shown in Figure 9 and DID records as shown in Figure 13 when a user registers with the system, call blocking records of the type shown in Figure 26, call forwarding records of the type shown in Figure 28 and voicemail records of the type shown in Figure 30 may be added to the database 18 when a new subscriber is added to the system.

Referring back to Figure 8A, after retrieving a dialing profile for the caller, such as shown at 276 in Figure 10, the RC processor circuit 200 is directed to block 256 which directs the processor circuit (200) to determine whether the contents of the concurrent call field 277 are less then the contents of the maximum concurrent call field 275 of the dialing profile for the caller and, if so, block 271 directs the processor circuit to increment the contents of the concurrent call field 277. If the contents of concurrent call field 277 are equal to or greater than the contents of the maximum concurrent call field 275, block 259 directs the processor circuit 200 to send an error message back to the call controller (14) to cause the call controller to notify the caller that the maximum number of concurrent calls has been reached and no further calls can exist concurrently, including the presently requested call.

Assuming block **256** allows the call to proceed, the RC processor circuit **200** is directed to perform certain checks on the callee identifier provided by the contents of the callee field **154** in Figure **6**, of the RC request message **150**. These checks are shown in greater detail in Figure **8**B.

Referring to Figure 8B, the processor (202 in Figure 7) is directed to a first block 257 that causes it to determine whether a digit pattern of the callee identifier (154) provided in the RC request message (150) includes a pattern that matches the contents of the international dialing digits (IDD) field 264 in the caller profile shown in

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Figure 10. If so, then block 259 directs the processor (202) to set a call type code identifier variable maintained by the processor to indicate that the call is an international call and block 261 directs the processor to produce a reformatted callee identifier by reformatting the callee identifier into a predefined digit format. In this embodiment, this is done by removing the pattern of digits matching the IDD field contents 264 of the caller dialing profile to effectively shorten the callee identifier. Then, block 263 directs the processor 202 to determine whether or not the callee identifier has a length which meets criteria establishing it as a number compliant with the E.164 Standard set by the ITU. If the length does not meet this criteria, block 265 directs the processor 202 to send back to the call controller (14) a message indicating the length is not correct. The process is then ended. At the call controller 14, routines (not shown) stored in the program memory 104 may direct the processor (102 of Figure 4) to respond to the incorrect length message by transmitting a message back to the telephone (12 shown in Figure 1) to indicate that an invalid number has been dialled.

Still referring to Figure 8B, if the length of the amended callee identifier meets the criteria set forth at block 263, block 269 directs the processor (202 of Figure 7) to make a database request to determine whether or not the amended callee identifier 20 is found in a record in the direct-in-dial bank (DID) table. Referring back to Figure 8B, at block 269, if the processor 202 receives a response from the database indicating that the reformatted callee identifier produced at block 261 is found in a record in the DID bank table, then the callee is a subscriber to the system and the call is classified as a private network call by directing the processor to block 279 which directs the processor to copy the contents of the corresponding user name field (281 in Figure 14) from the callee DID bank table record (300 in Figure 14) into the callee ID buffer (211 in Figure 7). Thus, the processor 202 locates a subscriber user name associated with the reformatted callee identifier. The processor 202 is then directed to point B in Figure 8A.

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## Subscriber to Subscriber Calls Between Different Nodes

Referring to Figure 8A, block 280 directs the processor (202 of Figure 7) to execute a process to determine whether or not the node associated with the reformatted callee identifier is the same node that is associated with the caller identifier. To do this, the processor 202 determines whether or not a prefix (e.g., continent code 61) of the callee name held in the callee ID buffer (211 in Figure 7), is the same as the corresponding prefix of the caller name held in the username field 258 of the caller dialing profile shown in Figure 10. If the corresponding prefixes are not the same, block 302 in Figure 8A directs the processor (202 in Figure 7) to set a call type flag in the buffer memory (207 in Figure 7) to indicate the call is a cross-domain call. Then, block 350 of Figure 8A directs the processor (202 of Figure 7) to produce a routing message identifying an address on the private network with which the callee identified by the contents of the callee ID buffer is associated and to set a time to live for the call at a maximum value of 99999, for example.

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Thus the routing message includes a caller identifier, a call identifier set according to a username associated with the located DID bank table record and includes an identifier of a node on the private network with which the callee is associated.

The node in the system with which the callee is associated is determined by using the callee identifier to address a supernode table having records of the type as shown at 370 in Figure 17. Each record 370 has a prefix field 372 and a supernode address field 374. The prefix field 372 includes the first n digits of the callee identifier. In this embodiment n=2. The supernode address field 374 holds a code representing the IP address or a fully qualified domain name of the node associated with the code stored in the callee identifier prefix field 372. Referring to Figure 18, for example, if the prefix is 20, the supernode address associated with that prefix is sp.yvr.digifonica.com.

Referring to Figure 15, a generic routing message is shown generally at 352 and includes an optional supplier prefix field 354, and optional delimiter field 356, a callee user name field 358, at least one route field 360, a time to live field 362 and other fields 364. The optional supplier prefix field 354 holds a code for identifying supplier traffic. The optional delimiter field 356 holds a symbol that delimits the supplier prefix code from the callee user name field 358. In this embodiment, the symbol is a number sign (#). The route field 360 holds a domain name or IP address of a gateway or node that is to carry the call, and the time to live field 362 holds a value representing the number of seconds the call is permitted to be active, based on subscriber available minutes and other billing parameters.

Referring to Figure **8**A and Figure **16**, an example of a routing message produced by the processor at block **350** for a caller associated with a different node than the caller is shown generally at **366** and includes only a callee field **359**, a route field **361** and a time to live field **362**.

Referring to Figure 8A, having produced a routing message as shown in Figure 16, block 381 directs the processor (202 of Figure 7) to send the routing message shown in Figure 16 to the call controller 14 shown in Figure 1.

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Referring back to Figure 8B, if at block 257, the callee identifier stored in the callee id buffer (211 in Figure 7) does not begin with an international dialing digit, block 380 directs the processor (202) to determine whether or not the callee identifier begins with the same national dial digit code as assigned to the caller. To do this, the processor (202) is directed to refer to the retrieved caller dialing profile as shown in Figure 10. In Figure 10, the national dialing digit code 262 is the number 1. Thus, if the callee identifier begins with the number 1, then the processor (202) is directed to block 382 in Figure 8B.

Block 382 directs the processor (202 of Figure 7) to examine the callee identifier to determine whether or not the digits following the NDD digit identify an area code that is the same as any of the area codes identified in the local area codes field 267 of the caller dialing profile 276 shown in Figure 10. If not, block 384 of Figure 8B directs the processor 202 to set the call type flag to indicate that the call is a national call. If the digits following the NDD digit identify an area code that is the same as a local area code associated with the caller as indicated by the caller dialing profile, block 386 directs the processor 202 to set the call type flag to indicate a local call, national style. After executing blocks 384 or 386, block 388 directs the processor 202 to format the callee identifier into a pre-defined digit format to produce a re-formatted callee identifier by removing the national dialled digit and prepending a caller country code identified by the country code field 266 of the caller dialing profile shown in Figure 10. The processor (202) is then directed to block 263 of Figure 8B to perform other processing as already described above.

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If at block 380, the callee identifier does not begin with a national dialled digit, block 390 directs the processor (202) to determine whether the callee identifier begins with digits that identify the same area code as the caller. Again, the reference for this is the retrieved caller dialing profile shown in Figure 10. The processor (202) determines whether or not the first few digits of the callee identifier identify an area code corresponding to the local area code field 267 of the retrieved caller dialing profile. If so, then block 392 directs the processor 202 to set the call type flag to indicate that the call is a local call and block 394 directs the processor (202) to format the callee identifier into a pre-defined digit format to produce a reformatted callee identifier by prepending the caller country code to the callee identifier, the caller country code being determined from the country code field 266 of the retrieved caller dialing profile shown in Figure 10. The processor (202) is then directed to block 263 for further processing as described above.

Referring back to Figure 8B, at block 390, the callee identifier does not start with the same area code as the caller, block 396 directs the processor (202 of Figure 7) to determine whether the number of digits in the callee identifier, i.e. the length of the callee identifier, is within the range of digits indicated by the caller minimum local number length field 268 and the caller maximum local number length field 270 of the retrieved caller dialing profile shown in Figure 10. If so, then block 398 directs the processor (202) to set the call type flag to indicate a local call and block 400 directs the processor (202) to format the callee identifier into a pre-defined digit format to produce a reformatted callee identifier by prepending to the callee identifier the caller country code (as indicated by the country code field 266 of the retrieved caller dialing profile shown in Figure 10) followed by the caller area code (as indicated by the local area code field 267 of the caller profile shown in Figure 10). The processor (202) is then directed to block 263 of Figure 8B for further processing as described above.

Referring back to Figure 8B, if at block 396, the callee identifier has a length that 15 does not fall within the range specified by the caller minimum local number length field (268 in Figure 10) and the caller maximum local number length field (270 in Figure 10), block 402 directs the processor 202 of Figure 7 to determine whether or not the callee identifier identifies a valid user name. To do this, the processor 202 searches through the database (18 of Figure 10 of dialing profiles to find a dialing 20 profile having user name field contents (258 in Figure 10) that match the callee identifier. If no match is found, block 404 directs the processor (202) to send an error message back to the call controller (14). If at block 402, a dialing profile having a user name field 258 that matches the callee identifier is found, block 406 directs the 25 processor 202 to set the call type flag to indicate that the call is a private network call and then the processor is directed to block 280 of Figure 8A. Thus, the call is classified as a private network call when the callee identifier identifies a subscriber to the private network.

From Figure 8B, it will be appreciated that there are certain groups of blocks of codes that direct the processor 202 in Figure 7 to determine whether the callee identifier has certain features such as an international dialing digit, a national dialing digit, an area code and a length that meet certain criteria, and cause the processor 202 to reformat the callee identifier stored in the callee id buffer 211, as necessary into a predetermined target format including only a country code, area code, and a normal telephone number, for example, to cause the callee identifier to be compatible with the E.164 number plan standard in this embodiment. This enables block 269 in Figure 8B to have a consistent format of callee identifiers for use in searching through the DID bank table records of the type shown in Figure 13 to determine how to route calls for subscriber to subscriber calls on the same system. Effectively, therefore blocks 257, 380, 390, 396 and 402 establish call classification criteria for classifying the call as a public network call or a private network call. Block 269 classifies the call, depending on whether or not the formatted callee identifier 15 has a DID bank table record and this depends on how the call classification criteria are met and block 402 directs the processor 202 of Figure 7 to classify the call as a private network call when the callee identifier complies with a pre-defined format, i.e. is a valid user name and identifies a subscriber to the private network, after the callee identifier has been subjected to the classification criteria of blocks 257, 380, 390 and 396.

#### Subscriber to Non-Subscriber Calls

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25 Not all calls will be subscriber to subscriber calls and this will be detected by the processor 202 of Figure 7 when it executes block 269 in Figure 8B, and does not find a DID bank table record that is associated with the callee, in the DID bank table. When this occurs, the call is classified as a public network call by directing the processor 202 to block 408 of Figure 8B which causes it to set the contents of the 30 callee id buffer 211 of Figure 7 equal to the newly formatted callee identifier, i.e., a

number compatible with the E.164 standard. Then, block 410 of Figure 8B directs the processor (202) to search a database of route or master list records associating route identifiers with dialing codes shown in Figure 19 to locate a router having a dialing code having a number pattern matching at least a portion of the reformatted callee identifier.

Referring to Figure 19, a data structure for a master list or route list record is shown. Each master list record includes a master list ID field 500, a dialing code field 502, a country code field 504, a national sign number field 506, a minimum length field 508, 10 a maximum length field 510, a national dialled digit field 512, an international dialled digit field **514** and a buffer rate field **516**.

The master list ID field 500 holds a unique code such as 1019, for example, identifying the record. The dialing code field 502 holds a predetermined number pattern that the processor 202 of Figure 7 uses at block 410 in Figure 8B to find the master list record having a dialing code matching the first few digits of the amended callee identifier stored in the callee id buffer 211. The country code field 504 holds a number representing the country code associated with the record and the national sign number field 506 holds a number representing the area code associated with 20 the record. (It will be observed that the dialing code is a combination of the contents of the country code field 504 and the national sign number field 506.) The minimum length field 508 holds a number representing the minimum length of digits associated with the record and the maximum length field 51 holds a number representing the maximum number of digits in a number with which the record may be compared. The national dialled digit (NDD) field 512 holds a number representing an access code used to make a call within the country specified by the country code, and the international dialled digit (IDD) field 514 holds a number representing the international prefix needed to dial a call from the country indicated by the country code.

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Thus, for example, a master list record may have a format as shown in Figure 20 with exemplary field contents as shown.

Referring back to Figure 8B, using the country code and area code portions of the reformatted callee identifier stored in the callee id buffer 211, block 410 directs the processor 202 of Figure 7 to find a master list record such as the one shown in Figure 20 having a dialing code that matches the country code (1) and area code (604) of the callee identifier. Thus, in this example, the processor (202) would find a master list record having an ID field containing the number 1019. This number may be referred to as a route ID. Thus, a route ID number is found in the master list record associated with a predetermined number pattern in the reformatted callee identifier.

After executing block **410** in Figure **8**B, the process continues as shown in Figure **8**D. Referring to Figure **8**D, block **412** directs the processor **202** of Figure **7** to use the route ID number to search a database of supplier records associating supplier identifiers with route identifiers to locate at least one supplier record associated with the route identifier to identify at least one supplier operable to supply a communications link for the route.

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Referring to Figure 21, a data structure for a supplier list record is shown. Supplier list records include a supplier ID field 540, a master list ID field 542, an optional prefix field 544, a specific route identifier field 546, a NDD/IDD rewrite field 548, a rate field 550, and a timeout field 551. The supplier ID field 540 holds a code identifying the name of the supplier and the master list ID field 542 holds a code for associating the supplier record with a master list record. The prefix field 544 holds a string used to identify the supplier traffic and the specific route identifier field 546 holds an IP address of a gateway operated by the supplier indicated by the supplier ID field 540. The NDD/IDD rewrite field 548 holds a code representing a rewritten value of the NDD/IDD associated with this route for this supplier, and the rate field

550 holds a code indicating the cost per second to the system operator to use the route provided by the gateway specified by the contents of the route identifier field 546. The timeout field 551 holds a code indicating a time that the call controller should wait for a response from the associated gateway before giving up and trying the next gateway. This time value may be in seconds, for example. Exemplary supplier records are shown in Figures 22, 23 and 24 for the exemplary suppliers shown at 20 in Figure 1, namely Telus, Shaw and Sprint.

Referring back to Figure 8D, at block 412 the processor 202 finds all supplier 10 records that identify the master list ID found at block **410** of Figure **8**B.

Referring back to Figure 8D, block 560 directs the processor 202 of Figure 7 to begin to produce a routing message of the type shown in Figure 15. To do this, the processor 202 loads a routing message buffer as shown in Figure 25 with a supplier prefix of the least costly supplier where the least costly supplier is determined from the rate fields 550 of Figure 21 of the records associated with respective suppliers.

Referring to Figures 22-24, in the embodiment shown, the supplier "Telus" has the lowest number in the rate field 550 and therefore the prefix 4973 associated with that supplier is loaded into the routing message buffer shown in Figure 25 first.

Block 562 in Figure 8D directs the processor to delimit the prefix 4973 by the number sign (#) and to next load the reformatted callee identifier into the routing message buffer shown in Figure 25. At block 563 of Figure 8D, the contents of the route identifier field 546 of Figure 21 of the record associated with the supplier "Telus" are added by the processor 202 of Figure 7 to the routing message buffer shown in Figure 25 after an @ sign delimiter, and then block 564 in Figure 8D directs the processor to get a time to live value, which in one embodiment may be 3600 seconds, for example. Block 566 then directs the processor 202 to load this time to 30 live value and the timeout value (551) in Figure 21 in the routing message buffer of

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Figure 25. Accordingly, a first part of the routing message for the Telus gateway is shown generally at 570 in Figure 25.

Referring back to Figure 8D, block 571 directs the processor 202 back to block 560 5 and causes it to repeat blocks 560, 562, 563, 564 and 566 for each successive supplier until the routing message buffer is loaded with information pertaining to each supplier identified by the processor at block 412. Thus, a second portion of the routing message as shown at 572 in Figure 25 relates to the second supplier identified by the record shown in Figure 23. Referring back to Figure 25, a third portion of the routing message as shown at 574 and is associated with a third supplier as indicated by the supplier record shown in Figure 24.

Consequently, referring to Figure 25, the routing message buffer holds a routing message identifying a plurality of different suppliers able to provide gateways to the 15 public telephone network (i.e. specific routes) to establish at least part of a communication link through which the caller may contact the callee. embodiment, each of the suppliers is identified, in succession, according to rate. Other criteria for determining the order in which suppliers are listed in the routing message may include preferred supplier priorities which may be established based 20 on service agreements, for example.

Referring back to Figure 8D, block 568 directs the processor 202 of Figure 7 to send the routing message shown in Figure 25 to the call controller 14 in Figure 1.

#### 25 Subscriber to Subscriber Calls Within the Same Node

Referring back to Figure 8A, if at block 280, the callee identifier received in the RC request message has a prefix that identifies the same node as that associated with the caller, block 600 directs the processor 202 to use the callee identifier in the callee id buffer 211 to locate and retrieve a dialing profile for the callee. The dialing profile may be of the type shown in Figure 11 or 12, for example. Block 602 of Figure

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8A then directs the processor 202 of Figure 7 to get call block, call forward and voicemail records from the database 18 of Figure 1 based on the user name identified in the callee dialing profile retrieved by the processor at block 600. Call block, call forward and voicemail records may be as shown in Figures 26, 27, 28 and 30 for example.

Referring to Figure 26, the call block records include a user name field 604 and a block pattern field 606. The user name field holds a user name corresponding to the user name in the user name field (258 in Figure 10) of the callee profile and the block pattern field 606 holds one or more E.164-compatible numbers or user names identifying PSTN numbers or system subscribers from whom the subscriber identified in the user name field 604 does not wish to receive calls.

Referring to Figure 8A and Figure 27, block 608 directs the processor 202 of Figure
7 to determine whether or not the caller identifier received in the RC request
message matches a block pattern stored in the block pattern field 606 of the call
block record associated with the callee identified by the contents of the user name
field 604 in Figure 26. If the caller identifier matches a block pattern, block 610
directs the processor to send a drop call or non-completion message to the call
controller (14) and the process is ended. If the caller identifier does not match a
block pattern associated with the callee, block 609 directs the processor to store the
username and domain of the callee, as determined from the callee dialing profile,
and a time to live value in the routing message buffer as shown at 650 in Figure 32.
Referring back to Figure 8A, block 612 then directs the processor 202 to determine
whether or not call forwarding is required.

Referring to Figure 28, the call forwarding records include a user name field 614, a destination number field 616, and a sequence number field 618. The user name field 614 stores a code representing a user with which the record is associated. The destination number field 616 holds a user name representing a number to which the

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current call should be forwarded, and the sequence number field **618** holds an integer number indicating the order in which the user name associated with the corresponding destination number field **616** should be attempted for call forwarding. The call forwarding table may have a plurality of records for a given user. The processor **202** of Figure **7** uses the contents of the sequence number field **618** to place the records for a given user in order. As will be appreciated below, this enables the call forwarding numbers to be tried in an ordered sequence.

Referring to Figure 8A and Figure 29, if at block 612, the call forwarding record for the callee identified by the callee identifier contains no contents in the destination number field 616 and accordingly no contents in the sequence number field 618, there are no call forwarding entries for this callee, and the processor 202 is directed to block 620 in Figure 8C. If there are entries in the call forwarding table 27, block 622 in Figure 8A directs the processor 202 to search the dialing profile table to find a dialing profile record as shown in Figure 9, for the user identified by the destination number field 616 of the call forward record shown in Figure 28. The processor 202 of Figure 7 is further directed to store the username and domain for that user and a time to live value in the routing message buffer as shown at 652 in Figure 32, to produce a routing message as illustrated. This process is repeated for each call forwarding record associated with the callee identified by the callee id buffer 211 in Figure 7 to add to the routing message buffer all call forwarding usernames and domains associated with the callee.

Referring back to Figure 8A, if at block 612 there are no call forwarding records, then at block 620 in Figure 8C the processor 202 is directed to determine whether or not the user identified by the callee identifier has paid for voicemail service. This is done by checking to see whether or not a flag is set in a voicemail record of the type shown in Figure 30 in a voicemail table stored in the database 18 shown in Figure 1.

Referring to Figure 30, voicemail records in this embodiment may include a user name field 624, a voicemail server field 626, a seconds to voicemail field 628 and an enable field 630. The user name field 624 stores the user name of the callee. The voicemail server field 626 holds a code identifying a domain name of a voicemail server associated with the user identified by the user name field 624. The seconds to voicemail field 628 holds a code identifying the time to wait before engaging voicemail, and the enable field 630 holds a code representing whether or not voicemail is enabled for the user. Referring back to Figure 8C, at block 620 if the processor 202 of Figure 7 finds a voicemail record as shown in Figure 30 having user name field 624 contents matching the callee identifier, the processor is directed to examine the contents of the enabled field 630 to determine whether or not voicemail is enabled. If voicemail is enabled, then block 640 in Figure 8C directs the processor 202 to Figure 7 to store the contents of the voicemail server field 626 and the contents of the seconds to voicemail field 628 in the routing message buffer, as shown at 654 in Figure 32. Block 642 then directs the processor 202 to get time to live values for each path specified by the routing message according to the cost of routing and the user's balance. These time to live values are then appended to corresponding paths already stored in the routing message buffer.

Referring back to Figure 8C, block 644 then directs the processor 202 of Figure 7 to store the IP address of the current node in the routing message buffer as shown at 656 in Figure 32. Block 646 then directs the processor 202 to send the routing message shown in Figure 32 to the call controller 14 in Figure 1. Thus in the embodiment described the routing controller will produce a routing message that will cause at least one of the following: forward the call to another party, block the call and direct the caller to a voicemail server.

Referring back to Figure 1, the routing message whether of the type shown in Figures 16, 25 or 32, is received at the call controller 14 and the call controller 30 interprets the receipt of the routing message as a request to establish a call.

Referring to Figure 4, the program memory 104 of the call controller 14 includes a routing to gateway routine depicted generally at 122.

Where a routing message of the type shown in Figure 32 is received by the call controller 14, the routing to gateway routine 122 shown in Figure 4 may direct the processor 102 cause a message to be sent back through the internet 13 shown in Figure 1 to the callee telephone 15, knowing the IP address of the callee telephone 15 from the user name.

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Alternatively, if the routing message is of the type shown in Figure **16**, which identifies a domain associated with another node in the system, the call controller may send a SIP invite message along the high speed backbone **17** connected to the other node. The other node functions as explained above, in response to receipt of a SIP invite message.

If the routing message is of the type shown in Figure 25 where there are a plurality of gateway suppliers available, the call controller sends a SIP invite message to the first supplier, in this case Telus, using a dedicated line or an internet connection to determine whether or not Telus is able to handle the call. If the Telus gateway returns a message indicating it is not able to handle the call, the call controller 14 then proceeds to send a SIP invite message to the next supplier, in this case Shaw. The process is repeated until one of the suppliers responds indicating that it is available to carry the call. Once a supplier responds indicating that it is able to carry the call, the supplier sends back to the call controller 14 an IP address for a gateway provided by the supplier through which the call or audio path of the call will be carried. This IP address is sent in a message from the call controller 14 to the media relay 9 which responds with a message indicating an IP address to which the caller telephone should send its audio/video, traffic and an IP address to which the gateway should send its audio/video for the call. The call controller conveys the IP

address at which the media relay expects to receive audio/video from the caller telephone, to the caller telephone 12 in a message. The caller telephone replies to the call controller with an IP address at which it would like to receive audio/video and the call controller conveys that IP address to the media relay. The call may then be conducted between the caller and callee through the media relay and gateway.

Referring back to Figure 1, if the call controller 14 receives a routing message of the type shown in Figure 32, and which has at least one call forwarding number and/or a voicemail number, the call controller attempts to establish a call to the callee telephone 15 by seeking from the callee telephone a message indicating an IP address to which the media relay should send audio/video. If no such message is received from the callee telephone, no call is established. If no call is established within a pre-determined time, the call controller 14 attempts to establish a call with the next user identified in the call routing message in the same manner. This process is repeated until all call forwarding possibilities have been exhausted, in 15 which case the call controller communicates with the voicemail server 19 identified in the routing message to obtain an IP address to which the media relay should send audio/video and the remainder of the process mentioned above for establishing IP addresses at the media relay 9 and the caller telephone is carried out to establish audio/video paths to allowing the caller to leave a voicemail message with the voicemail server.

When an audio/video path through the media relay is established, a call timer maintained by the call controller **14** logs the start date and time of the call and logs the call ID and an identification of the route (i.e., audio/video path IP address) for later use in billing.

### Time to Live

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Referring to Figures 33A and 33B, a process for determining a time to live value for any of blocks 642 in Figure 8C, 350 in Figure 8A or 564 in Figure 8D above is

described. The process is executed by the processor 202 shown in Figure 7. Generally, the process involves calculating a cost per unit time, calculating a first time value as a sum of a free time attributed to a participant in the communication session and the quotient of a funds balance held by the participant to the cost per unit time value and producing a second time value in response to the first time value and a billing pattern associated with the participant, the billing pattern including first and second billing intervals and the second time value being the time to permit a communication session to be conducted.

- Referring to Figure 33A, in this embodiment, the process begins with a first block 700 that directs the RC processor to determine whether or not the call type set at block 302 in Figure 8A indicates the call is a network or cross-domain call. If the call is a network or cross-domain call, block 702 of Figure 33A directs the RC processor to set the time to live equal to 99999 and the process is ended. Thus, the network or cross-domain call type has a long time to live. If at block 700 the call type is determined not to be a network or cross-domain type, block 704 directs the RC processor to get a subscriber bundle table record from the database 18 in Figure 1 and store it locally in the subscriber bundle record buffer at the RC 14.
- Referring to Figure **34**, a subscriber bundle table record is shown generally at **706**. The record includes a user name field **708** and a services field **710**. The user name field **708** holds a code identifying the subscriber user name and the services field **710** holds codes identifying service features assigned to the subscriber, such as free local calling, call blocking and voicemail, for example.

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Figure 35 shows an exemplary subscriber bundle record for the Vancouver caller. In this record the user name field 708 is loaded with the user name 2001 1050 8667 and the services field 710 is loaded with codes 10, 14 and 16 corresponding to free local calling, call blocking and voicemail, respectively. Thus, user 2001 1050 8667 has free local calling, call blocking and voicemail features.

Referring back to Figure 33A, after having loaded a subscriber bundle record into the subscriber bundle record buffer, block 712 directs the RC processor to search the database (18) determine whether or not there is a bundle override table record for the master list ID value that was determined at block 410 in Figure 8B. An exemplary bundle override table record is shown at 714 in Figure 36. The bundle table record includes a master list ID field 716, an override type field 718, an override value field 720 a first interval field 722 and a second interval field 724. The master list ID field 716 holds a master list ID code. The override type field 718 holds an override type code indicating a fixed, percent or cent amount to indicate the amount by which a fee will be increased. The override value field 720 holds a real number representing the value of the override type. The first interval field 722 holds a value indicating the minimum number of seconds for a first level of charging and the second interval field 724 holds a number representing a second level of charging.

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Referring to Figure 37, a bundle override record for the located master list ID code is shown generally at 726 and includes a master list ID field 716 holding the code 1019 which was the code located in block 410 of Figure 8B. The override type field 718 includes a code indicating the override type is a percentage value and the override value field 720 holds the value 10.0 indicating that the override will be 10.0% of the charged value. The first interval field 722 holds a value representing 30 seconds and the second interval field 724 holds a value representing 6 seconds. The 30 second value in the first interval field 722 indicates that charges for the route will be made at a first rate for 30 seconds and thereafter the charges will be made at a different rate in increments of 6 seconds, as indicated by the contents of the second interval field 724.

Referring back to Figure 33A, if at block 712 the processor finds a bundle override record of the type shown in Figure 37, block 728 directs the processor to store the bundle override record in local memory. In the embodiment shown, the bundle

override record shown in Figure 37 is stored in the bundle override record buffer at the RC as shown in Figure 7. Still referring to Figure 33A, block 730 then directs the RC processor to determine whether or not the subscriber bundle table record 706 in Figure 35 has a services field including a code identifying that the user is entitled to free local calling and also directs the processor to determine whether or not the call type is not a cross domain cell, i.e. it is a local or local/national style. If both of these conditions are satisfied, block 732 directs the processor to set the time to live equal to 99999, giving the user a long period of time for the call. The process is then ended. If the conditions associated with block 730 are not satisfied, block 734 of Figure 33B directs the RC processor to retrieve a subscriber account record associated with a participant in the call. This is done by copying and storing in the subscriber account record buffer a subscriber account record for the caller.

Referring to Figure 38, an exemplary subscriber account table record is shown generally at 736. The record includes a user name field 738, a funds balance field 740 and a free time field 742. The user name field 738 holds a subscriber user name, the funds balance field 740 holds a real number representing the dollar value of credit available to the subscriber and the free time field 742 holds an integer representing the number of free seconds that the user is entitled to.

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An exemplary subscriber account record for the Vancouver caller is shown generally at 744 in Figure 39, wherein the user name field 738 holds the user name 2001 1050 8667, the funds balance field 740 holds the value \$10.00, and the free time field 742 holds the value 100. The funds balance field holding the value of \$10.00 indicates the user has \$10.00 worth of credit and the free time field having the value of 100 indicates that the user has a balance of 100 free seconds of call time.

Referring back to Figure 33B, after copying and storing the subscriber account record shown in Figure 39 from the database to the subscriber account record buffer RC, block 746 directs the processor to determine whether or not the subscriber

account record funds balance field **740** or free time field **742** are greater than zero. If they are not greater than zero, block **748** directs the processor to set the time to live equal to zero and the process is ended. The RC then sends a message back to the call controller to cause the call controller to deny the call to the caller. If the conditions associated with block **746** are satisfied, block **750** directs the processor to calculate the call cost per unit time. A procedure for calculating the call cost per unit time is described below in connection with Figure **41**.

Assuming the procedure for calculating the cost per second returns a number representing the call cost per second, block **752** directs the processor **202** in Figure **7** to determine whether or not the cost per second is equal to zero. If so, block **754** directs the processor to set the time to live to **99999** to give the caller a very long length of call and the process is ended.

If at block 752 the call cost per second is not equal to zero, block 756 directs the processor 202 in Figure 7 to calculate a first time to live value as a sum of a free time attributed to the participant in the communication session and the quotient of the funds balance held by the participant to the cost per unit time value. To do this, the processor 202 of Figure 7 is directed to set a first time value or temporary time to live value equal to the sum of the free time provided in the free time field 742 of the subscriber account record shown in Figure 39 and the quotient of the contents of the funds balance field 740 in the subscriber account record for the call shown in Figure 39 and the cost per second determined at block 750 of Figure 33B. Thus, for example, if at block 750 the cost per second is determined to be three cents per second and the funds balance field holds the value \$10.00, the quotient of the funds balance and cost per second is 333 seconds and this is added to the contents of the free time field 742, which is 100, resulting in a time to live of 433 seconds.

Block **758** then directs the RC processor to produce a second time value in response to the first time value and the billing pattern associated with the participant as

established by the bundle override record shown in Figure 37. This process is shown in greater detail at 760 in Figure 40 and generally involves producing a remainder value representing a portion of the second billing interval remaining after dividing the second billing interval into a difference between the first time value and the first billing interval.

Referring to Figure 40, the process for producing the second time value begins with a first block 762 that directs the processor 202 in Figure 7 to set a remainder value equal to the difference between the time to live value calculated at block 756 in 10 Figure 33B and the contents of the first interval field 722 of the record shown in Figure 37, multiplied by the modulus of the contents of the second interval field 724 of Figure 37. Thus, in the example given, the difference between the time to live field and the first interval field is 433 minus 30, which is 403 and therefore the remainder produced by the mod of 403 divided by 6 is 0.17. Block 764 then directs the processor to determine whether or not this remainder value is greater than zero and, if so, block 766 directs the processor to subtract the remainder from the first time value and set the difference as the second time value. To do this the processor is directed to set the time to live value equal to the current time to live of 403 minus the remainder of 1, i.e., 402 seconds. The processor is then returned back to block 758 of Figure 33B.

Referring back to Figure 40, if at block 764 the remainder is not greater than zero, block 768 directs the processor 202 of Figure 7 to determine whether or not the time to live is less than the contents of the first interval field 722 in the record shown in 25 Figure 37. If so, then block 770 of Figure 40 directs the processor to set the time to live equal to zero. Thus, the second time value is set to zero when the remainder is greater than zero and the first time value is less than the free time associated with the participant in the call. If at block **768** the conditions of that block are not satisfied, the processor returns the first time to live value as the second time to live value.

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Thus, referring to Figure 33B, after having produced a second time to live value, block 772 directs the processor to set the time to live value for use in blocks 342, 350 or 564.

### 5 Cost per Second

Referring back to Figure 33B, at block 750 it was explained that a call cost per unit time is calculated. The following explains how that call cost per unit time value is calculated.

Referring to Figure 41, a process for calculating a cost per unit time is shown generally at 780. The process is executed by the processor 202 in Figure 7 and generally involves locating a record in a database, the record comprising a markup type indicator, a markup value and a billing pattern and setting a reseller rate equal to the sum of the markup value and the buffer rate, locating at least one of an override record specifying a route cost per unit time amount associated with a route associated with the communication session, a reseller record associated with a reseller of the communications session, the reseller record specifying a reseller cost per unit time associated with the reseller for the communication session and a default operator markup record specifying a default cost per unit time and setting as the cost per unit time the sum of the reseller rate and at least one of the route cost per unit time, the reseller cost per unit time and the default cost per unit time.

The process begins with a first set of blocks **782**, **802** and **820** which direct the processor **202** in Figure **7** to locate at least one of a record associated with a reseller and a route associated with the reseller, a record associated with the reseller, and a default reseller mark-up record. Block **782**, in particular, directs the processor to address the database **18** to look for a record associated with a reseller and a route with the reseller by looking for a special rate record based on the master list ID established at block **410** in Figure **8**C.

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Referring to Figure 42, a system operator special rate table record is shown generally at 784. The record includes a reseller field 786, a master list ID field 788, a mark-up type field 790, a mark-up value field 792, a first interval field 794 and a second interval field 796. The reseller field 786 holds a reseller ID code and the master list ID field 788 holds a master list ID code. The mark-up type field 790 holds a mark-up type such as fixed percent or cents and the mark-up value field 792 holds a real number representing the value corresponding to the mark-up type. The first interval field 794 holds a number representing a first level of charging and the second interval field 796 holds a number representing a second level of charging.

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An exemplary system operator special rate table for a reseller known as "Klondike" is shown at 798 in Figure 43. In this record, the reseller field 786 holds a code indicating the retailer ID is Klondike, the master list ID field 788 holds the code 1019 to associate the record with the master list ID code 1019. The mark-up type field 790 holds a code indicating the mark-up type is cents and the mark-up value field 792 holds a mark-up value indicating 1/10 of one cent. The first interval field 794 holds the value 30 and the second interval field 796 holds the value 6, these two fields indicating that the operator allows 30 seconds for free and then billing is done in increments of 6 seconds after that.

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Referring back to Figure 41, if at block 782 a record such as the one shown in Figure 43 is located in the system operator special rates table, the processor is directed to block 800 in Figure 41. If such a record is not found in the system operator special rates table, block 802 directs the processor to address the database 18 to look in a system operator mark-up table for a mark-up record associated with the reseller.

Referring to Figure **44**, an exemplary system operator mark-up table record is shown generally at **804**. The record includes a reseller field **806**, a mark-up type field **808**, a mark-up value field **810**, a first interval field **812** and a second interval field **814**. The reseller mark-up type, mark-up value, first interval and second interval fields are as

described in connection with the fields by the same names in the system operator special rates table shown in Figure 42.

Figure 45 provides an exemplary system operator mark-up table record for the 5 reseller known as Klondike and therefore the reseller field 806 holds the value "Klondike", the mark-up type field 808 holds the value cents, the mark-up value field holds the value 0.01, the first interval field 812 holds the value 30 and the second interval field 814 holds the value 6. This indicates that the reseller "Klondike" charges by the cent at a rate of one cent per minute. The first 30 seconds of the call are free and billing is charged at the rate of one cent per minute in increments of 6 seconds.

Figure 46 provides an exemplary system operator mark-up table record for cases where no specific system operator mark-up table record exists for a particular 15 reseller, i.e., a default reseller mark-up record. This record is similar to the record shown in Figure 45 and the reseller field 806 holds the value "all", the mark-up type field 808 is loaded with a code indicating mark-up is based on a percentage, the mark-up value field 810 holds the percentage by which the cost is marked up, and the first and second interval fields 812 and 814 identify first and second billing levels.

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Referring back to Figure 41, if at block 802 a specific mark-up record for the reseller identified at block 782 is not located, block 820 directs the processor to get the markup record shown in Figure 46, having the "all" code in the reseller field 806. The processor is then directed to block 800.

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Referring back to Figure 41, at block 800, the processor 202 of Figure 7 is directed to set a reseller rate equal to the sum of the mark-up value of the record located by blocks 782, 802 or 820 and the buffer rate specified by the contents of the buffer rate field 516 of the master list record shown in Figure 20. To do this, the RC processor sets a variable entitled "reseller cost per second" to a value equal to the sum of the contents of the mark-up value field (792, 810) of the associated record, plus the contents of the buffer rate field (516) from the master list record associated with the master list ID. Then, block 822 directs the processor to set a system operator cost per second variable equal to the contents of the buffer rate field (516) from the master list record. Block 824 then directs the processor to determine whether the call type flag indicates the call is local or national/local style and whether the caller has free local calling. If both these conditions are met, then block 826 sets the user cost per second variable equal to zero and sets two increment variables equal to one, for use in later processing. The cost per second has thus be calculated and the process shown in Figure 41 is ended.

If at block **824** the conditions of that block are not met, the processor **202** of Figure **7** is directed to locate at least one of a bundle override table record specifying a route cost per unit time associated with a route associated with the communication session, a reseller special destinations table record associated with a reseller of the communications session, the reseller record specifying a reseller cost per unit time associated with the reseller for the communication session and a default reseller global markup record specifying a default cost per unit time.

To do this block 828 directs the processor 202 of Figure 7 to determine whether or not the bundle override record 726 in Figure 37 located at block 712 in Figure 33A has a master list ID equal to the stored master list ID that was determined at block 410 in Figure 8B. If not, block 830 directs the processor to find a reseller special destinations table record in a reseller special destinations table in the database (18), having a master list ID code equal to the master list ID code of the master list ID that was determined at block 410 in Figure 8B. An exemplary reseller special destinations table record is shown in Figure 47 at 832. The reseller special destinations table record includes a reseller field 834, a master list ID field 836, a mark-up type field 838, a mark-up value field 840, a first interval field 842 and a second interval field 844. This record has the same format as the system operator

special rates table record shown in Figure 42, but is stored in a different table to allow for different mark-up types and values and time intervals to be set according to resellers' preferences. Thus, for example, an exemplary reseller special destinations table record for the reseller "Klondike" is shown at 846 in Figure 48. The reseller field 834 holds a value indicating the reseller as the reseller "Klondike" and the master list ID field holds the code 1019. The mark-up type field 838 holds a code indicating the mark-up type is percent and the mark-up value field 840 holds a number representing the mark-up value as 5%. The first and second interval fields identify different billing levels used as described earlier.

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Referring back to Figure **41**, the record shown in Figure **48** may be located at block **830**, for example. If at block **830** such a record is not found, then block **832** directs the processor to get a default operator global mark-up record based on the reseller ID.

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Referring to Figure 49, an exemplary default reseller global mark-up table record is shown generally at 848. This record includes a reseller field 850, a mark-up type field 852, a mark-up value field 854, a first interval field 856 and a second interval field 858. The reseller field 850 holds a code identifying the reseller. The mark-up type field 852, the mark-up value field 854 and the first and second interval fields 856 and 858 are of the same type as described in connection with fields of the same name in Figure 47, for example. The contents of the fields of this record 860 may be set according to system operator preferences, for example.

Referring to Figure **50**, an exemplary reseller global mark-up table record is shown generally at **860**. In this record, the reseller field **850** holds a code indicating the reseller is "Klondike", the mark-up type field **852** holds a code indicating the mark-up type is percent, the mark-up value field **854** holds a value representing **10**% as the mark-up value, the first interval field **856** holds the value **30** and the second interval

field **858** holds the values **30** and **6** respectively to indicate the first **30** seconds are free and billing is to be done in **6** second increments after that.

Referring back to Figure **41**, should the processor get to block **832**, the reseller global mark-up table record as shown in Figure **50** is retrieved from the database and stored locally at the RC. As seen in Figure **41**, it will be appreciated that if the conditions are met in blocks **828** or **830**, or if the processor executes block **832**, the processor is then directed to block **862** which causes it to set an override value equal to the contents of the mark-up value field of the located record, to set the first increment variable equal to the contents of the first interval field of the located record and to set the second increment variable equal to the contents of the second interval field of the located record. (The increment variables were alternatively set to specific values at block **826** in Figure **41**.)

It will be appreciated that the located record could be a bundle override record of the type shown in Figure 37 or the located record could be a reseller special destination record of the type shown in Figure 48 or the record could be a reseller global mark-up table record of the type shown in Figure 50. After the override and first and second increment variables have been set at block 862, the processor 202 if Figure 7 is directed to set as the cost per unit time the sum of the reseller rate and at least one of the route cost per unit time, the reseller cost per unit time and the default cost per unit time, depending on which record was located. To do this, block 864 directs the processor to set the cost per unit time equal to the sum of the reseller cost set at block 800 in Figure 41, plus the contents of the override variable calculated in block 862 in Figure 41. The cost per unit time has thus been calculated and it is this cost per unit time that is used in block 752 of Figure 33B, for example.

### Terminating the Call

In the event that either the caller or the callee terminates a call, the telephone of the terminating party sends a SIP bye message to the controller **14**. An exemplary SIP

bye message is shown at **900** in Figure **51** and includes a caller field **902**, a callee field **904** and a call ID field **906**. The caller field **902** holds a twelve digit user name, the callee field **904** holds a PSTN compatible number or user name, and the call ID field **906** holds a unique call identifier field of the type shown in the call ID field **65** of the SIP invite message shown in Figure **3**.

Thus, for example, referring to Figure **52**, a SIP bye message for the Calgary callee is shown generally at **908** and the caller field **902** holds a user name identifying the caller, in this case **2001 1050 8667**, the callee field **904** holds a user name identifying the Calgary callee, in this case **2001 1050 2222**, and the call ID field **906** holds the code FA**10** @ **192.168.0.20**, which is the call ID for the call.

The SIP bye message shown in Figure 52 is received at the call controller 14 and the call controller executes a process as shown generally at 910 in Figure 53. The process includes a first block 912 that directs the call controller processor 202 of Figure 7 to copy the caller, callee and call ID field contents from the SIP bye message received from the terminating party to corresponding fields of an RC stop message buffer (not shown). Block 914 then directs the processor to copy the call start time from the call timer and to obtain a call stop time from the call timer. Block 916 then directs the call controller to calculate a communication session time by determining the difference in time between the call start time and the call stop time. This session time is then stored in a corresponding field of the RC call stop message buffer. Block 917 then directs the processor to decrement the contents of the current concurrent call field 277 of the dialing profile for the caller as shown in Figure 10, to indicate that there is one less concurrent call in progress. A copy of the amended dialing profile for the caller is then stored in the database 18 of Figure 1. Block 918 then directs the processor to copy the route from the call log. An RC call stop message produced as described above is shown generally at 1000 in Figure 54. An RC call stop message specifically associated with the call made to the Calgary callee is shown generally at 1020 in Figure 55.

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Referring to Figure **54**, the RC stop call message includes a caller field **1002**, callee field **1004**, a call ID field **1006**, an account start time field **1008**, an account stop time field **1010**, a communication session time **1012** and a route field **1014**. The caller field **1002** holds a username, the callee field **1004** holds a PSTN-compatible number or system number, the call ID field **1006** hold the unique call identifier received from the SIP invite message shown in Figure **3**, the account start time field **1008** holds the date and start time of the call, the account stop time field **1010** holds the date and time the call ended, the communication session time field **1012** holds a value representing the difference between the start time and the stop time, in seconds, and the route field **1014** holds the IP address for the communications link that was established.

Referring to Figure 55, an exemplary RC stop call message for the Calgary callee is shown generally at 1020. In this example the caller field 1002 holds the user name 2001 1050 8667 identifying the Vancouver-based caller and the callee field 1004 holds the user name 2001 1050 2222 identifying the Calgary callee. The contents of the call ID field 1006 are FA10 @ 192.168.0.20. The contents of the account start time field 1008 are 2006-12-30 12:12:12 and the contents of the account stop time field are 2006-12-30 12:12:14. The contents of the communication session time field 1012 are 2 to indicate 2 seconds call duration and the contents of the route field are 72.64.39.58.

Referring back to Figure **53**, after having produced an RC call stop message, block **920** directs the processor **202** in Figure **7** to send the RC stop message compiled in the RC call stop message buffer to the RC **16** of Figure **1**. Block **922** directs the call controller **14** to send a "bye" message back to the party that did not terminate the call.

The RC 16 of Figure 1 receives the call stop message and an RC call stop message process is invoked at the RC, the process being shown at 950 in Figures 56A, 56B and 56C. Referring to Figure 56A, the RC stop message process 950 begins with a first block 952 that directs the processor 202 in Figure 7 to determine whether or not 5 the communication session time is less than or equal to the first increment value set by the cost calculation routine shown in Figure 41, specifically blocks 826 or 862 thereof. If this condition is met, then block 954 of Figure 56A directs the RC processor to set a chargeable time variable equal to the first increment value set at block 826 or 862 of Figure 41. If at block 952 of Figure 56A the condition is not met, block 956 directs the RC processor to set a remainder variable equal to the difference between the communication session time and the first increment value mod the second increment value produced at block 826 or 862 of Figure 41. Then, the processor is directed to block 958 of Figure 56A which directs it to determine whether or not the remainder is greater than zero. If so, block 960 directs the RC processor to set the chargeable time variable equal to the difference between the communication session time and the remainder value. If at block 958 the remainder is not greater than zero, block 962 directs the RC processor to set the chargeable time variable equal to the contents of the communication session time from the RC stop message. The processor is then directed to block 964. In addition, after 20 executing block **954** or block **960**, the processor is directed to block **964**.

Block 964 directs the processor 202 of Figure 7 to determine whether or not the chargeable time variable is greater than or equal to the free time balance as determined from the free time field 742 of the subscriber account record shown in 25 Figure 39. If this condition is satisfied, block 966 of Figure 56A directs the processor to set the free time field 742 in the record shown in Figure 39, to zero. If the chargeable time variable is not greater than or equal to the free time balance, block 968 directs the RC processor to set a user cost variable to zero and Block 970 then decrements the free time field 742 of the subscriber account record for the caller by the chargeable time amount determined by block 954, 960 or 962.

If at Block 964 the processor 202 of Figure 7 was directed to Block 966 which causes the free time field (742 of Figure 39) to be set to zero, referring to Figure 56B, Block 972 directs the processor to set a remaining chargeable time variable equal to the difference between the chargeable time and the contents of the free time field (742 of Figure 39). Block 974 then directs the processor to set the user cost variable equal to the product of the remaining chargeable time and the cost per second calculated at Block 750 in Figure 33B. Block 976 then directs the processor to decrement the funds balance field (740) of the subscriber account record shown in Figure 39 by the contents of the user cost variable calculated at Block 974.

After completing Block **976** or after completing Block **970** in Figure **56**A, block **978** of Figure **56**B directs the processor **202** of Figure **7** to calculate a reseller cost variable as the product of the reseller rate as indicated in the mark-up value field **810** of the system operator mark-up table record shown in Figure **45** and the communication session time determined at Block **916** in Figure **53**. Then, Block **980** of Figure **56**B directs the processor to add the reseller cost to the reseller balance field **986** of a reseller account record of the type shown in Figure **57** at **982**.

The reseller account record includes a reseller ID field **984** and the aforementioned reseller balance field **986**. The reseller ID field **984** holds a reseller ID code, and the reseller balance field **986** holds an accumulated balance of charges.

Referring to Figure **58**, a specific reseller accounts record for the reseller "Klondike" is shown generally at **988**. In this record the reseller ID field **984** holds a code representing the reseller "Klondike" and the reseller balance field **986** holds a balance of **\$100.02**. Thus, the contents of the reseller balance field **986** in Figure **58** are incremented by the reseller cost calculated at block **978** of Figure **56**B.

Still referring to Figure 56B, after adding the reseller cost to the reseller balance field as indicated by Block 980, Block 990 directs the processor to 202 of Figure 7 calculate a system operator cost as the product of the system operator cost per second, as set at block 822 in Figure 41, and the communication session time as determined at Block 916 in Figure 53. Block 992 then directs the processor to add the system operator cost value calculated at Block 990 to a system operator accounts table record of the type shown at 994 in Figure 59. This record includes a system operator balance field 996 holding an accumulated charges balance. Referring to Figure 60 in the embodiment described, the system operator balance field 996 may hold the value \$1,000.02 for example, and to this value the system operator cost calculated at Block 990 is added when the processor executes Block 992 of Figure 56B.

Ultimately, the final reseller balance **986** in Figure **58** holds a number representing an amount owed to the reseller by the system operator and the system operator balance **996** of Figure **59** holds a number representing an amount of profit for the system operator.

While specific embodiments have been described and illustrated, such embodiments should be considered illustrative only and not as limiting the accompanying claims.

# EMBODIMENTS IN WHICH AN EXCLUSIVE PROPERTY OR PRIVILEGE IS CLAIMED ARE DEFINED AS FOLLOWS:

A method for routing a communication in a communication system between an
 Internet-connected first participant device associated with a first participant and a second participant device associated with a second participant, the method comprising:

in response to initiation of the communication by the first participant device, receiving, by a controller comprising at least one processor, over an Internet protocol (IP) network a first participant identifier and a second

participant identifier;

causing the at least one processor to access at least one database comprising user profiles using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes;

comparing at least a portion of the second participant identifier, using the at least one processor, with at least one of the plurality of first participant attributes obtained from the user profile for the first participant;

causing the at least one processor to access the at least one database to search for a user profile for the second participant;

classifying the communication, based on the comparing, as a system communication or an external network communication, using the at least one processor;

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when the communication is classified as a system communication, producing a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant, using the at least one processor, wherein the system routing message causes the communication to be established to the second participant device; and

when the communication is classified as an external network communication, producing an external network routing message identifying an Internet address associated with a gateway to an external network, using the at least one processor, wherein the external network routing message causes the communication to the second participant device to be established using the gateway to the external network.

## 15 2. The method of claim 1,

wherein producing the system routing message causes the communication to be established over an Internet protocol (IP) network;

and wherein producing the external network routing message causes a portion of a path taken by the communication to be established over a circuit switched network.

## 3. The method of claim 1 or 2,

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wherein the Internet address associated with the second participant device comprises an IP address or domain name of the communication system node associated with the second participant device, the communication system node being one of a plurality of communication

system nodes each operably configured to provide communications services to a plurality of communication system subscribers.

4. The method of any one of claims 1 to 3, wherein:

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(a) the causing the at least one processor to access the at least one database to search for the user profile for the second participant is based on the comparing at least a portion of the second participant identifier with the at least one of the plurality of first participant attributes; and

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- (b) the classifying the communication is based on the causing the at least one processor to access the at least one database to search for the user profile for the second participant.
- 15 **5**. The method of any one of claims **1** to **4**, further comprising:

accessing the at least one database to locate communication blocking information for the second participant, using the at least one processor; and blocking the communication when the communication blocking information identifies the first participant identifier.

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6. The method of any one of claims 1 to 5, further comprising:

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accessing the at least one database to locate communication forwarding information for the second participant, using the at least one processor; and

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wherein classifying the communication is based on the communication forwarding information for the second participant.

- 7. The method of claim **6**, wherein the communication forwarding information for the second participant comprises a plurality of destination identifiers.
- 8. The method of any one of claims 1 to 7, further comprising:

5

processing a plurality of communications from the first participant device to a plurality of communication recipient devices to classify each of the plurality of communications as a system communication or an external network communication, wherein the plurality of communications are concurrent; and

10

producing a respective plurality of routing messages, based on the classifying of each respective one of the plurality of communications, each respective routing message identifying an Internet address associated with a recipient device or identifying an Internet address associated with a gateway to an external network.

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**9**. The method of any one of claims **1** to **8**, further comprising updating the first participant profile, via the at least one processor, to cause at least one of the plurality of first participant attributes to be modified.

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**10**. The method of any one of claims **1** to **9**, wherein the external network routing message comprises a code identifying a communication supplier associated with the gateway to the external network.

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11. The method of any one of claims 1 to 5,

wherein the Internet address of the communication system node comprises a first Internet address and wherein the second participant

device uses a second Internet address distinct from the first Internet address;

wherein the user profile for the second participant is associated with communication forwarding information identifying at least one other destination device for the communication; and

wherein the at least one processor is further operably configured to use the communication forwarding information associated with the second participant to cause the communication system node to establish the communication to the at least one other destination device using a third Internet address distinct from the first and second Internet addresses.

**12**. The method of any one of claims **1** to **11**, further comprising causing the at least one processor to:

determine whether the second participant device is operably configured to communicate via the Internet; and

- if the second participant device is not operably configured to communicate via the Internet, classify the communication as the external network communication.
- **13**. The method of any one of claims **1** to **12**, further comprising causing the at least one processor to:

determine whether the user profile for the second participant exists in the at least one database; and

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if the user profile for the second participant does not exist in the at least one database, classify the communication as the external network communication.

5 **14**. The method of any one of claims **1** to **13**, further comprising causing the at least one processor to:

if the user profile for the second participant exists in the at least one database, classify the communication as the system communication.

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- **15**. The method of any one of claims **1** − **14**, wherein the communication comprises a voice-over-IP (VoIP) communication.
- 16. The method of any one of claims 1 14, wherein the communication comprises
  at least one of video data traffic and message traffic.
  - 17. The method of any one of claims 1 16, wherein at least one of the system routing message and the external network routing message is communicated to a call controller to effect routing of the communication.

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- **18**. A non-transitory computer readable medium encoded with program code for directing the at least one processor to execute the method of any one of claims **1** to **17**.
- 25 19. An apparatus for routing communications in a communication system that includes an Internet-connected first participant device associated with a first participant, the first participant device operable to initiate a communication to a second participant device associated with a second participant, the apparatus comprising:

a controller comprising at least one processor in communication with at least one memory storing processor readable instructions, wherein the at least one processor is operably configured by the processor readable instructions to:

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in response to initiation of the communication by the first participant device, receive over an Internet protocol (IP) network a first participant identifier and a second participant identifier;

10

access at least one database comprising user profiles using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes;

15

compare at least a portion of the second participant identifier with at least one of the plurality of first participant attributes obtained from the user profile for the first participant to generate a comparison result;

20

access the at least one database to search for a user profile for the second participant;

25

classify the communication, based on the comparison result, as a system communication or an external network communication;

30

when the communication is classified as a system communication, produce a system routing message identifying an Internet address of a communication system node associated with the second participant device based on the user profile for the second participant, wherein the system routing message causes the

communication to be established to the second participant device; and

5

when the communication is classified as an external network communication, produce an external network routing message identifying an Internet address associated with a gateway to an external network, wherein the external network routing message causes the communication to the second participant device to be established using the gateway to the external network.

10

20. The apparatus of claim 19,

wherein the at least one processor is operably configured to cause the communication to be established over an Internet protocol (IP) network when the communication is classified as a system communication; and

15

wherein the at least one processor is operably configured to cause a portion of a path taken by the communication to be established over a circuit switched network when the communication is classified as an external network communication.

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21. The apparatus of claim 19 or 20, wherein the at least one processor is further operably configured to:

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access the at least one database to locate communication blocking information associated with the second participant; and

block the communication when the communication blocking information identifies the first participant identifier.

22. The apparatus of any one of claims 19 to 20, wherein the at least one processor is further operably configured to:

access the at least one database to locate communication forwarding information, associated with the second participant, comprising a plurality of destination identifiers; and

use the communication forwarding information associated with the second participant to classify the communication.

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23. The apparatus of any one of claims 19 to 22, wherein the at least one processor is further operably configured to:

process a plurality of communications from the first participant device to a plurality of communication recipient devices to classify each of the plurality of communications as a system communication or an external network communication, wherein the plurality of communications are concurrent; and

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produce a respective plurality of routing messages, based on the classifying of each respective one of the plurality of communications, each respective routing message identifying an Internet address associated with a recipient device or identifying an Internet address associated with a gateway to an external network.

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24. An apparatus for routing communications in a communication system that includes an Internet-connected first participant device associated with a first participant, the first participant device operable to initiate a communication to a second participant device associated with a second participant, the apparatus comprising:

a controller comprising at least one processor in communication with at least one memory storing processor-readable instruction codes, wherein the at least one processor is operably configured by the processor-readable instruction codes to:

5

in response to initiation of the communication by the first participant device over the Internet, receive a first participant identifier and a second participant identifier;

10

access at least one database comprising user profiles using the first participant identifier, each user profile comprising a respective plurality of attributes for a respective user, to locate a user profile for the first participant including a plurality of first participant attributes;

15

compare at least a portion of the second participant identifier with at least one of the plurality of first participant attributes obtained from the user profile for the first participant to generate a comparison result;

20

access the at least one database to search for a user profile for the second participant and communication blocking information for the second participant;

25

classify the communication, based on at least one of the comparison result and the communication blocking information for the second participant, as a system communication, an external network communication or a blocked communication;

30

when the communication is classified as a system communication, produce a system routing message identifying a first Internet address

associated with the second participant device, causing the communication to be established entirely over an Internet protocol (IP) network; and

when the communication is classified as an external network communication, produce an external routing message identifying an Internet address associated with a gateway to a network that is external to the communication system, causing a portion of a path taken by the communication to be established over a circuit switched network.

### 10 **25**. The apparatus of claim **24**,

wherein the first Internet address identifies a first communication system node operably configured to establish the communication to the second participant device, wherein the second participant device uses a second Internet address distinct from the first Internet address;

wherein the user profile for the second participant is associated with communication forwarding information identifying at least one other destination device for the communication; and

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wherein the at least one processor is further operably configured to use the communication forwarding information for the second participant to cause the first communication system node to establish the communication to the at least one other destination device using a third Internet address distinct from the first and second Internet addresses.

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**26**. The apparatus of claim **24** or **25**, wherein the at least one processor is further operably configured to update the first participant profile to cause at least one of the plurality of first participant attributes to be modified.

- 27. The apparatus of any one of claims 24 to 26, wherein the first Internet address associated with the second participant device comprises an IP address or domain name of a communication system node associated with the second participant device, the communication system node being one of a plurality of communication system nodes each operably configured to provide communications services to a plurality of communication system subscribers.
- 28. The apparatus of any one of claims 24 to 27, wherein the external routing message comprises a code identifying a communication supplier associated
  with the gateway to the network that is external to the communication system.
  - 29. The apparatus of any one of claims 24 to 28, wherein the at least one processor is further operably configured to:
- access the at least one database to locate communication forwarding information, associated with the second participant, comprising a plurality of destination identifiers; and
- use the communication forwarding information associated with the second participant to classify the communication.
  - **30**. The apparatus of any one of claims **24** to **29**, wherein the at least one processor is further operably configured to:
- determine whether the second participant device is operably configured to communicate via the Internet; and
  - if the second participant device is not operably configured to communicate via the Internet, classify the communication as the external network communication.

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**31**. The apparatus of any one of claims **24** to **30**, wherein the at least one processor is further operably configured to:

determine whether the user profile for the second participant exists in the at least one database; and

if the user profile for the second participant does not exist in the at least one database, classify the communication as the external network communication.

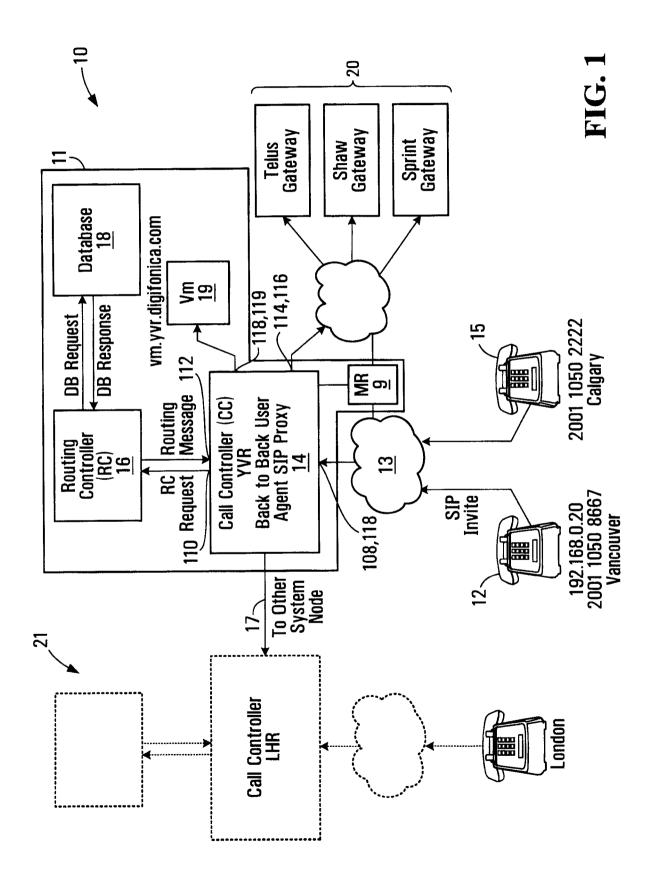
**32**. The apparatus of any one of claims **24** to **31**, wherein the at least one processor is further operably configured to:

if the user profile for the second participant exists in the at least one database, classify the communication as the system communication, wherein the first Internet address associated with the second participant device is determined based on the user profile for the second participant.

20 33. The apparatus of any one of claims 24 to 32, wherein producing the external routing message comprises causing the at least one processor to select the Internet address associated with the gateway from among a plurality of Internet addresses associated with a respective plurality of gateways to the network that is external to the communication system, and wherein the external routing message comprises a code identifying a communication supplier associated with the gateway.

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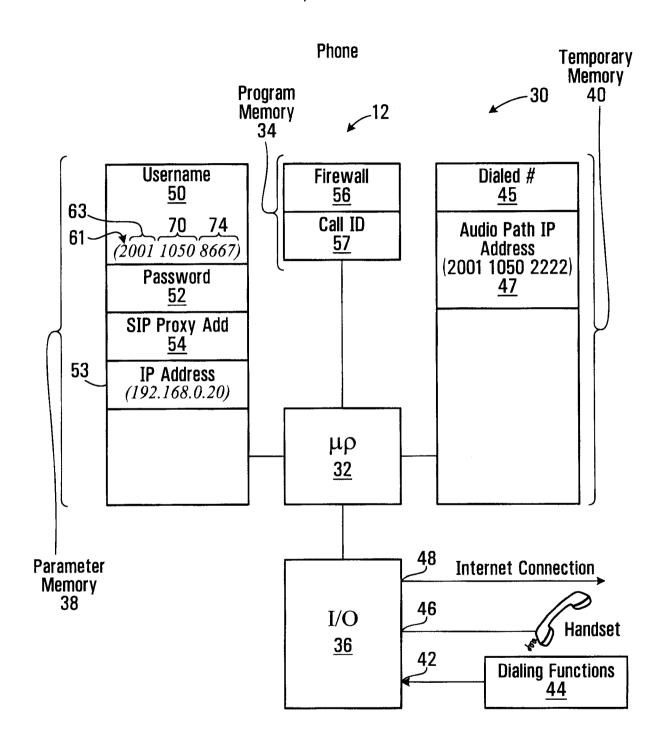
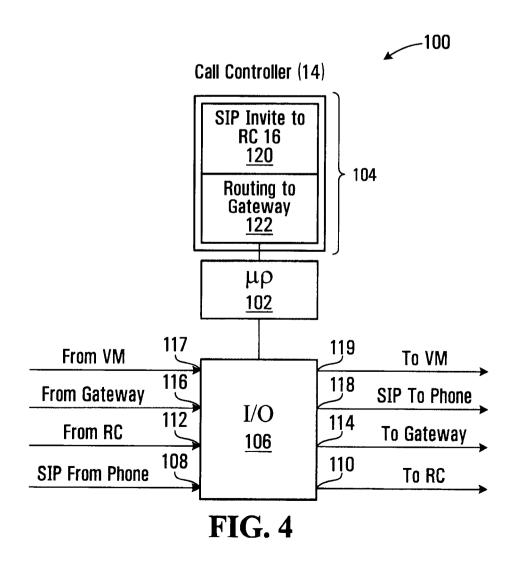


FIG. 2

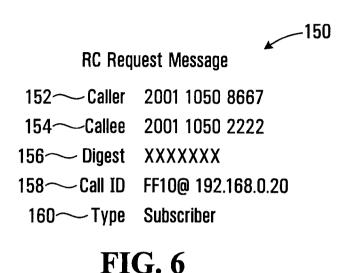
#### SIP Invite Message

60 Caller 2001 1050 8667
62 Callee 2001 1050 2222
64 Digest Parameters XXXXXXX
65 Call ID FF10@ 192.168.0.20
67 IP Address 192.168.0.20
69 Caller UDP Port 1



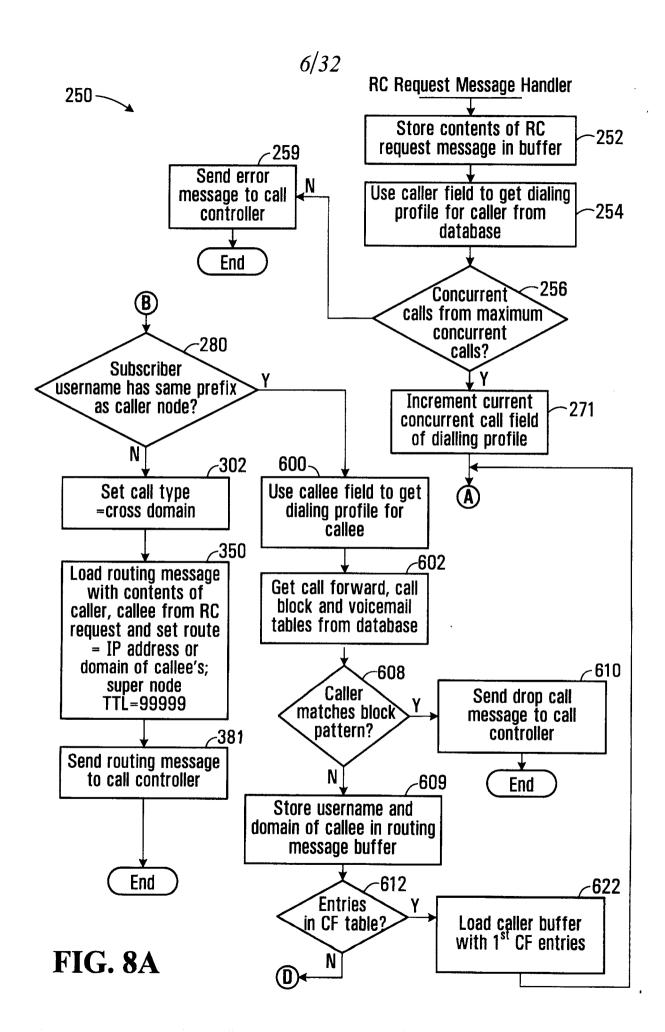
# **Call Controller Process** 120 **Receive SIP Invite** Message Authenticate **Error Handling** 122 124 Caller ID N Field Contents= IP Address? 121 Set Type = 3<sup>rd</sup> Party Invite Set Type = Regular Invite 125 123 **Establish Call ID** 126 Prepare RC Request Message 128 Send RC Request Message 129

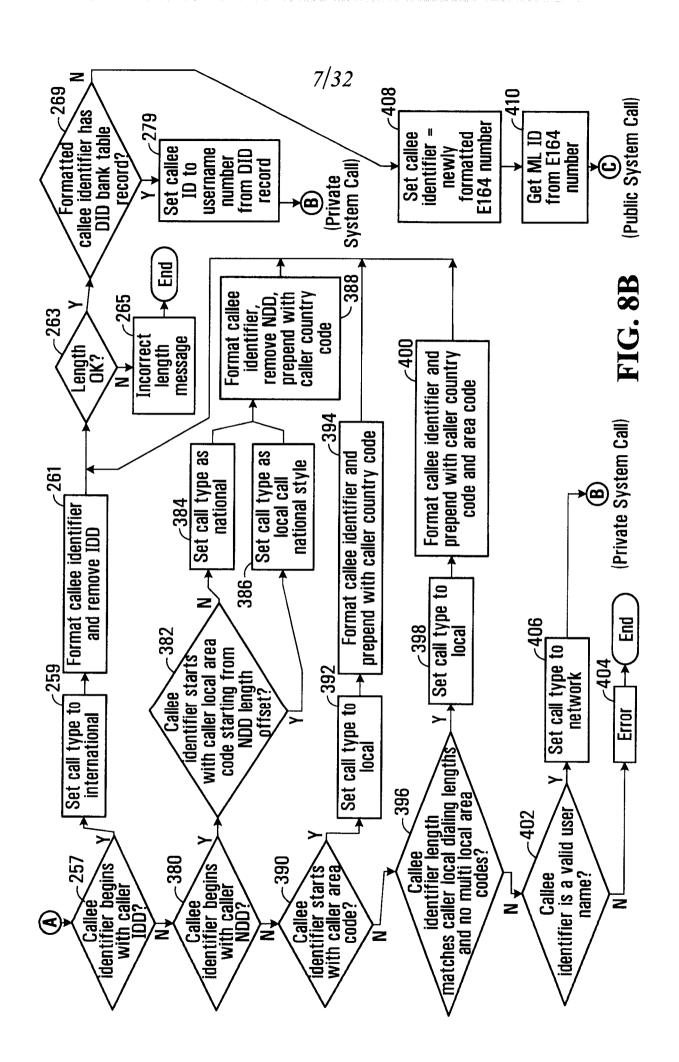
FIG. 5



**Routing Controller (16)** 207 **-200** 209 Caller ID **Program** 211 **RC** Request Callee ID Memory Message 204 Handler Retrieved 250 **Dialing Profile** 206 Call Type Flag Username To Domain Table **Subscriber** μρ **Bundle Record** 202 **Bundle Override** Record 210 **DB** Response **DB** Request I/O 214 216 **RC** Request 208 **Routing Message** 

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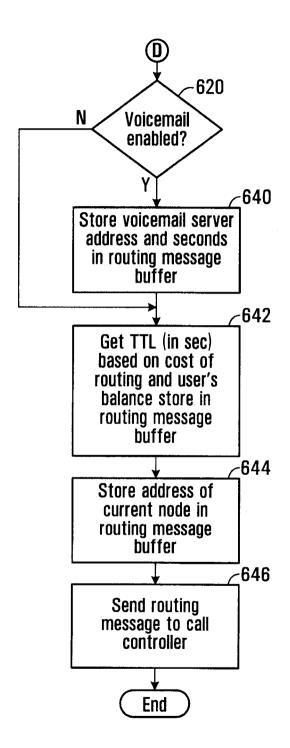


FIG. 8C



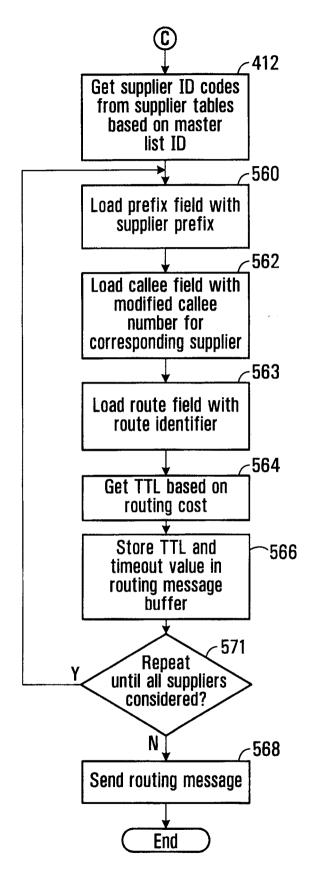
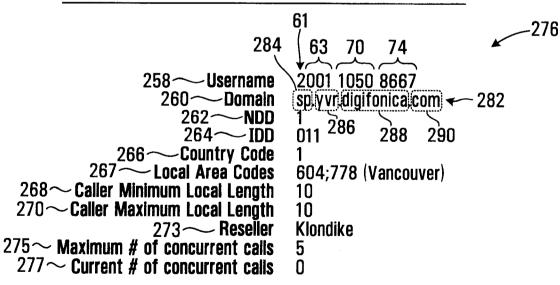


FIG. 8D

253 Dialing Profile for a User 258 — Username Assigned on Subscription 260 — Domain Domain Associated with User 262~NDD 264~IDD 011 266 Country Code 267 Local Area Codes 604:778 268 ~ Caller Minimum Local Length 10 270 ∼ Caller Maximum Local Length 10 273 ~ Reseller Retailer 275 ~ Maximum # of concurrent calls **Assigned on Subscription**  $277 \sim$  Current # of concurrent calls **Assigned on Subscription** 

# FIG. 9

### Dialing Profile for Caller (Vancouver Subscriber)



#### Callee Profile for Calgary Subscriber

Username 2001 1050 2222 Domain sp.yvr.digifonica.com NDD IDD 011 **Country Code** Local Area Codes 403 (Calgary) **Caller Minimum Local Length Caller Maximum Local Length** 10 Deerfoot Reseller Maximum # of concurrent calls Current # of concurrent calls U

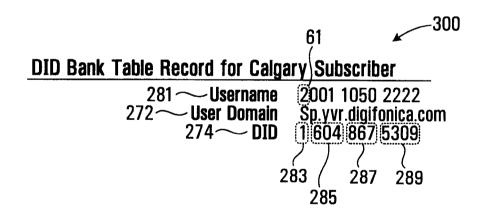
# **FIG. 11**

#### Callee Profile for London Subscriber

Username 4401 1062 4444 Domain sp.lhr.digifonica.com NDD IDD 00 **Country Code** 44 **Local Area Codes** 20 (London) **Caller Minimum Local Length** 10 **Caller Maximum Local Length** 11 Reseller Marble Arch Maximum # of concurrent calls Current # of concurrent calls 0

#### **DID Bank Table Record Format**

281 Username System subscriber
272 User Domain Host name of supernode E164#



**FIG. 14** 

#### **Routing Message Format**

354 Supplier Prefix (optional)
356 Delimiter (optional)
358 Callee
360 Route
362 Time to Live(TTL)

Code identifying supplier traffic
Symbol separating fields
PSTN compatible number or Digifonica number
Domain name or IP address
In seconds

364 Other TBD

# **FIG. 15**

Example of Routing Message - Different Node

440110624444@sp.lhr.digifonica.com;ttl=9999

359 361 363

### **FIG. 16**

370

### **Prefix to Supernode Table Record Format**

372 Prefix First n digits of callee identifier
374 Supernode Address IP address or fully qualified domain name

### **FIG. 17**

### Prefix to Supernode Table Record for Calgary Subscriber

Prefix 20 Supernode Address sp.yvr.digifonica.com

# **Master List Record Format**

500 ml_id 502 Dialing code 504 Country code	Alphanumeric Number Sequence The country code is the national prefix to be used when dialing TO a particular country FROM another country.
506 Nat Sign #(Area Code) 508 Min Length 510 Max Length 512 NDD	Number Sequence Numeric Numeric The NDD prefix is the access code used to make a call WITHIN that country from one city to another (when calling another city in the same vicinity, this may not be necessary).
514 ~ IDD	The IDD prefix is the international prefix needed to dial a call FROM the country listed TO another country.
516 — Buffer rate	Safe change rate above the highest rate charged by suppliers

# FIG. 19

# **Example: Master List Record with Populated Fields**

ml_id Dialing code Country code	1019 1604
Nat Sign #(Area Code)	604
Min Length	7
Max Length	7
NDD	1
IDD	011
Buffer rate	\$0.009/min

# **Suppliers List Record Format**

540~ Sup_id	Name code
542~ Ml_id	Numeric code
544~ Prefix (optional)	String identifying supplier's traffic #
546~ Specific Route	IP address
548~ NDD/IDD rewrite 550~ Rate 551~ Timeout	Cost per second to Digifonica to use this route Maximum time to wait for a response when requesting this gateway

# FIG. 21

# **Telus Supplier Record**

	<del></del>	
Sup_id	2010 (Telus)	
MI_id	1019	
Prefix (optional)	4973#	
Specific Route	72.64.39.58	
NDD/IDD rewrite	011	
Rate	\$0.02/min	
Timeout	20	
	FIG 22	

# **Shaw Supplier Record**

Sup_id	2011 (Shaw)	
Ml <sup>'</sup> id	1019	
Prefix (optional)	4974#	
Specific Route	73.65.40.59	
NDD/IDD rewrite	011	
Rate	\$0.025/min	
Timeout	30	
	FIG. 23	

# **Sprint Supplier Record**

	the state of the s	
Sup_id	2012 (Sprint)	
MI_id	1019	
Prefix (optional)	4975#	
Specific Route	74.66.41.60	
NDD/IDD rewrite	011	
Rate	\$0.03/min	
Timeout	40	
	TITO A	

#### **Routing Message Buffer for Gateway Call**

4973#0116048675309@72.64.39.58;ttl=3600;to=20 570 4974#0116048675309@73.65.40.59;ttl=3600;to=30 572 4975#0116048675309@74.66.41.60;ttl=3600;to=40 574

## **FIG. 25**

#### **Call Block Table Record Format**

604 Username Digifonica # PSTN compatible or Digifonica #

# **FIG. 26**

#### **Call Block Table Record for Calgary Callee**

604 Username of Callee 2001 1050 2222 606 Block Pattern 2001 1050 8664

### **FIG. 27**

#### **Call Forwarding Table Record Format for Callee**

614 Username of Callee Digifonica #
616 Destination Number Digifonica #
618 Sequence Number Integer indicating order to try this

### **FIG. 28**

## **Call Forwarding Table Record for Calgary Callee**

614 Username of Callee 2001 1050 2222 616 Destination Number 2001 1055 2223 618 Sequence Number 1

#### **Voicemail Table Record Format**

624 Username of Callee Digifonica # domain name

628 Seconds to Volcemall time to wait before engaging voicemail ves/no

FIG. 30

#### Voicemail Table Record for Calgary Callee

Username of Callee 2001 1050 2222
Vm Server vm.yvr.digifonica.com
Seconds to Voicemail 20
Fnabled 1

**FIG. 31** 

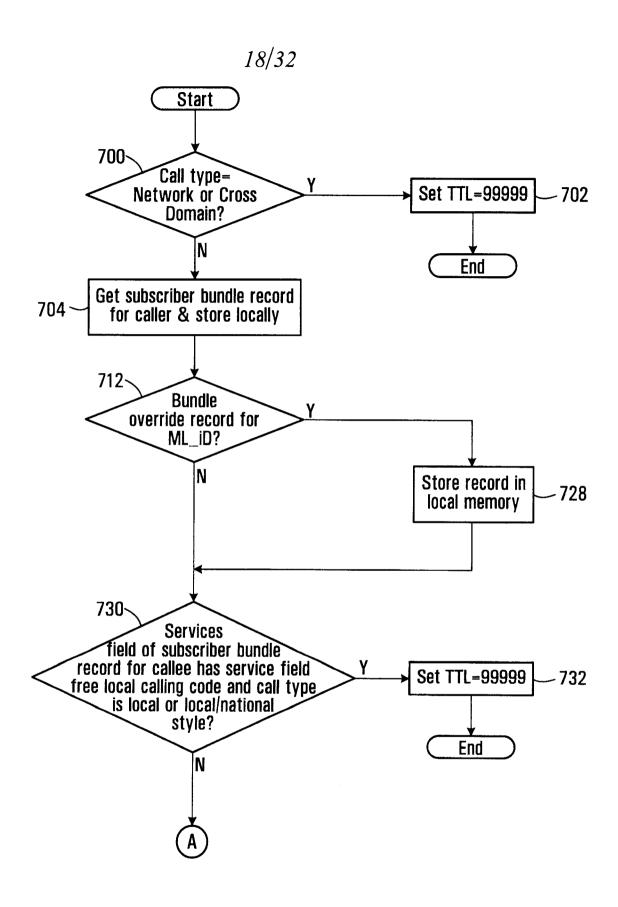
### Routing Message Buffer - Same Node

650 200110502222@sp.yvr.digifonica.com;ttl=3600

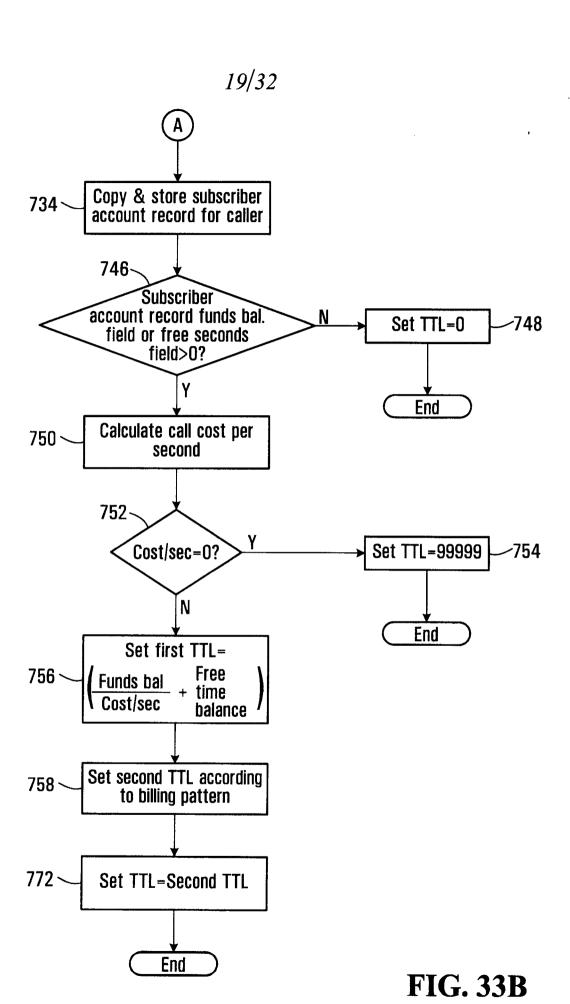
652 200110552223@sp.yvr.digifonica.com;ttl=3600

654 vm.yvr.digifonica.com;20;ttl=60

656 — sp.yvr.digifonica.com



**FIG. 33A** 



**Subscriber Bundle Table Record** 

706

708 Username
710 Services

Subscriber username Codes identifying service features

(e.g. Free local calling; call blocking, voicemail)

**FIG. 34** 

#### **Subscriber Bundle Record for Vancouver Caller**

708 Username 2001 1050 8667 710 Services 10; 14; 16

**FIG. 35** 

#### **Bundle Override Table Record**

714 ح

**726** 

716 ML\_Id
718 Override type

Master list ID code Fixed: percent; cents

720 Override value
722 Inc1

real number representing value of override type first level of charging (minimum # of seconds) charge

724 Inc2 second level of charging

**FIG. 36** 

#### Bundle Override Record for Located ML\_iD

716 ML\_Id 1019

718 Override type percent 720 Override value 10.0

722 **Inc1** 30 seconds 724 **Inc2** 6 seconds

### **Subscriber Account Table Record**

, 736

738 Username
740 Funds balance
742 Free time balance

Subscriber username real number representing \$ value of credit integer representing # of free seconds

**FIG. 38** 

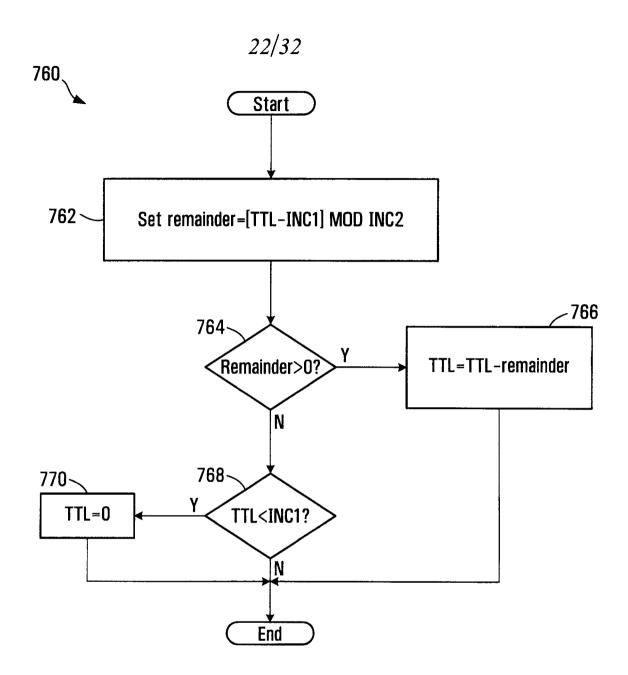
**Subscriber Account Record for Vancouver Caller** 

744

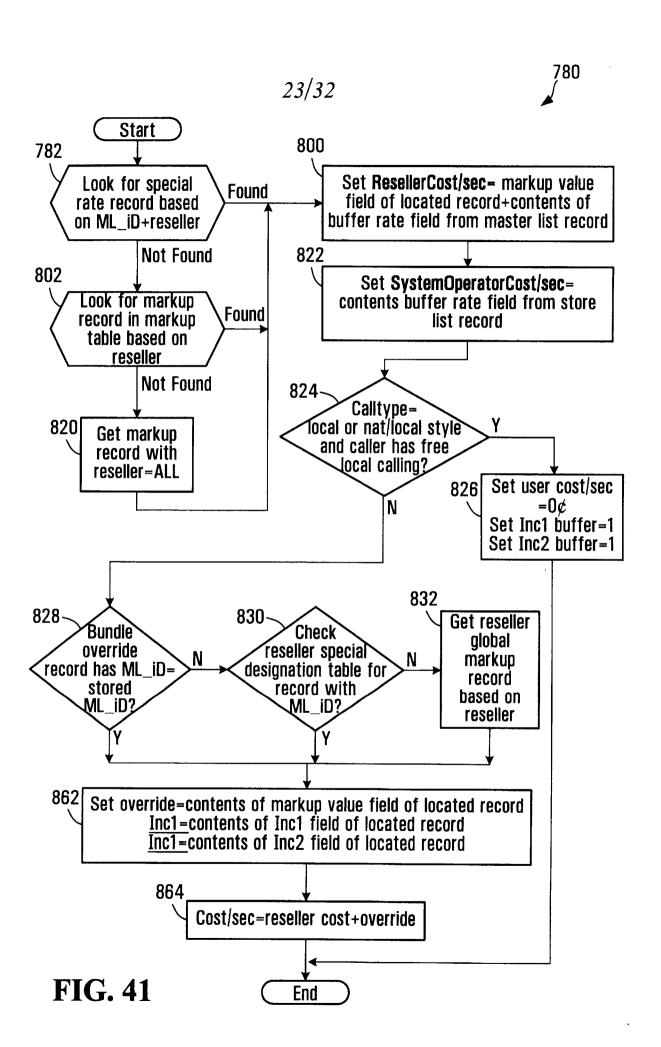
<u>r Vancouver Caller</u>
2001 1050 8667

738 Username
740 Funds balance
742 Free time balance

\$10.00 100



**FIG. 40** 



# System Operator Special Rates Table Record

786 ~ Reseller	retailer id
788~ ML_Id	master list id
790 <b>Markup Table</b>	fixed; percent; cents
792 <b>Markup Value</b>	real number representing value of markup type
794~ <b>Inc1</b>	first level of charging (minimum # of seconds) charge
796 ~ Inc2	second level of charging

# FIG. 42

<sub>\_\_\_</sub>798

# **System Operator Special Rates Table Record for Klondike**

786 Reseller	Klondike
788~ ML_Id	1019
790 <b>Markup Table</b>	cents
792 Markup Value	\$0.001
794~ Inc1	30
796~ Inc2	6

# System Operator Markup Table Record

804

806 ~ Reseller	reseller id code
808 <b>Markup Table</b>	fixed; percent; cents
810 <b>Markup Value</b>	real number representing value of markup type
812 <b>Inc1</b>	first level of charging (minimum # of seconds) charge
814~ Inc2	second level of charging

# **FIG. 44**

# System Operator Markup Table Record for the Reseller Klondike

806 ~ Reseller	Klondike
808 Markup Table	cents
810 Markup Value	\$0.01
812 <b>Inc1</b>	30
814~ Inc2	6

# FIG. 45

# **System Operator Markup Table Record**

806 ~ Reseller	all
808 <b>Markup Table</b>	percent
810 <b>Markup Value</b>	1.0
812 <b>Inc1</b>	30
814 ~ Inc2	6

2	6/	3	2
_	V/	J	_

Reseller	Snecial	<b>Destinations</b>	Table	Record

834 Reseller reseller id code 836 ML\_id Master List ID code 838 Markup Table fixed; percent; cents

840 Markup Value real number representing value of markup type 842 Inc1 first level of charging (minimum # of seconds) charge

844 Inc2 second level of charging

### **FIG. 47**

, 846

832

### Reseller Special Destinations Table Record for the Reseller Klondike

834 Reseller Klondike
836 ML\_id 1019
838 Markup Table percent
840 Markup Value 5%
842 Inc1 30
844 Inc2 6

### **FIG. 48**

### Reseller Global Markup Table Record

, 848

850 Reseller reseller id code
852 Markup Table fixed; percent; cents
854 Markup Value real number representing value of markup type
856 Inc1 first level of charging (minimum # of seconds) charge second level of charging

### **FIG. 49**

860

## Reseller Global Markup Table Record for the Reseller Klondike

850 Reseller Klondike 852 Markup Table percent 854 Markup Value 10% 856 Inc1 30 858 Inc2 6

900

### SIP Bye Message

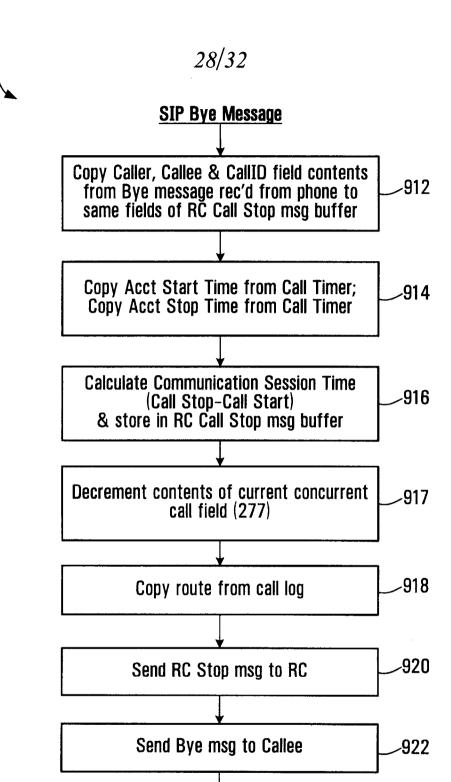
902 Caller Username
904 Callee PSTN compatible # or Username
906 Call ID unique call identifier (hexadecimal string@IP))

# **FIG. 51**

908

## SIP Bye Message

902 Caller 2001 1050 8667 904 Callee 2001 1050 2222 906 Call ID FA10@192.168.0.20



**FIG. 53** 

End

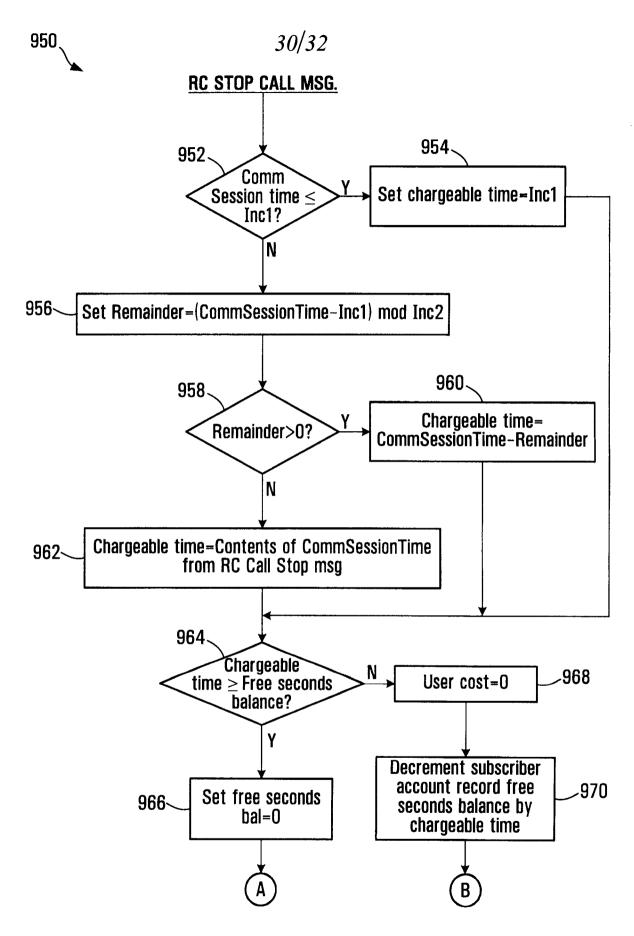
910

# RC Call Stop Message

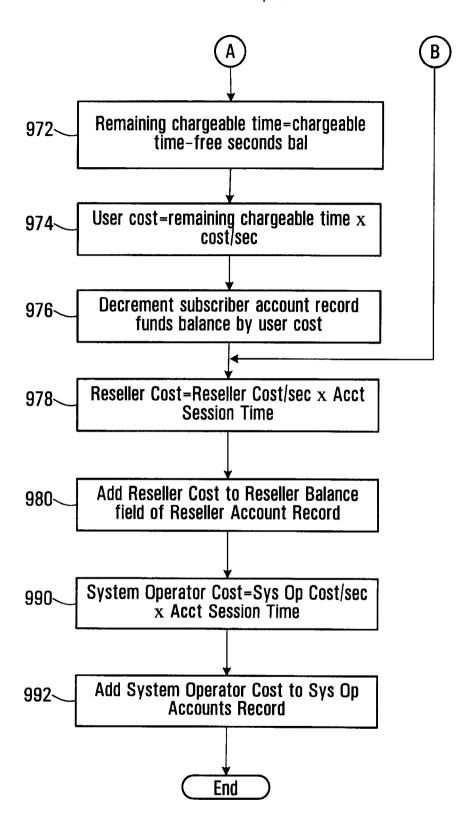
I string@IP) ns link that

# **FIG. 54**

RC Call Stop Message for Ca	lgary Callee	1020
1006 Call ID 1008 Acct Start Time 1010 Acct Stop Time 1012 Acct Session Time	2001 1050 8667 2001 1050 2222 FA10@192.168.0.20 2006-12-30 12:12:12 2006-12-30 12:12:14 2 72.64.39.58	



**FIG. 56A** 



**FIG. 56B** 

**Reseller Accounts Table Record** 

<sub>\_\_\_</sub>982

984 ~ Reseller ID 986 ~ Reseller balance

reseller id code

accumulated balance of charges

**FIG. 57** 

Reseller Accounts Table Record for Klondike

988

984 Reseller ID Klondike 986 Reseller balance \$100.02

**FIG. 58** 

**System Operator Accounts Table Record** 

994

996 System Operator balance accumulated balance of charges

**FIG. 59** 

System Operator Accounts Record for this System Operator

996 System Operator balance \$1000.02

