



(51) International Patent Classification:

H04L 1/00 (2006.01) H04L 29/06 (2006.01)
H03M 13/37 (2006.01) H04L 9/08 (2006.01)

(21) International Application Number:

PCT/EP2018/073650

(22) International Filing Date:

04 September 2018 (04.09.2018)

(25) Filing Language:

English

(26) Publication Language:

English

(30) Priority Data:

17190779.3 13 September 2017 (13.09.2017) EP

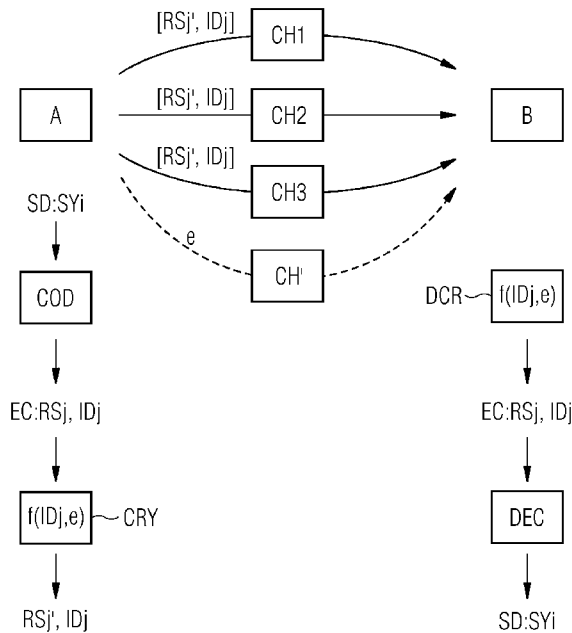
(71) Applicant: SIEMENS AKTIENGESELLSCHAFT
[DE/DE]; Werner-von-Siemens-Straße 1, 80333 München (DE).

(72) Inventors: CUELLAR, Jorge; Höllriegelskreuther Weg 14, 82065 Baierbrunn (DE). GASIBA, Tiago; Chiemgaustraße 94, 81549 München (DE).

(81) Designated States (unless otherwise indicated, for every kind of national protection available): AE, AG, AL, AM, AO, AT, AU, AZ, BA, BB, BG, BH, BN, BR, BW, BY, BZ, CA, CH, CL, CN, CO, CR, CU, CZ, DE, DJ, DK, DM, DO, DZ, EC, EE, EG, ES, FI, GB, GD, GE, GH, GM, GT, HN, HR, HU, ID, IL, IN, IR, IS, JO, JP, KE, KG, KH, KN, KP, KR, KW, KZ, LA, LC, LK, LR, LS, LU, LY, MA, MD, ME, MG, MK, MN, MW, MX, MY, MZ, NA, NG, NI, NO, NZ, OM, PA, PE, PG, PH, PL, PT, QA, RO, RS, RU, RW, SA, SC, SD, SE, SG, SK, SL, SM, ST, SV, SY, TH, TJ, TM, TN, TR, TT, TZ, UA, UG, US, UZ, VC, VN, ZA, ZM, ZW.

(84) Designated States (unless otherwise indicated, for every kind of regional protection available): ARIPO (BW, GH, GM, KE, LR, LS, MW, MZ, NA, RW, SD, SL, ST, SZ, TZ,

(54) Title: A METHOD FOR SENDING DIGITAL DATA OVER A NUMBER OF CHANNELS



(57) Abstract: The invention refers to a method for sending digital data over a number of channels (CH1, CH2, CH3) wherein a sender (A) performs the following steps: - encoding (COD) source data (SD) having a first number of source symbols (SYi), the encoding (COD) being such that an error correction code (EC) is generated from the source data (SD), the error correction code (EC) comprising a second number of repair symbols (RSj) higher than the first number as well as identifiers (IDj) where each identifier (IDj) is assigned to a corresponding repair symbol (RSj), the error correction code (EC) adding redundancy to the source data (SD); - encrypting each repair symbol (RSj) by an encryption process (CRY) which is based on a shared secret (e) between the sender (A) and a receiver (B), where the encryption process (CRY) for a respective repair symbol (RSj) depends on the identifier (IDj) assigned to the respective repair symbol (RSj); - feeding pairs of the encrypted repair symbols (RSj') and the assigned identifiers (IDj) to the number of channels



UG, ZM, ZW), Eurasian (AM, AZ, BY, KG, KZ, RU, TJ, TM), European (AL, AT, BE, BG, CH, CY, CZ, DE, DK, EE, ES, FI, FR, GB, GR, HR, HU, IE, IS, IT, LT, LU, LV, MC, MK, MT, NL, NO, PL, PT, RO, RS, SE, SI, SK, SM, TR), OAPI (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, KM, ML, MR, NE, SN, TD, TG).

Published:

— *with international search report (Art. 21(3))*

Description

A method for sending digital data over a number of channels

5 The invention refers to a method for sending digital data of a number of channels. Furthermore, the invention refers to a method for processing digital data sent by the method for sending digital data. Moreover, the invention refers to a method for transmitting digital data from a sender to a re-
10 ceiver as well as to a corresponding sender and receiver. Moreover, the invention refers to a transmission system comprising a sender and a receiver.

When transmitting digital data from a sender to a receiver,
15 data losses may occur due to transmission errors. In order to reconstruct the original data even in case of data losses, so-called error correction techniques, e.g. forward error correction, are used which encode the digital data to obtain an error correction code in which redundancy is added to the
20 data.

Although error correction codes enable a reconstruction of the original source information in case of data losses, those codes do not provide cryptographic protection against attack-
25 ers listening to the data transmission. To provide such a cryptographic protection, it is known from the prior art to encrypt encoded symbol identifiers transmitted together with corresponding repair symbols of the error correction code. However, this encryption is not very strong.

30 It is an object of the invention to provide a method for sending digital data over a number of channels enabling both an error correction as well as a strong encryption of the digital data being sent. Furthermore, it is an object of the
35 invention to provide a method to reconstruct source data from the digital data sent by the method of the invention.

This object is solved by the independent patent claims. Preferred embodiments of the invention are described in the dependent claims.

5 The method of the invention provides a mechanism for sending digital data over a number of channels, i.e. over at least one channel. The term "channel" is to be interpreted broadly and may refer to any continuous transmission path from a sender to a receiver. I.e., a channel may also have different
10 channel portions being based on different transmission techniques. In the method of the invention, a sender performs the steps as described in the following.

The sender encodes source data having a first number of
15 source symbols, where the encoding is such that an error correction code is generated from the source data, the error correction code comprising a second number of repair symbols higher than the first number as well as identifiers where each identifier is assigned to a corresponding repair symbol.
20 This error correction code adds redundancy to the source data. Any known technique for generating such an error correction code may be used in this step. Particularly, the error correction code may be a FEC code (FEC = Forward Error Correction). In a preferred embodiment, fountain coding is used,
25 i.e. the error correction code is a fountain code.

After having generated the error correction code, the sender encrypts each repair symbol of this code by an encryption process which is based on a shared secret between the sender
30 and a receiver by which the encrypted repair symbols shall be received, where the encryption process for a respective repair symbol depends on the identifier assigned to the respective repair symbol. The term "shared secret" refers to a secret value which shall only be known by the sender and the
35 receiver.

In a next step, the sender feeds pairs of the encrypted repair symbols and the assigned identifiers to the number of

channels. I.e., corresponding pairs, each comprising an encrypted repair symbol and the assigned identifier, are fed to the number of channels.

5 The method of the invention has the advantage that a strong encryption of the error correction code is provided by encrypting the repair symbols not only based on a shared secret, but also in dependency of the corresponding identifiers.

10

In a preferred embodiment of the invention, the source symbols and the repair symbols each are bit sequences having the same fixed bit length.

15 In another preferred variant of the invention, the encryption process uses a pseudo-random function, i.e. a function of a pseudo-random generator. The terms "pseudo-random function" as well as "pseudo-random generator" are well-known for a skilled person. The pseudo-random function is applied to the
20 repair symbols, where the pseudo-random function depends on the shared secret and on the identifier of the repair symbol to be encrypted. This embodiment provides an efficient realization of an encryption process depending on both the shared secret and the identifier of the repair symbol.

25

In a preferred implementation of the above variant, the pseudo-random function generates an intermediate bit sequence having the bit length of the repair symbol to be encrypted, the encrypted repair symbol being the XOR combination of the
30 intermediate bit sequence with the repair symbol to be encrypted.

In another, particularly preferred embodiment, the number of channels comprises several independent channels where the
35 pairs of the encrypted repair symbols and the assigned identifiers are fed to the independent channels such that each independent channel receives at least one pair of the encrypted repair symbols and the assigned identifiers. The term

"independent channels" means that an access to one channel does not permit an access to any other independent channel. The use of independent channels strongly enhances the security of the data transmission as information from several channels is needed in order to reconstruct the whole data.

In a preferred variant of the above embodiment, the independent channels are channels of a multipath TCP protocol. This kind of protocol is known from the prior art. However, the independent channels may also be defined differently, e.g. by different transmission technologies.

In a particularly preferred embodiment, the number of pairs of the encrypted repair symbols and the assigned identifiers in each independent channel is less than the first number referring to the number of source symbols. According to this embodiment, source symbols cannot be reconstructed by repair symbols in a single channel even in case that an attacker succeeds to decrypt the repair symbols.

In another variant of the method according to the invention, the encryption process further depends on the channel to which the encrypted repair symbol is fed, resulting in an enhanced security against attacks.

The above described shared secret used for encryption may be notified to the receiver before sending data based on the method of the invention. However, in a preferred embodiment of the invention, parallel to the feeding of the pairs of the encrypted repair symbols and the assigned identifiers to the number of channels, the shared secret is fed by the sender to an additional channel being connected to the receiver and being independent of the channels of the number of channels. Hence, the transmission of the shared secret needs not be performed in a separate step but may be part of the method for sending data according to the invention.

Besides the above method for sending digital data, the invention also refers to a method for processing digital data being sent by the method of the invention or by one or more preferred embodiments of the method according to the invention. To do so, a receiver performs the following steps:

- receiving the pairs of the encrypted repair symbols and the assigned identifiers over the number of channels;
- decrypting each encrypted repair symbol by a decryption process which is based on the shared secret between the sender and the receiver, resulting in the error correction code, where the decryption process for a respective encrypted repair symbol depends on the identifier assigned to the respective repair symbol;
- decoding the error correction code such that the source data are restored.

In case that the method for sending digital data uses an encryption process based on a pseudo-random function, the same pseudo-random function is used in the above method for processing digital data in order to decrypt the encrypted repair symbols. Particularly, in case that the pseudo-random function generates an intermediate bit sequence, the decryption will be such that a corresponding decrypted repair symbol is the XOR combination of the intermediate bit sequence with the encrypted repair symbol.

The invention also refers to a method for transmitting digital data from a sender to a receiver over a number of channels, where the sender performs a method for sending digital data according to the invention or one or more preferred embodiments thereof and where the receiver performs a method for processing digital data according to the invention or one or more preferred embodiments thereof.

Moreover, the invention refers to a sender for sending digital data, wherein the sender comprises means for performing the method for sending digital data according to the invention or one or more preferred embodiments thereof.

Furthermore, the invention refers to a receiver for processing digital data wherein the receiver comprises means for performing the method for processing digital data according to the invention or one or more preferred embodiments thereof.

Additionally, the invention refers to a transmission system comprising the above sender as well as the above receiver.

In the following, an embodiment of the invention will be described in detail with respect to the accompanying Fig. 1. This figure is a schematic diagram illustrating the transmission of digital data based on a variant of the invention.

Fig. 1 shows a scenario where digital data are transmitted between a sender A and a receiver B using three channels CH1, CH2 and CH3. Additionally to those three channels, a further channel CH' is used for transmitting a shared secret e , as will be explained further on.

As mentioned above, the term "channel" is to be interpreted broadly and refers to a continuous transmission path from sender A to sender B where the path can have portions with different transmission technologies. The channels CH1, CH2, CH3 and CH' are independent from each other. I.e., an attacker listening to one channel only receives information transmitted in this channel and not from the other channels.

In one variant of the invention, the channels CH1, CH2 and CH3 may refer to different UDP Internet streams or to different channels of a multipath TCP protocol. Contrary to that, the channel CH' may refer to an SMS channel. In another setting, where only two channels are used between sender A and receiver B, one channel may use an intermediate node, like a proxy to which the communication is secured either from sender A or from receiver B. Thus, sender A can communicate with a remote proxy over a secure channel portion which in turn

communicates with receiver B over a non-secure channel portion. For the other channel, the situation could be opposite. In this channel, sender A communicates with a proxy known to receiver B over a non-secure channel portion but this proxy
5 communicates with the receiver B over a secure tunnel. If the two proxies are not too close, it is almost impossible for an attacker to control both non-secure channels.

The embodiment shown in Fig. 1 results in a data transmission
10 providing an error correction as well as an efficient encryption for the data in the channels CH1 to CH3. The source data to be transmitted from sender A to sender B are designated in Fig. 1 by reference numeral SD. Those data are divided into k source symbols SY_i ($i = 0, 1, \dots, k-1$) where this dividing
15 step is performed by sender A. Each of the source symbols SY_i has a fixed and equal number of bits b , i.e. $b \cdot k$ corresponds to the total number of source bits of the source data SD.

According to Fig. 1, the sender performs at first an encoding
20 step COD which is based on a known error correction technique in order to generate an error correction code EC from the source data SD. This error correction code comprises n repair symbols RS_j ($j = 0, 1, \dots, n-1$) as well as so-called encoded symbol identifiers ID_j for each repair symbol. The repair
25 symbols have the same number of bits b as the source symbols. However, the number of repair symbols is larger than the number of source symbols, i.e. n is a number larger than k .

In a preferred embodiment, the encoding step COD generates
30 well-known fountain codes. However, other types of error correction codes may be generated as well, provided that the encoding technique COD results in a code adding redundancy to the original source data SD so that the source data can be restored after transmission even if some information got lost
35 during transmission.

When generating a fountain code as an error correction code EC, the respective repair symbols RS_j are typically deter-

mined by first computing, in dependence on the encoded symbol identifier ID_j , a degree d which is an integer in the range between 1 and k . Secondly, given the computed degree, d source symbols are selected from the source data. Those selected source symbols are then XORed in order to produce a repair symbol RS_j . An ideal fountain code can recover source information from any k different repair symbols.

After having computed the error correction code EC , the sender A applies an encryption method CRY on the respective repair symbols RS_j . This encryption uses a pseudo-random function $f(ID_j, e)$ which generates b random bits in dependence on the respective encoded symbol identifier ID_j and in dependence on the shared secret e between the sender A and the receiver B. Those random bits are XORed with the corresponding repair symbol RS_j in order to obtain the encrypted repair symbol RS_j' .

The sender A and the sender B may have received the shared secret e before the transmission of the source data SD . However, as shown in the embodiment of Fig. 1, the shared secret e is transmitted between sender A and sender B over the channel CH' simultaneously to the data transmission over the channels CH_1 to CH_3 .

An essential feature of the invention resides in the fact that the encryption CRY is not only based on the shared secret e but also depends on the encoded symbol identifier ID_j . This improves the security of the data transmission because the encryption process is different for different repair symbols RS_j .

The encryption CRY of Fig. 1 results in the encrypted repair symbols RS_j' where the respective identifiers ID_j of the non-encrypted repair symbols RS_j are assigned to the respective encrypted repair symbol RS_j' . In a next step, respective pairs $[RS_j', ID_j]$ of encrypted repair symbols and encoded symbol identifiers are sent by the sender A over the channels

CH1 to CH3, where each channel transmits a subset of those pairs. The encoded symbol identifier ID_j needs to be transmitted together with the encrypted repair symbol RS_j' in order to determine which repair symbols have been lost in case that not all repair symbols are received by the receiver B. This information is needed in order to reconstruct the source symbols in case of data losses.

In the embodiment described herein, the number of pairs of encrypted repair symbols and encoded symbol identifiers transmitted over each of the channels CH1 to CH3 is always less than the original number k of source symbols. This ensures that an attacker listening to one channel cannot reconstruct any source symbols from this channel even in case that he succeeds to decrypt the encrypted repair symbols RS_j' .

After having received the pairs $[RS_j', ID_j]$ over channels CH1 to CH3 as well as the shared secret e over channel CH', the receiver B performs a decryption method DCR in order to decrypt the encrypted repair symbols RS_j' . This decryption uses the same function $f(ID_j, e)$ as the encryption method. To obtain the decrypted repair symbols RS_j , the sequence of random bits being the value of the function $f(ID_j, e)$ is XORed with the respective encrypted repair symbol RS_j' , i.e. $RS_j = RS_j' \text{ XOR } f(ID_j, e)$.

As a result of the decryption step DCR, the original error correction code EC comprising the non-encrypted repair symbols RS_j and the corresponding encoded symbol identifiers ID_j is obtained. Thereafter, the receiver performs a decoding DEC being inverse to the encoding COD in order to reconstruct the original source data SD with source symbols SY_i . This decoding enables the restoration of the source data even in case that some of the repair symbols are lost.

In a modification of the embodiment of Fig. 1, the encryption of the repair symbols additionally depends on the channel over which the respective repair symbol is sent. This can be

achieved by a pseudo-random function f which additionally depends on the corresponding channel CH_x ($x = 1, 2, 3$). I.e., a pseudo-random function $f(ID_j, e, CH_x)$ is used. There are several options for the receiver B to correctly identify the channels. For instance, the channels can identify themselves to the receiver B (e.g. by static allocation, broadcast information, peer-to-peer data exchange, hashing, etc.). Alternatively, the receiver B is able to identify the channels. E.g., the receiver may know that the first channel uses UDP while the second one uses a mail channel and so on, or the sender may have sent some tags on the data to identify the channel.

When using a function depending on the channels, the sender A computes $RS_j' = RS_j \text{ XOR } f(ID_j, e, CH_x)$ for the channel CH_x and sends the pair $[RS_j', ID_j]$ over the channel CH_x . At the receiver side, repair symbols that are received through channel CH_x are reverted by computing $RS_j = RS_j' \text{ XOR } f(ID_j, e, CH_x)$. After having reverted the repair symbols based on the function $f(ID_j, e, CH_x)$, the standard error correction decoding DEC described with respect to Fig. 1 can be used in order to recover the source data SD.

In the following, an example of an encoding and encryption of source symbols and the corresponding decoding and decryption is described. In this example, a sender wishes to transmit the following three source symbols to a receiver:

SY0 = 00001010
30 SY1 = 10100000
SY2 = 10101010

The sender at first generates fountain codes in the form of seven repair symbols RS1 to RS6 which read as follows:

35
RS0 = SY0 = 00001010
RS1 = SY1 = 10100000
RS2 = SY2 = 10101010

RS3 = SY0 XOR SY1 = 10101010
 RS4 = SY0 XOR SY2 = 10100000
 RS5 = SY1 XOR SY2 = 00001010
 RS6 = SY0 XOR SY1 XOR SY2 = 00000000

5

An encoded symbol identifier ID_j is associated with each repair symbol RS_j (j = 0, 1, ..., 6). E.g., the identifiers may be chosen to be ID₀ = 0, ID₁ = 1, ID₂ = 2, ID₃ = 3, ID₄ = 4, ID₅ = 5, ID₆ = 6. The sender also generates the shared secret e, e.g. by computing the random number e = 33 by a pseudo-random generator. Based on the shared secret e and the encoded symbol identifiers ID_j, the function f(ID_j, e) generates the following intermediate bit sequences for the respective seven repair symbols RS₀ to RS₆:

10

15

f(ID₀, e) = 00010101
 f(ID₁, e) = 10111101
 f(ID₂, e) = 11011100
 f(ID₃, e) = 01000010
 f(ID₄, e) = 01010001
 f(ID₅, e) = 01001011
 f(ID₆, e) = 11010001

20

Based on those intermediate bit sequences, the sender computes the encrypted repair symbols RS_j as follows:

25

RS₀' = RS₀ XOR f(ID₀, e) = 00001010 XOR 00010101 = 00011111
 RS₁' = RS₁ XOR f(ID₁, e) = 10100000 XOR 10111101 = 00011101
 RS₂' = RS₂ XOR f(ID₂, e) = 10101010 XOR 11011100 = 01110110
 RS₃' = RS₃ XOR f(ID₃, e) = 10101010 XOR 01000010 = 11101001
 RS₄' = RS₄ XOR f(ID₄, e) = 10100000 XOR 01010001 = 11110001
 RS₅' = RS₅ XOR f(ID₅, e) = 00001010 XOR 01001011 = 01000001
 RS₆' = RS₆ XOR f(ID₆, e) = 00000000 XOR 11010001 = 11010001

30

In a next step, the pairs of encoded repair symbols RS_j and assigned identifiers ID_j are sent over at least one channel, e.g. the Internet. I.e., the pairs [RS₀', ID₀], [RS₁', ID₁], [RS₂', ID₂], [RS₃', ID₃], [RS₄', ID₄], [RS₅', ID₅] and [RS₆',

35

ID6] are sent through the at least one channel. Contrary to that, the shared secret $e = 33$ is sent through another different channel, e.g. through SMS. Upon receiving e and the pairs $[RSj', IDj]$ (where some of the pairs may be dropped due to channel errors), the receiver can recompute RSj from RSj' by using $f(IDj, e)$, i.e. by determining $RSj = RSj' \text{ XOR } f(IDj, e)$. Afterwards, a standard error correction decoding is performed on the correctly received repair symbols in order to recover the original source information comprising the source symbols $SY0, SY1$ and $SY2$.

In the foregoing, a pseudo-random function f is used in the process for encrypting and decrypting the error correction code. However, it may also be possible to use a normal function instead of a pseudo-random function. Nevertheless, in case that the output length of the function f is longer than the length of the shared secret e , it is preferable to use a pseudo-random function instead of a normal function.

The invention as described in the foregoing has several advantages. Particularly, a method for data transmission is provided achieving both an error correction by adding redundancy as well as a high security by encrypting the error correction code in dependency on the encoded symbol identifiers of the respective repair symbols in the error correction code.

Patent Claims

1. A method for sending digital data over a number of channels (CH1, CH2, CH3), wherein a sender (A) performs the following steps:
- 5 - encoding (COD) source data (SD) having a first number of source symbols (SY_i), the encoding (COD) being such that an error correction code (EC) is generated from the source data (SD), the error correction code (EC) comprising a second number of repair symbols (RS_j) higher than
10 the first number as well as identifiers (ID_j) where each identifier (ID_j) is assigned to a corresponding repair symbol (RS_j), the error correction code (EC) adding redundancy to the source data (SD);
 - 15 - encrypting each repair symbol (RS_j) by an encryption process (CRY) which is based on a shared secret (e) between the sender (A) and a receiver (B), where the encryption process (CRY) for a respective repair symbol (RS_j) depends on the identifier (ID_j) assigned to the respective
20 repair symbol (RS_j);
 - feeding pairs of the encrypted repair symbols (RS_j') and the assigned identifiers (ID_j) to the number of channels (CH1, CH2, CH3) which are connected to the receiver (B).
- 25 2. The method according to claim 1, wherein the error correction code (EC) is a fountain code.
3. The method according to claim 1 or 2, wherein the source symbols (SY_i) and the repair symbols (RS_j) each are bit sequences having the same fixed bit length.
- 30 4. The method according to one of the preceding claims, wherein the encryption process (CRY) uses a pseudo-random function (f(ID_j, e)) applied to the repair symbols (RS_j), the pseudo-random function (f(ID_j, e)) depending on the shared
35 secret (e) and on the identifier (ID_j) assigned the repair symbol (RS_j) to be encrypted.

5. The method according to claim 4, wherein the pseudo-random function ($f(\text{ID}_j, e)$) generates an intermediate bit sequence having the bit length of the repair symbol (RS_j) to be encrypted, the encrypted repair symbol (RS_j') being the XOR combination of the intermediate bit sequence with the repair symbol (RS_j) to be encrypted.
6. The method according to one of the preceding claims, wherein the number of channels ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) comprises several independent channels where the pairs of the encrypted repair symbols (RS_j') and the assigned identifiers (ID_j) are fed to the independent channels ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) such that each independent channel ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) receives at least one pair of the encrypted repair symbols (RS_j') and the assigned identifiers (ID_j).
7. The method according to claim 6, wherein the independent channels ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) are channels of a Multipath TCP protocol.
8. The method according to claim 6 or 7, wherein the number of pairs of the encrypted repair symbols (RS_j') and the assigned identifiers (RS_j) in each independent channel ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) is less than the first number.
9. The method according to one of claims 6 to 8, wherein the encryption process (COD) further depends on the channel ($\text{CH}_1, \text{CH}_2, \text{CH}_3$) to which the encrypted repair symbol (RS_j') is fed.
10. The method according to one of the preceding claim, wherein, parallel to the feeding of the pairs of the encrypted repair symbols (RS_j') and the assigned identifiers (ID_j) to the number of channels ($\text{CH}_1, \text{CH}_2, \text{CH}_3$), the shared secret (e) is fed by the sender (A) to an additional channel (CH') being connected to the receiver (B) and being independent of the channels of the number of channels ($\text{CH}_1, \text{CH}_2, \text{CH}_3$).

11. A method for processing digital data being sent by a method according to one of the preceding claims, wherein a receiver (B) performs the following steps:

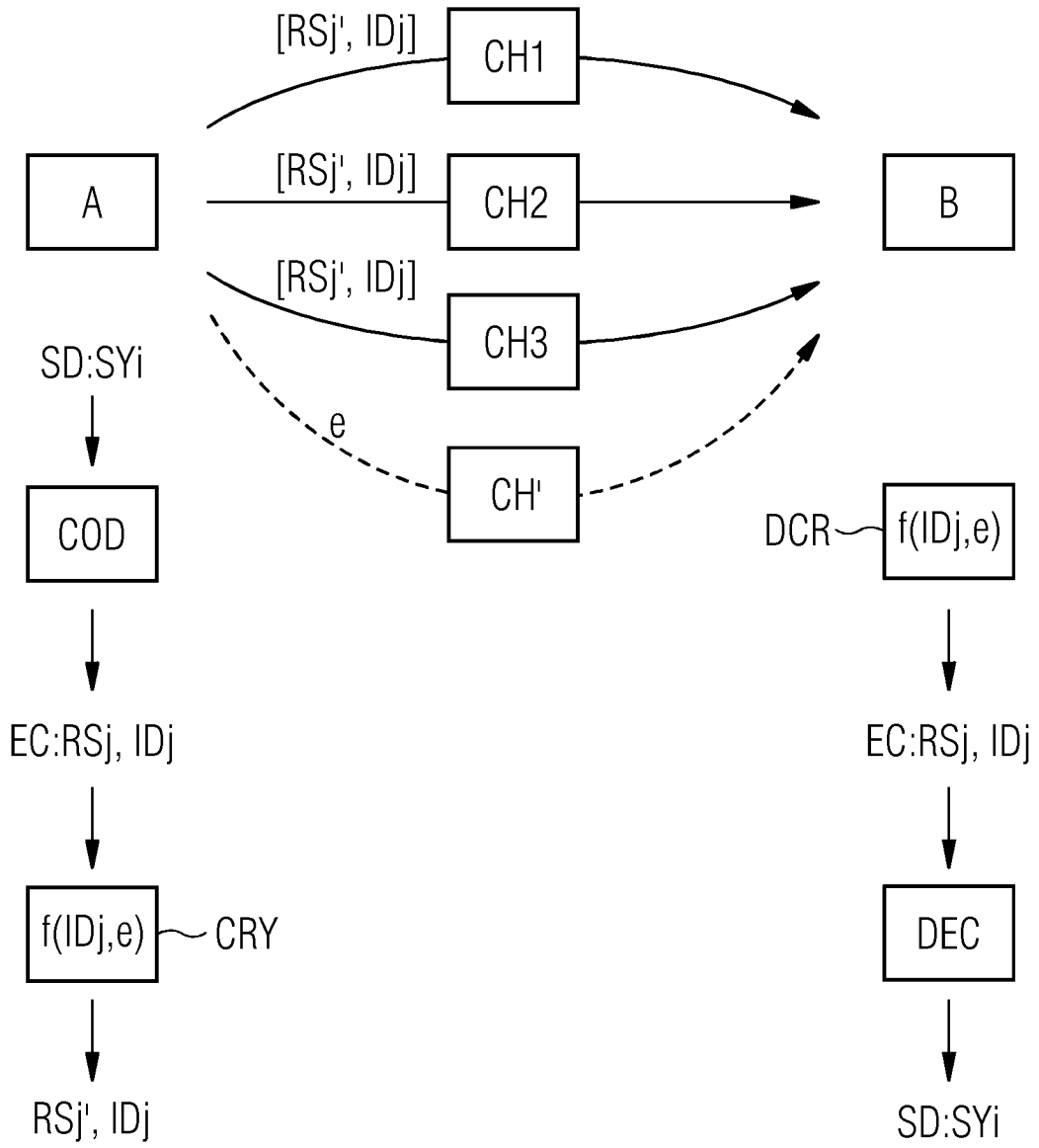
- receiving the pairs of the encrypted repair symbols (RSj') and the assigned identifiers (IDj) over the number of channels (CH1, CH2, CH3);
- decrypting each encrypted repair symbol (RSj') by a decryption process (DCR) which is based on the shared secret (e) between the sender (A) and the receiver (B), resulting in the error correction code (EC), where the decryption process (CRY) for a respective encrypted repair symbol (RSj') depends on the identifier (IDj) assigned to the respective repair symbol (RSj);
- decoding (DEC) the error correction code (EC) such that the source data (SD) are restored.

12. A method for transmitting digital data from a sender (A) to a receiver (B) over a number of channels (CH1, CH2, CH3), wherein the sender performs a method according to one of claims 1 to 10 and the receiver (B) performs a method according to claim 11.

13. A sender for sending digital data, wherein the sender (A) comprises means for performing a method according to one of claims 1 to 10.

14. A receiver for processing digital data wherein the receiver (B) comprises means for performing a method according to claim 11.

15. A transmission system, comprising a sender (A) according to claim 13 and a receiver (B) according to claim 14.



INTERNATIONAL SEARCH REPORT

International application No
PCT/EP2018/073650

A. CLASSIFICATION OF SUBJECT MATTER
INV. H04L1/00 H03M13/37 H04L29/06 H04L9/08
ADD.
According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED
Minimum documentation searched (classification system followed by classification symbols)
H04L H04M H03M
Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)
EPO-Internal, WPI Data

| C. DOCUMENTS CONSIDERED TO BE RELEVANT | | |
|--|---|-----------------------|
| Category* | Citation of document, with indication, where appropriate, of the relevant passages | Relevant to claim No. |
| A | US 9 496 897 B1 (TRIANDOPOULOS NIKOLAOS [US] ET AL) 15 November 2016 (2016-11-15) column 5, line 16 - column 8, line 24 column 13, line 12 - column 14, line 44 figures 1A-1B, 3-6, 9 claim 1 | 1-15 |
| A | US 7 548 556 B1 (WITTENSCHLAEGER THOMAS MICHAEL [US]) 16 June 2009 (2009-06-16) column 2, line 26 - line 43 column 7, line 9 - line 18 ----- -/-- | 1-15 |

Further documents are listed in the continuation of Box C.

See patent family annex.

* Special categories of cited documents :

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier application or patent but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art

"&" document member of the same patent family

| | |
|--|--|
| Date of the actual completion of the international search 26 November 2018 | Date of mailing of the international search report 20/12/2018 |
| Name and mailing address of the ISA/ European Patent Office, P.B. 5818 Patentlaan 2 NL - 2280 HV Rijswijk Tel. (+31-70) 340-2040, Fax: (+31-70) 340-3016 | Authorized officer Vucic, Nikola |

INTERNATIONAL SEARCH REPORT

International application No
PCT/EP2018/073650

| C(Continuation). DOCUMENTS CONSIDERED TO BE RELEVANT | | |
|--|--|-----------------------|
| Category* | Citation of document, with indication, where appropriate, of the relevant passages | Relevant to claim No. |
| A | <p>YONG CUI ET AL: "FMTCP: A Fountain Code-Based Multipath Transmission Control Protocol", IEEE / ACM TRANSACTIONS ON NETWORKING, IEEE / ACM, NEW YORK, NY, US, vol. 23, no. 2, 1 April 2015 (2015-04-01), pages 465-478, XP058071800, ISSN: 1063-6692, DOI: 10.1109/TNET.2014.2300140 Section IV; page 468 - page 470 -----</p> | 1-15 |

INTERNATIONAL SEARCH REPORT

Information on patent family members

International application No

PCT/EP2018/073650

| Patent document cited in search report | Publication date | Patent family member(s) | Publication date |
|--|------------------|-------------------------|------------------|
| US 9496897 | B1 | 15-11-2016 | NONE |
| ----- | | | |
| US 7548556 | B1 | 16-06-2009 | NONE |
| ----- | | | |