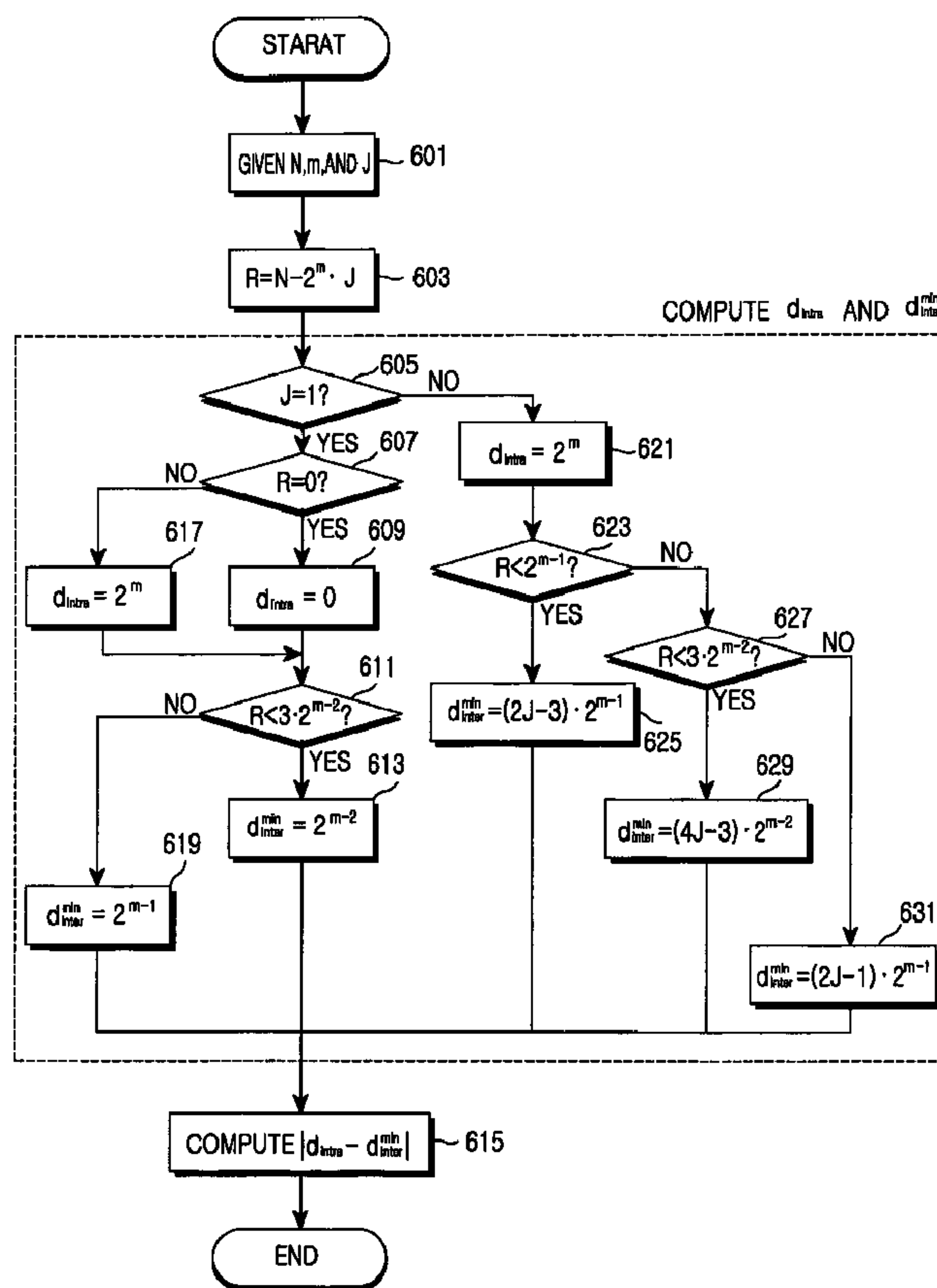




(86) Date de dépôt PCT/PCT Filing Date: 2003/02/06
 (87) Date publication PCT/PCT Publication Date: 2003/08/14
 (45) Date de délivrance/Issue Date: 2008/01/15
 (85) Entrée phase nationale/National Entry: 2003/10/02
 (86) N° demande PCT/PCT Application No.: KR 2003/000261
 (87) N° publication PCT/PCT Publication No.: 2003/067766
 (30) Priorité/Priority: 2002/02/06 (KR10-2002-0006890)

(51) Cl.Int./Int.Cl. *H03M 13/27* (2006.01),
H04L 1/22 (2006.01)
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(54) Titre : ENTRELACEUR ET PROCEDE D'ENTRELACEMENT DANS UN SYSTEME DE COMMUNICATION
 (54) Title: INTERLEAVER AND INTERLEAVING METHOD IN A COMMUNICATION SYSTEM



(57) **Abrégé/Abstract:**

A P-BRO interleaver and a method for optimizing parameters according to an interleaver size for the P-BRO interleaver. The P-BRO interleaver sequentially, by columns, arranges an input data stream of size N in a matrix having 2m rows and (J-1) columns, and R rows in a Jth column, P-BRO interleaves the arranged data, and reads the interleaved data by rows.

(12) INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(19) World Intellectual Property Organization
International Bureau(43) International Publication Date
14 August 2003 (14.08.2003)

PCT

(10) International Publication Number
WO 03/067766 A1(51) International Patent Classification⁷: H03M 13/27(72) Inventors: HA, Sang-Hyuck; 488, Kokbanjong-dong, Kwonsong-gu, Suwon-shi, 441-400 Kyonggi-do (KR).
KIM, Min-Goo; 968, Youngtong-dong, Paldal-gu, Suwon-shi, 442-470 Kyonggi-do (KR).

(21) International Application Number: PCT/KR03/00261

(22) International Filing Date: 6 February 2003 (06.02.2003)

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(25) Filing Language: English

(26) Publication Language: English

(81) Designated States (national): AU, BR, CA, CN, IN, JP, RU.

(30) Priority Data:

10-2002-0006890 6 February 2002 (06.02.2002) KR

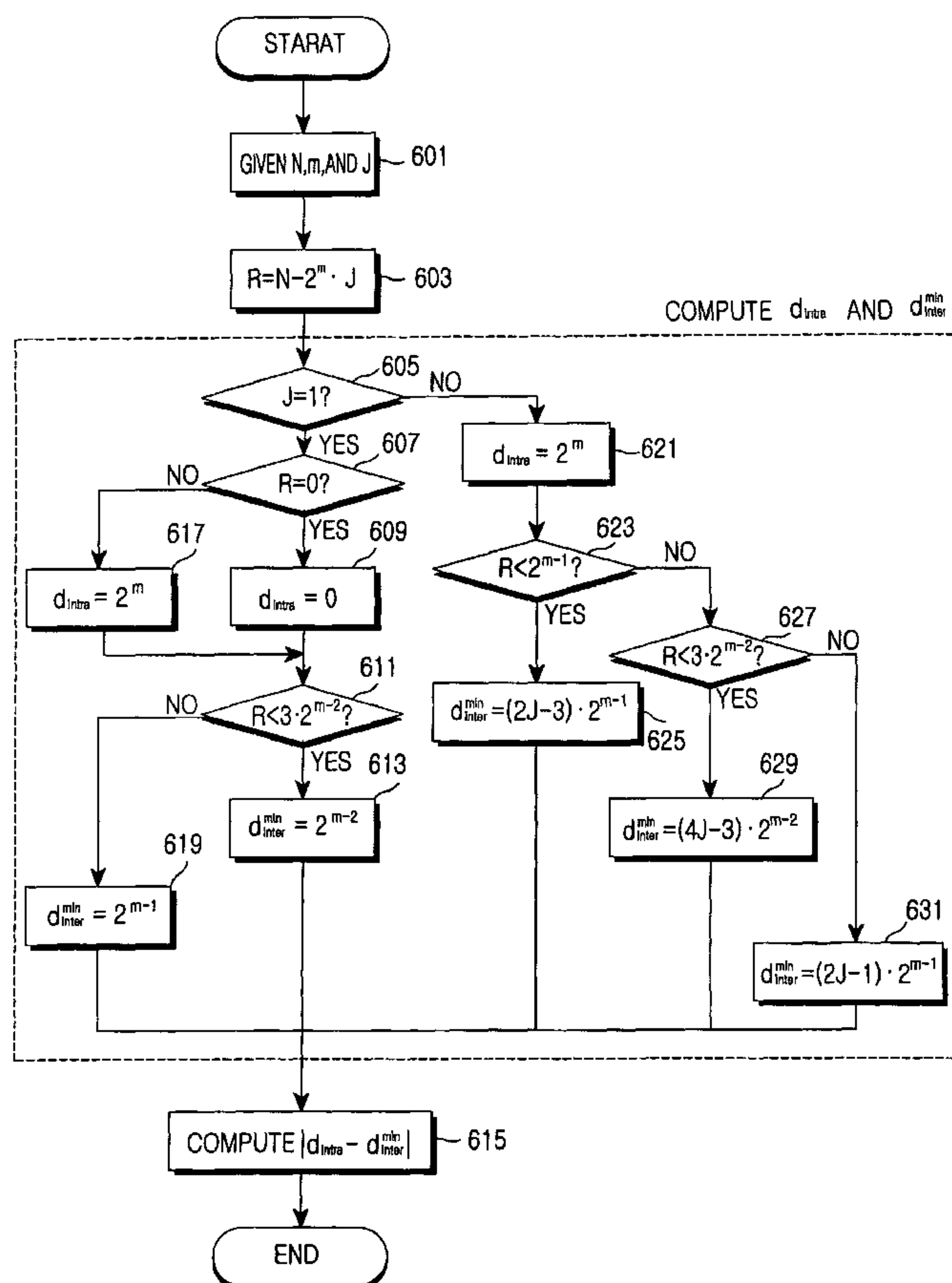
Published:

— with international search report

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For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

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INTERLEAVER AND INTERLEAVING METHOD IN A COMMUNICATION SYSTEM

BACKGROUND OF THE INVENTION

Field of the Invention:

5 The present invention relates generally to interleaving in a communication system, and in particular, to a method of optimizing parameters according to an interleaver size for partial bit reversal order (P-BRO) interleaving and an interleaver using the same.

Description of the Related Art:

10 While a sub-block channel interleaver designed in accordance with the IS-2000 Release C(1xEV-DV) F/L specification performs P-BRO operation for row permutation similarly to an existing channel interleaver designed in accordance with the IS-2000 Release A/B spec., the sub-block channel interleaver differs from the channel interleaver in that the former generates read addresses in a different manner and requires
15 full consideration of the influence of a selected interleaver parameter on Quasi-Complementary Turbo code (QCTC) symbol selection.

Hence, there is a need for analyzing the operating principles of the sub-block channel interleaver and the channel interleaver and creating criteria on which to generate
20 optimal parameters for the channel interleavers. The optimal parameters will offer the best performance in channel interleavers built in accordance with both the IS-2000 Release A/B and IS-2000 Release C.

SUMMARY OF THE INVENTION

25 An object of the present invention is to substantially solve at least the above problems and/or disadvantages and to provide at least the advantages described below. Accordingly, it is an object of the present invention to provide a method of optimizing parameters for P-BRO interleaving and an interleaver using the optimizing parameters.

30 It is another object of the present invention to provide a method of optimizing parameters m and J according to an interleaver size for P-BRO interleaving and an interleaver using the same

To achieve the above and other objects, there are provided a P-BRO
35 interleaver and a method for optimizing parameters according to an interleaver size for

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the P-BRO interleaver. The P-BRO interleaver sequentially, by columns, arranges an input data stream of size N in a matrix having 2^m rows, $(J-1)$ columns, and R rows in a J th column. The P-BRO interleaver interleaves the arranged data, and reads the interleaved data by rows. Here, N , m , J and R are given as follows:

N	m	J	R
408	7	4	24
792	8	4	24
1560	9	4	24
2328	10	3	280
3096	10	4	24
3864	11	2	1816

5

BRIEF DESCRIPTION OF THE DRAWINGS

The above and other objects, features and advantages of the present invention will become more apparent from the following detailed description of preferred embodiments thereof when taken in conjunction with the accompanying drawings, in which:

10

Fig. 1 illustrates P-BRO interleaving when $N=384$, $m=7$ and $J=3$ according to an embodiment of the present invention;

15

Fig. 2 illustrates distances between read addresses after P-BRO interleaving when $N=384$, $m=7$ and $J=3$ according to an embodiment of the present invention;

20

Fig. 3 illustrates P-BRO interleaving when $N=408$, $m=7$, $J=3$ and $R=24$ according to an embodiment of the present invention;

Fig. 4 illustrates the minimum intra-row distance after P-BRO interleaving when $N=408$, $m=7$ and $J=3$ according to an embodiment of the present invention;

25

Fig. 5 is a block diagram of an interleaver to which an embodiment of the

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present invention is applied;

Fig. 6 is a flowchart illustrating a first example of the optimal interleaver parameters determining operation according to an embodiment of the present invention;
5 and

Fig. 7 is a flowchart illustrating another example of the optimal interleaver parameters determining operation according to an embodiment of the present invention.

10 DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Several preferred embodiments of the present invention will be described in detail with reference to the accompanying drawings. In the drawings, the same or similar elements are denoted by the same reference numerals, even though they are depicted in different drawings. In the following description, a detailed description of known
15 functions or configurations incorporated herein have been omitted for conciseness.

Hereinbelow, a description will be made of P-BRO interleaving to which various embodiments of the present invention are applied, as well as the principle of determining parameters for optimal P-BRO interleaving in accordance with
20 embodiments of the present invention.

Fig. 5 is a block diagram of a P-BRO interleaver to which an embodiment of the present invention is applied. Referring to Fig. 5, an address generator 511 receives an interleaver size N , a first parameter m (i.e., Bit_Shift), a second parameter J (i.e.,
25 Up_Limit) and a clock signal Clock, and generates read addresses to read bit symbols from an interleaver memory 512. The parameters m and J are determined in an higher-layer controller (not shown) and provided to the address generator 511, or determined according to the interleaver size N in the address generator 511. The interleaver memory 512 sequentially stores input bit symbols at write addresses corresponding to count values of a counter 513 in a write mode, and outputs bit symbols from read addresses
30 received from the address generator 511 in a read mode. The counter 513 receives the clock signal Clock, generates a count value, and provides it as a write address Write ADDR to the interleaver memory 512.

35 As described above, the P-BRO interleaver writes input data sequentially in the interleaver memory 512 in the write mode and reads data from the interleaver memory 512 according to read addresses generated from the address generator 511. For

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details of the P-BRO interleaver, reference is made to Korea Patent Application No. 1998-54131, filed on December 10, 1998, the entire contents of which are expressly incorporated herein.

5 In operation, the address generator 511 generates a read address A_i for symbol permutation by

$$A_i = 2^m (i \bmod J) + \text{BRO}_m(\lfloor i/J \rfloor) \dots \dots (1)$$

10 where $i=0, 1, \dots, N-1$ and $N=2^m \times J$.

In Eq. (1), N denotes the size of an interleaver input sequence and m and J are interleaver parameters called *Up_Limit* and *Bit_Shift*, respectively.

15 Fig. 1 illustrates P-BRO interleaving when $N=384$, $m=7$ and $J=3$. Referring to Fig. 1, an interleaving matrix has 2^m rows starting from index 0 and J columns starting from index 0. After step 101, the row index and column index of a symbol in the resulting matrix are expressed as $\lfloor i/J \rfloor$ and $(i \bmod J)$, respectively. Therefore, after $2^m(i \bmod J) + \lfloor i/J \rfloor$, an i th symbol in an input sequence has a number corresponding to an
20 $\lfloor i/J \rfloor$ th row and an $(i \bmod J)$ column as its read address. J symbols are in each row and the distance between symbols is 2^m in the row.

The row index $\lfloor i/J \rfloor$ is BRO-operated in step 102. If the distance between symbols in adjacent rows of the same column is row distance d_{row} , the BRO operation of
25 the row indexes results in a row permutation such that two minimum row distances d_{row} are 2^{m-2} and 2^{m-1} , as illustrated in Fig. 2. Thus, after $2^m(i \bmod J) + \text{BRO}_m(\lfloor i/J \rfloor)$, the i th symbol in the input sequence has a number corresponding to a $\text{BRO}_m(\lfloor i/J \rfloor)$ th row and an $(i \bmod J)$ th column as its read address in the third matrix from the left.

In summary, a read address sequence is generated by row permutations of a $2^m \times J$ matrix
30 in the P-BRO interleaver. The row-permuted matrix is read first by rows from the top to the bottom, then subsequently reading each row from the left to the right.

For clarity of description, the distance between adjacent addresses in the same row is defined as "intra-row distance d_{intra} ". If $J \neq 1$, $d_{\text{intra}} = 2^m$. If $J=1$, there is no
35 intra-row distance.

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The distance between adjacent addresses in different rows, that is, the distance between the last address in a row and the first address in the next row is defined as "inter-row distance d_{inter} ". d_{inter} is one of a plurality of values calculated from a function of the parameters m and J . When m and J are determined, the resulting minimum inter-row distance d_{inter}^{min} is defined as d_{inter}^{min} .

Since two minimum rows distances d_{row} are 2^{m-2} and 2^{m-1} ,

$$\begin{aligned} \text{If } J = 1, \quad & d_{inter}^{min} = d_{row}^{min} = 2^{m-2}, \\ \text{Else,} \quad & d_{inter}^{min} = (J-1) \cdot 2^m - 2^{m-1} = (2 \cdot J - 3) \cdot 2^{m-1} \end{aligned}$$

10 (2)

The reason for computing d_{inter}^{min} by Eq. (2) when $J \neq 1$ is apparent in Fig. 2. If $J=1$, which implies that the interleaving matrix has only one column, d_{inter}^{min} is d_{row}^{min} , that is, 2^{m-2} .

As described above, the interleaver parameters m and J are used as the numbers of rows and columns in a read address sequence matrix and parameters for a function that determines distances between read addresses. Consequently, the characteristics of the P-BRO channel interleaver depend on the interleaver parameters m and J .

Before presenting a description of a method of determining sub-block channel interleaver parameters that ensure the best interleaving performance according to an embodiment of the present invention, the purposes of channel interleavers in the IS-2000 specifications, Releases A/B and C will first be described. Following that, the interleaver parameter determination will then be described separately in two cases: $N=2^m \times J$; and $N=2^m \times J + R$.

The purpose of channel interleaving in the IS-2000 specification, Release A/B, is to improve decoding performance, which is degraded when fading adversely influences successive code symbols, through error scattering resulting from symbol permutation. To improve decoding performance, interleaving must be performed such that the distance between adjacent addresses (inter-address distance) is maximized.

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5 Meanwhile, the purpose of sub-block channel interleaving as described in the IS-2000 specification, Release C, is to allow a QCTC symbol selector at the rear end of an interleaver to select appropriate code symbols according to a coding rate and thus ensure the best performance at the coding rate, as well as to scatter errors through symbol permutation. To achieve this purpose, interleaving must be performed such that inter-address distances are maximized and are uniform.

10 Accordingly, to satisfy the requirements of the channel interleaver of the IS-2000 specification, Release A/B, and the sub-block channel interleaver of the IS-2000 specification, Release C, an interleaver must be designed so that a read address sequence is uniformly permuted by interleaving. This is possible by determining the interleaver parameters m and j that maximize a minimum inter-address distance and minimize the difference between inter-address distances.

15 As stated before, the inter-address distances are categorized into the intra-row distance d_{intra} and the inter-row distance d_{inter} . The intra-row distance is a function of m and the inter-row distance is a function of m and J . Since there are a plurality of inter-row distances, a minimum inter-row distance d_{inter}^{min} is calculated. A minimum inter-address distance is always 2^{m-2} when J is 1, and the smaller of the minimum inter-row distance d_{inter}^{min} and the minimum intra-row distance d_{intra}^{min} when J is not 1. The difference between inter-address distances is 2^{m-2} when J is 1, since the intra-row distance d_{intra} is 0, and is equal to the difference between the intra-row distance d_{intra} and the minimum inter-row distance d_{inter}^{min} when J is not 1.

25 This can be expressed as follows:

$$\begin{aligned} \text{If } J = 1, \quad & |0 - 2^{m-2}| = 2^{m-2}, \\ \text{Else,} \quad & |d_{intra} - d_{inter}^{min}| = |2^m - (2 \cdot J - 3) \cdot 2^{m-1}| = |2 \cdot J - 5| \cdot 2^{m-1} \\ & \dots \dots (3) \end{aligned}$$

30 Since $N=2^m \times J$, 2^m is replaced by N/J in Eq. (3), it follows that

$$\begin{aligned} \text{If } J = 1, \quad & 2^{m-2} = \frac{1}{4} \cdot \frac{N}{J} = 0.25 \frac{N}{J}, \\ \text{Else,} \quad & |d_{intra} - d_{inter}^{min}| = |2 \cdot J - 5| \cdot 2^{m-1} = \left| J - \frac{5}{2} \frac{N}{J} \right| = \left| 1 - \frac{2.5}{J} \right| \cdot N \end{aligned}$$

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..... (4)

When $J=3$ in Eq. (4), the difference between inter-address distances is minimized. Thus $|d_{intra} - d_{inter}^{min}| = 0.166667N$.

5 Table 1 below illustrates changes in inter-read address distances as m increases when $N=384$. When $J=3$, a maximum difference between inter-address distances is minimized, 64 and a minimum inter-address distance d^{min} is maximized, 128.

Table 1

N	m	J	d_{intra}	d_{inter}^{min}	$ d_{intra} - d_{inter}^{min} $	d^{min}
384	4	24	16	360	344	16
	5	12	32	336	304	32
	6	6	64	288	224	64
	7	3	128	192	64	128

10 The method of determining optimal interleaver parameters when $N=2^m \times J$ has been described above. Now, a method of determining optimal interleaver parameters when $N=2^m \times J + R$ will be described. Here, R is the remainder of dividing N by 2^m . Thus R is a positive integer less than 2^m .

15 Fig. 3 illustrates P-BRO interleaving when $N=408$, $m=7$, $J=3$ and $R \neq 0$. Referring to Fig. 3, similarly to the case where $R=0$, numbers in a row-permuted matrix after step 302 are read as read addresses by rows from the top to the bottom, reading each row from the left to the right, as described in step 303. Since $R \neq 0$, the number of columns is $J+1$, and numbers are filled in only R rows of a $(J+1)$ th column with no
20 numbers in the other $(2^m - R)$ rows.

In summary, when $R \neq 0$, a read address sequence is generated by a row permutation of a $2^m \times J$ matrix, each row including J or $J+1$ elements in the P-BRO interleaver. The row-permuted matrix is read by rows from the top to the bottom,
25 reading each row from the left to the right.

Furthermore, when $R \neq 0$, the interleaver parameters m and J are determined such that a minimum inter-read address distance is maximized and the difference between inter-read address distances is minimized.

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An inter-row distance d_{inter} is a function of m , 2^m irrespective of whether $R=0$ or $R \neq 0$. However, while the minimum inter-row distance d_{inter}^{min} is a function of m and J when $R=0$, it is a function of m , J and R when $R \neq 0$.

5 The minimum inter-row distance is determined according to J by Eq. (5) and Eq. (6).

When $J = 1$,

For $0 \leq R < 3 \cdot 2^{m-2}$, $d_{inter}^{min} = 2^{m-2}$

For $3 \cdot 2^{m-2} \leq R < 2^m$, $d_{inter}^{min} = 2^{m-1}$

..... (5)

10

When $J \neq 1$,

For $0 \leq R < 2^{m-1}$, $d_{inter}^{min} = (J-1) \cdot 2^m - 2^{m-1} = (2J-3) \cdot 2^{m-1}$

For $2^{m-1} \leq R < 3 \cdot 2^{m-2}$, $d_{inter}^{min} = (J-1) \cdot 2^m - (-2^{m-2}) = (4J-3) \cdot 2^{m-2}$

For $3 \cdot 2^{m-2} \leq R < 2^m$, $d_{inter}^{min} = J \cdot 2^m - 2^{m-1} = (2J-1) \cdot 2^{m-1}$

... (6)

15 Fig. 4 illustrates how Eq. (6) is derived when $m=7$ and $J=3$. Referring to Fig. 4, when $0 \leq R < 2^{m-1}$, the inter-row distance between two adjacent rows having a row distance d_{row} of 2^{m-1} , the last column of the upper row being empty, is a minimum inter-row distance ($d_{inter}^{min} = (2J-3) \cdot 2^{m-1}$). When $2^{m-1} \leq R < 3 \cdot 2^{m-2}$, the inter-row distance between two adjacent rows having a row distance d_{row} of 2^{m-2} , the last column of the upper row being empty, is a minimum inter-row distance ($d_{inter}^{min} = (4J-3) \cdot 2^{m-2}$).

20

When $3 \cdot 2^{m-2} \leq R < 2^m$, the inter-row distance between two adjacent rows having a row distance d_{row} of 2^{m-2} and elements in the last columns, is a minimum inter-row distance ($d_{inter}^{min} = (2J-1) \cdot 2^{m-1}$). For example, if $R=0$, the minimum inter-row distance is 192, as indicated by reference numeral 401. If $R=64$ (2^{m-1}), the minimum inter-row distance is 288, as indicated by reference numeral 402. If $R=96$ ($3 \cdot 2^{m-2}$), the minimum inter-row distance is 320, as indicated by reference numeral 403. In the same manner, Eq. (5) can be derived when $J=1$.

30

Table 2 below illustrates changes in the interleaver parameters J and R , the intra-row distance d_{intra} , the minimum inter-row distance d_{inter}^{min} , and the minimum inter-read

- 9 -

address distance d_{\min} as m increases, with respect to six encoder packet (EP) sizes as described in the IS-2000 specification, Release C.

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Table 2

N	m	J	R	d_{intra}	d_{inter}^{min}	$ d_{intra} - d_{inter}^{min} $	d^{min}	$n(d^{min})$
408	3	51	0	8	396	388	8	400
	4	25	8	16	388	372	16	392
	5	12	24	32	368	336	32	376
	6	6	24	64	288	224	64	344
	7	3	24	128	192	64	128	280
	8	1	152	256	64	192	64	40
792	4	49	8	16	772	756	16	776
	5	24	24	32	752	720	32	760
	6	12	24	64	672	608	64	728
	7	6	24	128	576	448	128	664
	8	3	24	256	384	128	256	536
	9	1	280	512	128	384	128	104
1560	5	48	24	32	1520	1488	32	1528
	6	24	24	64	1440	1376	64	1496
	7	12	24	128	1344	1216	128	1432
	8	6	24	256	1152	896	256	1304
	9	3	24	512	768	256	512	1048
	10	1	536	1024	256	768	256	232
2328	6	36	24	64	2208	2144	64	2264
	7	18	24	128	2112	1984	128	2200
	8	9	24	256	1920	1664	256	2072
	9	4	280	512	1664	1152	512	1816
	10	2	280	1024	512	512	512	232
	11	1	280	2048	512	1536	512	512
3096	6	48	24	64	2976	2912	64	3032
	7	24	24	128	2880	2752	128	2968
	8	12	24	256	2688	2432	256	2840
	9	6	24	512	2304	1792	512	2584
	10	3	24	1024	1536	512	1024	2072
	11	1	1048	2048	512	1536	512	488
3864	6	60	24	64	3744	3680	64	3800
	7	30	24	128	3648	3520	128	3736

3864

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8	15	24	256	3456	3200	256	3608
9	7	280	512	3200	2688	512	3352
10	3	792	1024	2560	1536	1024	2840
<u>11</u>	<u>1</u>	<u>1816</u>	<u>2048</u>	<u>1024</u>	<u>1024</u>	<u>1024</u>	<u>1024</u>

As described above, similarly to the case where $R=0$, optimal interleaver parameters are selected which maximize a minimum inter-address distance and minimize the difference between inter-address distances.

5

In Table 2, the minimum inter-read address distance d^{\min} in the eighth column is the smaller of the intra-row distance d_{intra} and the minimum inter-row distance d_{inter}^{\min} . Hence, parameters that maximize the minimum inter-read address distance d^{\min} can be obtained by selecting a row having the maximum value in the eighth column. For EP sizes of 2328 and 3864, three rows and two rows satisfy this condition. In this case, rows that satisfy another condition of minimizing the difference between inter-read address $|d_{\text{intra}} - d_{\text{inter}}^{\min}|$ must be selected. They are shown in bold and underlined in Table 2. The validity of this condition is apparent by comparing the rows having the maximum d^{\min} in terms of $n(d^{\min})$ in the last column. Here, $n(d^{\min})$ indicates the number of address pairs having a minimum inter-address distance d^{\min} .

10

15

Rows marked in bold and underlined in Table 2 satisfy the above two conditions for selecting optimal interleaver parameters. As noted, once the second condition is satisfied, the first condition is naturally satisfied. For reference, it is made clear that the intra-row distances d_{intra} and the minimum inter-row distances d_{inter}^{\min} listed in Table 2 are equal to those computed on P-BRO-interleaved read addresses. Table 2 covers both cases of dividing N by 2^m or J with no remainder and of dividing N by 2^m or J with a remainder R (i.e., $N=2^m \times J + R$ ($0 \leq R < 2^m$)). Here, interleaver parameters shown in bold and underlined are optimal for each EP size.

20

25

When $N=2^m \times (J-1) + R$ ($0 \leq R < 2^m$), that is, N is divided by 2^m or J either with no remainder or with a remainder R , optimal interleaver parameters for each interleaver size N are listed in Table 3. The description made in the context of J is also applied when J is replaced by $(J-1)$.

30

Table 3

N	m	J	R
---	---	---	---

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408	7	4	24
792	8	4	24
1560	9	4	24
2328	10	3	280
3096	10	4	24
3864	11	2	1816

5 The above description has provided a method of selecting interleaver parameters expected to offer the best performance when, for example, a channel interleaver built in accordance with the IS-2000 Release A/B specification, and a sub-block channel interleaver built in accordance with the IS-2000 Release C specification are used.

10 As described above, the optimal interleaver parameters are those that maximize an inter-address distance and at the same time, minimize the difference between inter-address distances when generating read addresses in a channel interleaver. Consequently, interleaver parameters for sub-block channel interleaving in circumstances wherein a sub-block channel interleaver is built in accordance with the IS-2000 Release C specification are values in the rows in bold and underlined in Table 2. While interleaver parameters selection has been described for the sub-block channel
15 interleaver built in accordance with the IS-2000 Release C specification, it is obvious that the same thing can also be applied to systems of other standards.

20 Fig. 6 is a flowchart illustrating an optimal interleaver parameters determining operation according to an embodiment of the present invention. Particularly, this operation is concerned with the computation of $|d_{intra} - d_{inter}^{min}|$. An optimal (m, J) that minimizes $|d_{intra} - d_{inter}^{min}|$ is selected by computing $|d_{intra} - d_{inter}^{min}|$, changing (m, J).

25 Referring to Fig. 6, when an interleaver size N, and parameters m and J are given in step 601, a parameter R is calculated by subtracting $2^m \times J$ from N in step 603. In step 605, it is determined whether J is 1. This is a determination, therefore, of whether an interleaving matrix has a single column or not. If J is 1, the procedure goes to step 607 (“Yes” path from decision step 605) and if J is not 1, the procedure goes to step 621 (“No” path from decision step 605). In step 607, it is determined whether R is 0 (i.e., whether N is an integer multiple of 2^m). On the contrary, if R is 0 (“Yes” path from decision step 607), an intra-row distance d_{intra} is set to 0 in step 609. If R is not 0 (“No”
30 path from decision step 607), d_{intra} is set to 2^m in step 617.

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After d_{intra} is determined, it is determined whether R is less than $3 \times 2^{m-2}$ in step 611. If R is less than $3 \times 2^{m-2}$ ("Yes" path from decision step 611) a minimum inter-row distance d_{inter}^{min} is set to 2^{m-2} in step 613. If R is equal to or greater than $3 \times 2^{m-2}$ ("No" path from decision step 611) d_{inter}^{min} is set to 2^{m-1} in step 619. After d_{inter}^{min} is determined, $|d_{intra} - d_{inter}^{min}|$ is calculated in step 615.

Meanwhile, if J is not 1 in step 605, d_{intra} is set to 2^m in step 621 and it is determined whether R is less than 2^{m-1} in step 623. If R is less than 2^{m-1} ("Yes" path from decision step 623) d_{inter}^{min} is set to $(2J-3) \times 2^{m-1}$ in step 625 and then the procedure goes to step 615. If R is equal to or greater than 2^{m-1} ("No" path from decision step 623), it is determined whether R is less than $3 \times 2^{m-2}$ in step 627. If R is less than $3 \times 2^{m-2}$ ("Yes" path from decision step 627), d_{inter}^{min} is set to $(4J-3) \times 2^{m-2}$ in step 629. If R is equal to or greater than $3 \times 2^{m-2}$ ("No" path from decision step 627), d_{inter}^{min} is set to $(2J-1) \times 2^{m-1}$ in step 631 and then the procedure goes to step 615.

Optimal interleaver parameters m and J are achieved for a given N by computing $|d_{intra} - d_{inter}^{min}|$, changing (m, J). If J is one of 1, 2 and 3, a logical formula that facilitates selection of J without the repeated computation can be derived.

With a description of a logical equation deriving procedure omitted, the logical equation is

If $\log_2 N - \lfloor \log_2 N \rfloor < \log_2 3 - 1 = 0.5849625$,

$$\text{For } \left(\frac{3}{4}\right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 1 \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 3,$$

$$\text{For } 1 \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{3}{2}\right) \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 2,$$

$$\text{For } \left(\frac{3}{2}\right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 2 \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 1.$$

Else if $\log_2 N - \lfloor \log_2 N \rfloor \geq \log_2 3 - 1 = 0.5849625$,

$$\text{For } 1 \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{3}{2}\right) \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 2,$$

$$\text{For } \left(\frac{3}{2}\right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{7}{4}\right) \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 3,$$

$$\text{For } \left(\frac{7}{4}\right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 2 \cdot 2^{\lfloor \log_2 N \rfloor}, \quad J = 1.$$

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..... (7)

From an optimal J from Eq. (7), an optimal m is calculated by

$$m = \left\lceil \log_2 \left(\frac{N}{J} \right) \right\rceil \quad \dots (8)$$

The selection of optimal interleaver parameters by the simple logical equations is summarized below and illustrated in Fig. 7.

1. An optimal J is obtained by Eq. (7) for a given N ; and
2. m is calculated by computing Eq. (8) using N and J .

Fig. 7 is a flowchart illustrating an optimal interleaver parameters determining operation according to another embodiment of the present invention.

Referring to Fig. 7, when N is given, a variable α is calculated by $\log_2 N - \lfloor \log_2 N \rfloor$ and a variable β is calculated by $2^{\lfloor \log_2 N \rfloor}$ in step 701. Decision step 703, determines whether α is less than a first threshold, 0.5849625. If α is less than the first threshold ("Yes" path from decision step 703), another decision is made, whether N is less than β in decision step 705. If N is equal to or greater than β ("No" path from decision step 705), the procedure goes to step 707. On the contrary, if N is less than β ("Yes" path from decision step 705), J is determined to be 3 in step 713.

Meanwhile, decision step 707 determines whether N is less than $(3/2) \times \beta$. If N is less than $(3/2) \times \beta$ ("Yes" path from decision step 707), J is determined to be 2 in step 711. Otherwise, J is determined to be 1 in step 709 ("No" path from decision step 707).

If α is equal to or greater than the first threshold in step 703 ("No" path from decision step 703), a decision is made whether N is less than $(3/2) \times \beta$ in decision step 717. If N is less than $(3/2) \times \beta$ ("Yes" path from decision step 717), J is determined to be 2 in step 721. Otherwise, decision step 719 determines whether N is less than $(7/4) \times \beta$. If N is less than $(7/4) \times \beta$ ("Yes" path from decision step 719), J is determined to be 3 in step 723. Otherwise, J is determined to be 1 in step 725 ("No" path from decision step 719).

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As described above, optimal m and J can be calculated simply by the logical equations using N . The optimal m and J are equal to m and J resulting from repeated computation using different (m, J) values as illustrated in Table 2. This obviates the need for storing optimal m and J values according to N values.

When $N=2328$, for example, optimal m and J values are calculated in the procedure illustrated in Fig. 7 or by Eq. (8) to Eq. (10), as follows.

$$\alpha = \log_2 N - \lfloor \log_2 N \rfloor = \log_2 2328 - \lfloor \log_2 2328 \rfloor = 11.1848753 - 11 = 0.1848753.$$

$$\beta = 2^{\lfloor \log_2 N \rfloor} = 2^{\lfloor \log_2 2328 \rfloor} = 2^{11} = 2048.$$

$$\alpha \leq 0.5849625 \text{ and } \beta = 2048 \leq N = 2328 < \left(\frac{3}{2}\right) \cdot \beta = 3072. \text{ Thus } J = 2.$$

$$m = \left\lfloor \log_2 \frac{N}{J} \right\rfloor = \left\lfloor \log_2 \left(\frac{2328}{2}\right) \right\rfloor = \lfloor \log_2 1164 \rfloor = 10, R = N - 2^m \cdot J = 2328 - 2^{10} \cdot 2 = 280.$$

For reference, Eq. (7) is derived as follows.

In each case depicted in Fig. 6, Eq. (5) and Eq. (6), $|d_{intra} - d_{inter}^{min}|$ is determined by

15 A. When $J=1$,

$$\text{A-1. If } R=0, |d_{intra} - d_{inter}^{min}| = |0 - 2^{m-2}| = 2^{m-2}$$

$$\text{A-2. If } 0 < R < 3 \cdot 2^{m-2}, |d_{intra} - d_{inter}^{min}| = |2^m - 2^{m-2}| = 3 \cdot 2^{m-2}$$

$$\text{A-3. If } 3 \cdot 2^{m-2} \leq R < 2^m, |d_{intra} - d_{inter}^{min}| = |2^m - 2^{m-1}| = 2^{m-1}$$

20 B. When $J \neq 1$,

$$\text{B-1. If } 0 \leq R < 2^{m-1}, |d_{intra} - d_{inter}^{min}| = |2^m - (2J - 3) \cdot 2^{m-1}| = |2J - 5| \cdot 2^{m-1}$$

$$\text{B-2. If } 2^{m-1} \leq R < 3 \cdot 2^{m-2}, |d_{intra} - d_{inter}^{min}| = |2^m - (4J - 3) \cdot 2^{m-2}| = |4J - 7| \cdot 2^{m-2}$$

$$\text{B-3. If } 3 \cdot 2^{m-2} \leq R < 2^m, |d_{intra} - d_{inter}^{min}| = |2^m - (2J - 1) \cdot 2^{m-1}| = |2J - 3| \cdot 2^{m-1}$$

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Since $N=2^m \cdot J+R$ and $0 \leq R < 2^m$, $J \cdot 2^m \leq N < (J+1) \cdot 2^m$. When this is divided by J and then subject to a log base 2 operation,

$$m \leq \log_2 \left(\frac{N}{J} \right) < \log_2 \left(\left(\frac{J+1}{J} \right) \cdot 2^m \right) = m + \log_2 \left(1 + \frac{1}{J} \right) < m + 1$$

5 Thus, $m = \left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor$. Using $m = \left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor$, J can be expressed as a function of N for

all the cases of A and B.

A'. When $J=1$, since $m = \lfloor \log_2 N \rfloor$, $R = N - 2^m = N - 2^{\lfloor \log_2 N \rfloor}$. Then the cases A-1, A-2 and A-3 can be expressed as functions of N . It therefore follows that:

10 A'-1: If $N = 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = 2^{m-2} = \left(\frac{1}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$

A'-2: If $2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{7}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \left(\frac{3}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$

A'-3: If $\left(\frac{7}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 2 \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \left(\frac{1}{2} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$

B'. When $J \neq 1$, since $m = \left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor$, $R = N - J \cdot 2^m = N - J \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$.

15 Then the cases B-1, B-2 and B-3 can be expressed as functions of N instead of R . Therefore,

B'-1: If $J \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor} \leq N < \left(J + \frac{1}{2} \right) \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$,

$$|d_{intra} - d_{inter}^{min}| = \left| J - \frac{5}{2} \right| \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$$

20 B'-2: If $\left(J + \frac{1}{2} \right) \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor} \leq N < \left(J + \frac{3}{4} \right) \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$,

$$|d_{intra} - d_{inter}^{min}| = \left| J - \frac{7}{4} \right| \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$$

B'-3: If $\left(J + \frac{3}{4} \right) \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor} \leq N < (J+1) \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$,

$$|d_{intra} - d_{inter}^{min}| = \left| J - \frac{3}{2} \right| \cdot 2^{\left\lfloor \log_2 \left(\frac{N}{J} \right) \right\rfloor}$$

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B''. When J=2, since $\left\lfloor \log_2 \left(\frac{N}{2} \right) \right\rfloor = \lfloor \log_2 N - 1 \rfloor = \lfloor \log_2 N \rfloor - 1$,

B''-1: If $2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{5}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{1}{4} \cdot 2^{\lfloor \log_2 N \rfloor}$

5 B''-2: If $\left(\frac{5}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{11}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{1}{8} \cdot 2^{\lfloor \log_2 N \rfloor}$

B''-3: If $\left(\frac{11}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{3}{2} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{1}{4} \cdot 2^{\lfloor \log_2 N \rfloor}$

B'''. When J=3,

since $\left\lfloor \log_2 \left(\frac{N}{3} \right) \right\rfloor = \begin{cases} \lfloor \log_2 N \rfloor - 2, & \text{if } \log_2 N - \lfloor \log_2 N \rfloor < \log_2 3 - 1 \\ \lfloor \log_2 N \rfloor - 1, & \text{otherwise} \end{cases}$,

10

if $\log_2 N - \lfloor \log_2 N \rfloor < \log_2 3 - 1 = 0.5849625$,

B'''-1': If $\left(\frac{3}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{7}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{1}{8} \cdot 2^{\lfloor \log_2 N \rfloor}$

B'''-2': If $\left(\frac{7}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{15}{16} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{5}{16} \cdot 2^{\lfloor \log_2 N \rfloor}$

B'''-3': If $\left(\frac{15}{16} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{3}{8} \cdot 2^{\lfloor \log_2 N \rfloor}$

15

if $\log_2 N - \lfloor \log_2 N \rfloor \geq \log_2 3 - 1 = 0.5849625$,

B'''-1'': If $\left(\frac{3}{2} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{7}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{1}{4} \cdot 2^{\lfloor \log_2 N \rfloor}$

B'''-2'': If $\left(\frac{7}{4} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < \left(\frac{15}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{5}{8} \cdot 2^{\lfloor \log_2 N \rfloor}$

B'''-3'': If $\left(\frac{15}{8} \right) \cdot 2^{\lfloor \log_2 N \rfloor} \leq N < 2 \cdot 2^{\lfloor \log_2 N \rfloor}$, $|d_{intra} - d_{inter}^{min}| = \frac{3}{4} \cdot 2^{\lfloor \log_2 N \rfloor}$

20 If J is 4 or more, this case is neglected because $|d_{intra} - d_{inter}^{min}|$ cannot be less than $|d_{intra} - d_{inter}^{min}|$ in any of the cases where J=1, 2, and 3.

Eq. (7) is obtained by selecting a case having a minimum $|d_{intra} - d_{inter}^{min}|$ among the cases of A'-1, A'-2, A'-3, B''-1, B''-2, B''-3, B'''-1', B'''-2', and B'''-3'.

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Similarly, Eq. (8) is obtained by selecting a case having a minimum $|d_{intra} - d_{inter}^{min}|$ among the cases of A'-1, A'-2, A'-3, B''-1, B''-2, B''-3, B'''-1'', B'''-2'', and B'''-3''.

5 In accordance with the embodiments of the present invention as described above, interleaver parameters m and J are simply optimized according to an interleaver size N, for P-BRO interleaving.

10 While the invention has been shown and described with reference to certain preferred embodiments thereof, it will be understood by those skilled in the art that various changes in form and details may be made therein without departing from the spirit and scope of the invention as defined by the appended claims.

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What is claimed is:

- 5 1. An interleaver that sequentially arranges by columns an input data stream of size N in a matrix having 2^m rows, $(J-1)$ columns and R rows in a J th column, partial-bit reversal order (P-BRO) interleaves the arranged data and reads the interleaved data by rows wherein example of N , m , J and R are given as follows:

N	m	J	R
408	7	4	24
792	8	4	24
1560	9	4	24
2328	10	3	280
3096	10	4	24
3864	11	2	1816

- 10 2. A method of determining parameters for an interleaver in a communication system, comprising:

sequentially arranging by columns an input data stream of size N in a matrix having 2^m rows, $(J-1)$ columns and R rows in a J th column ($0 \leq R < 2^m$);

partial-bit reversal order (P-BRO) interleaving the arranged data and generating read addresses for reading the interleaved data by rows;

- 15 calculating a first minimum distance between addresses indicating adjacent columns in the same row in the generated read addresses;

calculating a second minimum distance between addresses indicating a last column of a row and an address indicating a first column of the next row in the generated read addresses; and

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repeating the steps of arranging, P-BRO interleaving, calculating the first minimum distance and calculating the second minimum distance until the m and J values, that minimize the difference between the first minimum distance and the second minimum distance, are determined.

5

3. The method of claim 2, wherein the parameters N , m , J , and R are determined to be

N	m	J	R
408	7	4	24
792	8	4	24
1560	9	4	24
2328	10	3	280
3096	10	4	24
3864	11	2	1816

10

4. The method of claim 2, wherein the second minimum distance is determined according to the following Equations.

When $J = 1$,

For $0 \leq R < 3 \cdot 2^{m-2}$, $d_{\text{inter}}^{\text{min}} = 2^{m-2}$

For $3 \cdot 2^{m-2} \leq R < 2^m$, $d_{\text{inter}}^{\text{min}} = 2^{m-1}$

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When $J \neq 1$,

$$\text{For } 0 \leq R < 2^{m-1}, \quad d_{inter}^{min} = (J-1) \cdot 2^m - 2^{m-1} = (2J-3) \cdot 2^{m-1}$$

$$\text{For } 2^{m-1} \leq R < 3 \cdot 2^{m-2}, \quad d_{inter}^{min} = (J-1) \cdot 2^m - (-2^{m-2}) = (4J-3) \cdot 2^{m-2}$$

$$\text{For } 3 \cdot 2^{m-2} \leq R < 2^m, \quad d_{inter}^{min} = J \cdot 2^m - 2^{m-1} = (2J-1) \cdot 2^{m-1}$$

5. A method of determining parameters for an interleaver in a communication system, comprising:

sequentially arranging by columns an input data stream of size N in a matrix having 2^m rows, (J-1) columns and R rows in a last column ($0 \leq R < 2^m$);

partial-bit reversal order (P-BRO) interleaving the arranged data, and generating read addresses for reading the interleaved data by rows;

calculating a first minimum distance between addresses indicating adjacent columns in the same row in the generated read addresses;

calculating a second minimum distance between addresses indicating a last column of a row and an address indicating a first column of the next row in the generated read addresses; and repeating the steps of arranging, P-BRO interleaving, calculating the first minimum distance and calculating the second minimum distance until the m and J values, that maximize one of the first minimum distance and the second minimum distance, are determined.

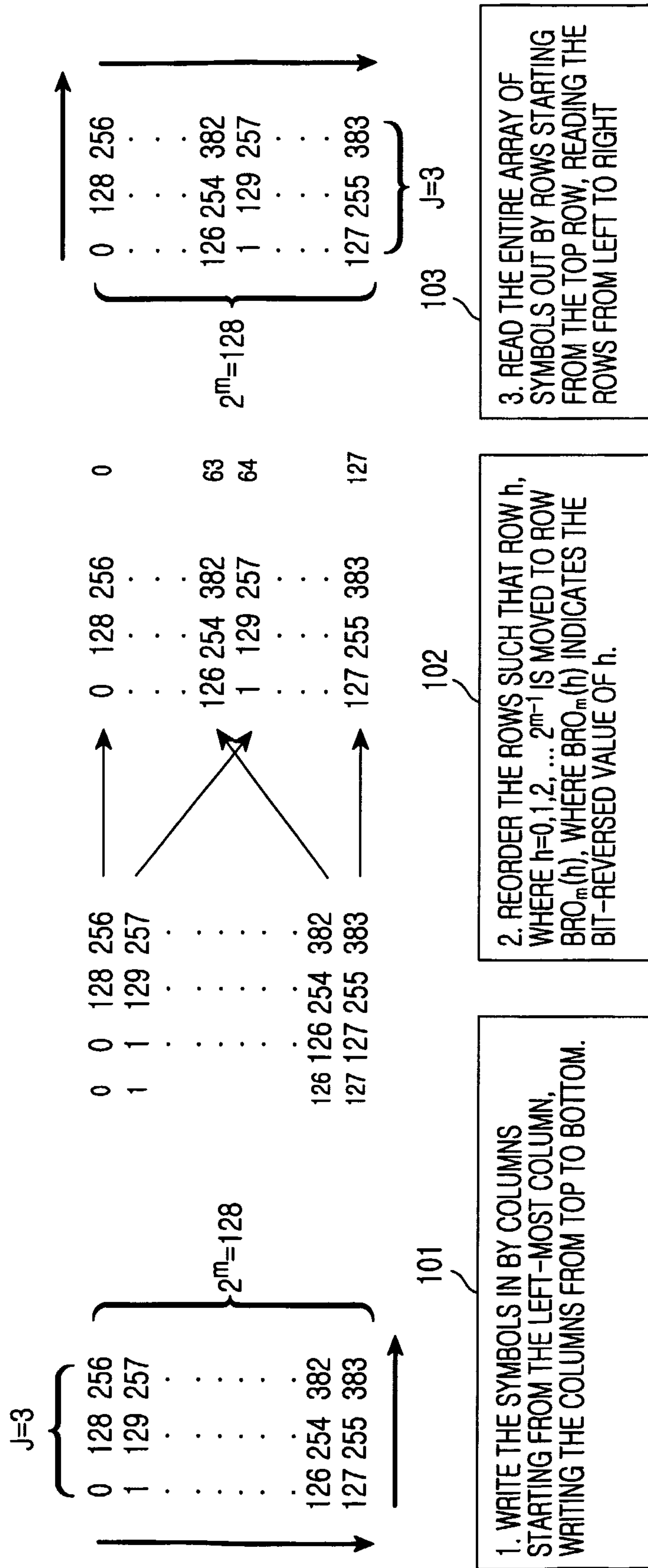


FIG. 1

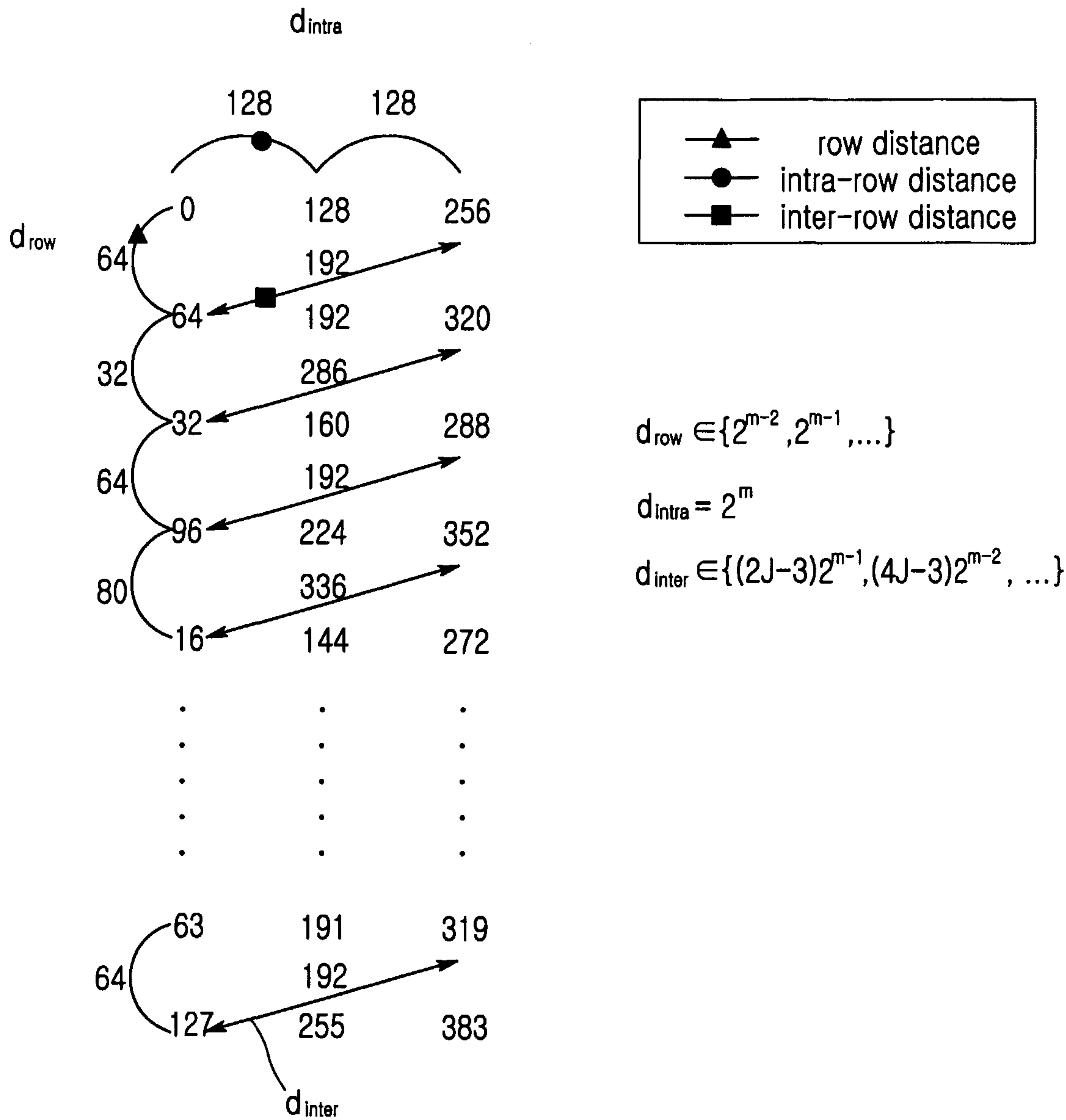
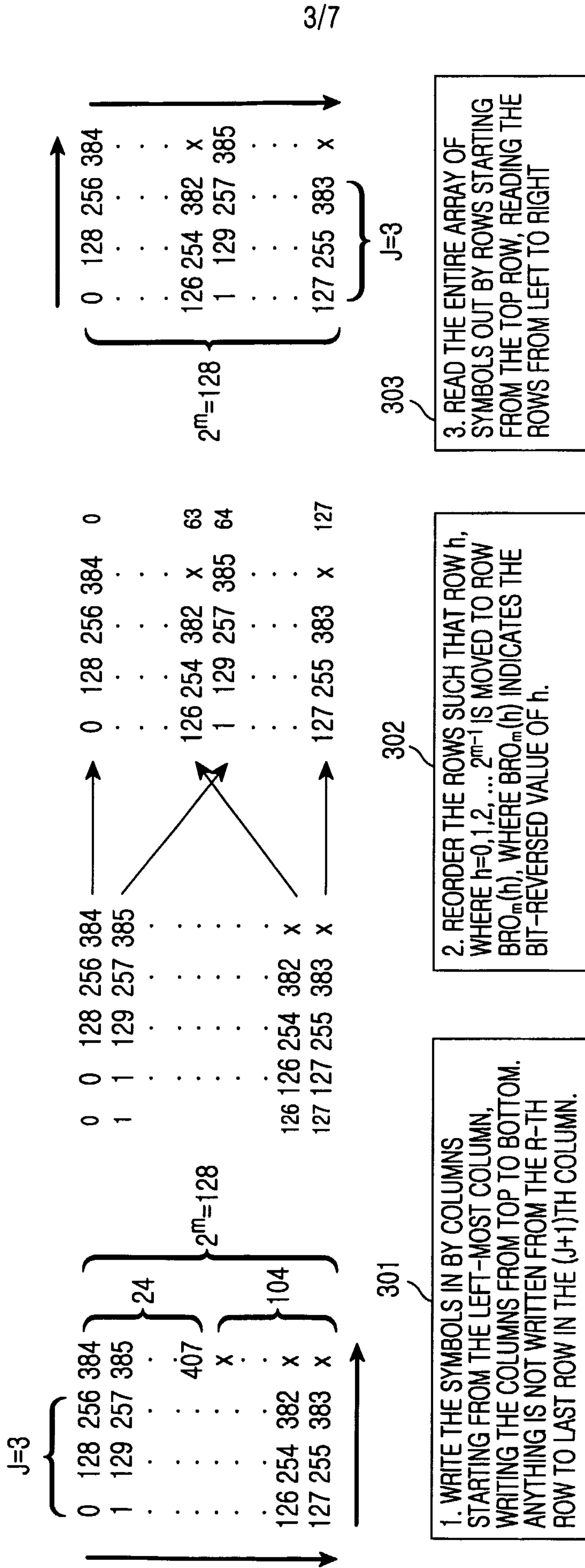


FIG. 2



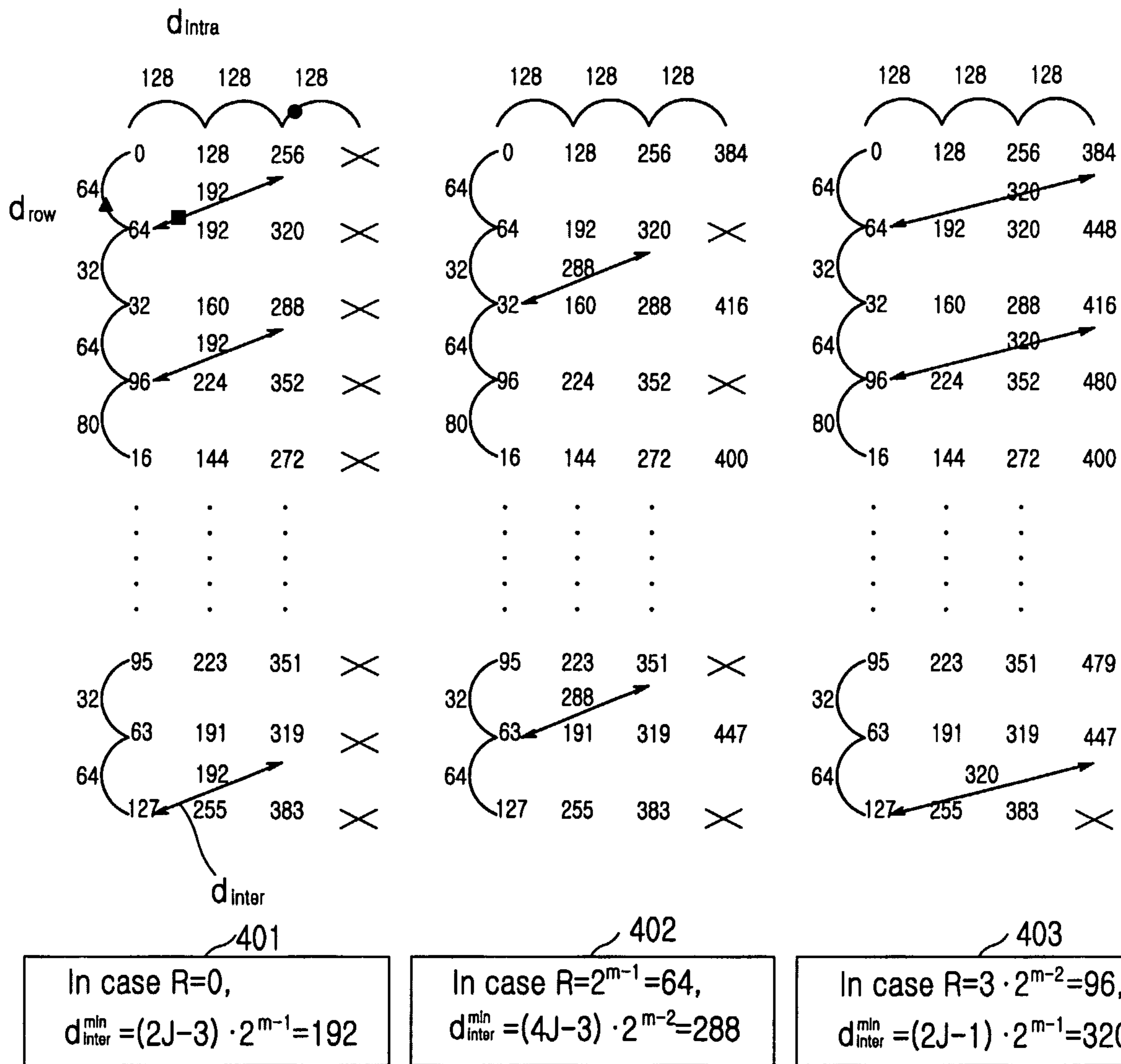


FIG. 4

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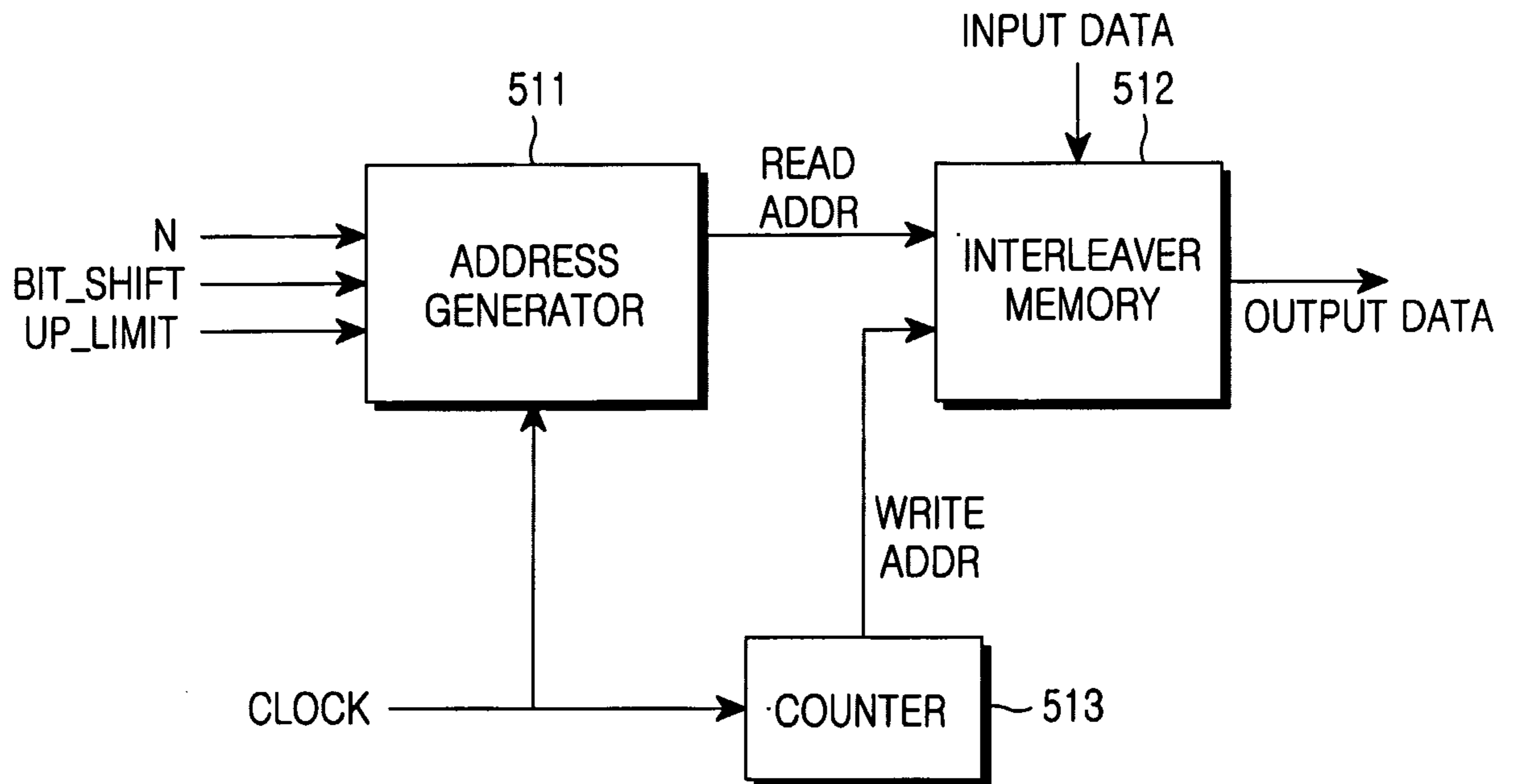


FIG. 5

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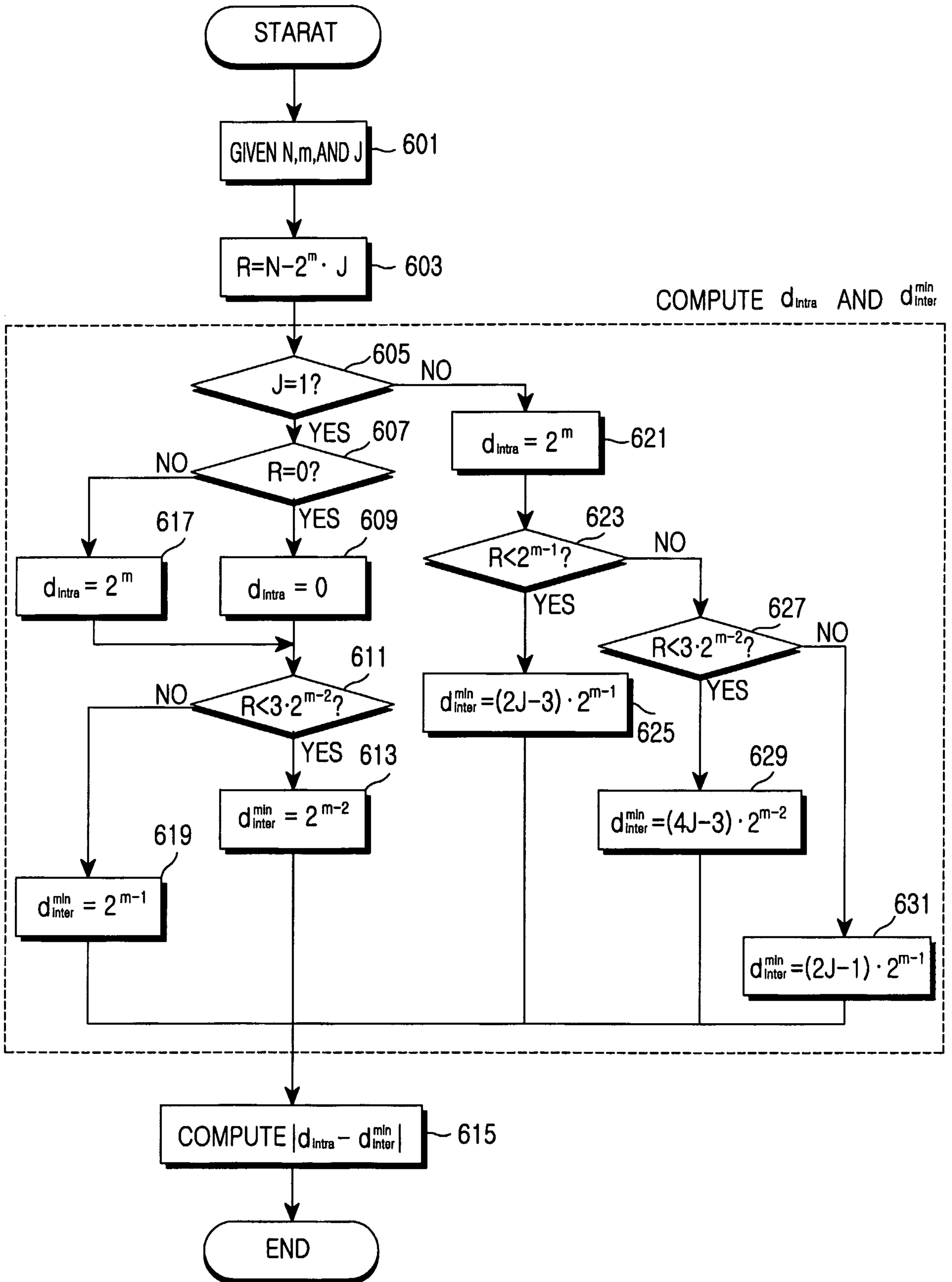


FIG. 6

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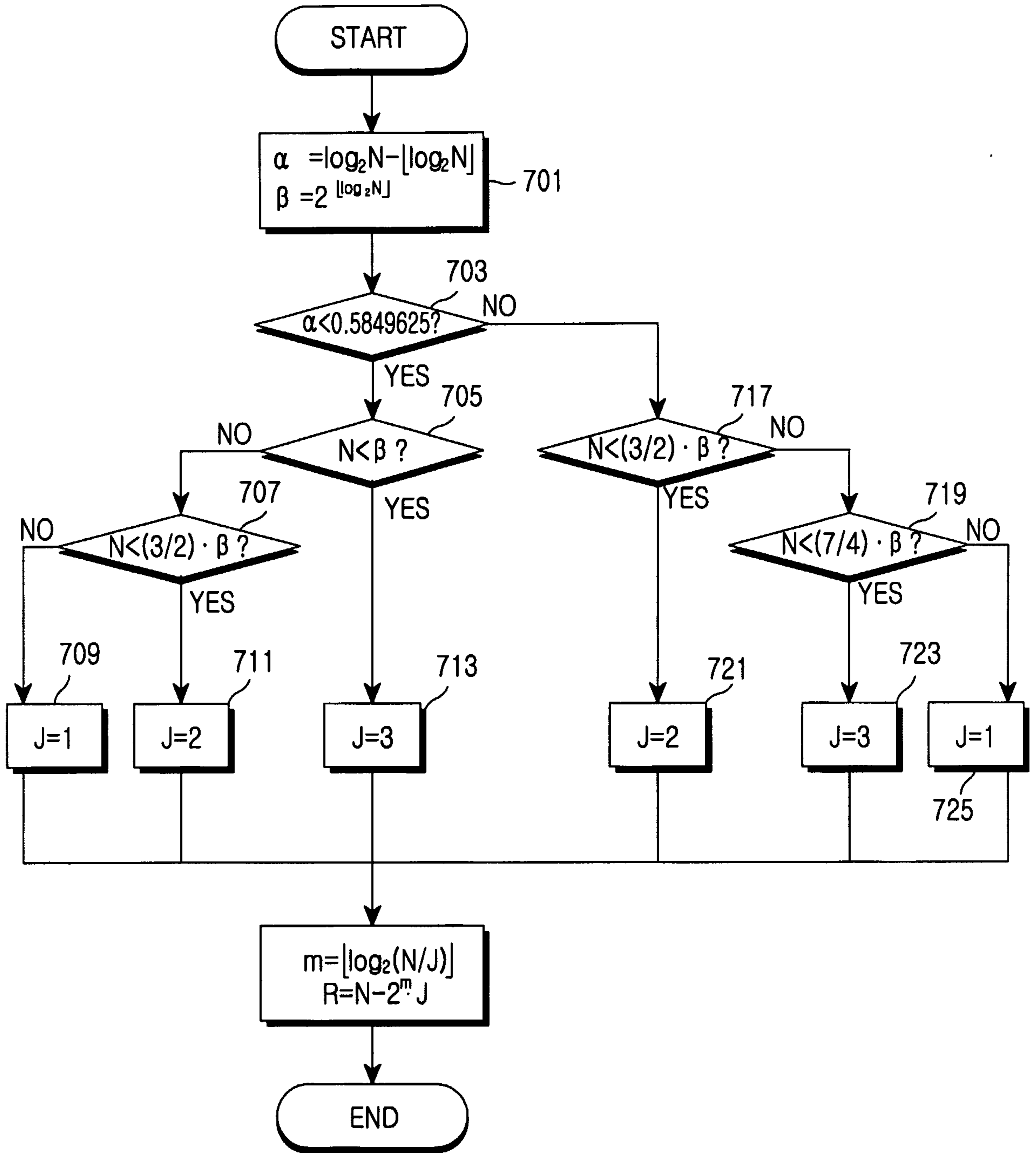


FIG. 7

STARAT

GIVEN N,m,AND J 601

$R = N - 2^m \cdot J$ 603

COMPUTE d_{intra} AND d_{inter}^{min}

J=1? 605

NO

YES 607

$d_{intra} = 2^m$ 621

R=0? 609

NO

$d_{intra} = 2^m$ 617

$d_{intra} = 0$ 609

$R < 2^{m-1}$? 623

NO

YES

$d_{inter}^{min} = (2J-3) \cdot 2^{m-1}$ 625

$R < 3 \cdot 2^{m-2}$? 627

NO

YES

$d_{inter}^{min} = (4J-3) \cdot 2^{m-2}$ 629

$R < 3 \cdot 2^{m-2}$? 611

NO

YES 613

$d_{inter}^{min} = 2^{m-2}$ 613

619

$d_{inter}^{min} = 2^{m-1}$ 619

$d_{inter}^{min} = (2J-1) \cdot 2^{m-1}$ 631

COMPUTE $|d_{intra} - d_{inter}^{min}|$ 615

END