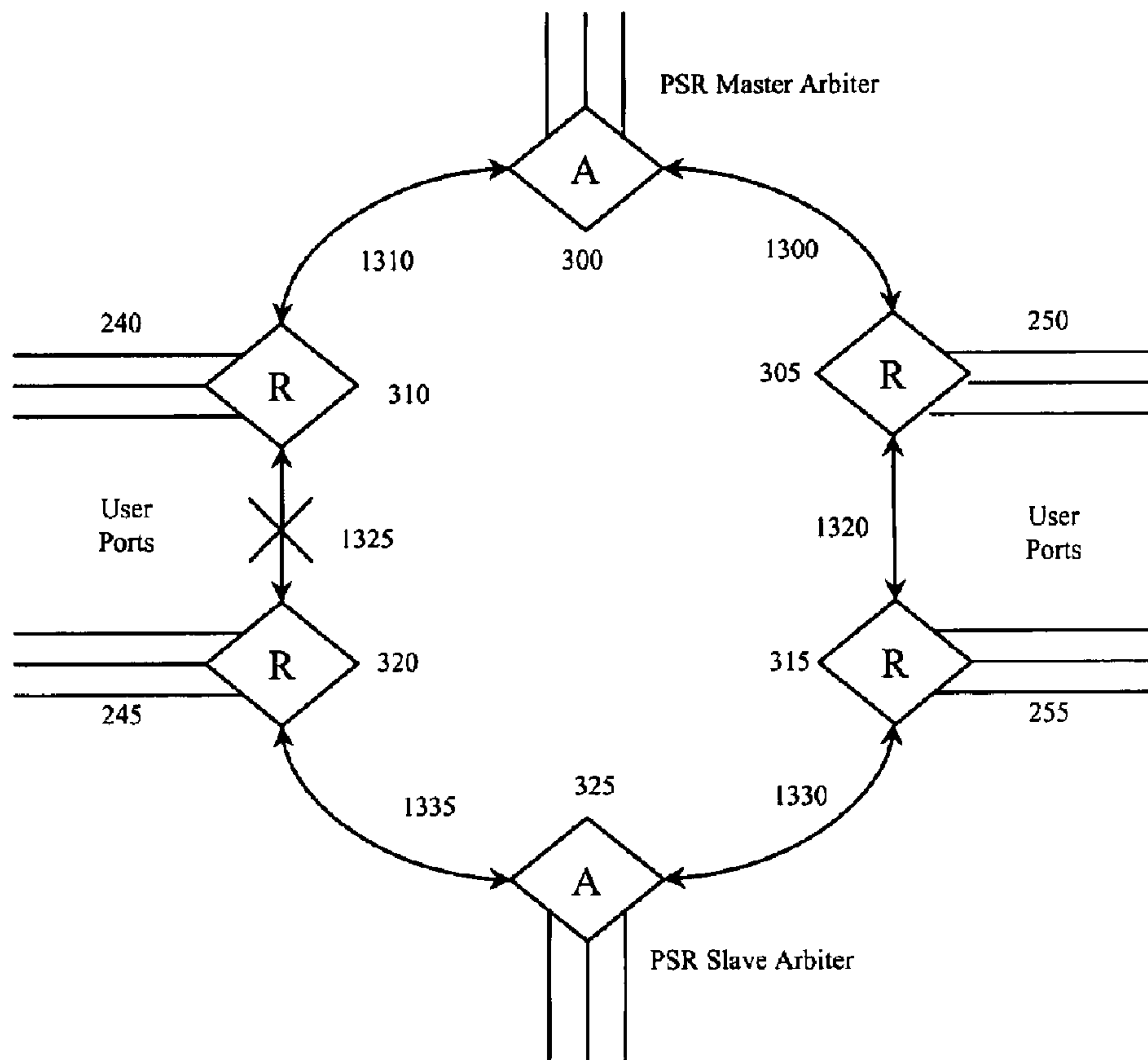




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(54) Titre : ANNEAU DE COMMUTATION PROTEGE
(54) Title: PROTECTED SWITCHING RING



(57) Abrégé/Abstract:

The protected switching ring architecture has distributed processing architecture enables ethernet to be switched between nodes over a loop. The architecture forces a tree structure by removing a link between the nodes in the loop. The architecture reconfigures links between nodes as shown in Fig 5.



ABSTRACT

Normal 802.3 Ethernet requires a tree topology. If a ring or a loop exists, then packets will be forwarded around the ring indefinitely. If the ring is broken, then there is no possibility of packets being propagated forever. This invention shows how to quickly impose a virtual break in the ring such that all nodes can communicate with each other, and how to remove the virtual break when a real failure occurs. This is accomplished by placing intelligent nodes on the ring that work together to virtually break and restore the ring. An embodiment is disclosed that handles a unidirectional break in a communication link.

(01) **PROTECTED SWITCHING RING**

FIELD OF THE INVENTION

(02) This invention relates to communication networks, and more particularly, to an
5 apparatus and method for Ethernet equipment in a ring topology.

BACKGROUND OF THE INVENTION

(03) As Ethernet is deployed in Metro and Access networks, and services are offered
on these networks, there is a desire to maintain SONET-like resiliency (i.e. recover from
10 a fault in less than 50 ms). One common means of providing resiliency is with a ring
topology. However, Ethernet switches will not work properly if there is a ring or loop in
the topology. Protocols such as IEEE 802.1d Spanning Tree Protocol (STP) or IEEE
802.1w Rapid Reconfiguration were invented to detect and remove loops. However, they
are slow and cannot achieve path restoral within 50 ms as desired.

15 (04) To solve this problem, the IEEE is working on 802.17 Resilient Packet Ring
(RPR). Others are looking at Multiprotocol Label Switching (MPLS) with Fast Reroute
capabilities. Both of these approaches are quite complex. RPR requires a new Media
Access Control (MAC) Layer, and MPLS requires extensive signaling. Because of the
complexities, these approaches will drive up the cost of the nodes on the ring.

20 (05) This invention introduces a new way (Protected Switching Ring or "PSR") of
providing protection for Ethernet deployed in a ring topology with resiliency that does
not require a new MAC layer, and that can be built using low cost Ethernet chips and
methods.

25 (06) This invention differs from some previous inventions. One of interest is described
in US patent number 6,430,151, granted on August 6, 2002. PSR is similar to '151 in
that:

- Both are based on nodes arranged in a ring topology.
- Both aim to overcome the limitations of STP.
- Both describe making or breaking a ring based on the passage or blockage of test
30 messages.
- Both have two classes of nodes on the ring, one of which is a controller or master.

(07) Some of the differences between PSR and the teachings of the '151 include:

- '151 is composed of bridging nodes that do dynamic layer 2 learning, while PSR is composed of nodes that are configured to switch (add and/or drop) packets based on a VLAN tag.
- 5 • '151 has a single redundancy manager (RM), while PSR can support dual redundancy Ring Arbiters (RA). PSR can operate in the presence of a failed RA, thus providing a higher level of availability.
- The nodes in '151 learn an association between ports and MAC addresses for ring traffic. When the topology changes, the RM of the '151 must send a "flush" message to tell the nodes to clear their databases. In contrast, the Ring Relay ("RR") nodes in PSR always send messages (both data and control) around the ring in both directions, thus removing half of the propagation delay from the recovery time. In this way, a flush command is not needed to redirect traffic on the ring, thus reducing the recovery time.
- 10 • '151 can cause packets to be duplicated during a restoral as there will be a ring upon restoral. Duplication of packets violates the IEEE 802.3 specifications. The state machines in PSR prevent this.
- Since nodes in PSR are not performing learning for ring traffic, there is less overhead and a higher packet rate can be sustained for a given amount of processing power.
- 20

(08) Another approach to the problem is described in US patent 4,354,267. The '267 patent describes a set of homogeneous layer 2 devices arranged in a ring. Each node in '267 forwards packets around the ring, and the originator removes the packet.

(09) Some of the differences between PSR and the teachings of the '267 include:

- 25 • '267 assumes that data sent that is sent one way around the ring makes it all the way around. In layer 2 systems, each node may pick off packets addressed to it, so this assumption is not valid.
- '267 assumes that each node can repair a fault. See claim 1 in column 10, starting at line 34, and claim 5, in column 12, starting at line 38. In contrast, PSR concentrates the recovery mechanism in just two nodes.
- 30

SUMMARY OF THE DISCLOSURE

(10) Normal 802.3 Ethernet requires a tree topology. If a ring or a loop exists, then packets will be forwarded around the ring indefinitely. STP was created to solve this

problem by detecting and breaking any rings. If the ring is broken, then there is no possibility of packets being propagated forever.

(11) This invention shows how to virtually break the ring such that all nodes can communicate with each other, and how to remove the virtual break when a real failure occurs. This is accomplished by placing intelligent nodes on the ring that work together to virtually break and restore the ring.

(12) In PSR, the nodes communicate between and among themselves to determine when and where a break occurs. The relevant state machines for a preferred embodiment of the present invention are contained within this disclosure.

10 BRIEF DESCRIPTION OF THE DRAWINGS

Figure 1 shows an example of prior art.

Figure 2 shows an example Protected Switching Ring in the Full Ring mode in normal operation.

Figure 3 shows an example Protected Switching Ring in the Full Ring mode during a failure.

Figure 4 shows an example Protected Switching Ring in the High Availability mode in normal operation.

Figure 5 shows an example Protected Switching Ring in the High Availability mode during a failure.

Figure 6 shows the state machine for a Ring Arbiter node in the Full Ring mode.

Figure 7 shows the state machine for a Ring Relay node in the Full Ring mode.

Figure 8 shows the state machine for the Ring Side of a Ring Relay node in the High Availability mode.

Figure 9 shows the state machine for the Extension Side of a Ring Relay node in the High Availability mode.

Figure 10 illustrates a unidirectional ring break.

Figure 11 shows the "Dual Homing" embodiment providing User Ports 1140 with redundant links to the existing network.

INDUSTRIAL APPLICABILITY AND

30 MODES FOR CARRYING OUT THE INVENTION

(13) The present invention will now be described more fully hereinafter with reference to the accompanying drawings, in order to disclose selected embodiments. This invention

may, however, be embodied in many different forms and should not be construed as limited to the embodiments set forth herein; rather, these embodiments are provided so that this disclosure will be thorough and complete, and will fully convey the scope of the invention to those skilled in the art. Like numbers refer to like elements throughout.

5 **Overview**

(14) The Protected Switching Ring (PSR) enables building of partial or full ring topologies from low-cost Ethernet equipment, while providing for sub-50 millisecond recovery from equipment or link faults. PSR nodes support the transport of point-to-point port-switched connections across the ring topology. During normal, non-fault operation,
10 one port in the ring will be blocked to user traffic, thus preventing a loop. In the event of a fault in the ring, the blocked port will be unblocked, allowing access to all nodes on the ring.

(15) Two topologies using the present invention are described below. The first topology is the PSR Full Ring ("FR") configuration that consists of a full ring of PSR
15 nodes. Port-switch connections can be configured between any two subscriber ports on the ring. The second topology is the PSR High-Availability (HA) configuration. This configuration provides a partial-ring extension of a SONET or RPR ring, or a partial-ring addition to existing layer-2 switching equipment. In either case, a path is engineered through the existing equipment to complete the path for the PSR protocol traffic and user
20 data.

(16) PSR nodes are designated as Ring Arbiters or Relays. Each ring contains two Ring Arbiters. The Ring Arbiters communicate with a "hello" protocol to coordinate the blocking or forwarding of user traffic. In a preferred embodiment, the PSR Ring Arbiter ports take on the role of master or slave on the ring according to their relative node
25 priority. In a preferred embodiment, the priority could be a unique identifier, such as a MAC address. In a highly preferred embodiment, the priority can be the concatenation of an operator-configurable priority with the MAC address (or other unique identifier) such that the priorities of two nodes would never be equal. In either of these preferred embodiments, the reception of a HELLO with the same priority would indicate a ring
30 with only one arbiter, where that arbiter was receiving its own HELLOs.

(17) In general, during normal fault-free operation of two Ring Arbiters, the slave Ring Arbiter will block one of its ring ports in order to terminate the ring loop. A ring may contain one or more Relay nodes. The Relay nodes may be distributed in any fashion

around the ring, although some benefit is provided by distributing approximately equal numbers of Relays on each “side” of a full-ring configuration.

(18) In addition to the “hello” protocol, each node performs a “discovery” protocol that allows each node to know about all the other nodes on the ring. The discovery protocol is also used to detect persistent ring faults and to generate the associated alarms. Both protocols operate at layer 2, employing reserved multi-cast MAC addresses.

(19) IP connectivity between all ring nodes is accomplished over a control VLAN used only for that purpose. This allows Telnet and a UDP-based signaling protocol to operate between any nodes on the ring. (An explanation of Telnet is not critical to the understanding of the present invention but Telnet is a terminal emulation program used with TCP/IP networks that allows remote entry of commands that are treated as if input at the network device.) Bridging techniques are used to provide the connectivity for these IP-based applications; all user traffic is transported across the ring using port-switching. As such, all user traffic is point-to-point across the ring; traffic from a subscriber-port/VLAN on one node is connected to a subscriber-port/VLAN on another node.

(20) An additional embodiment of the present invention addresses a partial failure of a network link so that the communication link is lost in only one direction. Yet, another embodiment uses a single arbiter to provide a high reliability connection of user ports to an existing network ring by creating a switching ring with the arbiter and two network ring access points.

Ring Nomenclature

(21) When the PSR is configured, two ports are designated as the ring ports and may be called East and West ports. Also the node type is given to distinguish Ring Arbiter types and Relays (also called Ring Relays or Relay Nodes). The Ring Arbiter type may be High-Availability (HA) or Full-Ring (FR). The two Ring Arbiters on the ring must be of the same type. When speaking of a specific ring port, the partner port refers to the other port of the pair of ring ports on that Ring Arbiter or Ring Relay.

(22) An additional distinction is made in the case of a HA Ring Arbiter. The port of the HA Ring Arbiter connected to the existing SONET or RPR ring is designated the “extension side” (ES) port. This port interfaces with the existing equipment for which we wish to extend a ring segment. The other Ring Arbiter port is referred to as the “ring side” port. It is connected to a string of one or more Ring Relays or directly to the other Ring Arbiter.

HELLO Protocol

(23) Each PSR Ring Arbiter periodically issues a “HELLO” protocol packet out each ring port. In a preferred embodiment each PSR Ring Arbiter issues a “HELLO” protocol packet out each ring port every 10 milliseconds. The packet uses a special multicast MAC address as the destination address. The Relay nodes are configured to have the data plane pass the packet from one ring port to the other, so a Relay node adds only a small amount of delay as the packet moves from one Ring Arbiter to the other. The remote Ring Arbiter node will terminate the packet and send the packet to the control plane. The control plane uses the presence of the new packet and some control information to drive its state machine for the Ring Arbiter ports. The absence of a new HELLO message for 30 milliseconds constitutes a *ring timeout*. If the timeout persists for 1.5 seconds, a *ring failure* is declared and the appropriate alarm is issued.

(24) The significantly longer period used as a trigger for a ring failure keeps a short intermittent problem from being deemed ring failures though the problems may be handled by the declaration of ring timeouts. In one embodiment, the ring failure is detected by loss of Discovery messages, described below. One of skill in the art could implement the ring failure to be based on the absence of HELLO messages rather than Discovery messages. One of skill in the art would appreciate that the HELLOs are not processed at the RR nodes, whereas the Discovery messages are. HELLOs therefore propagate around the ring faster than Discovery messages. A timeout threshold for loss of HELLOs can be set lower than an equivalent threshold for Discovery messages.

(25) A ring timeout causes the state machines to transition a slave Ring Arbiter port to a FORWARDING state. This response ensures that any loss in connectivity due to a single failure across the ring will only persist for 50 milliseconds or less.

(26) In a preferred embodiment the sequence number in the HELLO PDU is used at the receiving Ring Arbiter to distinguish the arrival of a new HELLO PDU. Those of skill in the art will recognize that other methods could be employed to detect the arrival of a new HELLO PDU. The Relay nodes do not process the HELLO PDUs; they only forward them between ring ports.

Discovery Protocol

(27) The discovery protocol is an optional protocol that can be implemented in order to add functionality. Note since the discovery protocol is not a necessary requirement of the state machines for any of the Ring Arbiters, Protected Switching Rings in accordance

with the teachings of the present invention could be implemented without implementing the discovery protocol.

(28) The discovery protocol also uses a special multicast destination MAC, but runs every 500 milliseconds. The discovery PDU is originated by the Ring Arbiters, appended to by intervening Relay nodes, and terminated at the remote Ring Arbiter. As the discovery PDU traverses the path between Ring Arbiters, each node in the path appends its management IP address, egress port for the PDU, and node type to the PDU. Since the discovery messages are flowing in both directions on the ring, each node on the ring can see the path of nodes to each Ring Arbiter on each of its ring ports. For example in Figure 2, the Ring port 210 will receive a discovery message on one port directly from the RA 200 and will receive the other discovery message from the RA 225 after that discovery message passes through the ring port 220. Thus after receiving the two discovery PDUs, each ring port knows the identity of all devices between the ring port and each RA.

(29) Additionally, as each Ring Arbiter constructs the discovery message to send out a ring port, the Ring Arbiter adds the completed node list received at its partner port. This allows every node in the PSR to know all the IP addresses of the nodes in the ring.

(30) In the event of a ring or node failure, the Relay nodes closest to the point of failure will originate the discovery message. In other words, if a relay fails to receive a discovery PDU from its upstream neighbor (due to a link or node failure), then the relay will create and send a discovery PDU in the downstream direction. All downstream nodes will detect that the Ring Arbiter is no longer the originator of the discovery message and declare a fault alarm. If a node either does not receive a Discovery message or receives a Discovery message without a ring Arbiter as the originator, a ring failure is declared after 1.5 seconds. The fault is cleared when the node receives a Discovery message with a ring Arbiter as the originator.

PSR Data Plane for User Traffic

(31) User traffic may enter and leave the PSR at any Ring Arbiter or relay node. A PSR connection defines the entry and exit points for a full-duplex flow of user traffic across the ring. This flow is defined by a pair of port/VLAN ID/PSR Node Address tuples. The connection defines a path through the ring between 2 user ports, each residing on a PSR node, configured to carry the user traffic for specific or all VLAN IDs on that port.

(32) As the user traffic enters the ring, a *ring tag* is added to the packet. The ring tag is a VLAN tag and is unique on the ring. The ring tag defines a given connection between two ring nodes. At the egress node of the PSR connection, the ring tag is removed from the frame before forwarding to the user port. In this way, the VLAN tags present in the user data are transparently transported across the ring. VLAN IDs used on one user port do not interfere with IDs used on another user port.

(33) A PSR node is either an *endpoint* of a given connection or a *transit node* for that connection. A PSR node is an endpoint for a connection if one of its user ports is specified in the definition of the given connection. The node is a transit node if neither endpoint of the connection resides on the node. In either case, a *switch table* used by the data plane is configured on each PSR node to either terminate one end of a given connection or to act as a transit node for that connection. When a node is a transit node for a given connection, the node simply transfers frames from one ring port to the other, based on the ring tag, without modification. When a node is an endpoint node for a given connection, the data plane directs the data arriving on a ring port to the correct user port and removes the ring tag. Conversely, the node's data plane directs user packets from the given user port with the given VLAN ID to the ring ports, adding the correct ring tag in the process.

PSR Control Plane for Control Traffic

(34) A PSR requires a mechanism to transport HELLO PDUs, discovery PDUs, and IP traffic for ring control applications between PSR nodes. While user traffic transport is transported using switching techniques, in a preferred embodiment the control functions are transported using bridging techniques. By using bridging techniques, full PSR node control connectivity is attained with all nodes appearing on the same IP subnet. This makes configuration much simpler.

(35) One ring tag is reserved for PSR control traffic. The data plane uses learning procedures and forwarding table lookups to direct control traffic to the correct PSR node. Note that the use of learning procedures and forwarding table lookups for the direct control traffic imposes an overhead burden that is orders of magnitude smaller than the overhead needed to use learning techniques for user data traffic. In the preferred embodiment the HELLO and discovery messages use known multicast MAC addresses and thus do not add additional learned database entries to be forwarded. Flushing is not

needed for the control traffic upon failure, recovery, or reconfiguration of the ring, as the new port entries are learned from bidirectional traffic after a short period of time.

(36) While the use of bridging for control traffic is preferred, it is not required in order to implement the present invention. The present invention could be implemented to use switching techniques for data packets and some or all types of control traffic. Care must be taken in creating this variation that the control traffic described in this application as passing when data packets are blocked, must be allowed to pass.

Example Recovery for Full Ring

Fault Detection

(37) Figure 3 shows a full ring where the link 1325 fails between nodes 310 and 320. This means that RA nodes 300 and 325 are unable to communicate with each other via the left hand side of the ring. Prior to the failure, assume that RA node 325, the slave Ring Arbiter, is blocking traffic on link 1330 (thus no counterclockwise communication on 1330) and forwarding traffic on link 1335. Also, any user traffic arriving on link 1330 is discarded. So clockwise traffic on 1330 is discarded at the 1330 side of RA 325. Communications to subscriber ports connected to RA 325 reach those ports through counterclockwise communication over link 1335 to RA 325.

(38) Assuming RA node 300 was the master Ring Arbiter, when RA node 325 detects the loss of communication; RA node 325 will start forwarding traffic to the right hand side of the ring onto link 1330 and accepting user traffic arriving on link 1330 and relaying the traffic to link 1335 and to the subscriber ports of RA 325. This will restore communications between all of the nodes on the ring. At this point, RA 325 is forwarding traffic on both ring ports. The ring port that is facing link 1335 is in MASTER FORWARDING state, and the ring port that is facing link 1330 is in SLAVE FORWARDING state.

Link Restoral

(39) When link 1325 is restored, RA node 325 needs to block one of its ring ports to prevent a loop in the ring. When RA node 325 receives the first HELLO on link 1335 (due to the restoration of link 1325), RA node 325 will see that the partner port to the port that is facing link 1335 is in SLAVE FORWARDING state. RA node 325 will move the port that is facing link 1335 to the BLOCKING state. Assuming that the Ring ports of nodes 310 and 320 connected to link 1325 went to an OPER DOWN state during the failure, the TIMING state in the relay nodes 310 and 320 will prevent forwarding of

traffic until the slave Ring Arbiter has time to switch from MASTER FORWARDING to BLOCKING on the 325 side of the Ring Arbiter. OPER DOWN is an indication from the physical or transport layer that a link is not operational. It is normally based on the detection of loss or corruption of the incoming electrical or optical signal.

- 5 (40) The advance to the TIMING state is triggered by the reception of a HELLO message. This TIMING state delay in the resumption of operation of relay nodes 310 and 320 prevents duplicate packets from reaching a given destination when the failed link is restored. To illustrate the value of this delay in the Ring Relay ports, consider a message coming to Ring Relay 305 to a subscriber port connected with Ring Relay 310 just before
10 link 1325 is restored. Ring Relay 305 operating normally will send the same message onto link 1300 and link 1320. The message traveling counterclockwise reaches Ring Relay 310 in a conventional way. The message traveling clockwise to Ring Relay 310 will pass through Ring Arbiter 325 onto link 1335 as the West Port is operating in MASTER FORWARDING. When link 1325 is restored, there is a path for a duplicate
15 message to cross link 1325 to Ring Relay 310. This potential is eliminated if the Ring Relay ports undergo a suitable delay between receipt of the first HELLO and the ultimate state of FORWARDING. Note that the HELLO messages from Ring Arbiter 300 to Ring Arbiter 325 and from Ring Arbiter 325 to Ring Arbiter 300 will pass over link 1325 as soon as it is restored as the HELLO messages are not blocked at any port in any state.
- 20 (41) The preferred embodiments disclose using a timing delay to ensure that a port progressing from OPER DOWN to operational delays sending data packets long enough for the slave arbiter to impose a virtual break. One of skill in the art will recognize that the use of the timer could be replaced by a control signal sent by the slave arbiter after it has successfully imposed the virtual break. In either case, the port does not go to fully
25 operational until after the virtual break has been imposed to preclude the creation of a ring for data packets.

Example Recovery for HA Ring

Fault Detection

- 30 (42) Figure 4 shows a HA ring under normal fault-free operation. The ES slave port 1440 is in the BLOCKING state to prevent a ring loop. Figure 5 shows a HA ring where the link 1520, between nodes 510 and 520, fails. As for the full ring case, the bidirectional failure of link 1520 means that the Ring Arbiter nodes 500 and 530 are unable to communicate over the left side (Ring Side) portion of the HA ring. Assuming

Ring Arbiter node 530 is the slave, its ES port (the facing link 1540) would be unblocked when the failure is detected. Fault detection and subsequent un-blocking of the slave Ring Arbiter port is fundamentally the same as for the full ring case described above.

5 **Link Restoral**

(43) In a preferred embodiment, the HA ring favors the Ring Side once the fault is removed. Instead of leaving the slave Ring Arbiter ES port (the port facing link 1540) in the forwarding state and blocking the Ring Side port (the port facing link 1530), the HA slave Ring Arbiter 530 always returns to a FORWARDING state on the Ring Side
10 segment and blocks the ES port.

(44) The Ring Side segment of the HA ring is favored in order to minimize HA ring traffic on the existing SONET or RPR ring as this will cut some of the user traffic on the SONET ring segment between the Ring Access Equipment as one direction will be blocked (thus counterclockwise traffic from port 1440 will be blocked while clockwise
15 traffic from 1400 will continue to travel on the SONET Ring.

Nomenclature for State Diagrams

(45) In the following descriptions, "isMaster" is based on the last received HELLO. If no HELLO was ever received on the port, then isMaster is based on the partner's last HELLO. If no HELLOs have ever been received by either port, then isMaster is "true".
20 The Boolean variable "isSlave" is the logical negation of "isMaster".

(46) The term "PartnerHelloTimeout" indicates that the partner port's age timer has timed out.

The designation "RxHello < Node" means a HELLO message has been received whose priority is less than that of the receiving node. This event would cause the receiving node
25 to consider itself a master.

Full Ring Mode – Ring Arbiter

(47) Figure 6 shows the state diagram for a RA node. Each of the two ports on an RA node runs a copy of this state machine.

Description of States

30 (48) The state machine of Figure 6 has the following states.

TABLE A

Number	State	Description
600	PORT DOWN	The port is operationally down or has just been initialized Entered from any state.
610	BLOCKING	The node is sending HELLOs, but not forwarding data traffic.
620	SLAVE TIMING	Node knows that it is a slave, but port is waiting for a timer to expire before moving to a forwarding state.
630	MASTER TIMING	Node knows that it is a master, but port is waiting for a timer to expire before moving to a forwarding state.
640	SLAVE FORWARDING	The port on a Slave Node is forwarding packets
650	MASTER FORWARDING	The port on a Master Node is forwarding packets

Description of State Transitions

5 (49) The table below describes the transitions of the state machine shown in Figure 6.

(50) Note the fd timer reference below runs using a time that is a small fraction of the time used for the age timers in the RA and Relay nodes. This ensures that the relays are timed for a period long enough after the expiration of the fd timer for the loop to be broken. For example, one embodiment uses a 10 millisecond timer for the RA and Relay
10 nodes and the fd timer at just one "tick" (a single 10 millisecond delay). This 10 millisecond interval is a small fraction of the 30 millisecond interval used to detect a ring timeout when a new HELLO message does not arrive within that interval.

(51) Note that the state machine for ring arbiters in the preferred embodiment does not wait indefinitely to see a HELLO as long as the ports of the ring arbiter are operationally
15 UP. The goal is to let the parts of the network ring operate even if other portions of the network ring cannot operate.

TABLE B

Num	Event	Action
1610	port operationally down OR init	block user traffic, cancel all timers
1615	port operationally up	start age timer
1620	age timer expires OR RxHello < Node	start fd timer
1625	fd timer expires	restart age timer, forward user traffic
1630	RxHello > Node	start fd timer
1635	fd timer expires AND partner not SLAVE FORWARD	restart age timer, forward user traffic
1640	age timer expires OR RxHello < Node	restart fd, age timer
1645	RxHello > Node	restart fd timer
1650	Age timer expires OR RxHello < Node	restart age timer
1655	RxHello > Node AND partner not SLAVE FORWARD	restart age timer
1660	RxHello > Node AND partner SLAVE FORWARD	restart age timer

(52) The following table shows the complete state transitions for the Full-Ring Arbiter as shown in Figure 6.

TABLE C

Current State	PORT DOWN 600	BLOCKED 610	SLAVE TIMING 620	MASTER TIMING 630	SLAVE FORWARDING 640	MASTER FORWARDING 650
Event						
Oper Down	N/A	PORT DOWN	PORT DOWN	PORT DOWN	PORT DOWN	PORT DOWN
Oper Up	BLOCKED	N/A	N/A	N/A	N/A	N/A
Age Timer Expires	N/A	MASTER TIMING	MASTER TIMING	MASTER TIMING	MASTER FORWARDING	MASTER FORWARDING
fd Timer Expires	N/A	N/A	N/A	MASTER FORWARDING	N/A	N/A
fd timer Expires AND Partner SLAVE FORWARDING	N/A	N/A	SLAVE TIMING	N/A	N/A	N/A
fd Timer Expires	N/A	N/A	SLAVE	N/A	N/A	N/A

Current State						
AND Partner not SLAVE FORWARDING			FORWARDING			
RxHello < Node	N/A	MASTER TIMING	MASTER TIMING	MASTER TIMING	MASTER FORWARDING	MASTER FORWARDING
RxHello ≥ Node	N/A	SLAVE TIMING	SLAVE TIMING	SLAVE TIMING	N/A	N/A
RxHello ≥ Node AND Partner SLAVE FORWARDING	N/A	N/A	N/A	N/A	N/A	BLOCKED
RxHello ≥ Node AND Partner not SLAVE FORWARDING	N/A	N/A	N/A	N/A	SLAVE FORWARDING	SLAVE FORWARDING

(53) In a preferred embodiment, every 10 milliseconds, the two ports are checked in the same order. The combination of variations in when the HELLOs were generated plus transit delays may cause one HELLO on one port to arrive before the other HELLO on the other port. In any case, since one port is checked before the other then the other, it always appears as though one HELLO arrives before the other. The order that the ports are checked does affect which slave port is set to BLOCKING on the full ring.

(54) One of skill in the art will recognize that any embodiment that does not check one port before the other would need to address the case of two HELLOs arriving essentially simultaneously at the two ports.

TABLE D

Time	Port A	Input	State Change	Port B	Input	State Change
1	Port Down	Port up	1615 to Blocking	Port Down	Port up	1615 to Blocking
HELLOs generated by other RA and sent towards ports A and B of this RA. One HELLO arrives slightly before the other.						
2	Blocking	RxHello >node	1630 to Slave Timing	Blocking		
3	Slave Timing			Blocking	RxHello >node	1630 to Slave Timing
4	Slave Timing	fd timer expires and partner not SLAVE FORWARD	1635 to SLAVE FORWARD	Slave Timing		
5	Slave forward			Slave Timing	[cannot advance to Slave Forward as partner is in Slave Forward]	

Time	Port A	Input	State Change	Port B	Input	State Change
6	Link breaks					
7	Slave Forward	Link breaks, age timer expires	1650 to Master Forwarding	Slave Timing		
8	Master Forward			Slave Timing	fd timer expires and partner not SLAVE FORWARD	1635 to SLAVE FORWARD
9	Master Forward			Slave Forward		
10	Link Restored					
11	Master Forward	RxHello > Node and partner Slave Forward	1660 to Blocking	Slave Forward		
12	Blocking	RxHello > Node	1630 to Slave Timing	Slave Forward		
This continues until a port goes down, a link goes down, or the node number of the other RA changes to become less than Node (normally this would take a reconfiguration from the operator or the substitution of another RA unit).						

Full Ring Mode – Ring Relay

(55) Figure 7 shows the state machine for a Ring Relay node.

Description of States

5 (56) The state machine of Figure 7 has the following states.

TABLE E

Number	State	Description
700	PORT DOWN	The port is operationally down or has just been initialized. Entered from any state on an indication of the port going down due to a loss of signal or other similar alarm.
710	AWAITING HELLO	Port is operationally up, but no HELLO has been received
720	TIMING	The port is waiting for the fd timer to expire
730	FORWARDING	Normal forwarding.

Description of State Transitions

(57) The table below describes the transitions of the state machine shown in Figure 7.

TABLE F

Number	Event	Action
1710	port operationally down OR init	block user traffic, cancel all timers
1715	port operationally up	start age timer
1720	age timer expires OR RxHello	start fd timer
1725	fd timer expires	forward user traffic

High Availability Mode – Ring Arbiter – Ring Side

- 5 (58) Figure 8 shows the state machine for the Ring Side (RS) of a Ring Arbiter in HA mode.

Description of States

- (59) The state machine of Figure 8 has the following states.

TABLE G

10

Number	State	Description
800	PORT DOWN	The port is operationally down or has just been initialized. Entered from any state.
810	BLOCKING	The port is sending HELLOs, but is not forwarding data traffic.
820	SLAVE FORWARDING	The port on a Slave Node is forwarding packets
830	MASTER FORWARDING	The port on a Master Node is forwarding packets

Description of State Transitions

- (60) The table below describes the transitions of the state machine shown in Figure 8.

TABLE H

15

Number	Event	Action
1810	port operationally down OR init	block user traffic, cancel age timer
1815	port operationally up	start age timer
1820	(age timer expires AND isMaster) OR RxHello < Node	forward user traffic

Number	Event	Action
1825	(age timer expires AND isSlave) OR RxHello > Node	forward user traffic
1830	RxHello < Node	No action
1835	RxHello > Node	No action

High Availability Mode – Ring Arbiter – Extension Side

(61) Figure 9 shows the state machine for the Extension Side (ES) of a Ring Arbiter in HA mode.

5 Description of States

(62) The state machine of Figure 9 has the following states.

TABLE I

Number	State	Description
900	PORT DOWN	The port is operationally down or has just been initialized. Entered from any state.
910	BLOCKING	The node is sending HELLOs, but not forwarding data traffic.
920	SLAVE FORWARDING	The port on a Slave Node is forwarding packets
930	MASTER FORWARDING	The port on a Master Node is forwarding packets

Description of State Transitions

10 (63) The table below describes the transitions of the state machine shown in Figure 9.

TABLE J

Number	Event	Action
1910	port operationally down OR init	block user traffic, cancel age timer
1915	port operationally up	start age timer
1920	(age timer expires AND isMaster) OR RxHello < Node	forward user traffic
1925	(age timer expires AND isSlave)	forward user traffic
1930	RxHello < Node	continue forwarding user traffic
1935	RxHello > Node AND PartnerHelloTimeout	continue forwarding user traffic
1940	RxHello > Node AND NOT PartnerHelloTimeout	block user traffic, start age timer
1945	RxHello > Node AND NOT PartnerHelloTimeout	block user traffic, start age timer

(64) As shown in the sequence of events reported in the table below, the RS ports of the Arbiters are always forwarding, unless the ports are OPER DOWN. The point of interest is the state of the ES port of the slave Arbiter. In essence, the ES slave port is FORWARDING if there is a HELLO timeout on either the RS or ES.

TABLE K

T I M E	Port Status (before trigger)				Trigger	Reaction
	500 RS	500 ES	530 RS	530 ES		
1	800	900	800	900	500 initialized	500 RS goes Blocking, 500 ES Goes to Blocking
2	810	910	800	900	530 initialized	530 RS goes Blocking, 530 ES Goes to Blocking
3	810	910	810	910	500 receives HELLO from 530 and RxHello < node	500 RS state change 1820 to Master Forwarding 500 ES state change 1920 to Master Forwarding
4	810	910	830	930	530 received HELLO from 500 and RxHello > node	530 RS state change 1825 to Slave Forwarding 530 ES does not leave Blocking unless RS or ES has HELLO timeout

T I M E	Port Status (before trigger)				Trigger	Reaction
	500 RS	500 ES	530 RS	530 ES		
5	820	910	830	930		Continues operation with the virtual break in the HA ring at the ES of the slave (RA 500).
6	820	910	830	930	Break in link 1520 (ring side)	
7	820	910	830	930	RxHellos stop coming on RS	500 RS no change 530 ES state change 1925 to Slave Forwarding
8	820	920	830	930		All four ports forward traffic while there is a physical break
9	820	920	830	930	Break fixed	
10	820	920	830	930	RxHello received at 530 RS and > node	530 ES state change 1945 to blocking
11	820	920	830	930		Continues operation with virtual break.
12	820	910	830	930	Link break ES	
13	820	910	830	930	HELLOs stop on ES side of both RA units	530 ES notes that its age timer expires and it is Slave and has state change 1925 to slave forwarding
14	820	920	830	930		All four ports forward traffic while there is a physical break
15	820	920	830	930	Break fixed	
16	820	920	830	930	530 ES receives RxHellos > node and not PartnerHelloTimeout	530 ES moves along state transition 1945 to Blocking
17	820	910	830	930		Until next break, port down, or switch in node numbers sufficient to change master/slave relationship.

ALTERNATIVE EMBODIMENTS

Unidirectional Break

5 (65) The control system described above assumes that a break in a network ring will be a bidirectional break as it connects both the clockwise and counterclockwise virtual breaks upon failure to receive a HELLO. This bidirectional response could cause a loop in the event of a unidirectional failure.

10 (66) Figure 10, adds additional detail to the drawing shown in Figure 2. More specifically, the links are shown in their unidirectional components rather than as bidirectional links.

15 (67) For example, when the network ring is fully operational, Master Arbiter 1000 can receive HELLOs from Slave Arbiter 1025 via link 11037, relay 1020, link 11027, relay 1010, and link 11012. Likewise, Slave Arbiter 1025 can receive HELLOs from Master Arbiter 1000 via link 11010, relay 1010, link 11025, relay 1020, and link 11035.

- (68) If link 11027 was cut but link 11025 was left in service, then the West port on Master Arbiter 1000 would soon stop receiving HELLOs from Slave Arbiter 1025, while Slave Arbiter 1025 continued to receive HELLOS from Master Arbiter 1000. In the previously described embodiment, this unidirectional cut at link 11027 would not trigger the Slave Arbiter 1025 to unblock as it continues to receive HELLOs from Master Arbiter 1000 across intact link 11025. Thus ring relay 1020 as well as connected subscriber ports 1045 would be cut off from the east side of the ring as Slave Arbiter 1025 is still blocking data, including data that would otherwise travel from Slave Arbiter 1025 to ring relay 1020.
- (69) One alternative embodiment is to react to a port going to an OPER DOWN state by stopping the transmission of HELLOs and all data from that port in the opposite direction, effectively creating a virtual unidirectional break in the other direction. Hence when ring relay 1010 observes an OPER DOWN associated with link 11027, ring relay 1010 would stop sending HELLOs and all data on link 11025. After Slave Arbiter 1025 fails to receive HELLOs in an allotted time, the Slave Arbiter 1025 would remove the virtual break on its west side to allow data traffic from link 11035 to proceed towards link 11030 or the user ports 1060 and to allow traffic from link 11032 or user ports 1060 to proceed onto link 11037..

Dual Homing Using a Single Node Ring

- (70) Figure 11 shows an application of a particular embodiment of the present invention that is referred to as "dual homing". Dual homing allows a Slave Arbiter Node 1130 to provide protected access for User Ports 1140 to network via Ring Access Equipment nodes 1110 and 1120 using redundant links 1160 and 1170.

- (71) In this alternative embodiment, the Slave Arbiter node 1130 would see its own HELLOs. As described in Table C, one side of 1130 (for example the West side of the Slave Arbiter connected to link 1160) would go to the SLAVE FORWARDING state and one side (for example, the East side of the Slave Arbiter connected to link 1170) would go to the BLOCKED state.

- (72) Now, in response to a fault on the Ring Access Equipment 1110 or the link 1160, the East side of the Slave Arbiter 1130 would unblock, and the User Ports 1140 would continue to have access to the network. The network access for User Ports 1140 is therefore protected against faults in either the access links (1160 and 1170) as well as in the Ring Access Equipment nodes (1110 and 1120).

(73) The preferred embodiments disclose using a timing delay to ensure that a port progressing from OPER DOWN to operational delays sending data packets long enough for the slave arbiter to impose a virtual break. One of skill in the art will recognize that the use of the timer could be replaced by a control signal sent by the slave arbiter after it
5 has successfully imposed the virtual break.

(74) One of skill in the art will recognize that alternative embodiments set forth above are not universally mutually exclusive and that in some cases alternative embodiments can be created that implement two or more of the variations described above.

(75) Those skilled in the art will recognize that the methods and apparatus of the
10 present invention have many applications and that the present invention is not limited to the specific examples given to promote understanding of the present invention. Moreover, the scope of the present invention covers the range of variations, modifications, and substitutes for the system components described herein, as would be known to those of skill in the art.

15 (76) The legal limitations of the scope of the claimed invention are set forth in the claims that follow and extend to cover their legal equivalents. Those unfamiliar with the legal tests for equivalency should consult a person registered to practice before the patent authority which granted this patent..

ACRONYMS

20 ES Extension Side
FR Full Ring
HA High Availability
IP Internet Protocol
MAC Media Access Control
25 MPLS Multiprotocol Label Switching
PDU Packet Data Unit
PSR Protected Switching Ring
RA Ring Arbiter
RPR Resilient Packet Ring
30 RR Ring Relay
RS Ring Side
TCP Transmission Control Protocol
UDP User Datagram Protocol

We claim:

1. A network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network carrying data packets, the network ring comprising:

A) a Master Arbiter comprising:

a MA East Port;

a MA West Port; and

B) a Slave Arbiter comprising:

a SA East Port normally connected to the MA East Port by a full duplex communication connection;

a SA West port normally connected to the MA West Port by a full duplex communication connection;

a SA control system;

the SA control system comprising controls over the introduction and removal of a virtual break to prevent the formation of a communication ring for data packets within the network requiring a tree topology through the use of at least one non-forwarding state wherein control traffic is forwarded but data packets are not forwarded; the SA control system comprising:

- at least one rule causing the operation of a SA Port in a slave-forwarding state wherein the port forwards data packets and control traffic received at the port;
- at least one rule causing the operation of a SA Port in one of at least one non-forwarding states wherein control traffic is forwarded but data packets are not forwarded;
- at least one rule causing the Slave Arbiter to operate when the ring is fully operational such that one SA Port is in a slave-forwarding state and the other SA Port is in a non-forwarding state wherein control traffic is forwarded but data packets are not forwarded;
- at least one rule for changing the state of the Slave Arbiter to have both the SA East Port and the SA West Port acting to forward data packets after a functional break in the network ring is discerned; and

- at least one rule for changing the state of one of the two SA Ports to a non-forwarding state wherein control traffic is forwarded but data packets are not forwarded after detection that the functional break has been removed.

2. The network ring of claim 1 further comprising:

C) at least one East Side Ring Relay on the ring segment between the MA East Port and the SA East Port;

the at least one East Side Ring Relay comprising two network ring ports and at least one user port in data communication with the network ring through the East Side Ring Relay;

D) a First Ring Relay Port on an East Side Ring Relay directly linked to the MA East Port, the First Ring Relay Port having a partner port on the Slave Arbiter Side of the East Side Ring Relay with the First Ring Relay Port; and

E) a Second Ring Relay port on an East Side Ring Relay directly linked to the SA East Port; the Second Ring Relay Port having a partner port on the Master Arbiter Side of the East Side Ring Relay with the Second Ring Relay Port; wherein the First Ring Relay Port is the partner port to the Second Ring Relay Port when there is only one East Side Ring Relay;

the operation of the network ring further characterized by:

the First Ring Relay Port relaying to the Master Arbiter's MA East Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

the Second Ring Relay Port relaying to the Slave Arbiter's SA East Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

each of the at least one East Side Ring Relay relaying data packets received from the at least one Subscriber Port connected to that East Side Ring Relay out both ring ports;

the First Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the First Ring Relay Port and the Master Arbiter's MA East Port;

the First Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the First Ring Relay Port and the Master Arbiter's MA East Port; the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent formation of a communication ring;

the Second Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the Second Ring Relay Port and the Slave Arbiter's SA East Port; and

the Second Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the Second Ring Relay Port and the Slave Arbiter's SA East Port, the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent the formation of a communication ring.

3. The network ring of claim 1 further comprising:

C) at least one West Side Ring Relay on the ring segment between the MA West Port and the SA West Port;

the at least one West Side Ring Relay comprising two network ring ports and at least one user port in data communication with the network ring through the West Side Ring Relay;

D) a First Ring Relay Port on a West Side Ring Relay directly linked to the MA West Port, the First Ring Relay Port having a partner port on the Slave Arbiter Side of the West Side Ring Relay with the First Ring Relay Port; and

E) a Second Ring Relay port on a West Side Ring Relay directly linked to the SA West Port; the Second Ring Relay Port having a partner port on the Master Arbiter Side of the West Side Ring Relay with the Second Ring Relay Port; wherein the First Ring Relay Port is the partner port to the Second Ring Relay Port when there is only one West Side Ring Relay;

the operation of the network ring further characterized by:

the First Ring Relay Port relaying to the Master Arbiter's MA West Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

the Second Ring Relay Port relaying to the Slave Arbiter's SA West Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

each of the at least one West Side Ring Relay relaying data packets received from the at least one Subscriber Port connected to that West Side Ring Relay out both ring ports;

the First Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the First Ring Relay Port and the Master Arbiter's MA West Port;

the First Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the First Ring Relay Port and the Master Arbiter's MA West Port; the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent formation of a communication ring;

the Second Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the Second Ring Relay Port and the Slave Arbiter's SA West Port; and

the Second Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the Second Ring Relay Port and the Slave Arbiter's SA West Port, the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent the formation of a communication ring.

4. The network ring of claim 2 further comprising:
 - F) at least one West Side Ring Relay on the ring segment between the MA West Port and the SA West Port;

the at least one West Side Ring Relay comprising two network ring ports and at least one user port in data communication with the network ring through the West Side Ring Relay;

G) a First Ring Relay Port on a West Side Ring Relay directly linked to the MA West Port, the First Ring Relay Port having a partner port on the Slave Arbiter Side of the West Side Ring Relay with the First Ring Relay Port; and

H) a Second Ring Relay Port on a West Side Ring Relay directly linked to the SA West Port; the Second Ring Relay Port having a partner port on the Master Arbiter Side of the West Side Ring Relay with the Second Ring Relay Port; wherein the First Ring Relay Port is the partner port to the Second Ring Relay Port when there is only one West Side Ring Relay;

the operation of the network ring further characterized by:

the First Ring Relay Port relaying to the Master Arbiter's MA West Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

the Second Ring Relay Port relaying to the Slave Arbiter's SA West Port data packets received from its partner port that are not directed to exit over one of the at least one user ports are sent out the first port;

each of the at least one West Side Ring Relay relaying data packets received from the at least one Subscriber Port connected to that West Side Ring Relay out both ring ports;

the First Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the First Ring Relay Port and the Master Arbiter's MA West Port;

the First Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the First Ring Relay Port and the Master Arbiter's MA West Port; the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent formation of a communication ring;

the Second Ring Relay Port ceasing forwarding data packets upon detection of a functional break between the Second Ring Relay Port and the Slave Arbiter's SA West Port; and

the Second Ring Relay Port resuming forwarding data packets after the detection of the removal of the functional break between the Second Ring Relay Port and the Slave Arbiter's SA West Port, the resumption delayed by a delay period sufficient to allow the Slave Arbiter to create a virtual break to prevent the formation of a communication ring.

5. The network ring of claim 1 wherein the Master Arbiter can operate as a slave arbiter and the Slave Arbiter can operate as a master arbiter and a set of at least one control system rules operate to limit the network ring to a single master arbiter at any one time.

6. The network ring of claim 1 wherein the Master Arbiter further comprises at least one Subscriber Port and the data packets received from the at least one Subscriber Port are transmitted out both the MA East Port and the MA West Port.

7. The network ring of claim 1 wherein:

the MA East Port sends a control traffic message of a first type that is conveyed to the SA East Port;

and the SA control system is further characterized by comprising:

at least one rule for changing the state of the Slave Arbiter to have both the SA East Port and the SA West port acting to forward data packets after a functional break in the network ring is discerned; and

a process of discernment comprising monitoring the receipt of the control traffic messages of the first type received at the SA East Port and detecting a gap in the receipt of an incoming control traffic message of the first type at the SA East Port for a period in excess of a first time interval.

8. The network ring of claim 1 wherein:
the data packets within the network ring are forwarded using a switching technique; and
at least one type of control traffic is forwarded using a bridging technique.
9. The network ring of claim 1 wherein:
the data packets within the network ring are forwarded using a switching technique; and
at least one type of control traffic is forwarded using a switching technique.
10. The network ring of claim 1 wherein the control system further comprises an extended delay so that a port recovering from a port down state does not forward data packets until after the execution of the at least one rule for changing the state of one of the two SA Ports to a non-forwarding state wherein control traffic is forwarded but data packets are not forwarded whereby the removal of a functional break does not create a temporary loop for data packets and thus provides a safeguard against the duplication of data packets.
11. A network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network carrying data packets, the network ring comprising:
 - A) a Slave Arbiter comprising:
 - a SA East Port normally connected to a first Ethernet device by a full duplex communication connection;
 - a SA West port normally connected to a second Ethernet device by a full duplex communication connection;
 - a SA control system;
 - the SA control system comprising controls over the introduction and removal of a virtual break to prevent the formation of a communication ring for data packets within the network requiring a tree topology through the use of at least one

non-forwarding state wherein control traffic is forwarded but data packets are not forwarded; the SA control system comprising:

- at least one rule causing the operation of a SA Port in a slave-forwarding state wherein the port forwards data packets and control traffic received at the port;
- at least one rule causing the operation of a SA Port in one of at least one non-forwarding states wherein control traffic is forwarded but data packets are not forwarded;
- at least one rule causing the Slave Arbiter to operate when the network ring is fully operational such that one SA Port is in a slave-forwarding state and the other SA Port is in a non-forwarding state wherein control traffic is forwarded but data packets are not forwarded;
- at least one rule for changing the state of the Slave Arbiter to have both the SA East Port and the SA West port acting to forward data packets after a functional break in the network ring is discerned; and
- at least one rule for changing the state one of the two SA Ports to a non-forwarding state wherein control traffic is forwarded but data packets are not forwarded after detection that the functional break has been removed.

12. The network ring of claim 11 wherein the first Ethernet device is also the second Ethernet device.

13. A network ring that is manipulated to have at least one virtual break for data packets; the network ring comprising:

a first arbiter acting as Master Arbiter, the Master Arbiter having at least a MA West Port, and a MA East Port;

a second arbiter acting as Slave Arbiter, the Slave Arbiter having at least a SA West Port, and a SA East Port;

a First full duplex communication path between the Master Arbiter and the Slave Arbiter;

a Second full duplex communication path between the Master Arbiter and the Slave Arbiter;

the Master Arbiter and the Slave Arbiter serving as nodes on the network ring;

the Master Arbiter sending out a HELLO protocol packet frequently onto the First full duplex path and on the Second full duplex path;

the Slave Arbiter noting the arrival of the HELLO protocol packets and noting a ring timeout when the interval between arrivals of two sequential HELLO protocol packets on a particular port exceeds a first threshold value;

the Slave Arbiter reacting to a ring timeout by removing at least one virtual break within the Slave Arbiter and communicating data packets received from the First Communication Path to the Second Communication Path and from the Second Communication Path to the First Communication Path; and

the Slave Arbiter reacts to receipt of a HELLO protocol packet from the communication path deemed to have a ring timeout by introducing at least one virtual break so that data packets received from the First Communication Path are not communicated to the Second Communication Path and the data packets received from the Second Communication Path are not communicated to the First Communication Path so that the network ring is not operated to provide a data communication loop.

14. A network ring that is manipulated to have at least one virtual break for data packets; the network ring comprising:

a first arbiter acting as Master Arbiter, the Master Arbiter having at least a MA West Port, and a MA East Port;

a second arbiter acting as Slave Arbiter, the Slave Arbiter having at least a SA West Port, and a SA East Port;

a First full duplex communication path between the Master Arbiter and the Slave Arbiter comprising at least an "A" Ring Relay and a "B" Ring Relay;

a Second full duplex communication path between the Master Arbiter and the Slave Arbiter comprising at least an Initial Ring Relay;

the A Ring Relay, the B Ring Relay and the Initial Ring Relay all comprising a Master Arbiter Side Port, a Slave Arbiter Side Port, and a Subscriber Port;

the Master Arbiter; the Slave Arbiter; and the ring relays serving as nodes on the network ring;

the Master Arbiter sending out a HELLO protocol packet frequently onto the First full duplex path and on the Second full duplex path;

the ring relays passing the HELLO protocol packet from the Master Arbiter Side Port to the Slave Arbiter Side Port;

the Slave Arbiter noting the arrival of the HELLO protocol packets and noting a ring timeout when the interval between arrivals of two sequential HELLO protocol packets on a particular port exceeds a first threshold value;

the Slave Arbiter reacting to a ring timeout by removing at least one virtual break within the Slave Arbiter and communicating data packets received from the First Communication Path to the Second Communication Path and from the Second Communication Path to the First Communication Path; and

the Slave Arbiter reacts to receipt of a HELLO protocol packet from the communication path deemed to have a ring timeout by introducing at least one virtual break so that data packets received from the First Communication Path are not communicated to the Second Communication Path and the data packets received from the Second Communication Path are not communicated to the First Communication Path so that the network ring is not operated to provide a data communication loop.

15. The network ring of claim 14 wherein the Slave Arbiter notes a ring failure when the interval between arrivals of two sequential HELLO protocol packets exceeds a second threshold value, larger than the first threshold value.

16. The network ring of claim 14 wherein the first threshold value is selected to allow the declaration of a ring timeout and removal of at the least one virtual break within the Slave Arbiter in less than 50 milliseconds.

17. The network ring of claim 14 wherein the Slave Arbiter has at least one Subscriber Port and the Slave Arbiter continues to transmit data packets received from at least one Subscriber Port onto the First Communication Path and the Second Communication Path after noting the ring timeout, but after receipt of a HELLO protocol packet from the communication path deemed to have a ring timeout, the Slave Arbiter acts to prevent subsequent retransmission of these data packets by inserting at least one virtual break within the Slave Arbiter.

18. The network ring of claim 17 wherein the nodes closest to a failure in the network ring stop transmitting received data packets towards the failure in the network ring and do not resume transmission of data packets towards the failure in the network ring after correction of the failure until after the Slave Arbiter has an opportunity to insert at least one virtual break.

19. The network ring of claim 14 wherein the HELLO protocol packets are transported within the network ring through use of bridging procedures to direct the HELLO protocol packets to particular nodes on the network ring, but data packets received from user ports are transported around the network ring through the use of a switching technique.

20. The network ring of claim 14 wherein the HELLO protocol packets and data packets are transported within the network ring through use of a switching technique to direct the packets to particular nodes on the network ring.

21. A network node device for selectively creating and removing at least one virtual break in a network ring; the node device comprising:

a first ring connection and a second ring connection for use in connecting the node device to two ends of a network ring;

the node device operation within a network ring characterized by:

initializing so that the node device contains a virtual break to prevent data packets received on the first ring connection from continuing on the

second ring connection and data packets on the second data connection from continuing on the first ring connection;

creating and transmitting a series of HELLO protocol packets out the first ring connection and the second ring connection wherein the HELLO protocol packets are addressed in such a way that the HELLO protocol packets pass through any relay nodes located between the network node device and a device intended to process HELLO protocol packets while monitoring for actual breaks in the network ring;

receiving HELLO protocol packets from a device other than itself on the first ring connection and monitoring the interval since the last HELLO protocol packet was received on the first ring connection;

receiving HELLO protocol packets from a device other than itself on the second ring connection and monitoring the interval since the last HELLO protocol packet was received on the second ring connection;

discerning the existence of a break in the network ring by detecting an interval in excess of a first value since the receipt of the last HELLO packet on the first ring connection;

discerning the existence of a break in the network ring by detecting an interval in excess of a second value since the receipt of the last HELLO packet on the second ring connection where the second value can equal the first value;

responding to the discernment of the existence of a break in the network ring by removing the at least one virtual break within the node device;

detecting the resumed receipt of a HELLO protocol packet on a ring connection used to discern a break in the network ring, and concluding the discerned break is gone; and

introducing at least one virtual break within the node device after concluding that the discerned break is gone.

22. The node device of claim 21 wherein the node device inspects incoming HELLO protocol packets to discern whether the node device has the duty of selectively imposing at least one virtual break or whether that duty is given to a device other than itself.
23. The node device of claim 21 wherein the node device does use dynamic layer 2 learning tables for transmitting data packets.
24. The node device of claim 21 further comprising at least one user port wherein:
during a period when the node device has imposed a virtual break, there is only one path for data packets to travel from the network ring to egress the network ring to the at least one user port;
during the period when the node device has imposed a virtual break, there is only one path for data packets to travel from the at least one user port to the network ring; and
during a period when the node device has removed the virtual break, data packets arriving at the node device from one of the at least one user port are sent out both the first ring connection and the second ring connection independent of any information learned from previous transmissions.
25. The node device of claim 24 wherein:
the data packets received at the node device from the at least one user port are modified to add a ring tag specifying a connection between two ring nodes; and
the data packets received at the node device from the network ring where the data packets are destined for one of the at least one user ports on the node device are stripped of the ring tag added to the data packet at the node where the data packet entered the network ring.
26. The node device of claim 21 wherein the at least one virtual break is always placed on the first ring connection so that:

data packets destined for the at least one user port on the node device reach the user port through the second ring connection when the at least one virtual break is imposed;

data packets from the at least one user port are transmitted exclusively on the second ring connection when the at least one virtual break is imposed; and

at least one type of network control packet is sent out of the first ring connection even when the at least one virtual break is imposed in the node device.

27. A ring relay for use in a network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network ring carrying data packets; the ring relay comprising the following attributes:

having a first ring port and a second ring port connected to the network ring;

detecting a type of network ring failure and moving the relevant ring port to a port down state; and

delaying resumption of transmission of data packets to the ring port previously moved to a port down state after detecting the removal of the ring failure so that the transition from the port down state to operating is delayed sufficiently to allow another node in the network ring to impose virtual breaks to prevent the formation of a communication loop for data packets.

28. The ring relay of claim 27 further comprising at least one user port connected to the network ring through the ring relay;

the characteristics of the ring relay further comprising:

adding a network ring VLAN tag, unique on the network ring, to any data packet entering the network ring through the ring relay; and

removing a network ring VLAN tag, previously added by a network ring node, from any data packet leaving the network ring through the ring relay.

29. The ring relay of claim 28 where data packets received on the ring first port are simply transferred to the second ring port, without modification, if a switch table, based

on the network ring VLAN tag, indicates that the ring relay is a transit node for that data packet, rather than an endpoint for that data packet.

30. The ring relay of claim 27 where data packets received on the first ring port are simply transferred to the second ring port, without modification, if a switch table indicates that the ring relay is a transit node for that data packet rather than an endpoint for that data packet.

31. The ring relay of claim 27 wherein the ring relay periodically receives discovery messages on the first ring port and adds information about the ring relay to the discovery message before sending it out the second ring port.

32. The ring relay of claim 31 wherein the ring relay will create a discovery message indicating that the ring relay is the originator of the discovery message and send that discovery message out the second ring port if a discovery message is not received on the first ring port within a specified time interval.

33. The ring relay of claim 31 wherein the ring relay periodically receives HELLO messages on the first ring port and sends the HELLO message out the second ring port without processing so that HELLO messages travel around a portion of the network ring faster than discovery messages travel around that same portion.

34. A ring relay for use in a network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network ring carrying data packets; the ring relay comprising the following attributes:

- having a first ring port and a second ring port connected to the network ring;
- detecting a type of network ring failure and moving the relevant ring port to a port down state; and
- delaying resumption of transmission of data packets to the ring port previously moved to a port down state until after receiving a control signal message

from another node indicating the insertion of at least one virtual break to prevent the formation of a communication loop for data packets.

35. A network having a redundant path for increasing the availability of a segment of a network ring, the network comprising:

an existing network ring;

a Master Arbiter at a first access point on the existing network ring;

a Slave Arbiter at a second access point, different from the first access point, on the existing network ring;

a second network ring formed with a ring side and an extension side;

the ring side comprising the Master Arbiter, the Ring Side port on the Master Arbiter, the Ring-Side Port on the Slave Arbiter, and the Slave Arbiter;

the extension side comprising the Master Arbiter, the Extension Side Port on the Master Arbiter; the first access point, the portion of the existing network ring between the first access point and the second access point, the second access point, the Extension Side Port on the Slave Arbiter, and the Slave Arbiter;

the Master Arbiter passing data packets received on the Extension Side Port to the Ring Side Port and the data packets received on the Ring Side Port to the Extension Side Port;

the Slave Arbiter when acting in a blocking mode, blocking the passage of data packets from the Extension Side Port to the Ring Side Port and from the Ring Side Port to the Extension Side Port but forwarding at least one type of control packet from the Extension Side Port to the Ring Side Port and from the Ring Side Port to the Extension Side Port;

the Slave Arbiter when acting in a forwarding mode, forwarding data packets from the Extension Side Port to the Ring Side Port and from the Ring Side Port to the Extension Side Port;

the Master Arbiter and the Slave Arbiter monitoring the performance of the second network ring to detect a failure in the second network ring within a first specified period and switching the Slave Arbiter from acting in a blocking mode to acting in a forwarding mode within a second specified period; and

the Slave Arbiter reacting to an indication that the failure in the second network ring has been remedied by switching from acting in forwarding mode to acting in blocking mode; and

at least one port on the second network adjacent to the failure in the second network ring noting the failure in the second network ring and switching to a port down mode, the at least one port delaying recovery from a port down condition long enough for the Slave Arbiter to switch from acting in a forwarding mode to acting in a blocking mode to prevent the formation of a communication loop for data packets.

36. The network of claim 35 wherein the total time for:

the actions of the Master Arbiter and the Slave Arbiter monitoring the performance of the second network ring to detect a failure in the network ring within a first specified period and

the time required for switching the Slave Arbiter from acting in a blocking mode to acting in a forwarding mode is less than 50 milliseconds.

37. A method of providing a redundant connection for at least one user port for increasing the availability access between the at least one user port and an existing network; the method comprising:

connecting the at least one user port to a slave arbiter node;

connecting a first ring port on the slave arbiter node in bidirectional data communication with a first network access device connected to the existing network;

connecting a second ring port on the slave arbiter node in bidirectional data communication with a second network access device connected to the existing network;

sending HELLO messages out the first ring port to the second ring port via the first network access device and the second network access device;

operating in a blocking mode when the ring comprising the slave arbiter node, the first access device, and the second access device is functioning, wherein blocking mode comprises:

blocking one ring port on the slave arbiter node, by blocking the transmission of data packets from the at least one user port from leaving the blocked ring port to travel to the corresponding access device; and

blocking data packets received on the blocked ring port from reaching the at least one user port;

operating in a forwarding mode when the slave arbiter detects a problem with a communication path between the unblocked ring port and the existing network, wherein forwarding mode comprises:

sending data packets received from the at least one user port out the previously blocked ring port; and

sending data packets received from the previously blocked ring port to the at least one user port.

38. The method of claim 37 wherein the frequency of HELLO messages is sufficient to allow both the detection of a problem with the communication path between the unblocked port and the existing network and switching from blocking mode to forwarding mode in less than 50 milliseconds.

39. The method of claim 37 wherein HELLO messages are sent out the second ring port to the first ring port via the second network access device and the first network access device.

40. A ring relay for use in a bidirectional network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network ring carrying data packets; the ring relay comprising the following attributes:

having a first ring port in data communication with a first incoming link and a first outgoing link and a second ring port in data communication with a second incoming link and a second outgoing link connected to the network ring; and

after detecting a type of network ring failure on the second incoming link, reacting to the detection of the network failure on the second incoming link by blocking the second outgoing link such that a virtual break is imposed on the second outgoing link after detection of a type of network failure on the second incoming link such that HELLO protocol packets received on the first incoming link are not passed to the second outgoing link which then triggers corrective action by the bidirectional network ring to remove an existing virtual break from a device other than the ring relay.

41. A ring relay for use in a bidirectional network ring that is manipulated to provide a virtual break in order to operate within a network requiring a tree topology, the network ring carrying data packets; the ring relay comprising the following attributes:

having a first ring port in data communication with a first incoming link and a first outgoing link and a second ring port in data communication with a second incoming link and a second outgoing link connected to the network ring;

receiving HELLO messages on the first incoming link and passing the HELLO message out the second outgoing link;

receiving data packets on the first incoming link destined for the second outgoing link and passing these data packets to the second outgoing link; and

after detecting a type of network ring failure on the second incoming link, reacting to the detection of the network failure on the second incoming link by blocking the HELLO messages received on the first incoming link from passing onto the second outgoing link and blocking data packets received on the first incoming link destined for the second outgoing link from passing to the second outgoing link such that a virtual break is imposed on the second outgoing link after detection of a type of network failure on the second incoming link.

42. A network node device for selectively creating and removing at least one virtual break in a network ring; the node device comprising:

a first ring connection and a second ring connection for use in connecting the node device to two ends of a network ring;

the node device operation within a network ring characterized by:

- initializing so that the node device contains a virtual break to prevent data packets received on the first ring connection from continuing on the second ring connection and data packets on the second data connection from continuing on the first ring connection;

- creating and transmitting a series of HELLO protocol packets out the first ring connection and the second ring connection;

- receiving HELLO protocol packets from a device other than itself on the first ring connection and monitoring the interval since the last HELLO protocol packet was received on the first ring connection;

- receiving HELLO protocol packets from a device other than itself on the second ring connection and monitoring the interval since the last HELLO protocol packet was received on the second ring connection;

- discerning the existence of a break in the network ring by detecting an interval in excess of a first value since the receipt of the last HELLO packet on the first ring connection;

- discerning the existence of a break in the network ring by detecting an interval in excess of a second value since the receipt of the last HELLO packet on the second ring connection where the second value can equal the first value;

- responding to the discernment of the existence of a break in the network ring by removing the at least one virtual break within the node device;

- detecting the resumed receipt of a HELLO protocol packet on a ring connection used to discern a break in the network ring, and concluding the discerned break is gone; and

- introducing at least one virtual break within the node device after concluding that the discerned break is gone;

further characterized by:

during a period when the node device has imposed a virtual break, there is only one path for data packets to travel from the network ring to egress the network ring to the at least one user port;

during the period when the node device has imposed a virtual break, there is only one path for data packets to travel from the at least one user port to the network ring; and

during a period when the node device has removed the virtual break, data packets arriving at the node device from one of the at least one user port are sent out both the first ring connection and the second ring connection;

still further characterized by:

the data packets received at the node device from the at least one user port are modified to add a ring tag specifying a connection between two ring nodes; and

the data packets received at the node device from the network ring where the data packets are destined for one of the at least one user ports on the node device are stripped of the ring tag added to the data packet at the node where the data packet entered the network ring.

43. The ring relay of claim 34 further comprising at least one user port connected to the network ring through the ring relay;

the characteristics of the ring relay further comprising:

adding a network ring VLAN tag, unique on the network ring, to any data packet entering the network ring through the ring relay; and

removing a network ring VLAN tag, previously added by a network ring node, from any data packet leaving the network ring through the ring relay.

44. The ring relay of claim 43 where data packets received on the first ring port are simply transferred to the second ring port, without modification, if a switch table, based on the network ring VLAN tag, indicates that the ring relay is a transit node for that data packet, rather than an endpoint for that data packet.

45. The ring relay of claim 34 where data packets received on the first ring port are simply transferred to the second ring port, without modification, if a switch table indicates that the ring relay is a transit node for that data packet rather than an endpoint for that data packet.

46. The ring relay of claim 34 wherein the ring relay periodically receives discovery messages on the first ring port and adds information about the ring relay to the discovery message before sending it out the second ring port.

47. The ring relay of claim 46 wherein the ring relay will create a discovery message indicating that the ring relay is the originator of the discovery message and send that discovery message out the second ring port if a discovery message is not received on the first ring port within a specified time interval.

48. The ring relay of claim 46 wherein the ring relay periodically receives HELLO messages on the first ring port and sends the HELLO message out the second ring port without processing so that HELLO messages travel around a portion of the network ring faster than discovery messages travel around that same portion.

Figure 1
Prior Art

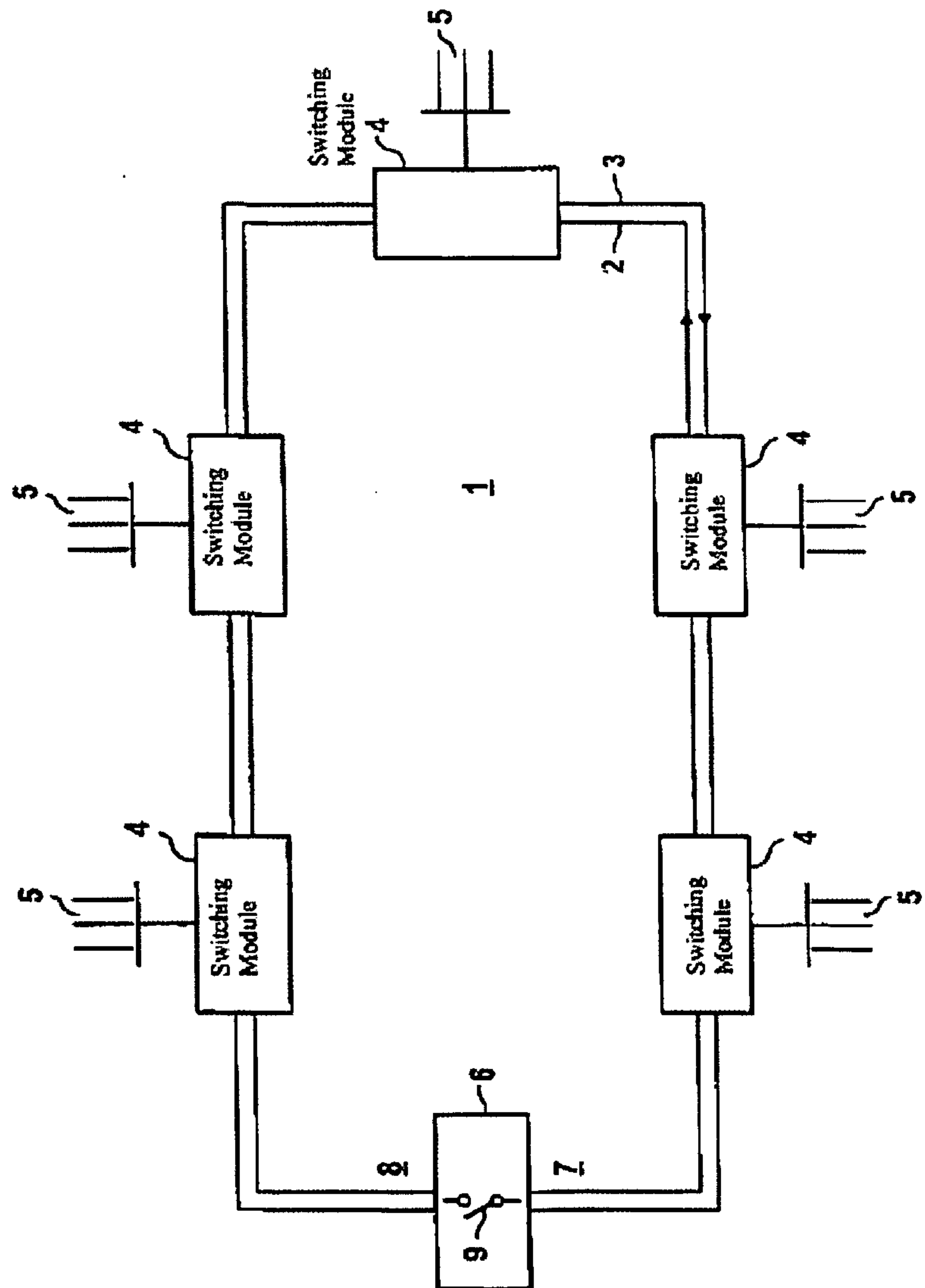


Figure 2

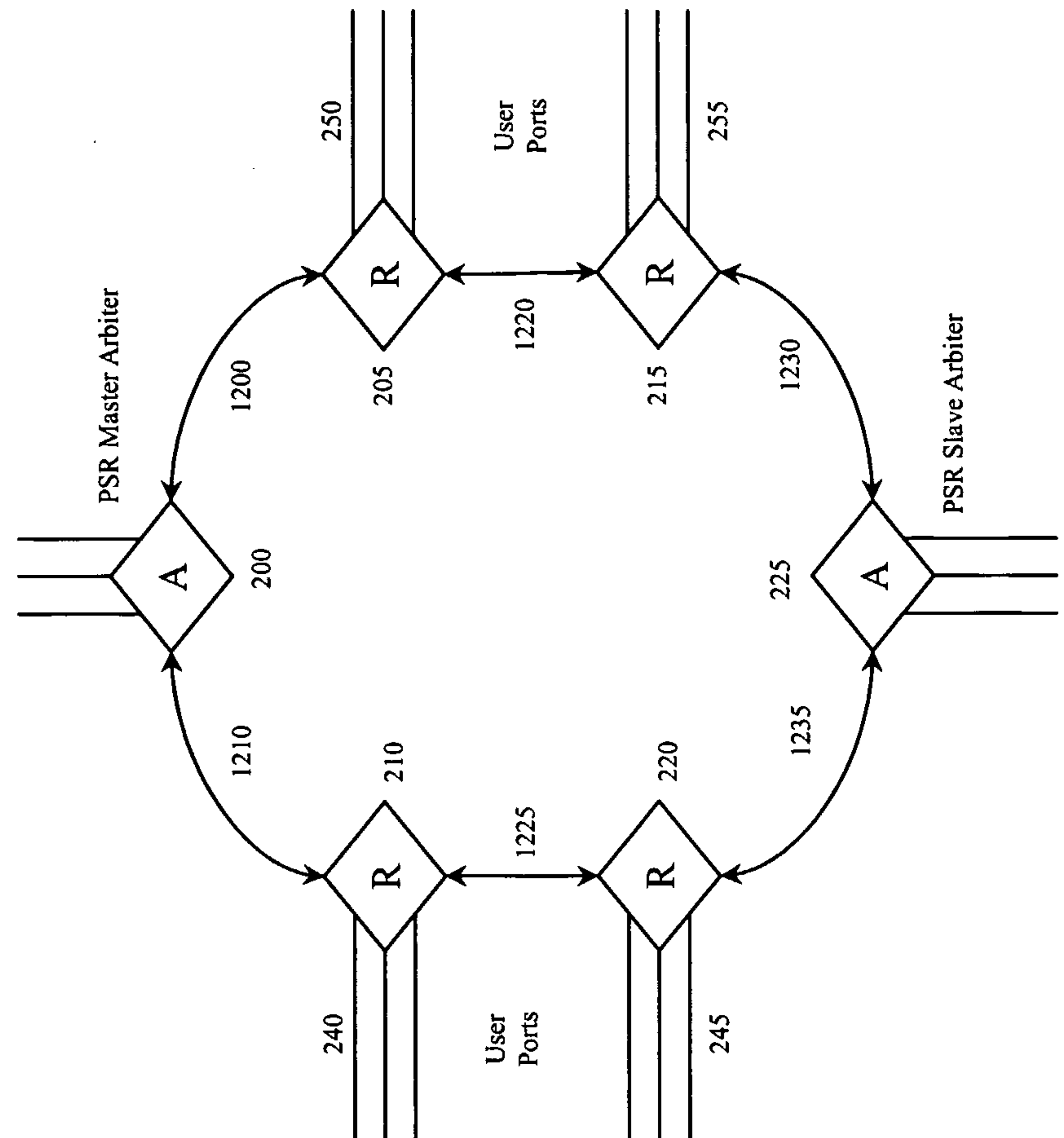


Figure 3

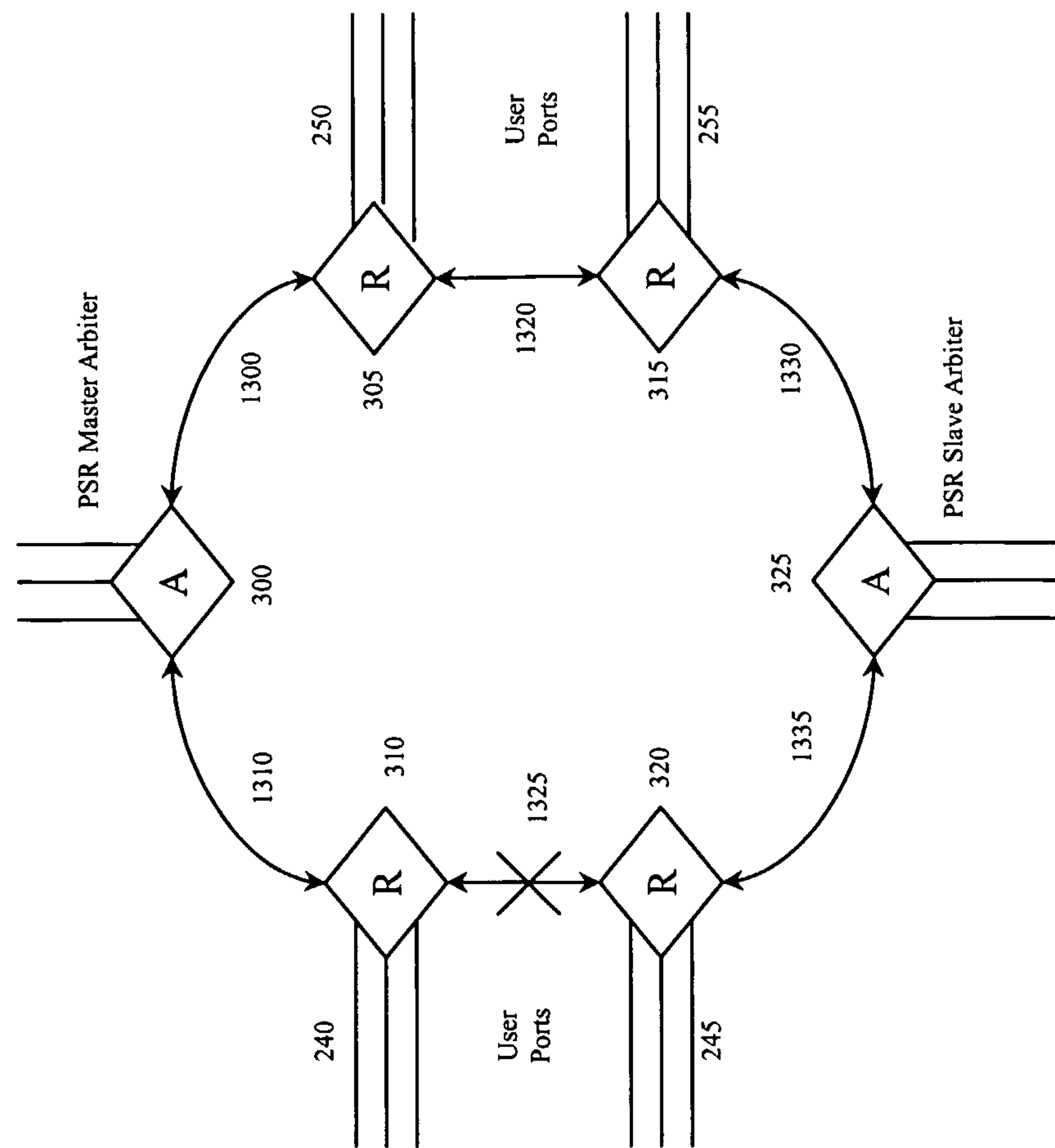


Figure 4

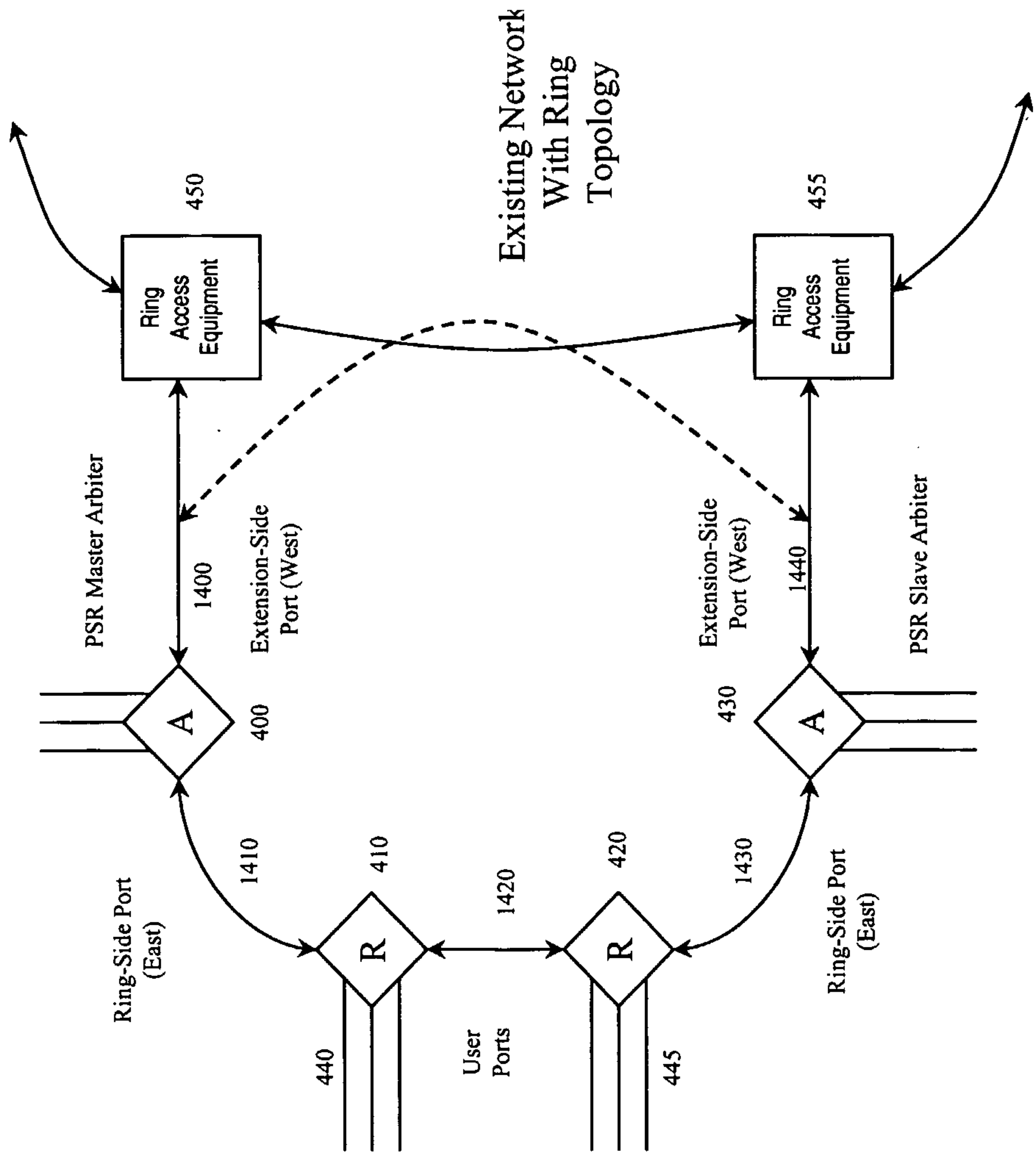


Figure 5

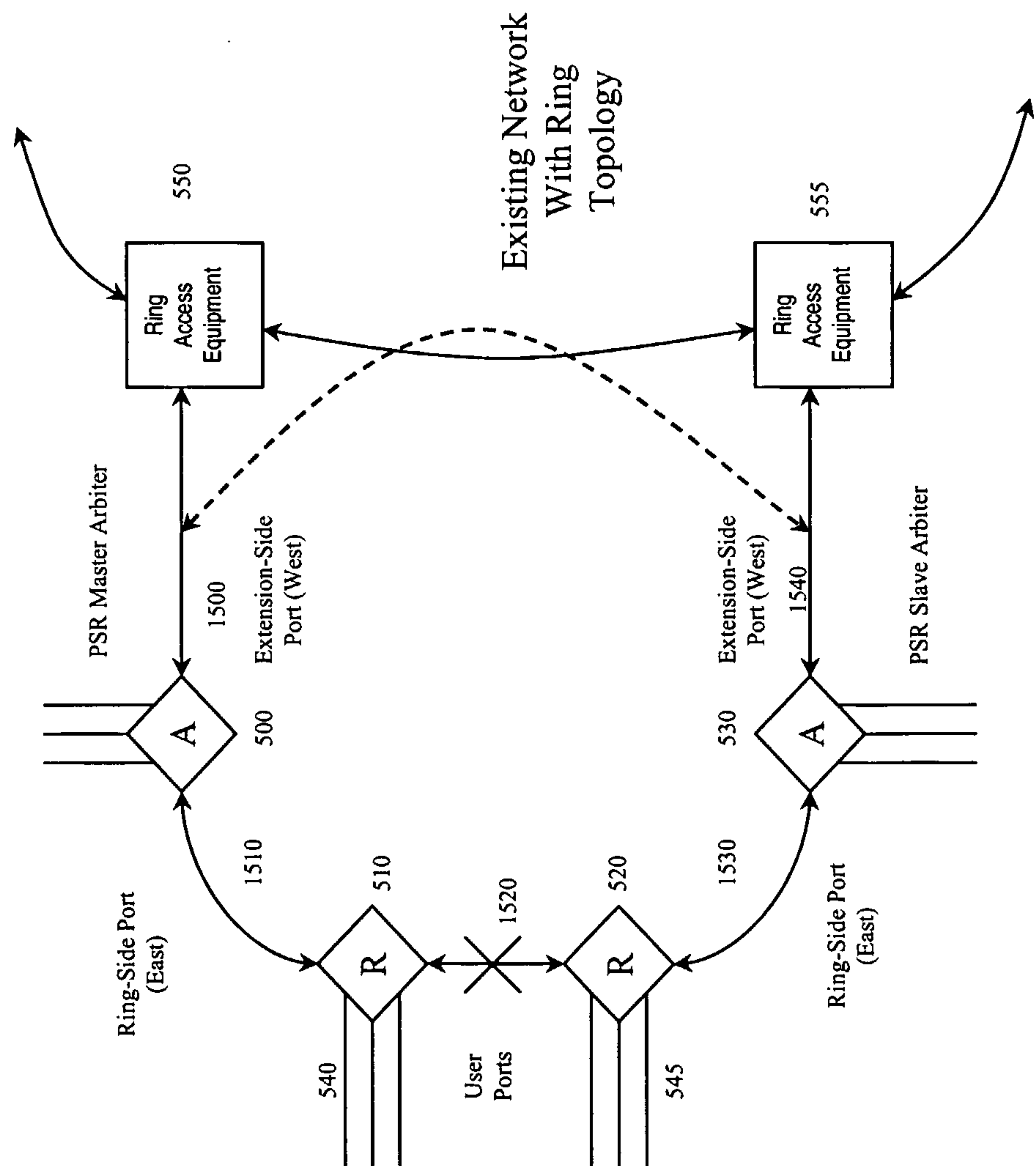


Figure 6

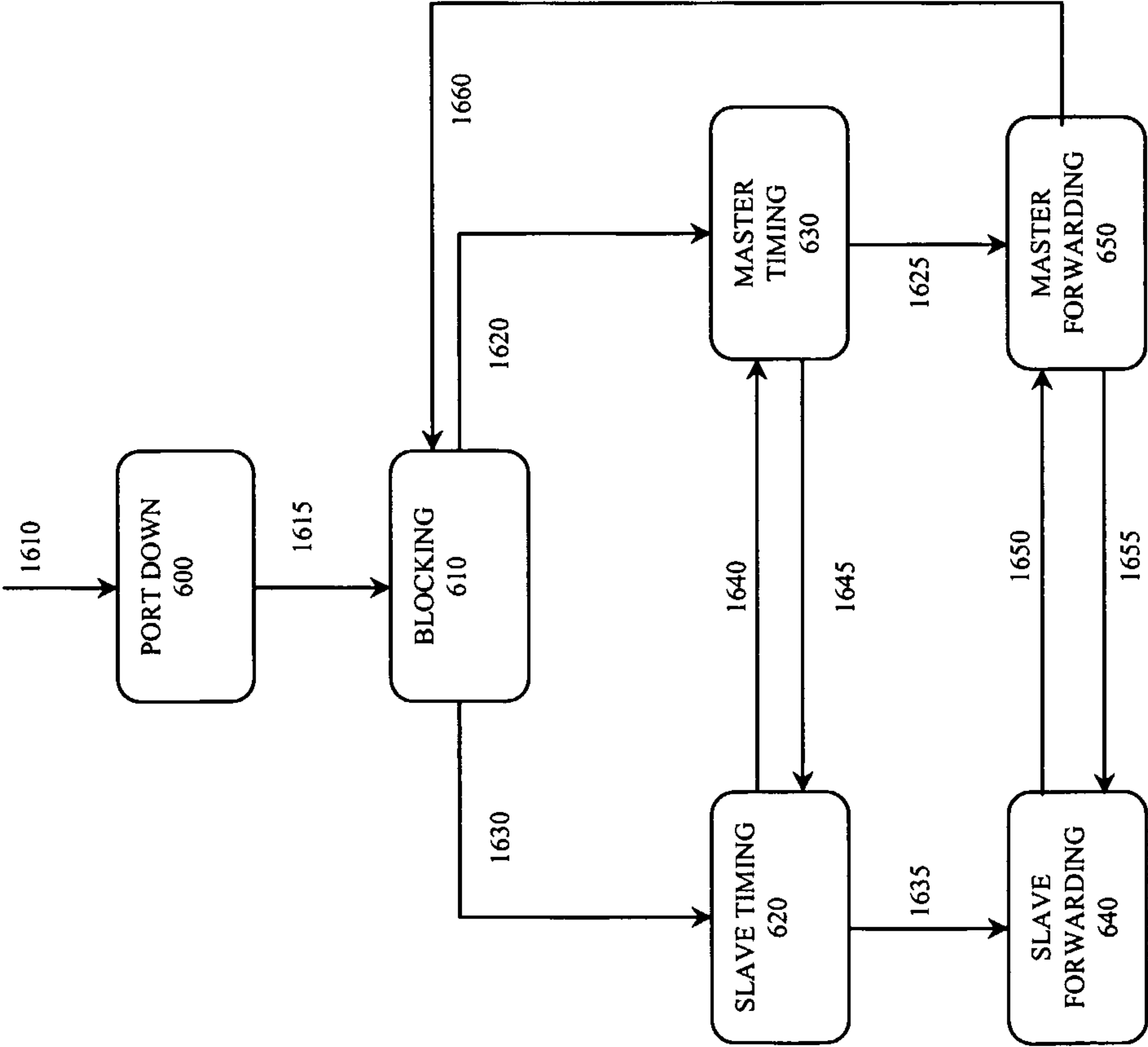


Figure 7

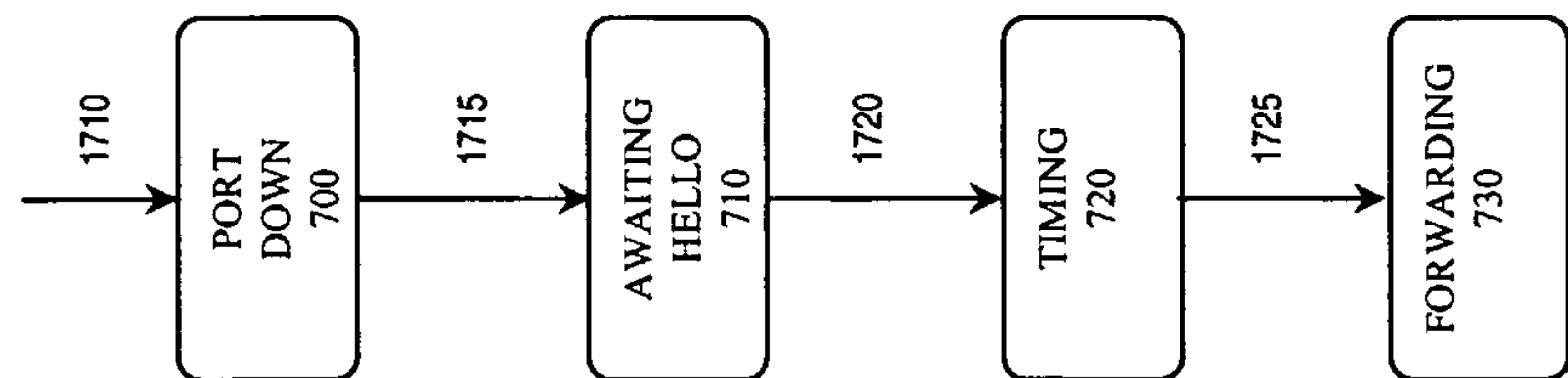


Figure 8

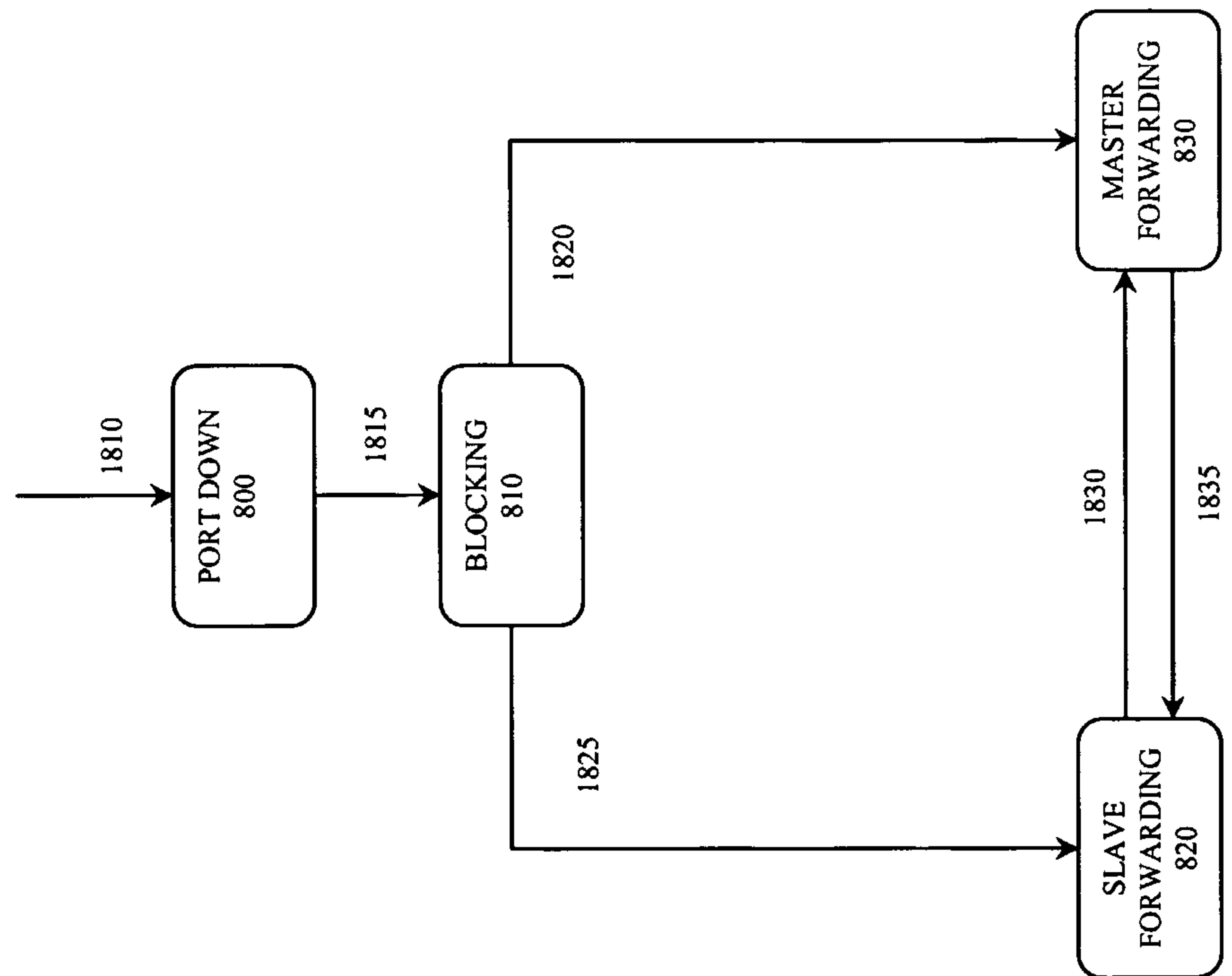


Figure 9

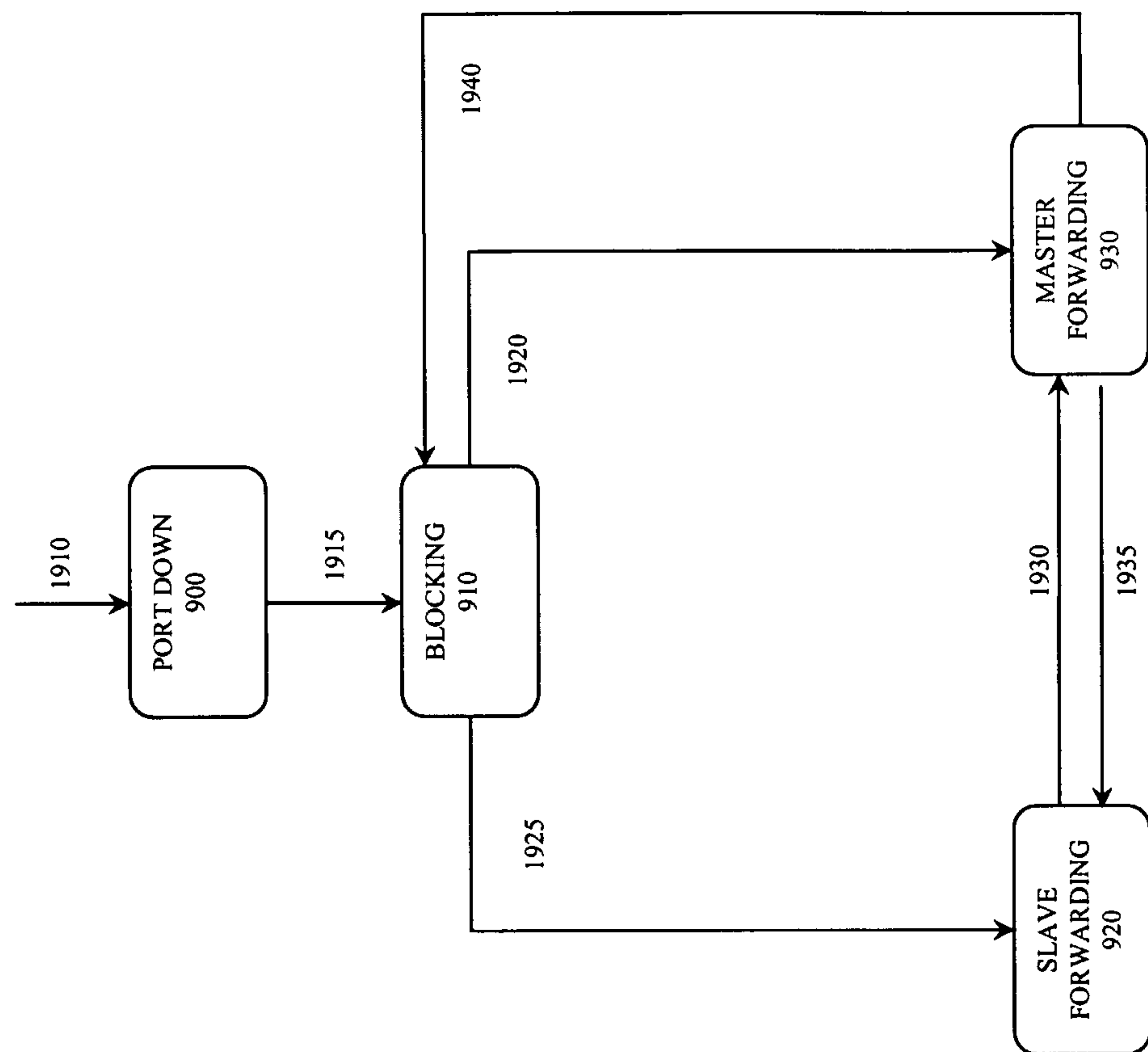


Figure 10

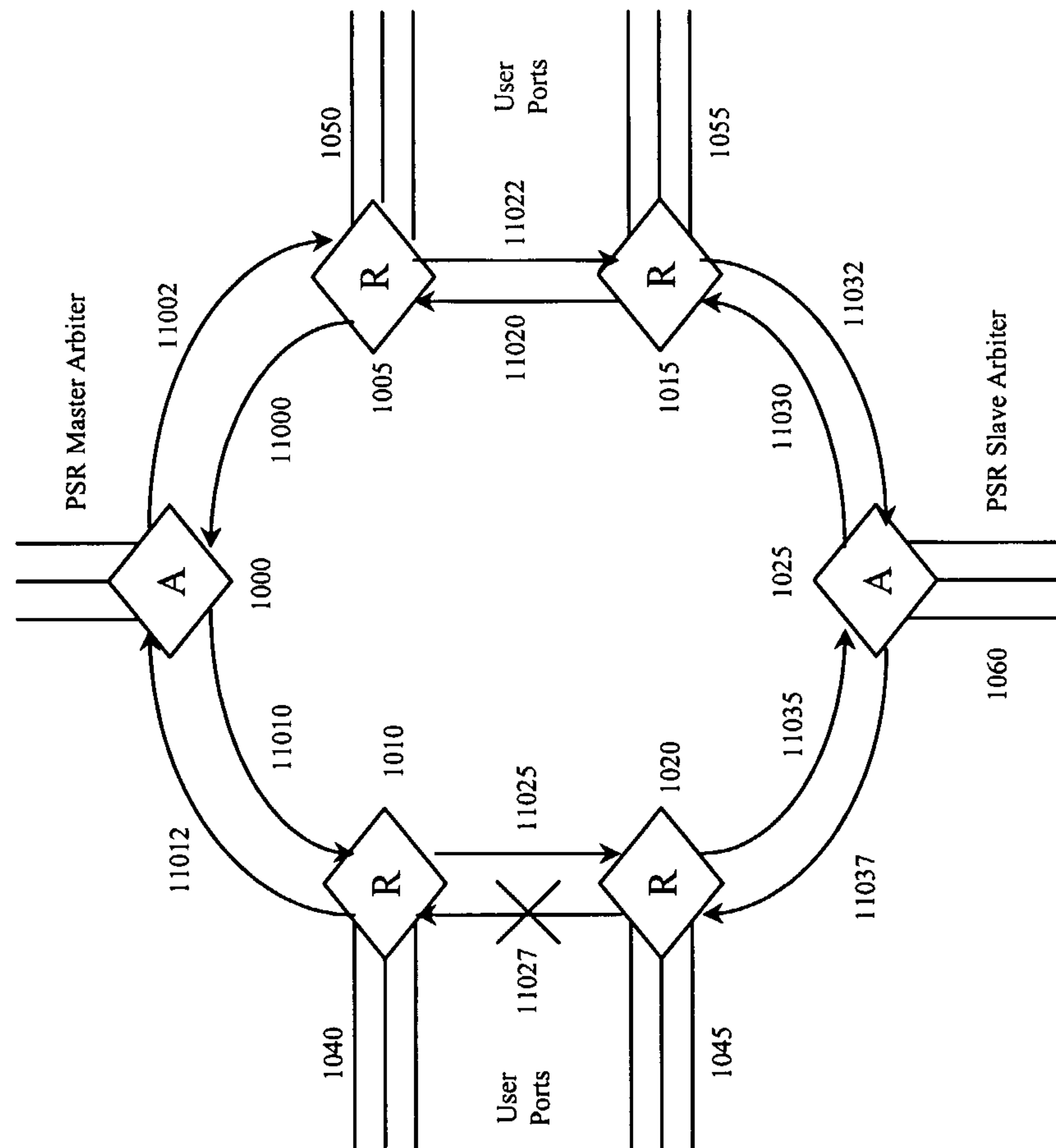


Figure 11

