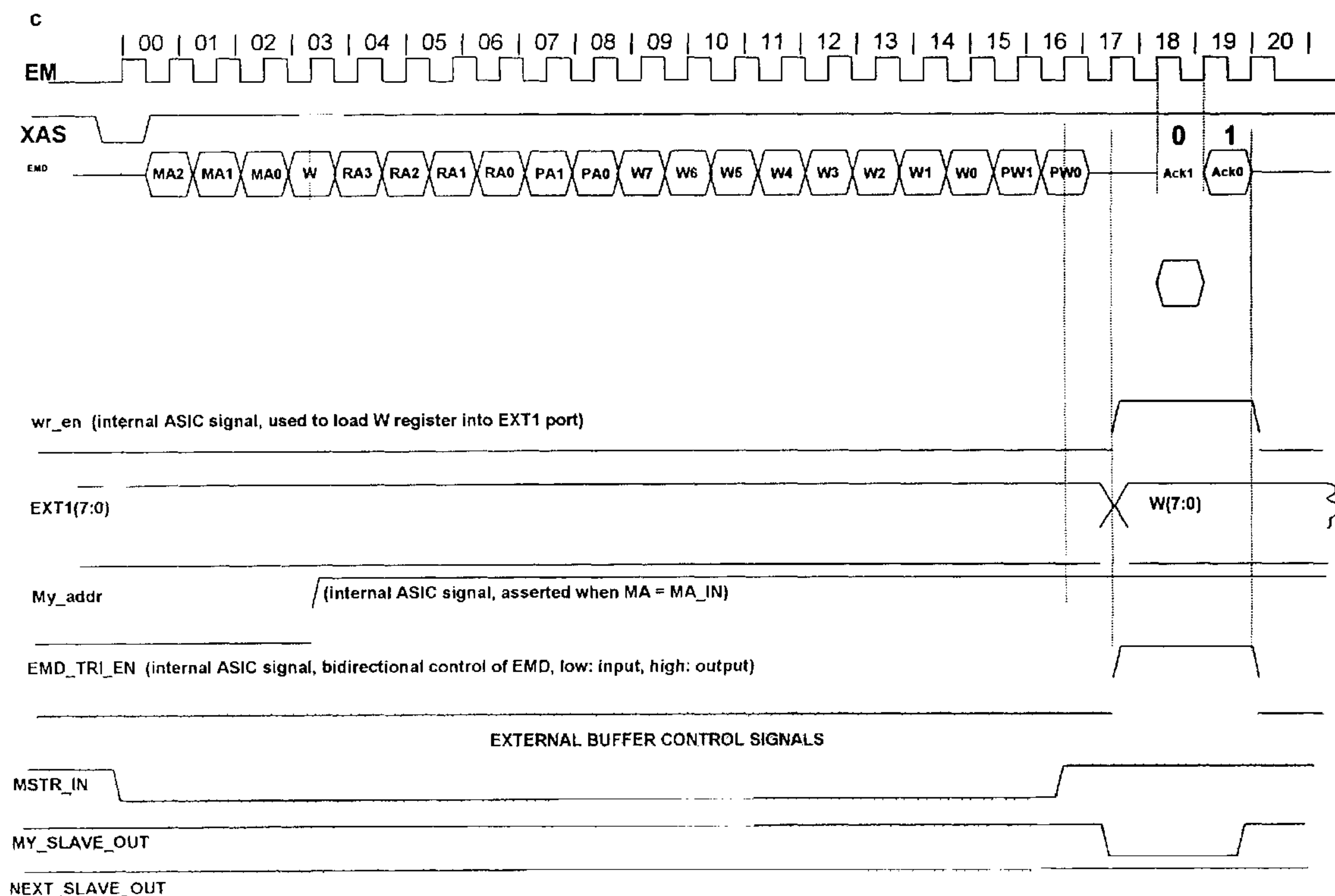




(86) Date de dépôt PCT/PCT Filing Date: 2001/03/21
(87) Date publication PCT/PCT Publication Date: 2001/09/27
(85) Entrée phase nationale/National Entry: 2002/09/18
(86) N° demande PCT/PCT Application No.: US 2001/008926
(87) N° publication PCT/PCT Publication No.: 2001/071514
(30) Priorité/Priority: 2000/03/21 (60/190,815) US

(51) Cl.Int.⁷/Int.Cl.⁷ G06F 13/38
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(54) Titre : SYSTEME, PROCEDE ET APPAREIL RELATIFS A UNE INTERFACE DE COMMUNICATION
(54) Title: A COMMUNICATION INTERFACE SYSTEM, METHOD AND APPARATUS



(57) **Abrégé/Abstract:**

A communication interface responds to a communication protocol for interfacing a controller and any of a plurality of discrete I/O devices. Each discrete I/O device has a different configuration. The interface has a plurality of modes of operation to accommodate the discrete I/O devices. In a first mode of operation, the interface accommodates a first discrete I/O device wherein a plurality of input pins input signals from a particular discrete I/O and a plurality of output pins output signals to the particular discrete I/O device. In a second mode of operation, the interface accommodates a second discrete I/O device wherein the input pins form a bidirectional input/output port and the output pins form a control and address line for controlling the second discrete I/O device and other discrete I/O devices.



(12) INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

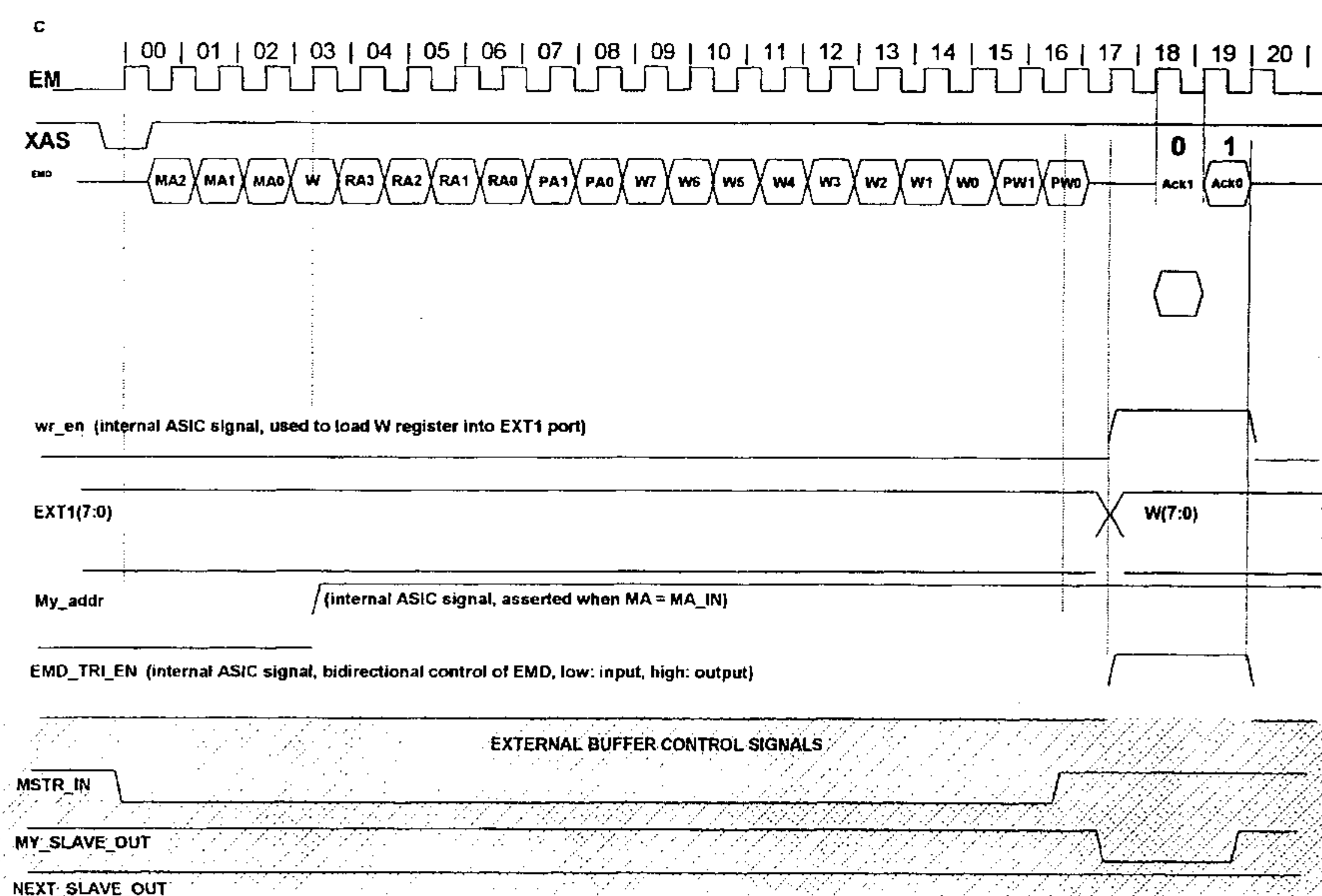
(19) World Intellectual Property Organization
International Bureau(43) International Publication Date
27 September 2001 (27.09.2001)

PCT

(10) International Publication Number
WO 01/071514 A3

- (51) International Patent Classification⁷: **G06F 13/38** (74) Agents: ASPERAS, I., Marc et al.; Siemens Corporation, Intellectual Property Department, 186 Wood Avenue South, Iselin, NJ 08830 (US).
- (21) International Application Number: **PCT/US01/08926**
- (22) International Filing Date: **21 March 2001 (21.03.2001)** (81) Designated States (*national*): CA, MX.
- (25) Filing Language: **English** (84) Designated States (*regional*): European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE, TR).
- (26) Publication Language: **English**
- (30) Priority Data: **60/190,815** **21 March 2000 (21.03.2000)** **US** Published:
— *with international search report*
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- (72) Inventors: **MASSIE, Michael, Ross; 440 Springfield Drive, Bristol, TN 37620 (US). MCNUTT, Alan, D.; 2902 Newbern Drive, Johnson City, TN 37604 (US).** *For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.*

(54) Title: A COMMUNICATION INTERFACE SYSTEM, METHOD AND APPARATUS



(57) Abstract: A communication interface responds to a communication protocol for interfacing a controller and any of a plurality of discrete I/O devices. Each discrete I/O device has a different configuration. The interface has a plurality of modes of operation to accommodate the discrete I/O devices. In a first mode of operation, the interface accommodates a first discrete I/O device wherein a plurality of input pins input signals from a particular discrete I/O and a plurality of output pins output signals to the particular discrete I/O device. In a second mode of operation, the interface accommodates a second discrete I/O device wherein the input pins form a bidirectional input/output port and the output pins form a control and address line for controlling the second discrete I/O device and other discrete I/O devices.

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TITLE OF THE INVENTION

A Communication Interface System, Method and Apparatus

BACKGROUND

Field Of The Invention.

The invention relates to a communication interface system, method and apparatus and, more particularly, to a universal integrated module for interfacing a control module to other modules.

Related Information.

In the past, interfaces have been introduced that interface a control module to other modules. However, there has been no universal interface for the interfacing a plurality of communications protocol as particularly contemplated by the present invention.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a table;
Figs. 2a and b are timing diagrams;
Fig. 3 is a table;
Fig. 4 is a timing diagram;
Fig. 5 is a timing diagram;
Fig. 6 is a table;
Fig. 7 is a timing diagram;
Fig. 8 is a timing diagram;
Fig. 9 is a timing diagram;
Fig. 10 is a timing diagram;
Fig. 11 is a timing diagram;
Fig. 12 is a timing diagram;
Fig. 13 is a timing diagram;

Fig. 14 is a timing diagram; and
Figs. 15a and 15b are block diagrams.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

The interface of the present invention provides serial communication to expansion modules (EM). A CPU of the interface controls all communications to and from the EM's and will be referred to as the "Master Function". The Expansion Module will include an ASIC to achieve this serial communication protocol and will be referred to as the "Slave Function" or "Slave ASIC".

The I/O Expansion bus signals as viewed by the Expansion Module are described in the following table. The connection from the PLC to the EM and from EM to EM will be 1 to 1 using a 10 pin header type connection. Refer to Figure 1 for Expansion I/O Bus connector pin assignments.

The following diagrams in Figures 2a – 2b illustrate an Expansion Bus Read cycle and Write cycle sequences. A bus transaction will be initiated by short active low pulse on XA_OD signal. In Fig. 2a:

MA[2:0]	- Module Address that CPU is addressing
W	- Transaction Type requested by CPU, Read/Write Bit
(1=>Write, 0=>Read)	
RA[3:0]	- Register Address that CPU is addressing
CP[1:0]	- Control Register Parity Bits generated by CPU
W[7:0]	- Data written to Expansion Module by CPU
DP[1:0]	- Data parity generated by CPU
Ack[1:0]	- Acknowledge bits returned to CPU by the Slave
	- A1 : 0 indicates a successful write cycle (no parity errors);
	1 indicates an invalid write transaction, data parity error encountered
	- A0 : defaults to a 1

In Fig. 2b:

MA[2:0]	- Module Address that CPU is addressing
W	- Transaction Type requested by CPU, Read/Write Bit (1=>Write, 0=>Read)
RA[3:0]	- Register Address that CPU is addressing
CP[1:0]	- Control Parity Bits generated by CPU
R[7:0]	- Data written to CPU from Expansion Module
DP[1:0]	- Data parity generated by Expansion Module

Each EM implements an electrical interface to the expansion I/O bus consisting of termination circuits and bus driver circuits. This will allow the addition of +5V power to be introduced anywhere in the daisy chained I/O bus and provides some protection of the SLAVE ASIC I/O. Given that EMD signal is a bi-directional signal, control circuitry will be provided on the bus driver circuits to avoid bus contention errors. Three control signals (MSTR_IN, MY_SLAVE_OUT, and NEXT_SLAVE_OUT) are used to enable/disable the EMD bus driver circuits. For details of the termination circuitry refer to Figure 3. Figure 3 describes EMD bus driver circuit configuration for three expansion modules.

A bus driver circuit is provided which has an active low enable line. The Slave ASIC generates 3 control signals to properly enable/disable the bus drivers. Control signal MSTR_IN enables bus driver circuits A and B. This allows the EMD signal to be input into the expansion module ASIC and to be input into the next expansion module ASIC down stream. MSTR_IN becomes active when XAS is detected and becomes inactive prior to any response from an EM. The signals MY_SLAVE_OUT and NEXT_SLAVE_OUT control EMD responses from the EM's according to the EM's physical address position. For example, MY_SLAVE_OUT controls the EMD signal as an output of the EM that was addressed by the CPU. The signal NEXT_SLAVE_OUT passes the EMD response through the EM if the EM addressed by the CPU has a higher address or is downstream from that EM. Neither MY_SLAVE_OUT or NEXT_SLAVE_OUT are asserted when the CPU

addresses an EM with a lesser address. All three control signals immediately become inactive on the occurrence of XOD.

In one particular embodiment of the invention, the initial Slave design may be developed in a 128 Macrocell CPLD using VHDL as a design instrument. The CPLD design is migrated into an ASIC design. The ASIC is designed to have 44 total pins with 36 usable pins for I/O and 8 pins for power and ground. The operating frequency of the SLAVE ASIC may be 4.125 MHz maximum, which is well within the ASIC capability. The following table defines the required inputs and outputs of the Slave ASIC. Refer to Figure 4 (SLAVE ASIC Suppliers Specification) for details of the ASIC.

The I/O described above accommodates all Expansion Module I/O configurations. That is, the invention is a universal interface. This is accomplished by implementing two modes of operation within the ASIC. The differences in ASIC operation is the implementation of EXT0 and EXT1 data ports.

The slave ASIC has two modes of operation to accommodate all Expansion Module I/O configurations. In Mode 0, the EXT0 bus is an 8 bit input register and EXT1 bus will be an 8 bit, active low, output register. EXT0 data bus and EXT1 data bus interfaces directly to the Expansion Module digital I/O. Mode 0 is used for Expansion Modules of 8IN / 8OUT or less. In Mode 1, EXT0 bus is used as an 8 bit, active high, bi-directional data bus and EXT1 bus is used as 8 bits of address and control. External registers and decode circuitry are required for Mode 1 operation. Mode 1 is used when the Expansion Module I/O configuration is greater than 8IN / 8OUT or for an intelligent module. The ID_REG is decoded by each EM at power up to determine its mode of operation. The ID_REG is also be read by the CPU to determine the Expansion Module type.

The polarity of EXT0 data port is "active high" for both Mode 0 and 1. EXT1 data port is "active low" while in mode 0 operation and in mode 1 the control lines is "active low" and the address lines is "active high."

In operation, the Slave ASIC implements a state machine architecture to provide proper communication and control. At initial power up the CPU issues an active XOD signal. Detection of XOD places the state machine into its home state and the EXT1 data port is cleared if in mode 0 or the external output registers is cleared if in mode 1. Also at initial power up, the Slave ASIC determines its Module Address (MA_IN), Mode of operation (Mode 0 or 1) and propagates the next Module Address (MA_OUT) by incrementing its Module Address by 1. Once XOD is released, the Slave ASIC state machine continuously monitors the XAS signal from its home state. A bus transaction is initiated when XAS becomes active and transitions to state 0 on the first rising edge of EMC0 clock. At state 0 the state machine is placed into a known state and propagates to state 1 on the next rising edge of EMC0 clock. If at any time the XOD signal becomes active, then the EXT1 data port is asynchronously reset and the state machine returns to its home state. If in mode 1 operation, the external output registers is asynchronously cleared and the state machine returns to its home state. The next occurrence of XAS synchronously places the state machine into a known state.

Fig. 5 illustrates a Mode 0 Write Transaction. As shown in Fig. 5, once a valid XAS is detected the Slave ASIC propagates to state 0 of the control state machine on the first rising edge of EMC0 clock. At state 0, the MSTR_IN signal is asserted and the EMD signal is enabled as an input to the Slave ASIC. The control state machine begins to shift in the control register data beginning on the rising edge of EMC0 clock 1 (state 1) and ending on the rising edge of EMC0 clock 10 (state 10). At state 4 (EMC0 clock 4) the Slave ASIC determines if the Module Address (MA) shifted in from the CPU equals the Module Address (MA_IN) propagated in at power up and if true "My_addr" is asserted. In the case that the addresses do not match, My_addr is not asserted, and the state machine continues through the remaining states to account for proper EMD bus driver control and control register parity checks. At State 10 the Slave ASIC determine the type of transaction to occur and enters either the write state machine or read state machine on the next EMC0 clock. During a write transaction the Slave ASIC shifts in Write data beginning at state 11 (EMC0 clock 11) and ending at state 20 (EMC0 20). Also, during state 11 the control register parity is checked and in the event an error is detected on the control register the write state machine returns to an idle state and the control state machine returns

to its home state on the next EMC0 clock. EXT1 data port is not disturbed and the MSTR_IN bus control signal becomes nactive. If no control register parity error is detected then the write data is shifted in accordingly. At state 20 the MSTR_IN bus control signal is released and at state 21 the state machine prepares the EMD bus control signals for a response back to the CPU. If "My_addr" is valid then MY_SLAVE_OUT bus control signal is asserted. If "My_addr" is not valid then the NEXT_SLAVE_OUT bus control signal is asserted only if the CPU has addressed a module of a greater address. At state 22 the Slave ASIC checks parity on the write data. When a write data parity error is detected the Slave ASIC returns an invalid (11) Acknowledge to the CPU and does not present new data to EXT1 data port. If no parity error is detected, then the Slave ASIC returns a valid (01) Acknowledge to the CPU, decodes the register address (RA) and enables new data onto the EXT1 data port if register address 'C' (hex) has been decoded. If any other register is decoded the state machine will return a valid Acknowledge to the CPU, but does not present new data to the EXT1 data port. This is only true while in mode 0 operation. On the rising edge of EMC0 clock 24 the write state machine returns to an idle state and the control state machine returns to its home state.

Figure 6 illustrates a mode 0 read transaction with the CPU addressing RA 8(hex). The control state machine operates the same as the previous write bus transaction, except at state 10 it now enters the read state machine. At state 11 the read machine releases the EMD bus control signal MSTR_IN, selects the EMD (EMD_TRI_EN) bi-directional port as an output, and checks the control register parity. If a parity error occurs the read state machine returns to an idle state and the control state machine returns to a home state on the next rising edge of EMC0 clock. If no error was detected then the read state machine propagates to state 12 on the next rising edge of EMC0 clock. At state 12 the state machine prepares the EMD bus control signals for a response back to the CPU by asserting bus control signal MY_SLAVE_OUT. Also at state 12 the register address (RA) is decoded to determine the data source. If RA decodes to 0 then the ID_BUF data is enabled, if RA decodes to 8 then EXT0 data bus is enabled, and if any other address is decoded then the hex value FF is enabled. On the falling edge of EMC0 12 the "R" register is loaded with the appropriate data either from the EXT0 data bus, the ID_BUF, or the default value of FF. For the case in Figure 2 the EXT0 data is loaded

into the "R" register. On the next rising edge of EMC0 clock (state 13) the read state machine shifts the read data bit 7 onto the EMD line and the last read bit 0 is shifted in on the rising edge of EMC0 clock 20 (state 20). The Slave ASIC generates 2 parity bits, PR1 and PR0, on the 8 bits of read data and shift this data onto the EMD line at states 21 and 22. At state 23 all EMD bus control signals are released, the read machine returns to an idle state and the control machine returns to its home state.

Mode 1 Operation will now be discussed with reference to Fig. 7. The EXT1 bus will be used as a control port in ASIC Mode 1. The table in Fig. 7 describes each bit for EXT1 data port when in ASIC Mode 1.

In a Mode 1 Write Transaction, the control register state machine and the write state machine function the same as in mode 0, however the external port usage and the available registers differs from mode 0. In mode 1, EXT1 data port is used as a control port for external decode circuitry and EXT0 data port is a bi-directional data port. Write data is enabled onto EXT0 data port on the rising edge of EMC0 clock 21 and is valid for 3 clock periods. The Register Address (RA[3:0]) is clock'd onto EXT1 data port on the falling edge of EMC0 clocks 5, 6, 7, & 8 respectively. All 16 register addresses is available for external decode. The WRSTRB is asserted on the falling edge of ECM0 clock 22 and is cleared on the falling edge of EMC0 clock 23. The signal Busy is asserted on the falling edge of EMC0 clock 16 and is cleared on the falling edge of EMC0 clock 24. If a parity error is detected on the WRITE byte, then both the ASIC registers and the external registers retain their last received value. Refer to the figures 8 and 9 for detailed timing information for a time exploded view of Mode 1, Write Transaction Timing.

In a Mode 1 Read Transaction, the control register state machine and the read state machine function the same as in mode 0, however external port usage and the available registers differ from mode 0. In mode 1, EXT1 data port is used as a control port for external decode circuitry and EXT0 data port is a bi-directional data port. Also, the CPU can access all 16 registers in Mode 1 with register 0 still the ID register. Refer to Figures 10 and 11 for timing details of the control port EXT1 and the data port EXT0. Figs. 10 and 11 illustrate Mode 1, Read Transaction Timing.

9.0 ID Register Definition:

The ID Register is addressed from Register Address (RA) 00 hex and is defined in Figure 11. The Slave ASIC filsl bit 7 with a 0 and the remaining 7 bits are hardwired according to the Module type. According to the ID Register definition, the Slave ASIC operate in ASIC Mode 0 only for ID Register values of 01, 04, and 05 hex. All other ID Register values operate in Mode 1.

Figs. 12-14 contains mode 1 Read and Write Bus transactions that display various parity errors. The figures illustrate the Bus operation/response under these conditions. Mode 0 bus transactions respond to these errors in the same manner. ^{Figure 6} Fig. 12 illustrates Mode 1, Control Register Parity Error during a Write Transaction. Fig. 13 illustrates Mode 1, Write Register Parity Error during a Write Transaction. Fig. 14 illustrates MODE 1, Control Register Parity Error During a READ Transaction.

Figs. 15a and 15b illustrate the two modes in block diagram form. As shown in Fig. 15a, the slave ASIC is in mode 1. In this state, a first bus line is dedicated as an output and the second bus provides control lines for sending control signals to the attached modules. In Fig. 15b, the same slave ASIC is switched into mode 0. In this state, the first bus remains a bused output. The second bus, however, is switched to a bused input. Thereby, the invention realizes the universal interface for a plurality of modes as contemplated in the description of the modes above.

While the present invention is described with reference to particular embodiments, it will be appreciated that the invention is not so limited to a specific

embodiment, but may encompass all modifications and permutations that are within the scope of the invention.

We claim:

1. A communication interface apparatus that responds to a communication protocol for interfacing a controller and any of a plurality of discrete I/O devices, each discrete I/O device having a different configuration, said apparatus having a plurality of modes of operation to accommodate said discrete I/O devices, said apparatus comprising:

a first mode of operation that accommodates a first discrete I/O device wherein a plurality of input pins input signals from a particular discrete I/O and a plurality of output pins output signals to said particular discrete I/O device; and

a second mode of operation that accommodates a second discrete I/O device wherein said input pins form a bidirectional input/output port and said output pins form a control and address line for controlling said second discrete I/O device and other discrete I/O devices.

2. The apparatus of claim 1, wherein said input pins are the same input pins for said first mode as for said second mode.

3. The apparatus of claim 1, further comprising data structures for inputting and outputting signals between the communication interface and said discrete I/O modules, wherein a format for a data structures for both of said first and second modes is the same.

4. The apparatus of claim 1, further comprising multiple read and write transactions in said first mode that provide an extended I/O bit protocol.

Signal Name	Description
EMD	Expansion Module Data: Bi-directional signal used to communicate address and data to and from the module.
EMC0	Expansion Module Clock 0: Used to access expansion module I/O. Active only during bus transaction. Frequency = 4.125 MHz
XA_OD	Address / Output Disable: Active Low, controlled by the Master Function. Dual function signal: XOD - A long pulse, 10 times the clock period or greater will reset the expansion modules. XAS - A short pulse, approximately one clock period, initiates an expansion module bus transaction.
EMA[2:0]	Expansion Module Address: Daisy chained from PLC to module to module. Value input to a module becomes that modules address. The module will output its address + 1 to the next module. The module closest to the PLC will have an address of 0. A total of 7 EM can be connected (0:6)
+5V	5 volt power supply – two signal lines
GND	Power supply return – two signal lines

Figure 1

Figure 2a

EM Bus Write Transaction Sequence

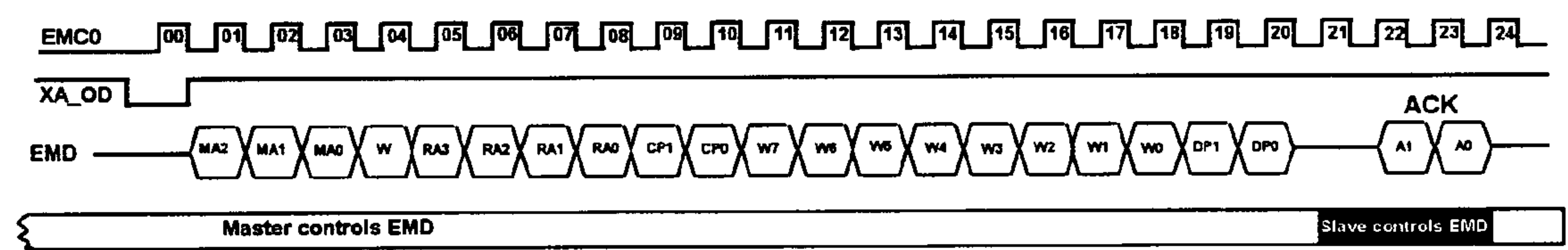


Figure 2b

EM Bus Read Transaction Sequence

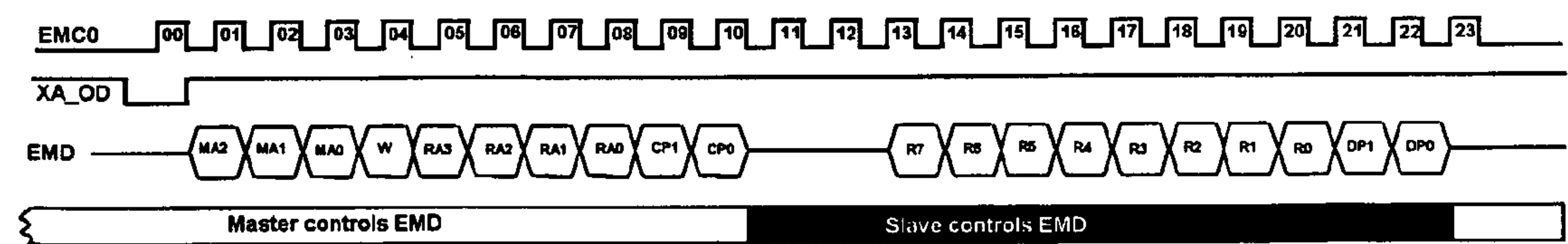


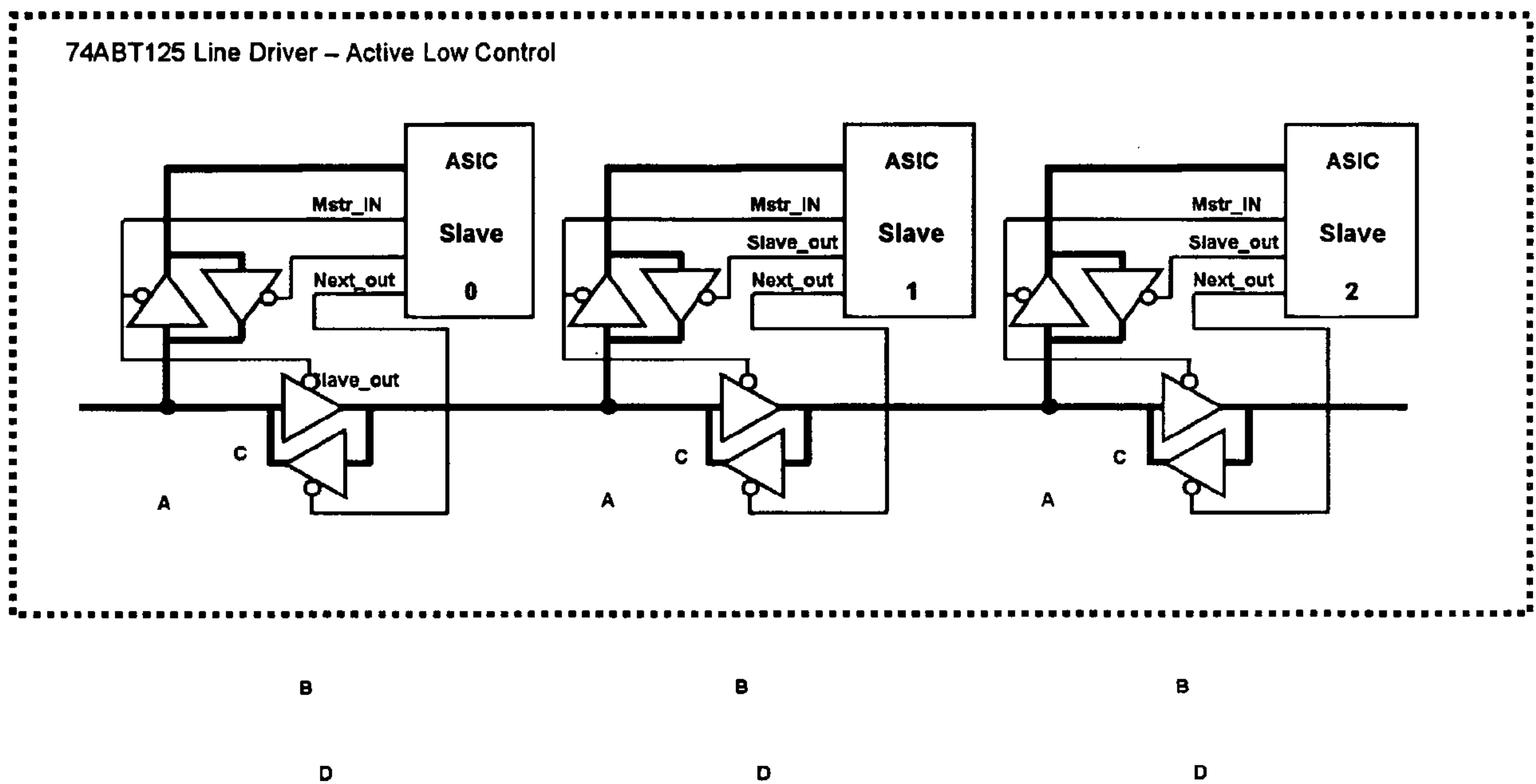
Figure 3

Figure 4

Signal Name	Description	# of Signal
MA_IN [2:0]	Input, Module Address	3
MA_OUT [2:0]	Output, Module Address + 1	3
XAS	Input: Active Low, initiates bus transaction	1
XOD	Input: Active Low, EM Reset	1
EMC0	Input: Gated Clock, 4.125 MHz Maximum Read => 24 clocks, Write => 25 clocks	1
EMD	Bi-directional Data Line	1
ID_BUF [6:0]	Input (Register ID), Identifies EM Type, Decodes ASIC Mode of operation (Mode 0 or 1)	7
EXT0 [7:0]	External Port 0 Dual Function: Mode 0. Input Data Bus Mode 1. Bi-directional Data Bus	8
EXT1 [7:0]	External Port 1 Dual Function: Mode 0. Output Data Bus Mode 1. Output Address & Control	8
MSTR_IN	Bus Driver Enable, Active Low	1
MY_SLAVE_OUT	Bus Driver Enable, Active Low	1
NEXT_SLAVE_OUT	Bus Driver Enable, Active Low	1
		36 Total I/O Pins

Figure 5

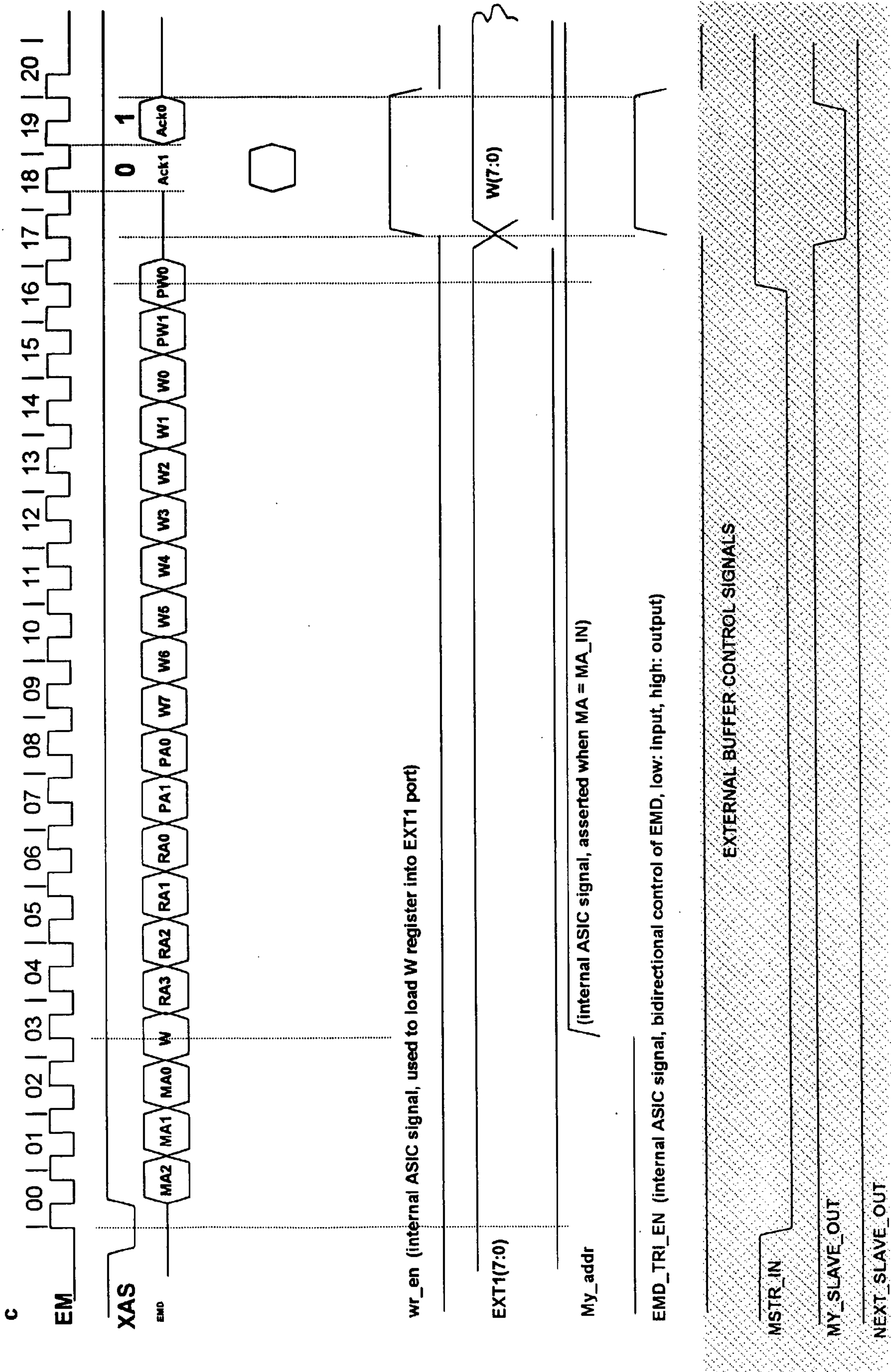
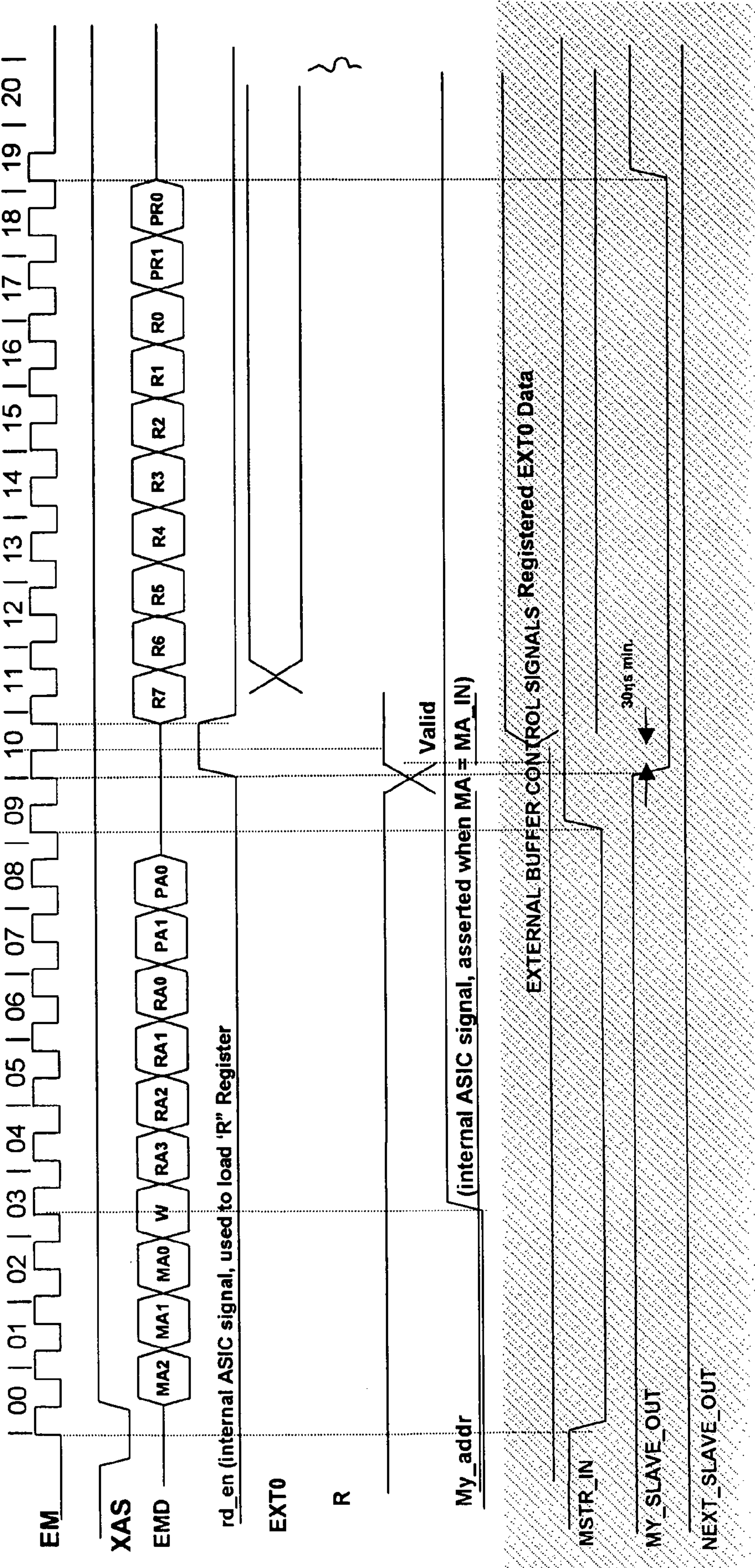


Figure 6



EXT 1	Mode 1 Function	Description
bit 7	CS	Chip Select, Active Low, OR'd condition of WRSTRB & RDSTRB
bit 6	Busy	Active Low, indicates activity on the EXT0 & EXT1
bit 5	WRSTRB	Write Strobe: Active Low, external decode circuitry
bit 4	RDSTRB	Read Strobe: Active Low, external decode circuitry
bit 3	RA3	Register Address Bit 3, Active High
bit 2	RA2	Register Address Bit 2, Active High
bit 1	RA1	Register Address Bit 1, Active High
bit 0	RA0	Register Address Bit 0, Active High

Figure 7

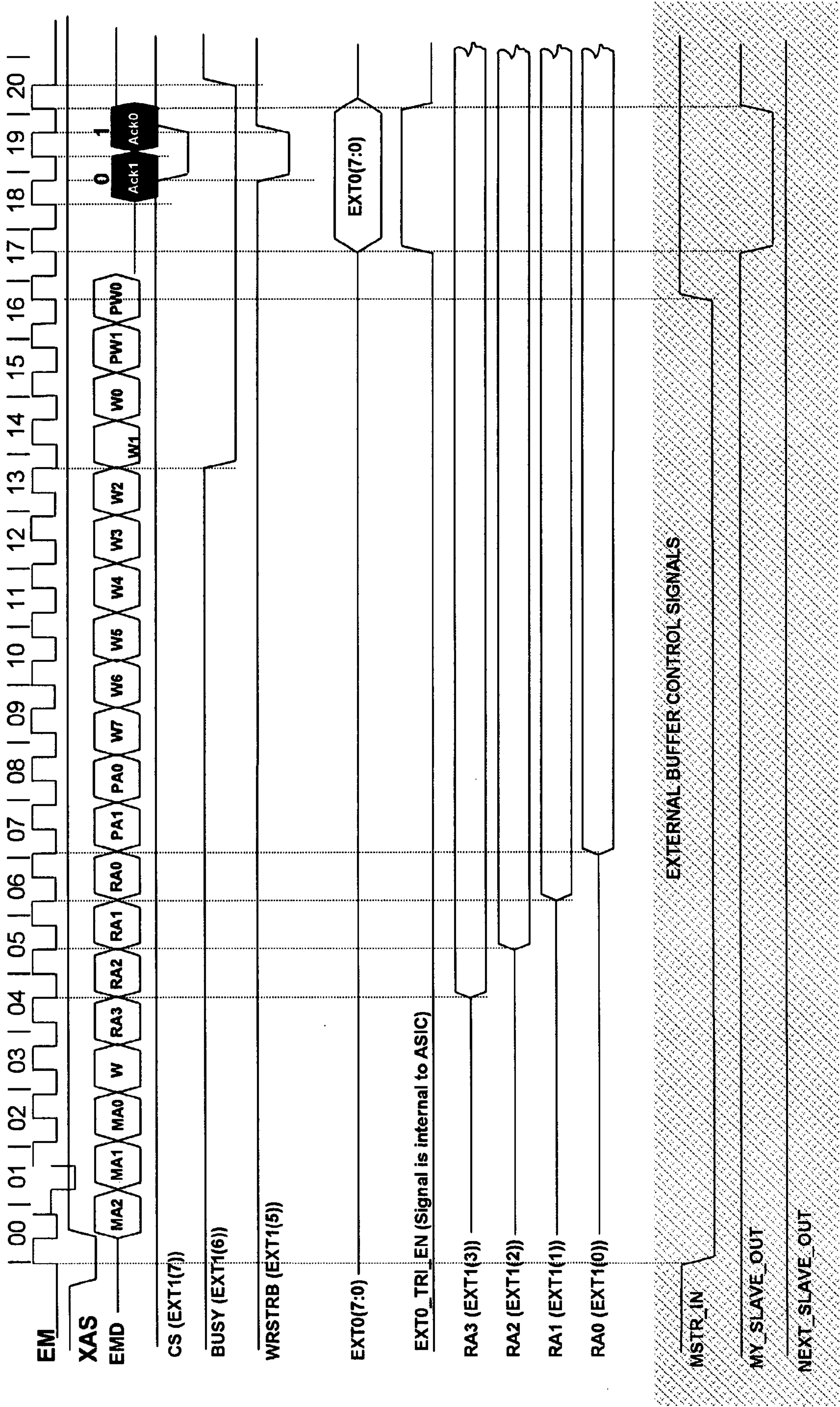


Figure 8

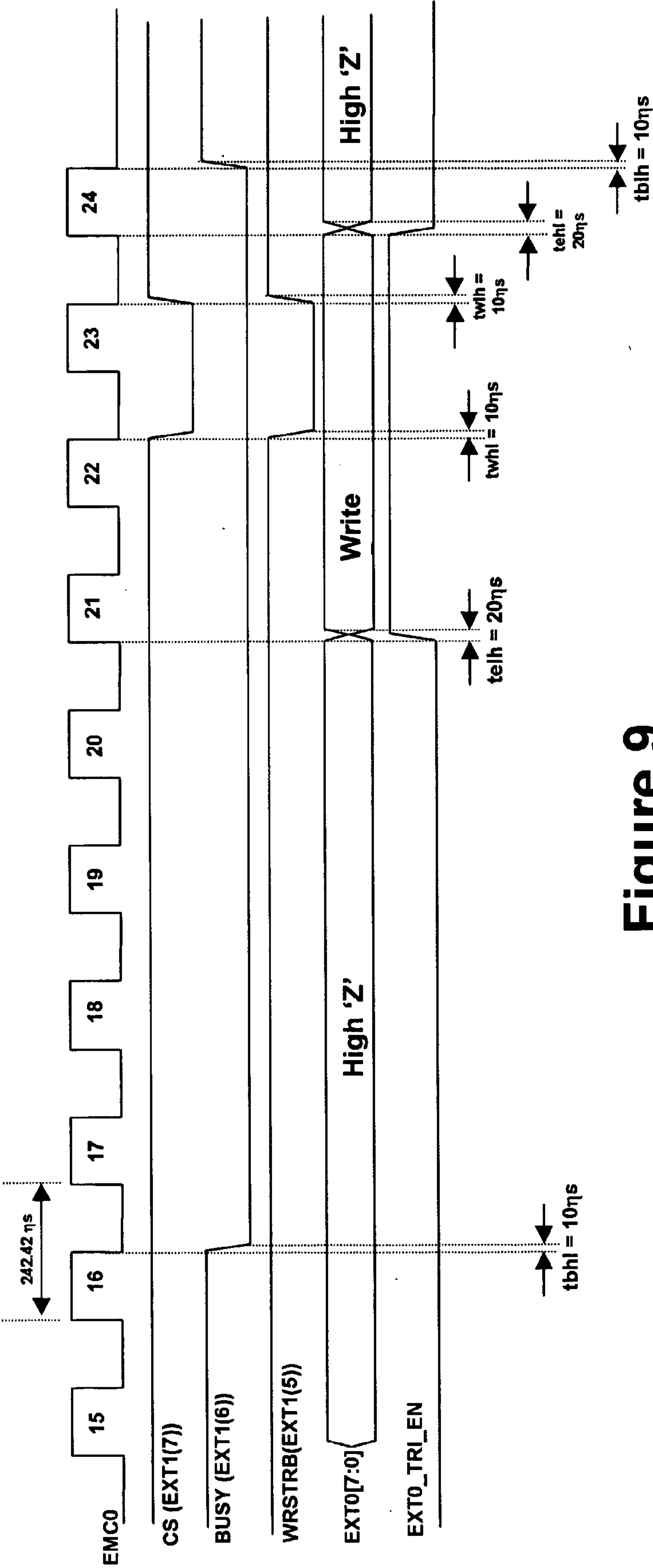


Figure 9

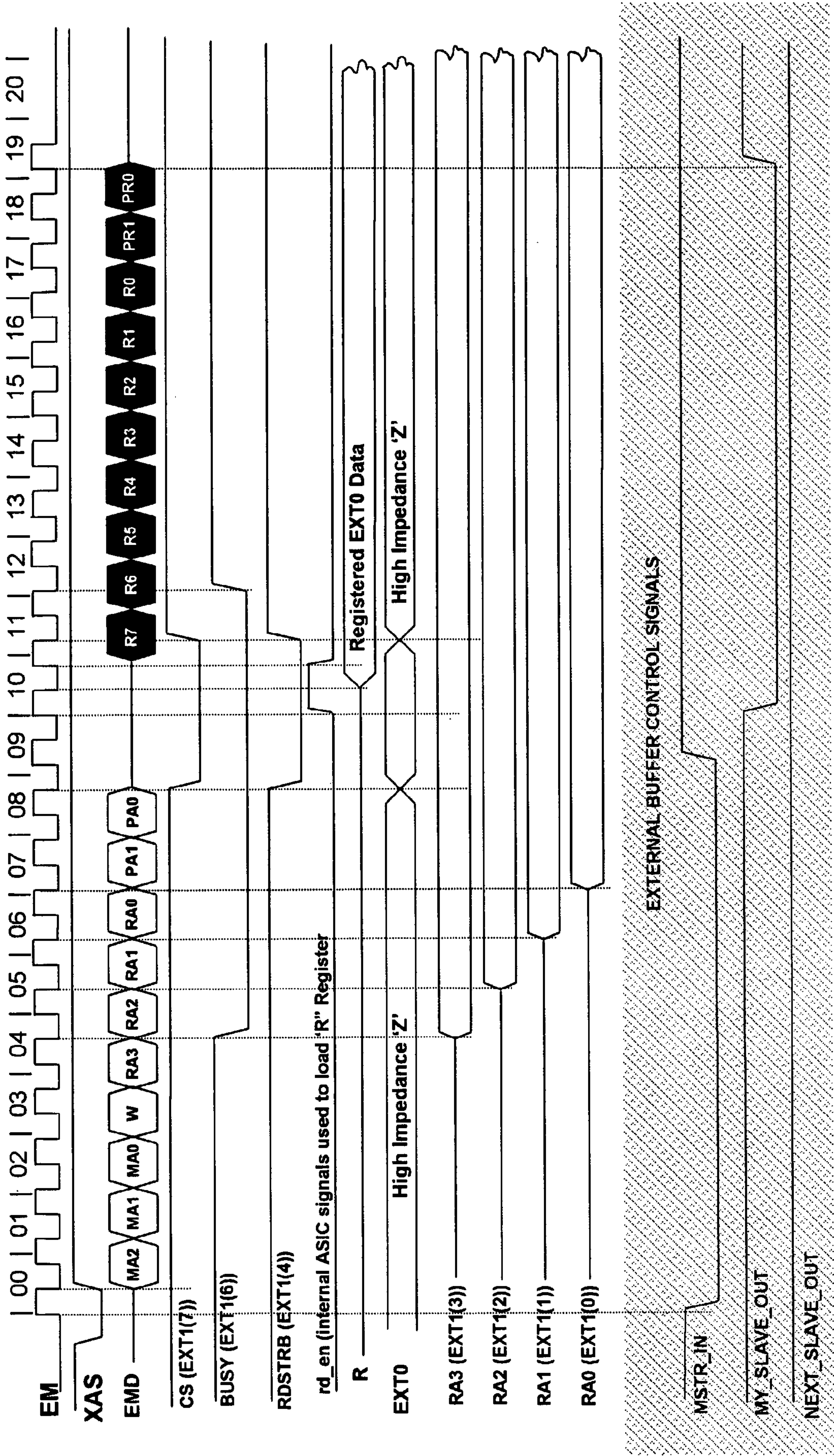


Figure 10

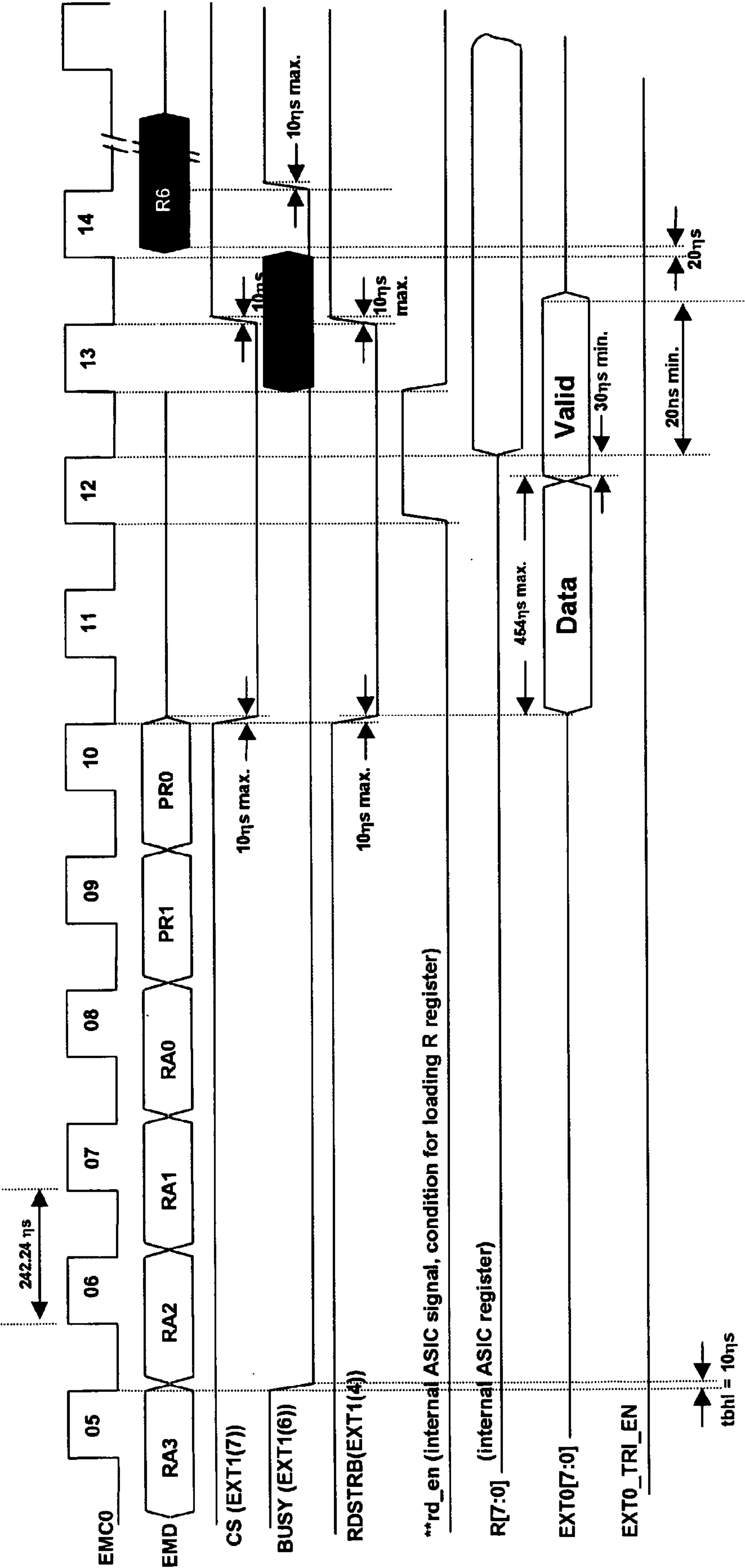


Figure 11

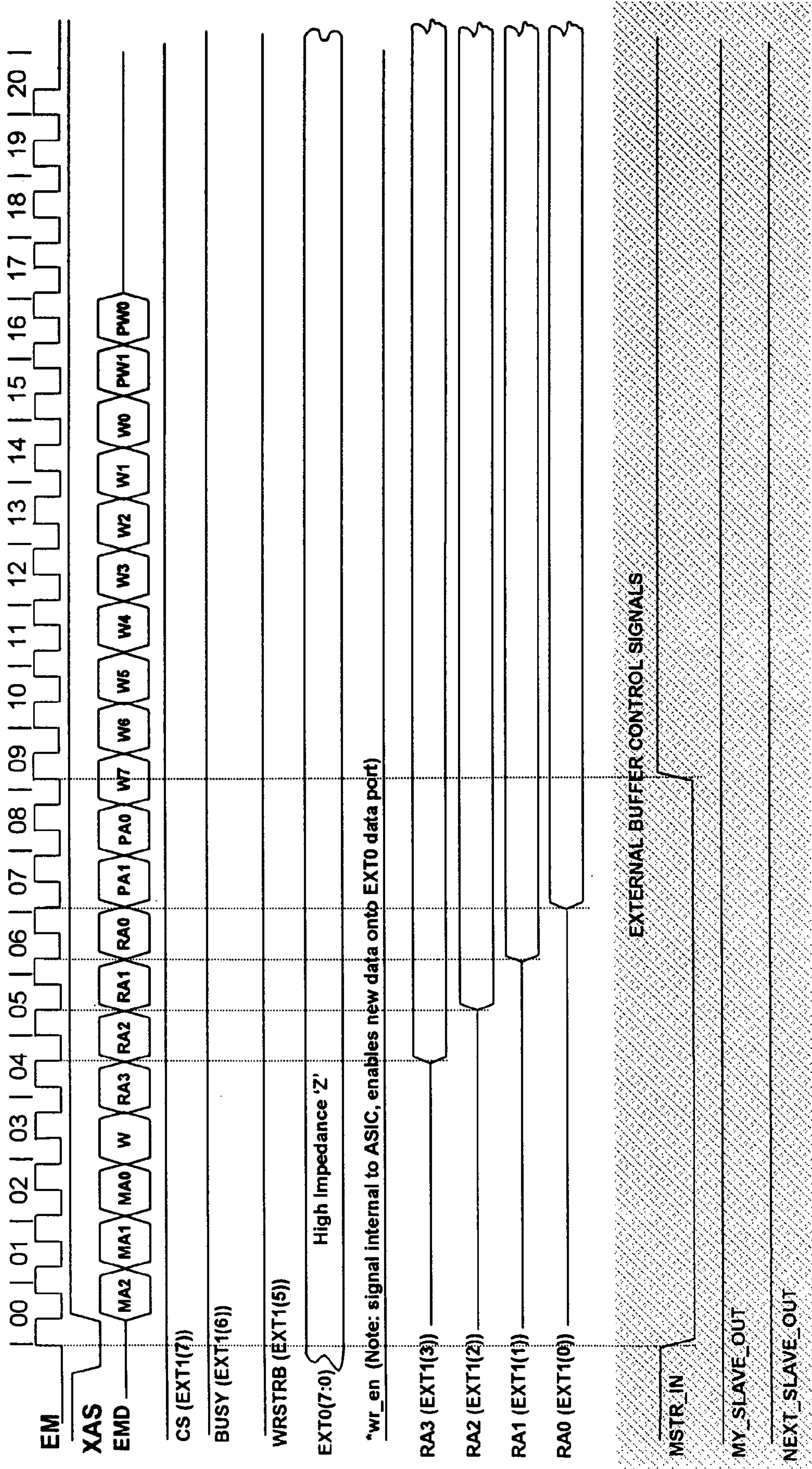


Figure 12

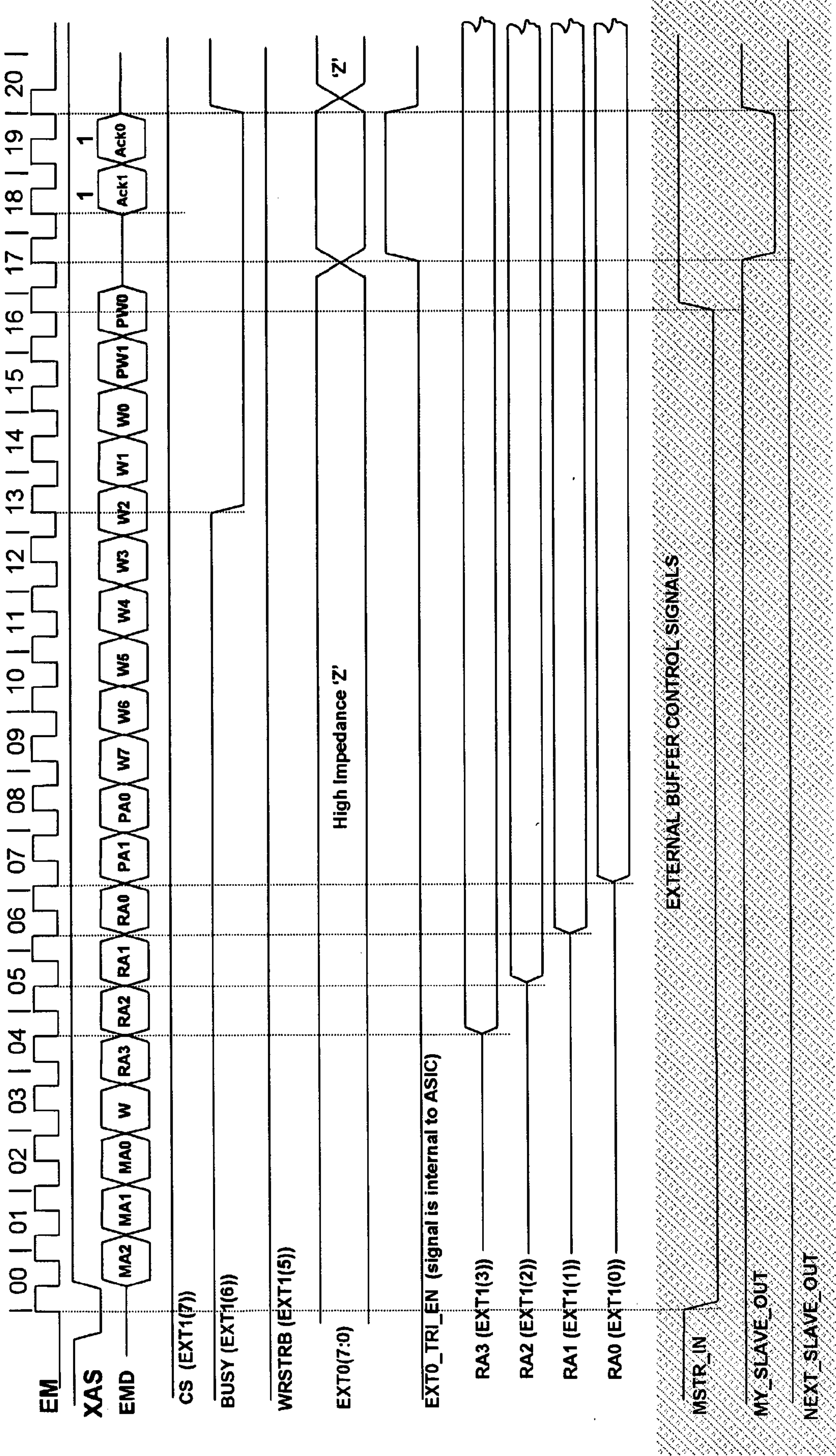


Figure 13

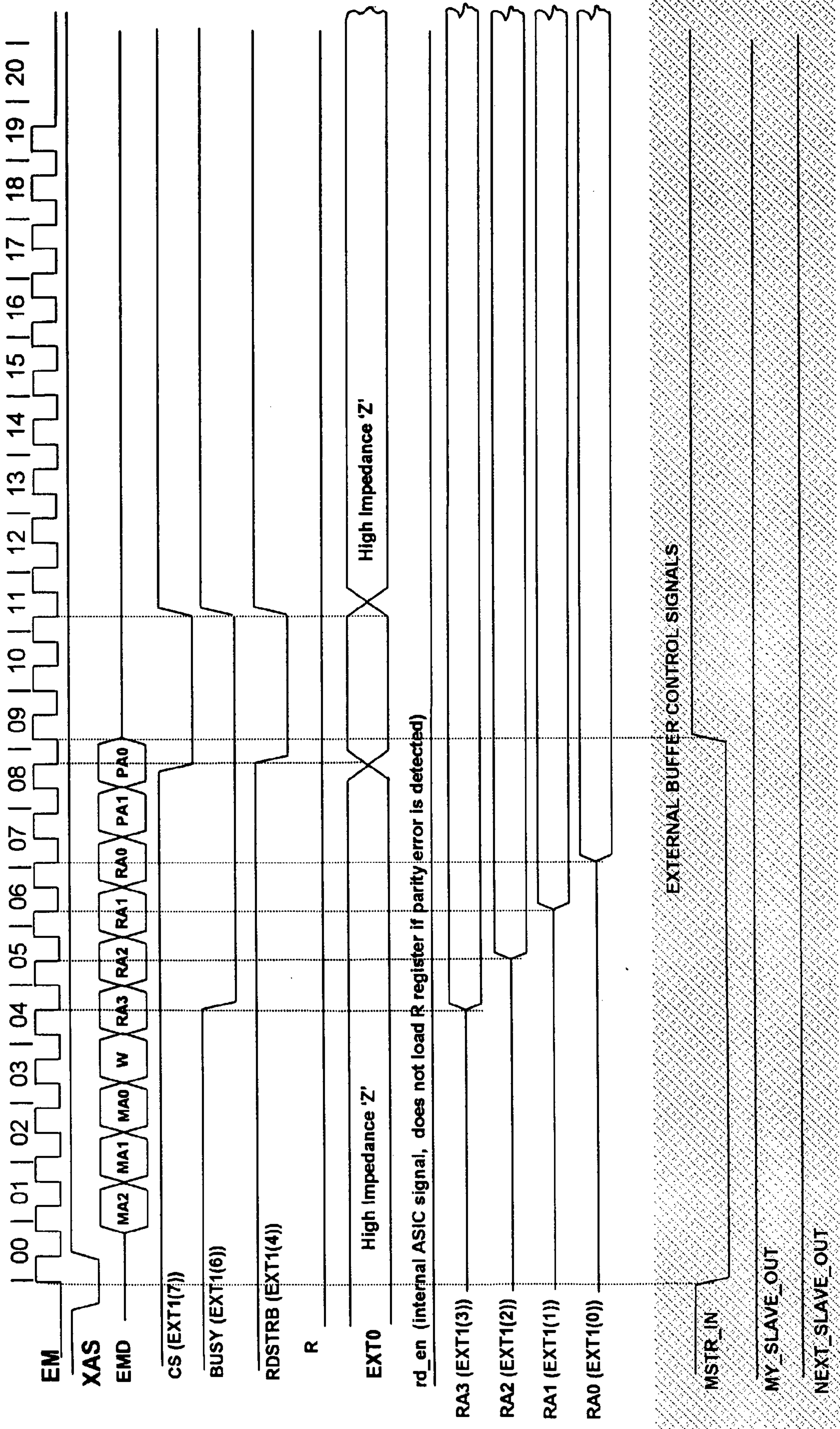


Figure 14

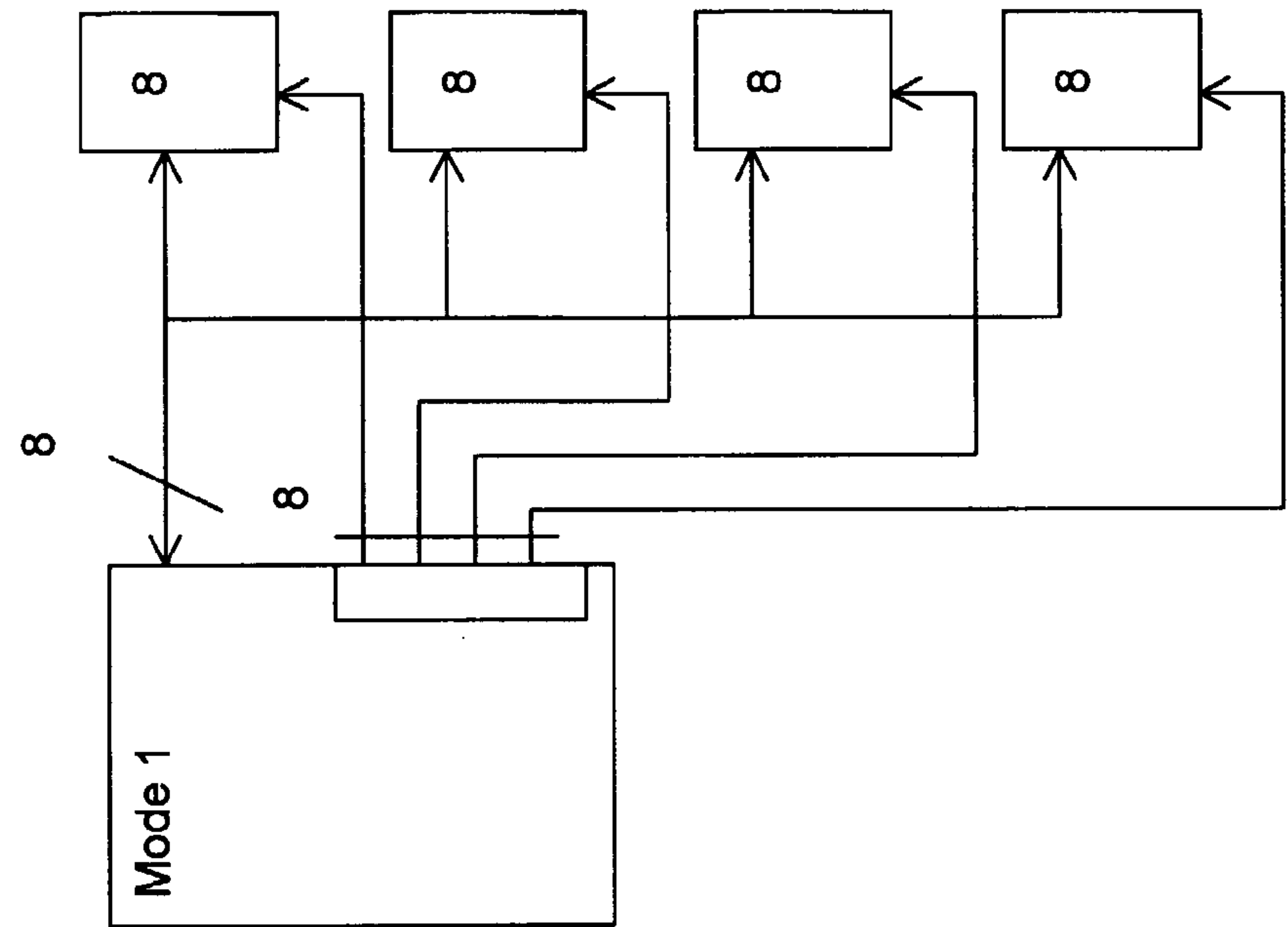


Figure 15a

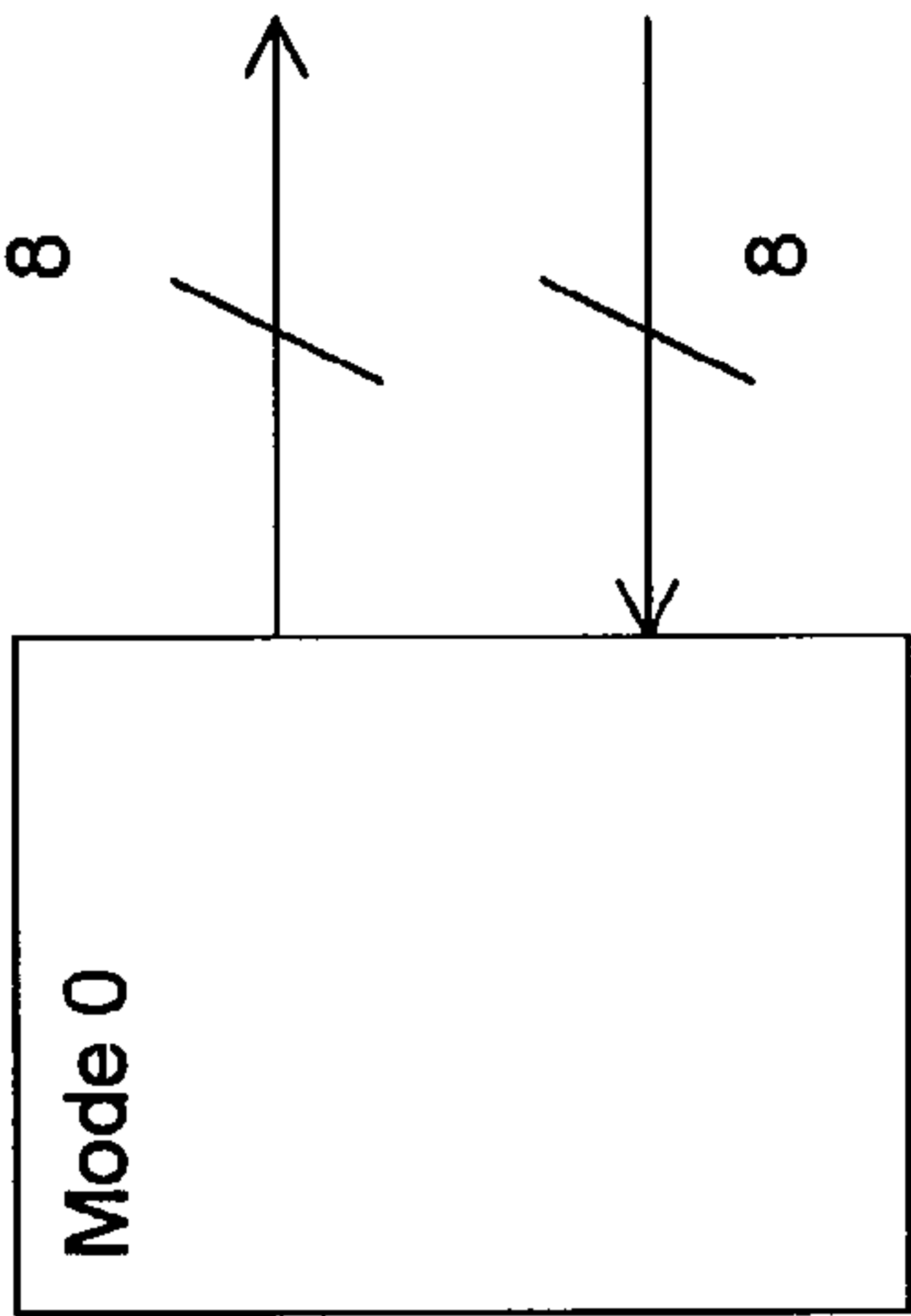


Figure 15b

