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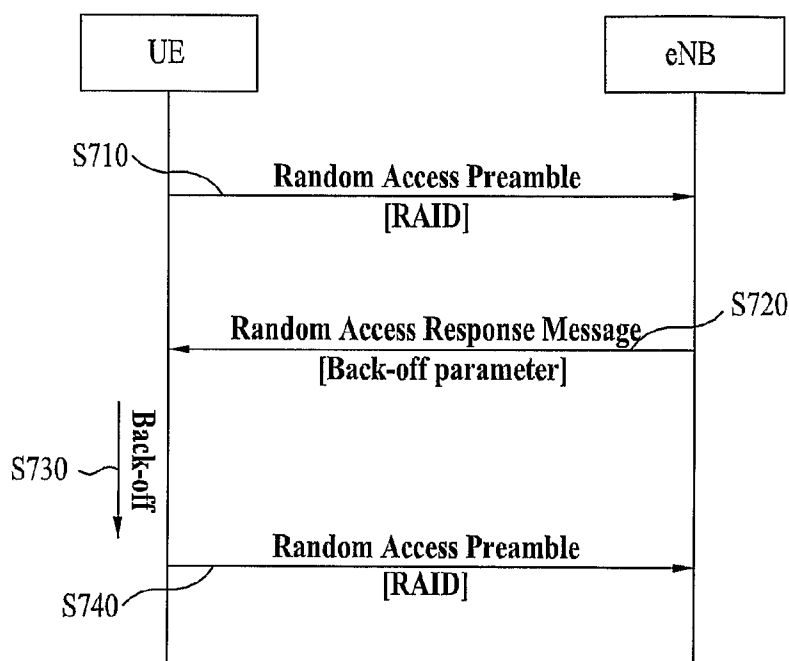
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(54) Title: METHOD FOR SIGNALING BACK-OFF INFORMATION IN RANDOM ACCESS

FIG. 7



(57) Abstract: A method for performing random access in a wireless communication system is provided. The method includes transmitting a preamble for random access in uplink, receiving a random access response message including back-off information as a response to the preamble, and performing back-off using the back-off information when the random access has failed.



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METHOD FOR SIGNALING BACK-OFF INFORMATION IN RANDOM ACCESS

BACKGROUND OF THE INVENTION5 Field of the Invention

The present invention relates to a wireless communication system, and more particularly, to a method for performing random access in a wireless communication system.

10 Discussion of the Related Art

The E-UMTS system is an evolved version of the conventional WCDMA UMTS system and basic standardization thereof is in progress under the 3rd Generation Partnership Project (3GPP). The E-UMTS is also referred to as a Long
15 Term Evolution (LTE) system. For details of the technical specifications of the UMTS and E-UMTS, refer to Release 7 and Release 8 of "3rd Generation Partnership Project; Technical Specification Group Radio Access Network".

The E-UMTS mainly includes a User Equipment (UE), a base
20 station (or eNB or eNode B), and an Access Gateway (AG) which is located at an end of a network (E-UTRAN) and which is connected to an external network. Generally, an eNB can simultaneously transmit multiple data streams for a broadcast service, a multicast service and/or a unicast service. The
25 AG can be divided into a part that handles processing of user

traffic and a part that handles control traffic. Here, the AG part for processing new user traffic and the AG part for processing control traffic can communicate with each other using a new interface. One or more cells may exist for one
5 eNB. An interface for transmitting user traffic or control traffic can be used between eNBs. A Core Network (CN) may include the AG and a network node or the like for user registration of the UE. An interface for discriminating between the E-UTRAN and the CN can be used. The AG manages
10 mobility of a UE on a Tracking Area (TA) basis. One TA includes a plurality of cells. When the UE has moved from a specific TA to another TA, the UE notifies the AG that the TA where the UE is located has been changed.

FIG. 1 illustrates a network structure of an Evolved
15 Universal Terrestrial Radio Access Network (E-UTRAN) system which is a mobile communication system to which the embodiment of the present invention is applied. The E-UTRAN system is an evolved version of the conventional UTRAN system. The E-UTRAN includes a base station that will also be
20 referred to as "eNode B" or "eNB". The eNBs are connected through an X2 interface. Each eNB is connected to the User Equipment (UE) through a radio interface and is connected to an Evolved Packet Core (EPC) through a S1 interface.

FIGS. 2 and 3 illustrate the configurations of a control
25 plane and a user plane of a radio interface protocol between

a UE and a UMTS Terrestrial Radio Access Network (UTRAN) based on the 3GPP radio access network standard. The radio interface protocol is divided horizontally into a physical layer, a data link layer and a network layer, and vertically
5 into a user plane for data transmission and a control plane for signaling. The protocol layers of FIGs. 2 and 3 can be divided into an L1 layer (first layer), an L2 layer (second layer) and an L3 layer (third layer) based on the lower three layers of the Open System Interconnection (OSI) reference
10 model widely known in communication systems.

The control plane is a passage through which control messages that a UE and a network use in order to manage calls are transmitted. The user plane is a passage through which data (e.g., voice data or Internet packet data) generated at
15 an application layer is transmitted. The following is a detailed description of the layers of the control and user planes in a radio interface protocol.

The physical layer, which is the first layer, provides an information transfer service to an upper layer using a
20 physical channel. The physical layer is connected to a Media Access Control (MAC) layer, located above the physical layer, through a transport channel. Data is transferred between the MAC layer and the physical layer through the transport channel. Data transfer between different physical layers,
25 specifically between the respective physical layers of

transmitting and receiving sides, is performed through the physical channel. The physical channel is modulated according to the Orthogonal Frequency Division Multiplexing (OFDM) method, using time and frequencies as radio resources.

5 The MAC layer of the second layer provides a service to a Radio Link Control (RLC) layer, located above the MAC layer, through a logical channel. The RLC layer of the second layer supports data transmission with reliability. The functions of the RLC layer may also be implemented through internal
10 functional blocks of the MAC layer. In this case, the RLC layer need not be existed. A PDCP layer of the second layer performs a header compression function to reduce unnecessary control information in order to efficiently transmit IP packets such as IPv4 or IPv6 packets in a radio interface
15 with a relatively narrow bandwidth.

A Radio Resource Control (RRC) layer located at the lowermost of the third layer is defined only in the control plane and is responsible for control of logical, transport, and physical channels in association with configuration, re-
20 configuration and release of Radio Bearers (RBs). The RB is a service that the second layer provides for data communication between the UE and the UTRAN. To accomplish this, the RRC layer of the UE and the RRC layer of the network exchange RRC messages. The UE is in RRC connected
25 mode if RRC connection has been established between the RRC

layer of the radio network and the RRC layer of the UE. Otherwise, the UE is in RRC idle mode.

A Non-Access Stratum (NAS) layer located above the RRC layer performs functions such as session management and mobility management.

One cell of the eNB is set to use a bandwidth such as 1.25, 2.5, 5, 10 or 20MHz to provide a downlink or uplink transmission service to UEs. Here, different cells may be set to use different bandwidths.

Downlink transport channels for transmission of data from the network to the UE include a Broadcast Channel (BCH) for transmission of system information, a Paging Channel (PCH) for transmission of paging messages and a downlink Shared Channel (SCH) for transmission of user traffic or control messages. User traffic or control messages of a downlink multicast or broadcast service may be transmitted through a downlink SCH and may also be transmitted through a downlink multicast channel (MCH). Uplink transport channels for transmission of data from the UE to the network include a Random Access Channel (RACH) for transmission of initial control messages and an uplink SCH for transmission of user traffic or control messages.

Logical channels, which are located above the transport channels and are mapped to the transport channels, include a Broadcast Control Channel (BCCH), a Paging Control Channel

(PCCH), a Common Control Channel (CCCH), a Multicast Control Channel (MCCH) and a Multicast Traffic Channel (MTCH).

FIG. 4 is a view showing an example of a physical channel structure used in an E-UMTS system. A physical channel includes several subframes on a time axis and several subcarriers on a frequency axis. Here, one subframe includes a plurality of symbols on the time axis. One subframe includes a plurality of resource blocks and one resource block includes a plurality of symbols and a plurality of subcarriers. In addition, each subframe may use certain subcarriers of certain symbols (e.g., a first symbol) of a subframe for a physical downlink control channel (PDCCH), that is, an L1/L2 control channel. In FIG. 4, an L1/L2 control information transmission area (PDCCH) and a data area (PDSCH) are shown. In one embodiment, a radio frame of 10 ms is used and one radio frame includes 10 subframes. In addition, one subframe includes two consecutive slots. The length of one slot may be 0.5 ms. In addition, one subframe includes a plurality of OFDM symbols and a portion (e.g., a first symbol) of the plurality of OFDM symbols may be used for transmitting the L1/L2 control information. A transmission time interval (TTI) which is a unit time for transmitting data is 1 ms.

A base station and a UE mostly transmit/receive data via a PDSCH, which is a physical channel, using a DL-SCH which is a

transmission channel, except a certain control signal or certain service data. Information indicating to which UE (one or a plurality of UEs) PDSCH data is transmitted and how the UE receive and decode PDSCH data is transmitted in a state of
5 being included in the PDCCH.

For example, in one embodiment, a certain PDCCH is CRC-masked with a radio network temporary identity (RNTI) "A" and information about data is transmitted using a radio resource "B" (e.g., a frequency location) and transmission format
10 information "C" (e.g., a transmission block size, modulation, coding information or the like) via a certain subframe. Then, one or more UEs located in a cell monitor the PDCCH using its RNTI information. And, a specific UE with RNTI "A" reads the PDCCH and then receive the PDSCH indicated by B and C in the
15 PDCCH information.

FIG. 5 is a process flow diagram illustrating a contention-based random access procedure.

The random access procedure is used to transmit short-length data in uplink. For example, the random access
20 procedure is performed upon initial access in an RRC idle mode, upon initial access after radio link failure, upon handover requiring the random access procedure, and upon the occurrence of uplink/downlink data requiring the random access procedure during an RRC connected mode. Some RRC
25 messages such as an RRC connection request message, a cell

update message, and an URA update message are transmitted using a random access procedure. Logical channels such as a Common Control Channel (CCCH), a Dedicated Control Channel (DCCH), or a Dedicated Traffic Channel (DTCH) can be mapped to a transport channel (RACH). The transport channel (RACH) can be mapped to a physical channel (e.g., Physical Random Access Channel (PRACH)). When a UE MAC layer instructs a UE physical layer to transmit a PRACH, the UE physical layer first selects an access slot and a signature and transmits a PRACH preamble in uplink. The random access procedure is divided into a contention-based procedure and a non-contention-based procedure.

As shown in FIG. 5, a UE receives and stores information regarding random access from an eNB through system information. Thereafter, when random access is needed, the UE transmits a random access preamble (message 1) to the eNB (S510). After transmitting the random access preamble (message 1), the UE monitors a PDCCH during a predetermined period of time in order to receive a random access response message. After receiving the random access preamble from the UE, the eNB transmits a random access response (message 2) to the UE (S520). Specifically, downlink scheduling information for the random access response message can be CRC-masked with a Random Access-RNTI and can be transmitted through an L1/L2 control channel (PDCCH). Upon receiving the downlink

scheduling signal masked with the RA-RNTI, the UE can receive and decode a random access response message from a PDSCH. Thereafter, the UE checks whether or not a random access response corresponding to the UE is present in the received
5 random access response message. Whether or not a random access response corresponding to the UE is present can be determined based on whether or not a RAID for the preamble that the UE has transmitted is present. After receiving response information, the UE transmits an uplink message
10 (Message 3) through an uplink SCH according to information regarding radio resources included in the response information(S530). After receiving the uplink message from the UE, the eNB transmits a contention resolution message (Message 4) (S540).

15 When random access has failed, the UE performs back-off. Here, the term "back-off" refers to delaying, by a UE, an access attempt by an arbitrary or predetermined time. If the UE makes an access attempt immediately after random access has failed, the access attempt is likely to fail again for
20 the same or similar reason. Accordingly, when random access has failed, the UE delays an access attempt by a predetermined time to prevent waste of radio resources due to failure of the access attempt and to increase the probability that the random access is successful.

FIG. 6 illustrates a method for signaling back-off information according to a conventional technology.

As shown in FIG. 6, an eNB transmits a back-off parameter to all UEs in the cell through system information (S610).

5 Thereafter, the UE performs its own back-off setting using a back-off parameter obtained from the system information. When random access is needed, the UE transmits a preamble for random access to the eNB (S620). The preamble may include a Random Access IDentity (RAID). Thereafter, when a random
10 access procedure has failed for some reason, the UE performs back-off (S630). Thereafter, the UE retransmits a preamble for random access to the eNB (S640).

In the conventional technology, the UE should receive and store a back-off parameter through system information before
15 making random access since the back-off parameter was broadcast through system information. Since a back-off parameter should be periodically broadcast through system information, a downlink overhead is always broadcast even when back-off is not performed since random access is
20 successful. In addition, it may also be necessary to apply a different back-off parameter due to a cause such as load. However, since a back-off parameter is broadcast through system information, each UE in the cell cannot perform different back-off.

SUMMARY OF THE INVENTION

Accordingly, the present invention is directed to a method for performing random access in a wireless communication system that substantially obviates one or more problems due to limitations and disadvantages of the related art.

An object of the present invention is to provide a method for reducing overhead and increasing the efficiency of radio resources in association with a random access procedure performed in a wireless communication system.

Another object of the present invention is to provide a method for efficiently signaling back-off information in random access.

Another object of the present invention is to provide a message format used for signaling back-off information in random access.

Additional advantages, objects, and features of the invention will be set forth in part in the description which follows and in part will become apparent to those having ordinary skill in the art upon examination of the following or may be learned from practice of the invention. The objectives and other advantages of the invention may be realized and attained by the structure particularly pointed out in the written description and claims hereof as well as the appended drawings.

To achieve these objects and other advantages and in accordance with the purpose of the invention, as embodied and broadly described herein, a method for performing random access in a wireless communication system includes
5 transmitting a preamble for random access in uplink; receiving a random access response message including back-off information as a response to the preamble; and performing back-off using the back-off information when the random access has failed. The method may further include obtaining
10 the back-off information from the random access response message. The method may further include retransmitting a preamble for random access in uplink.

In another aspect of the present invention, a method for performing random access in a wireless communication system
15 includes receiving a preamble for random access; and transmitting a random access response message including back-off information as a response to the preamble in downlink.

The embodiments of the present invention have the following advantages.

20 First, it is possible to reduce overhead and increase the efficiency of radio resources in association with a random access procedure performed in a wireless communication system.

Second, it is possible to signal back-off information only when random access is needed, thereby reducing overhead.

Third, it is possible to apply different back-off information to each UE in random access.

Fourth, it is possible to provide a message format for signaling back-off information in random access.

5 It is to be understood that both the foregoing general description and the following detailed description of the present invention are exemplary and explanatory and are intended to provide further explanation of the invention as claimed.

10

BRIEF DESCRIPTION OF THE DRAWINGS

The accompanying drawings, which are included to provide a further understanding of the invention and are incorporated in and constitute a part of this application, illustrate
15 embodiment(s) of the invention and together with the description serve to explain the principle of the invention. In the drawings:

FIG. 1 illustrates a schematic structure of an Evolved Universal Terrestrial Radio Access Network (E-UTRAN);

20 FIGS. 2 and 3 illustrate the configurations of a radio interface protocol between a UE and a UTRAN;

FIG. 4 illustrates an example physical channel structure used in an E-UMTS system;

25 FIG. 5 illustrates an example method for performing a contention-based random access procedure;

FIG. 6 illustrates a method for signaling a back-off parameter regarding random access according to a conventional technology;

FIG. 7 illustrates a method for signaling back-off information regarding random access according to an embodiment of the present invention;

FIG. 8 is a flow chart illustrating a random access procedure according to an embodiment of the present invention;

FIG. 9 illustrates a MAC PDU structure of a random access response message applied to an embodiment of the present invention;

FIG. 10 illustrates a structure of an E/R/RAID MAC sub-header;

FIG. 11 illustrates a MAC RAR structure;

FIG. 12 illustrates a MAC sub-header and a MAC RAR structure when a MAC PDU includes a dedicated back-off (BO) parameter for each RAID according to an embodiment of the present invention;

FIG. 13 illustrates a structure of a MAC sub-header including common back-off information according to an embodiment of the present invention;

FIG. 14 illustrates a MAC RAR structure when a MAC sub-header includes common back-off information according to an embodiment of the present invention;

FIG. 15 illustrates a random access response message when back-off information is included in a MAC sub-header according to an embodiment of the present invention; and

FIGs. 16 and 17 illustrate a MAC sub-header and a MAC RAR structure in the case where a special RAID is used according to an embodiment of the present invention.

DETAILED DESCRIPTION OF THE INVENTION

The above and other configurations, operations, and features of the present invention will be easily understood from the embodiments of the invention described below with reference to the accompanying drawings. The embodiments described below are examples wherein technical features of the invention are applied to an Evolved Universal Mobile Telecommunications System (E-UMTS).

Embodiment: Back-off Information Signaling using Random Access Response Message

FIG. 7 illustrates a method for performing random access according to an embodiment of the present invention.

As shown in FIG. 7, a UE transmits a random access preamble including a RAID. The random access preamble is transmitted through a Random Access Channel (RACH) (S710).

Thereafter, an eNB receives the random access preamble and transmits a random access response message including back-off information, as a response to the random access preamble, to

the UE. The UE obtains the back-off information from the received random access response message (S720). The format of the random access response message may vary according to a protocol layer that makes a response to the random access preamble. For example, the protocol layer may be a MAC layer. In this case, the random access response message may include a MAC header. The random access response message may further include a MAC Random Access Response (RAR) for one or more UEs. In this case, a MAC RAR for a specific UE may be indicated using a RAID included in a random access preamble transmitted by the UE. The back-off information may be included in a MAC header or a MAC RAR.

The back-off information indicates information required to perform back-off. For example, the back-off information may include a back-off parameter. The back-off parameter may include a delay time for retransmission of a random access preamble or an upper delay time limit. The back-off information may further include information regarding a UE to which the back-off parameter is applied. The back-off information may include a back-off parameter or an index, an identity (or identification), an indicator, simplified information, or the like for indicating a specific UE. The size of the back-off information is not limited to a specific size. For example, the size of the back-off information may

be 6 bits or less. The following table illustrates an example of the back-off parameter.

[TABLE 1]

Index	Backoff Parameter value (ms)
0	0
1	10
2	20
3	30
4	40
5	60
6	80
7	120
8	160
9	240
10	320
11	480
12	960

When the back-off parameter is indicated by an index, the index can be expressed by 4 bits.

A method for signaling back-off information using the random access response message will be described later in detail with reference to FIGs. 9 to 17. One important feature of this method is that back-off information is included in the random access response message so that back-off information is provided to the UE only when the UE has requested random access, thereby increasing the efficiency of radio resources and reducing a downlink overhead for signaling back-off.

Thereafter, when random access has failed, the UE performs back-off using the obtained back-off information (for example, a back-off parameter) (S730). Here, the term "back-off" refers to delaying, by a UE, an access attempt by an

arbitrary or predetermined time. Specifically, the UE delays transmission of a random access preamble according to the back-off parameter. More specifically, the UE may delay an access attempt by a time indicated by the back-off parameter.

5 In addition, the UE may delay an access attempt randomly within the time indicated by the back-off parameter. That is, the UE can delay an access attempt by a time that is selected with equal probability from a range of 0ms to the back-off parameter.

10 The following are the cases where random access has failed. The first case is where the UE has failed to receive a random access response message. The second case is where, although the UE has received a random access response message, no random access response corresponding to the UE is present in
15 the random access response message. The third case is that the UE has failed to perform uplink transmission using radio resources allocated through a random access response message. For example, the uplink transmission includes PUSCH transmission. The fourth case is where the UE has failed to
20 pass through a contention resolution procedure for random access. The contention resolution procedure includes a MAC contention resolution procedure or an RLC contention resolution procedure.

Depending on the cause of the failure of random access,
25 the UE may also fail to obtain back-off information from the

random access response message. In this case, the UE may use a default back-off parameter or a back-off parameter obtained from a previous random access response message. The default back-off parameter includes 0ms.

5 After performing back-off, the UE retransmits a random access preamble including a RAID (S740).

FIG. 8 is a flow chart illustrating a random access procedure according to an embodiment of the present invention.

As shown in FIG. 8, a UE receives RACH information (S802).

10 When no separate uplink transmission channel is present, the UE performs a random access procedure in the following manner (S804).

The UE initializes a random access preamble counter M to "1" (S806). Thereafter, the UE decides whether or not to
15 perform a random access procedure (S808). Thereafter, when the UE has decided to continue the random access procedure, the UE determines whether or not the value of the counter M is equal to or less than a maximum allowable value Mmax (S810). When the value of the counter M exceeds the maximum
20 allowable value Mmax, the UE terminates the random access procedure (S832). Otherwise, the UE continues the random access procedure. As needed, the UE updates RACH information (S812) and selects a preamble signature/resource (S814). Thereafter, the UE determines preamble transmission power
25 (S816) and transmits a PRACH preamble to an eNB (S818).

After transmitting the preamble, the UE determines whether or not an initial access procedure using the preamble is successful while monitoring an L1/L2 control region in order to receive a random access response message during a predetermined time (S820). For example, the UE determines that the initial access procedure has failed when the UE has failed to receive a downlink scheduling signal masked with an RA-RNTI within the predetermined time. The UE also determines that the initial access procedure has failed when information regarding the UE is not included in the random access response message although the UE has received a downlink scheduling signal masked with an RA-RNTI and has decoded the random access response message. The UE increases the counter value "M" by 1 (S828) and performs back-off when needed (①, S830). In this case, the back-off parameter is a default value including 0ms or a value obtained through a previous random access response message.

When the UE has successfully received the random access response message, the UE adjust transmission timing using information obtained from the random access response and transmits an uplink message through an uplink SCH (S822). When the eNB has not successfully received the uplink message, the UE increases the counter value "M" by 1 (S828) and performs back-off using a back-off parameter obtained from the random access response message when needed (②, S830).

When the uplink message has been successfully transmitted to the eNB, the eNB transmits a contention resolution message to the UE in downlink. Thereafter, when the UE has passed through the contention resolution procedure, the UE
5 terminates the random access procedure. On the other hand, when the UE has passed through the contention resolution procedure, the UE increases M by 1 (S828) and performs back-off using a back-off parameter obtained from the random access response message when needed (③, S830).

10 After performing back-off, the UE repeats the random access procedure, starting from step S808.

FIG. 9 illustrates a MAC PDU structure of a random access response message applied to an embodiment of the present invention. A MAC PDU is transmitted to a UE through a
15 downlink SCH channel.

As shown in FIG. 9, one MAC PDU includes a MAC header and MAC payload. The MAC header includes one or more MAC sub-headers (for example, E/R/RAID sub-headers) and the MAC payload includes one or more MAC RARs. The MAC sub-headers
20 correspond to the MAC RARs, respectively.

FIG. 10 illustrates a structure of an E/R/RAID MAC sub-header.

As shown in FIG. 10, the E/R/RAID MAC sub-header includes an E field, an R field, and a RAID field. The fields of the

MAC sub-header are arranged on an octet (byte) basis. The following is a description of each of the fields.

- E field: Extension field. This field serves as a flag indicating whether or not an additional field is included in the MAC header. When the E field is set to "1", this indicates that another E/R/RAID field (sub-header) is present. When the E field is set to "0", this indicates that a MAC RAR starts from the next byte.

- R field: Field including a reserved bit. This field currently has no usage and can be defined and used as needed at a later time. The size of the R field is 1 bit.

- RAID field: Field identifying a random access preamble transmitted from a UE. The size of the RAID field is 6 bits.

Back-off information signaling using a MAC sub-header will be described later in detail.

FIG. 11 illustrates a MAC RAR structure. As shown in FIG. 11, a MAC RAR includes a TA field, a UL grant field, and a T-CRNTI field. The fields of the MAC RAR are arranged on an octet (byte) basis. The following is a description of each of the fields.

- TA field: Timing Advance field. The TA field indicates uplink transmission timing required for timing synchronization. The size of the TA field is 11 bits.

- UL Grant field: Uplink grant field. The UL grant field indicates allocation information of radio resources used in uplink. The size of the UL grant field is 21 bits.

- T-CRNTI field: Temporary CRNTI field. The T-CRNTI field indicates a temporary identification used by the UE before another random access procedure is initiated or before the contention resolution procedure has passed. The size of the T-CRNTI field is 16 bits.

The following is a detailed description of signaling of back-off information using a random access response message (for example, a MAC PDU).

Back-off information included in a random access response message may be dedicated back-off information or common back-off information depending on implementations. In this specification, the term "dedicated back-off information" refers to back-off information that is individually applied to each of a plurality of UEs that use different random access preambles. For example, the dedicated back-off information may be defined for each RAID used to identify a random access preamble transmitted by the UE. In this case, the same back-off information is applied to UEs that use the same RAID while different back-off information is applied to UEs that use different RAIDs. The term "common back-off information" refers to back-off information that is commonly applied to UEs that use different random access preambles.

For example, common back-off information may be defined for a specific RAID group or for each random access response message. Thus, the same back-off information is applied to UEs even when the UEs use different RAIDs.

- 5 The following are descriptions of three possible options regarding signaling back-off information through a MAC PDU.

Option 1: Inclusion of Dedicated Back-off Information

Back-off information can be included in a random access response message for each RAID. Specifically, a back-off parameter can be included in a respective MAC RAR corresponding to each RAID. Since a back-off parameter is mapped one-to-one to a RAID, the back-off parameter is RAID-specific. That is, the back-off parameter is commonly applied to one or more UEs that have accessed a random access response message using the same RAID. On the other hand, different back-off parameters are applied to one or more UEs that have accessed a random access response message using different RAIDs. Since a different back-off parameter can be applied to each RAID in this manner, the back-off parameter can be used to control random access load.

FIG. 12 illustrates a MAC sub-header and a MAC RAR structure when a MAC PDU includes a dedicated back-off (BO) parameter for each RAID according to an embodiment of the present invention.

As shown in FIG. 12, the MAC sub-header has the same structure as the general E/R/RAID sub-header. That is, the MAC sub-header includes an E field, an R field, and a RAID field. On the other hand, a MAC RAR corresponding to the RAID includes a back-off parameter, unlike the general MAC RAR. Since the back-off parameter is provided for each RAID, a MAC RAR having a back-off parameter may also include a TA field. When a UE retransmits a random access preamble after performing back-off, the UE can control transmission timing of the random access preamble using a value included in the TA field of the MAC RAR. The MAC RAR may include only the back-off parameter although "TA" is illustrated as being additionally included in a MAC RAR having a back-off parameter in FIG. 12.

Option 2: Common Back-off Information

A random access response message may include a back-off parameter commonly used by a plurality of UEs. In this case, the common back-off parameter in the random access response message is applied to UEs that have simultaneously accessed the random access response message using different RAIDs. In this regard, the following two sub-options can be considered.

Option 2a: Replacement of RAID with Back-off Parameter

FIG. 13 illustrates a structure of a MAC sub-header including common back-off information according to an embodiment of the present invention. As shown in FIG. 13,

when an R bit in the E/R/RAID sub-header is set to "0", a normal RAID is included in the E/R/RAID sub-header (see FIG. 13(a)). On the other hand, when the R bit is set to "1", a back-off parameter is included in the RAID field (see FIG. 13(b)). Accordingly, when a UE confirms that an R bit in a specific sub-header is set to "0", the UE assumes that a RAID field in the specific sub-header is a back-off parameter. This analysis of the R bit is illustrative and the R bit can be interpreted in the reverse manner.

For this reason, when a back-off parameter is included in a specific MAC sub-header, a RAID is not included in the specific MAC sub-header. On the other hand, when a RAID is included in a specific MAC sub-header, a back-off parameter is not included in the specific MAC sub-header. In FIG. 13, the R bit is illustrated as indicating whether or not a back-off parameter field is included in the MAC sub-header. However, FIG. 13 is only illustrative and whether or not the back-off parameter field is included in the MAC sub-header can be indicated by any field having a flag indicating the type of the back-off parameter field or indicating presence or absence of the back-off parameter field. For example, whether or not a back-off parameter field is included in a MAC sub-header can be indicated by an E field or an additional field included in the MAC sub-header. In this

case, the size of the field indicating the type or presence of the back-off parameter field may be 1 bit.

The specific MAC sub-header including the back-off parameter may or may not correspond to a specific MAC RAR.

5 The case where the specific MAC sub-header corresponds to a specific MAC RAR will be described later with reference to FIG. 14. In the case where the specific MAC sub-header does not correspond to a specific MAC RAR, a MAC PDU for a random access response may be constructed as follows.

10 (1) Include only a specific sub-header including back-off parameter

The random access response message may include only a specific MAC sub-header including a back-off parameter. In this case, the back-off parameter is applied to all UEs that
15 receive the specific MAC sub-header. Accordingly, when a UE has received a random access response message including only one specific MAC sub-header including a back-off parameter, the UE obtains and stores the back-off parameter without the need to confirm whether or not a RAID of a preamble
20 transmitted by the UE is included in the random access response message. In this case, a corresponding MAC RAR is not included in the random access response message. Accordingly, all UEs that have received the random access response message assume that random access has failed and
25 perform back-off according to the back-off parameter.

(2) Additionally include normal sub-header without back-off parameter

The random access response message may include a sub-header including a back-off parameter and further include one or more normal sub-headers including no back-off parameter. The random access response message also includes normal MAC RARs corresponding respectively to the normal MAC sub-headers. The structures of the normal MAC sub-header and the normal MAC RAR are identical to those described above with reference to FIGs. 10 and 11. In this case, the back-off parameter is commonly applied to a plurality of UEs that receive the MAC sub-header. However, since the random access response message includes a MAC sub-header having a RAID and a corresponding MAC RAR, the scope of UEs to which the back-off parameter is applied is more restricted.

Specifically, when a UE has received a random access response message, the UE determines whether or not a RAID of the UE is included in a MAC header of the message. When the RAID of the UE is included in the MAC header, the UE can obtain information regarding uplink transmission timing, uplink radio resource allocation information for uplink SCH transmission, and temporary UE ID information from a corresponding MAC RAR. On the other hand, when the UE has failed to find a RAID of the UE in the MAC header, the UE obtains and stores a back-off parameter from a specific sub-

header including the back-off parameter. In this case, it is assumed that the random access procedure has failed since the UE has failed to find a RAID of the UE in the random access response message. Accordingly, the UE performs back-off
5 using the back-off parameter obtained from the specific sub-header. When the random access response message includes both the normal MAC sub-header and the specific MAC sub-header, the order of the two sub-headers is not limited. However, taking into consideration the fact that it is
10 desirable that back-off be applied when the UE has failed to find a RAID of the UE, it is more preferable that the specific MAC sub-header having a back-off parameter be located at the end of the MAC header.

FIG. 14 illustrates a MAC RAR structure when a MAC sub-
15 header includes common back-off information according to an embodiment of the present invention. As shown in FIG. 14, the structure of the MAC RAR is determined according to the type of the corresponding MAC sub-header. When the corresponding MAC sub-header includes a RAID (i.e., does not
20 include a back-off parameter), the MAC RAR has a normal structure illustrated in FIG. 11 (see FIG. 14(a)). That is, the MAC RAR includes a Timing Advance (TA) field indicating uplink transmission timing, a UL grant field including uplink radio resource allocation information for uplink SCH

transmission, and a T-CRNTI field that serves as a temporary UE identification.

As shown in FIG. 13, the random access response message in Option 2a may not include a MAC RAR corresponding to a MAC sub-header having a back-off parameter. However, the random access response message may further include a specific MAC RAR in order to more selectively control the scope of UEs to which the back-off parameter is applied. In this case, the specific MAC RAR may include UE information (e.g., a UE list) regarding one or more UEs to which the back-off parameter is applied. For example, the UE information may be a RAID for a random access preamble transmitted by a UE. Specifically, the specific MAC RAR may include a list of one or more different RAIDs. More specifically, a list of RAIDs in the specific MAC RAR may include an E field, an R field, and a RAID field. Here, the E field and the R field are identical to those described above with reference to the MAC E/R/RAID sub-header. By introducing a specific MAC RAR which corresponds to a MAC sub-header having a back-off parameter and includes a RAID list, the eNB can more selectively transmit a positive ACK to specific UEs and allow some other specific UEs to perform back-off.

Specifically, when a UE has received a random access response message, the UE determines whether or not a RAID of the UE is included in a MAC header of the message. When a

RAID of the UE is included in a MAC header of the message, the UE can obtain information regarding uplink transmission timing, uplink radio resource allocation information for uplink SCH transmission, and temporary UE ID information from a corresponding MAC RAR. On the other hand, when the UE has failed to find a RAID of the UE in the MAC header but instead has found only a specific sub-header including a back-off parameter, the UE determines whether or not a specific MAC RAR corresponding to the specific MAC sub-header is present. When no specific MAC RAR is present, the UE performs back-off using the back-off parameter included in the specific MAC sub-header. On the other hand, when a specific MAC RAR is present, the UE determines whether or not a RAID of the UE is included in the specific MAC RAR. When the UE has found a RAID of the UE in the specific MAC RAR, the UE obtains a back-off parameter from the MAC sub-header. It is determined that the random access procedure has failed since the UE has not received a positive ACK although the UE has obtained the back-off parameter. Accordingly, the UE performs back-off using the back-off parameter. On the other hand, a UE, which has not found a RAID of the UE in the specific MAC RAR, performs back-off using a default back-off parameter or a back-off parameter obtained from a previous random access response message. The default back-off parameter includes 0ms. The random access response message may include two or

more MAC sub-headers having different back-off parameters and specific MAC RARs corresponding respectively to the MAC sub-headers. In this case, even when the back-off parameters have common back-off parameter characteristics, UEs that have transmitted different RAIDs can be configured to use different back-off parameters. Accordingly, the eNB can selectively apply different back-off parameters to different UEs.

When the random access response message includes both the normal MAC RAR and the specific MAC RAR, the order of the two MAC RARs is not limited. In this regard, it is more preferable that the specific MAC sub-header having a back-off parameter be located at the end of the MAC header as described above. Thus, it is also preferable that the specific MAC RAR be located last among the MAC RARs.

FIG. 15 illustrates a random access response message when back-off information is included in a MAC sub-header according to an embodiment of the present invention. The random access response message of FIG. 15 is a combination of the MAC sub-header structure of FIG. 13 and the MAC RAR structure of FIG. 14. The random access response message structure shown in FIG. 15 can be represented by the following table.

[TABLE 2]

	MAC Header	MAC Payload
1	One 2nd MAC sub-header	-

2	One or more 1st MAC sub-headers ^{A)} + One 2nd MAC sub-header ^{B)}	One or more 1st MAC RARs ^{C)}
3	One or more 1st MAC sub-headers + One 2nd MAC sub-header	One or more 1st MAC RARs + One 2nd MAC RAR ^{D)}
4	One or more 1st MAC sub-headers + Two or more 2nd MAC sub-headers	One or more 1st MAC RARs + Two or more 2nd MAC RARs

A) Indicates a normal E/R/RAID sub-header (see FIG. 13A)

B) Indicates a MAC sub-header including a back-off parameter (see FIG. 13(b))

C) Indicates a normal MAC RAR corresponding to an E/R/RAID sub-header (see FIG. 14(a))

D) Indicates a MAC RAR which corresponds to a MAC sub-header including a back-off parameter and includes a RAID list (see FIG. 14(b))

Option 2b: Special RAID for Common Back-off Parameter

10 A certain RAID can be used only to indicate common back-off information. In this specification, the certain RAID is referred to as a "special RAID". The following is a detailed description thereof.

FIGs. 16 and 17 illustrate a MAC sub-header and a MAC RAR structure in the case where a special RAID is used according to an embodiment of the present invention. The MAC sub-header and the MAC RAR structure according to Option 2b are similar to those of Option 2a. The difference is that a special RAID rather than a back-off parameter is included in a MAC sub-header and a back-off parameter is included in a MAC RAR corresponding to the MAC sub-header. Although a list

of RAIDs is illustrated as being additionally included in the MAC RAR, the MAC RAR may include only the back-off parameter. Other basic operations or features are similar to those of Option 2a.

5 Specifically, when a UE has received a random access response message, the UE determines whether or not a RAID of the UE is included in a MAC header of the message. When a RAID of the UE is included in a MAC header of the message, the UE can obtain information regarding uplink transmission
10 timing, uplink radio resource allocation information for uplink SCH transmission, and temporary UE ID information from a corresponding MAC RAR. On the other hand, when the UE has failed to find a RAID of the UE in the MAC header but instead has found only a special RAID, the UE refers to a specific
15 MAC RAR corresponding to a sub-header having the special RAID. When the specific MAC RAR includes a back-off parameter alone, all UEs which have referred to the specific MAC RAR perform back-off after obtaining the back-off parameter. On the other hand, when a list of RAIDs is included in the specific
20 MAC RAR, the UE determines whether or not a RAID of the UE is included in the specific MAC RAR. When the UE has found a RAID of the UE in the specific MAC RAR, the UE obtains a back-off parameter from the MAC RAR and then performs back-off. On the other hand, a UE, which has not found a RAID of
25 the UE in the specific MAC RAR, performs back-off using a

default back-off parameter or a back-off parameter obtained from a previous random access response message. The default back-off parameter includes 0ms. Accordingly, the eNB can selectively apply different back-off parameters to different
5 UEs.

When the random access response message includes both the normal MAC RAR and the specific MAC RAR, the order of the two MAC RARs is not limited. However, taking into consideration the fact that it is desirable that back-off be applied when
10 the UE has failed to find a RAID of the UE in the MAC sub-header, it is preferable that a sub-header having a special RAID and a corresponding specific MAC RAR be located last among the MAC RARs.

The above embodiments are provided by combining components
15 and features of the present invention in specific forms. The components or features of the present invention should be considered optional if not explicitly stated otherwise. The components or features may be implemented without being combined with other components or features. The embodiments
20 of the present invention may also be provided by combining some of the components and/or features. The order of the operations described above in the embodiments of the present invention may be changed. Some components or features of one embodiment may be included in another embodiment or may be
25 replaced with corresponding components or features of another

embodiment. It will be apparent that claims which are not explicitly dependent on each other can be combined to provide an embodiment or new claims can be added through amendment after this application is filed.

5 The embodiments of the present invention have been described focusing mainly on the data communication relationship between a terminal and a Base Station (BS). Specific operations which have been described as being performed by the BS may also be performed by an upper node as
10 needed. That is, it will be apparent to those skilled in the art that the BS or any other network node may perform various operations for communication with terminals in a network including a number of network nodes including BSs. The term "base station (BS)" may be replaced with another term such as
15 "fixed station", "Node B", "eNode B (eNB)", or "access point". The term "terminal" may also be replaced with another term such as "user equipment (UE)", "mobile station (MS)", or "mobile subscriber station (MSS)".

 The embodiments of the present invention can be
20 implemented by hardware, firmware, software, or any combination thereof. In the case where the present invention is implemented by hardware, an embodiment of the present invention may be implemented by one or more application specific integrated circuits (ASICs), digital signal
25 processors (DSPs), digital signal processing devices (DSPDs),

programmable logic devices (PLDs), field programmable gate arrays (FPGAs), processors, controllers, microcontrollers, microprocessors, or the like.

In the case where the present invention is implemented by
5 firmware or software, the embodiments of the present invention may be implemented in the form of modules, processes, functions, or the like which perform the features or operations described above. Software code can be stored in a memory unit so as to be executed by a processor. The
10 memory unit may be located inside or outside the processor and can communicate data with the processor through a variety of known means.

As is apparent from the above description, the embodiments of the present invention have the following advantages.

15 First, it is possible to reduce overhead and increase the efficiency of radio resources in association with a random access procedure performed in a wireless communication system.

Second, it is possible to signal back-off information only when random access is needed, thereby reducing overhead.

20 Third, it is possible to apply different back-off information to each UE in random access.

Fourth, it is possible to provide a message format for signaling back-off information in random access.

Those skilled in the art will appreciate that the present
25 invention may be embodied in other specific forms than those

set forth herein without departing from the spirit and essential characteristics of the present invention. The above description is therefore to be construed in all aspects as illustrative and not restrictive. The scope of the invention
5 should be determined by reasonable interpretation of the appended claims and all changes coming within the equivalency range of the invention are intended to be embraced in the scope of the invention.

The present invention can be applied to a wireless
10 communication system, and more particularly, to a method for performing random access in a wireless communication system.

WHAT IS CLAIMED IS:

1. A method for performing random access in a wireless communication system, the method comprising:

5 transmitting a preamble for random access in uplink;
receiving a random access response message including back-off information as a response to the preamble; and
performing back-off using the back-off information when the random access has failed.

10 2. The method according to claim 1, wherein the back-off information is common back-off information or dedicated back-off information.

15 3. The method according to claim 1, wherein the back-off information includes a back-off parameter.

4. The method according to claim 3, wherein the back-off information further includes information regarding a User
20 Equipment (UE) to which the back-off parameter is applied.

5. The method according to claim 1, wherein the back-off information is included in a Medium Access Control (MAC) header.

6. The method according to claim 5, wherein the MAC header includes one or more sub-headers and the back-off information is included in at least one specific sub-header among the one or more sub-headers.

5

7. The method according to claim 6, wherein the specific sub-header includes a first field indicating the back-off information and a second field indicating presence of the first field.

10

8. The method according to claim 6, wherein the random access response message includes one specific sub-header and a MAC Random Access Response (RAR) corresponding to the specific sub-header is not present in the random access response message.

15

9. The method according to claim 6, wherein the random access response message further includes one or more MAC RARs and at least one specific MAC RAR among the one or more MAC RARs indicates one or more UEs to which the back-off information is applied.

20

10. The method according to claim 9, wherein the specific MAC RAR includes a Random Access IDentity (RAID)

corresponding to a UE to which the back-off information is applied.

11. The method according to claim 9, wherein the
5 specific MAC RAR is a last MAC RAR among the one or more MAC RARs.

12. A method for performing random access in a wireless communication system, the method comprising:

10 receiving a preamble for random access; and

transmitting a random access response message including back-off information as a response to the preamble in downlink.

15 13. The method according to claim 12, wherein the back-off information includes a back-off parameter.

14. The method according to claim 12, wherein the back-off information further includes information regarding a User
20 Equipment (UE) to which the back-off parameter is applied.

15. The method according to claim 12, wherein the back-off information is included in a Medium Access Control (MAC) header.

FIG. 1

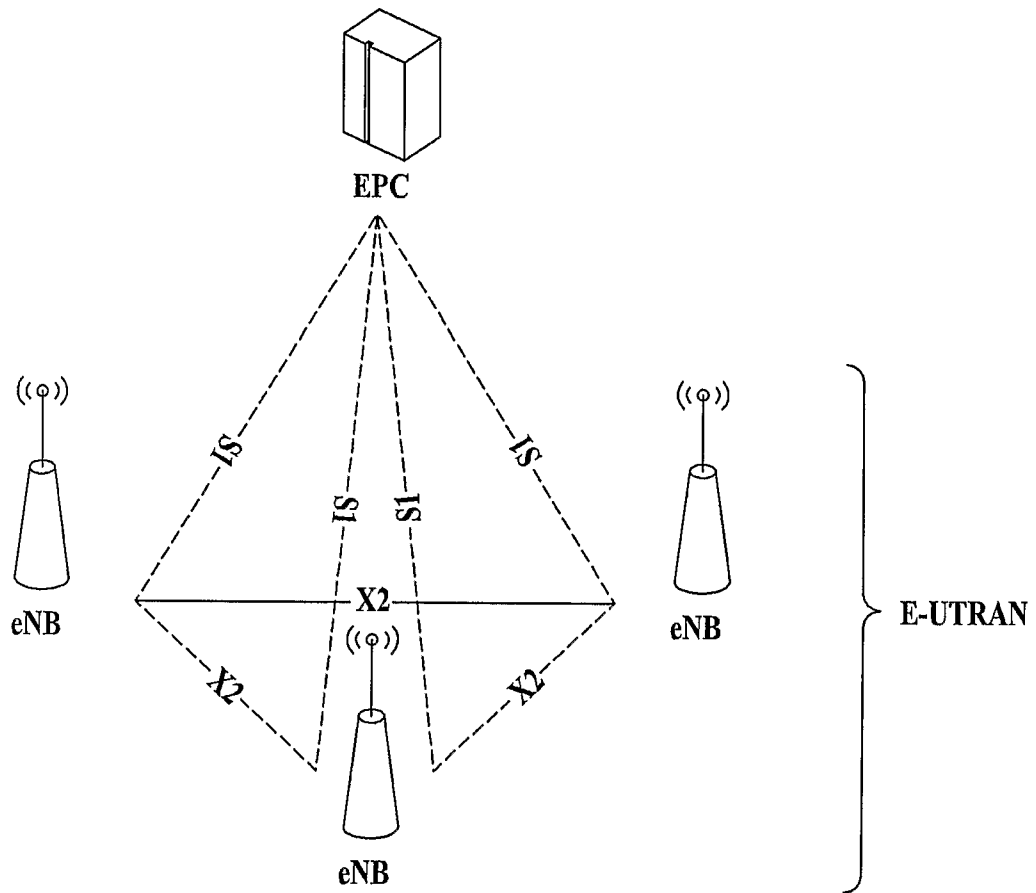


FIG. 2

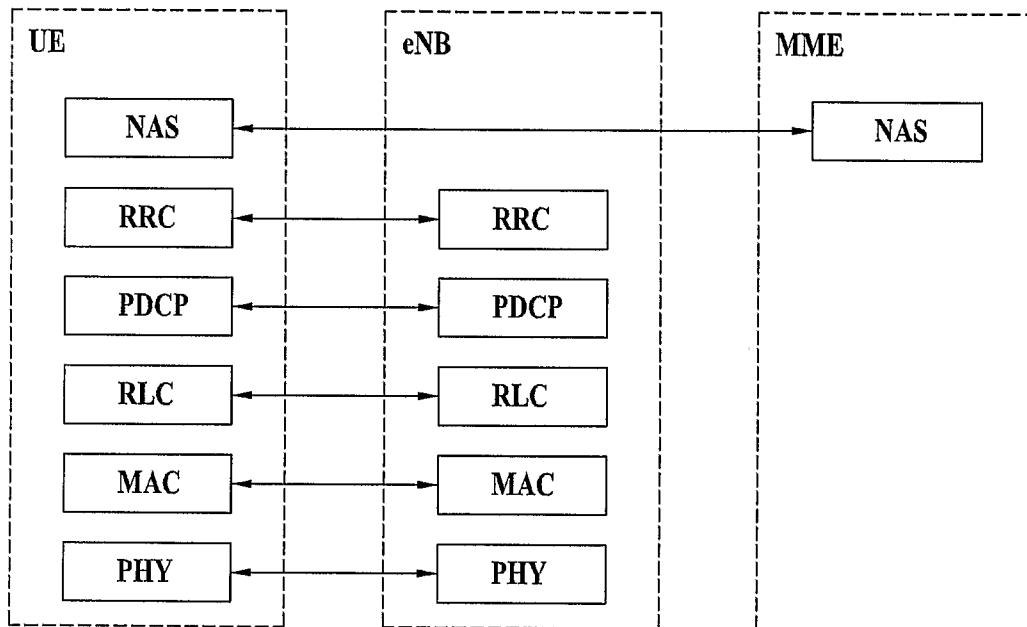


FIG. 3

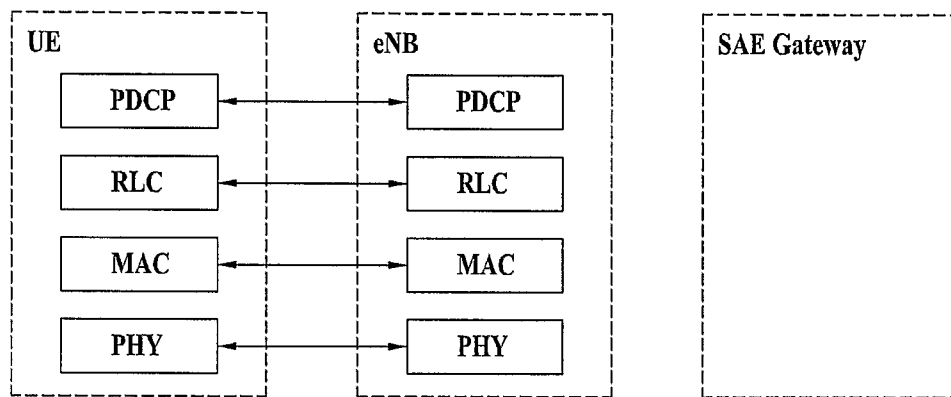


FIG. 4

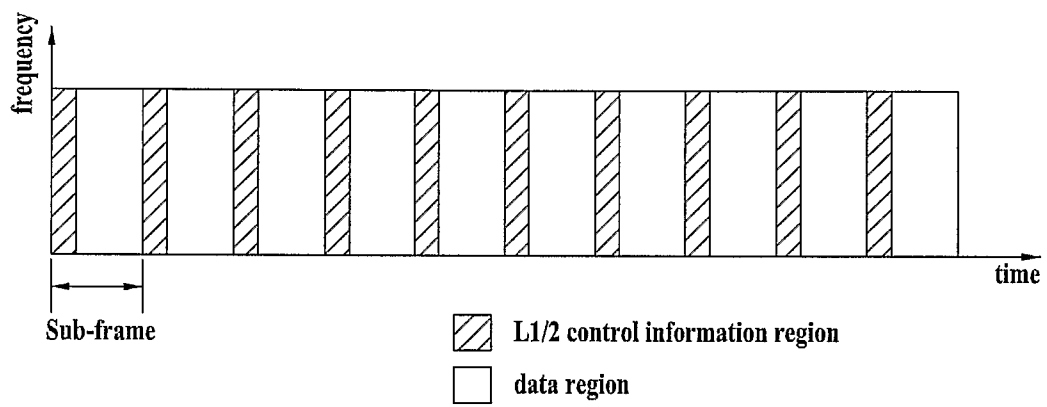


FIG. 5

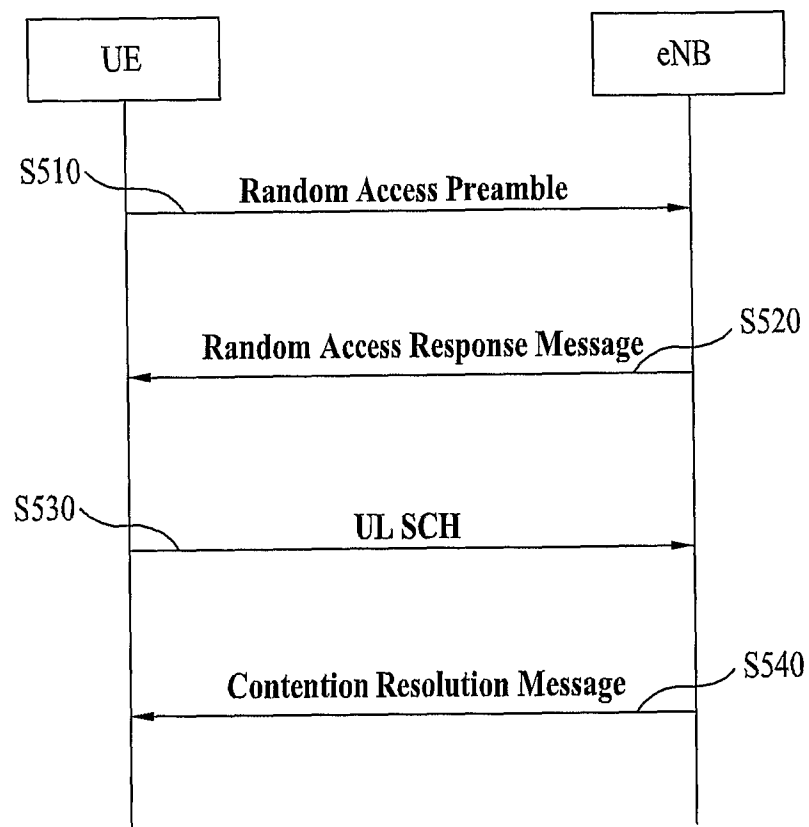


FIG. 6

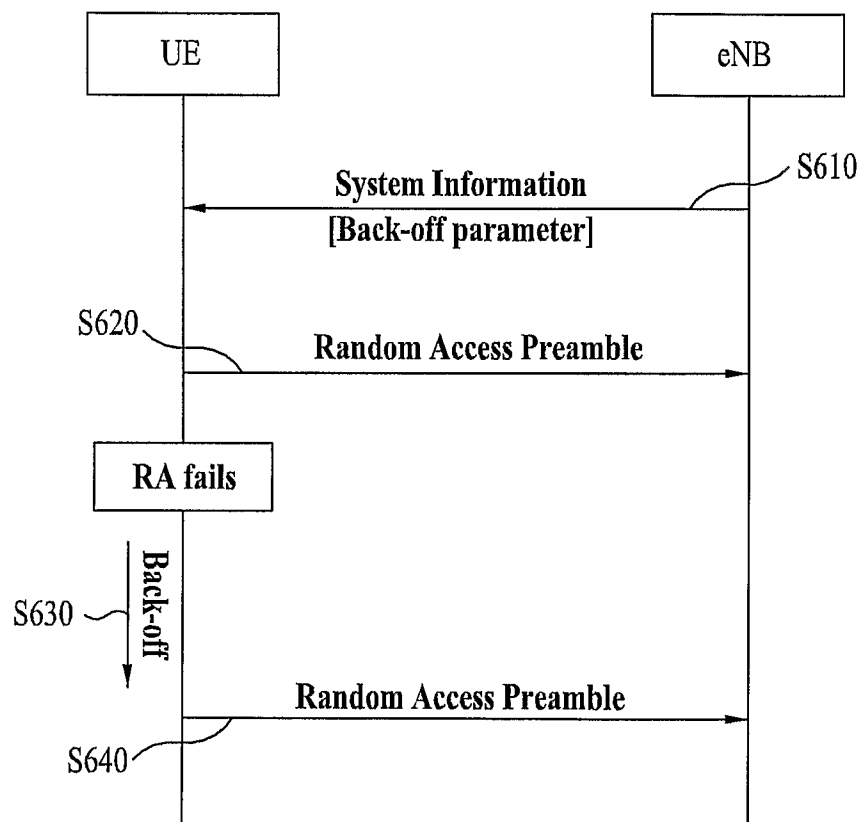


FIG. 7

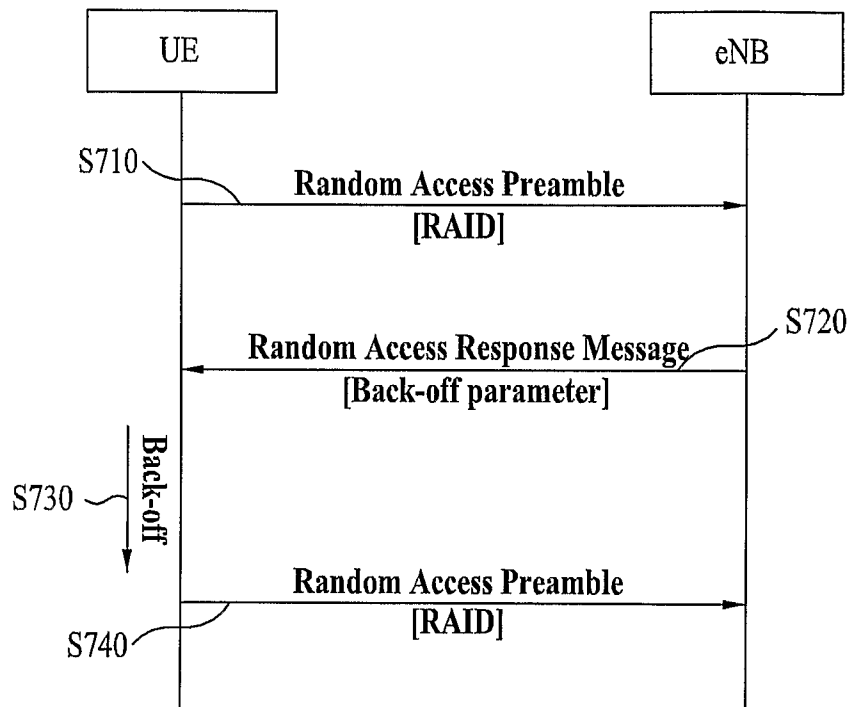


FIG. 8

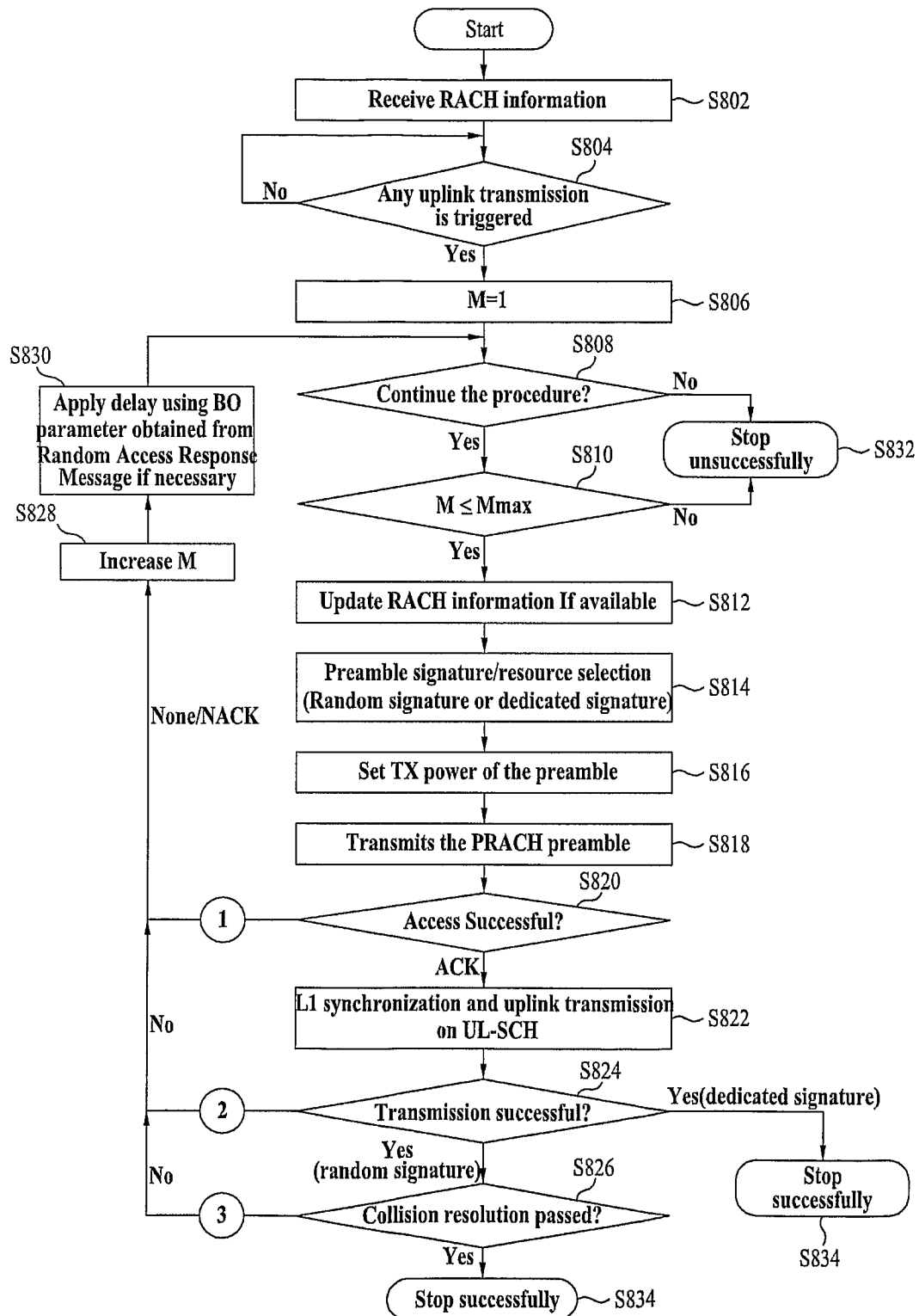


FIG. 9

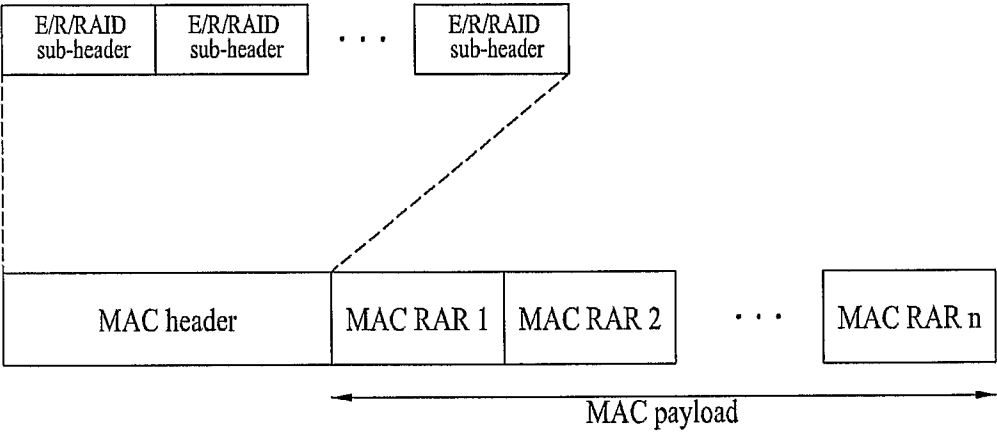


FIG. 10

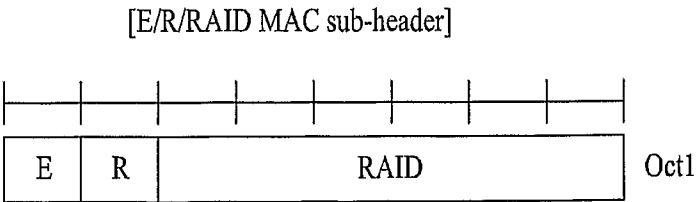


FIG. 12

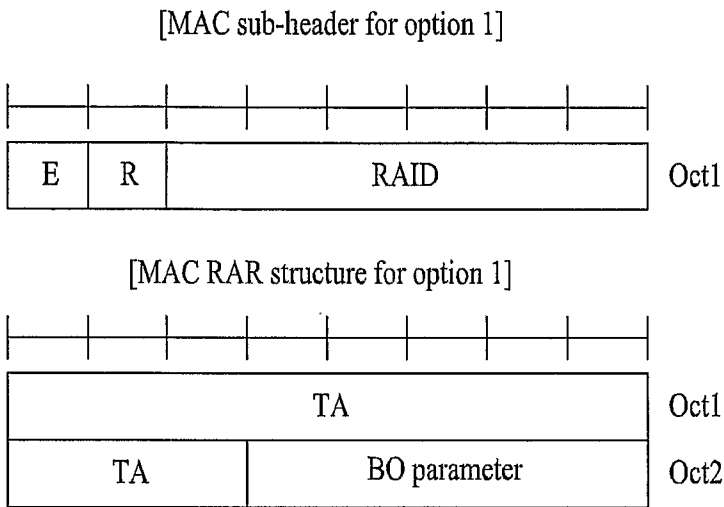


FIG. 13

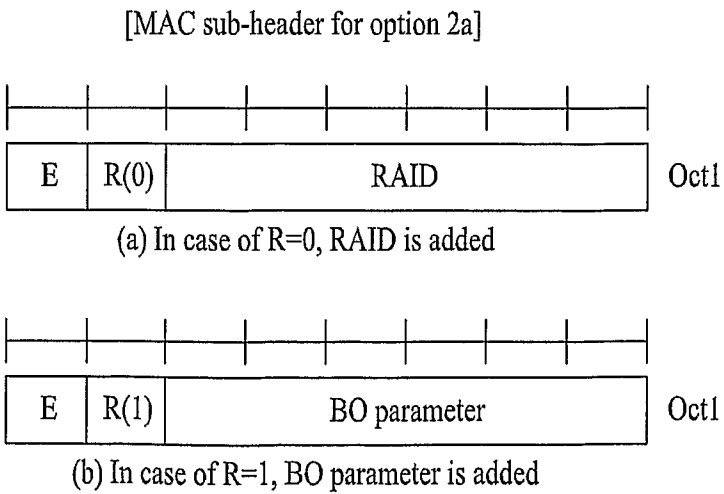


FIG. 14

[MAC RAR structure for option 2a]

TA								Oct1
TA			UL Grant					Oct2
UL Grant								Oct3
UL Grant								Oct4
T-CRNTI								Oct5
T-CRNTI								Oct6

(a) In case RAID is added in subheader

E	R	RAID #1						Oct1
E	R	RAID #2						Oct2
E	R	RAID #3						Oct3
E	R	RAID #4						Oct4
⋮								
E	R	RAID #N						Oct N

(b) In case BO parameter is added in subheader

FIG. 15

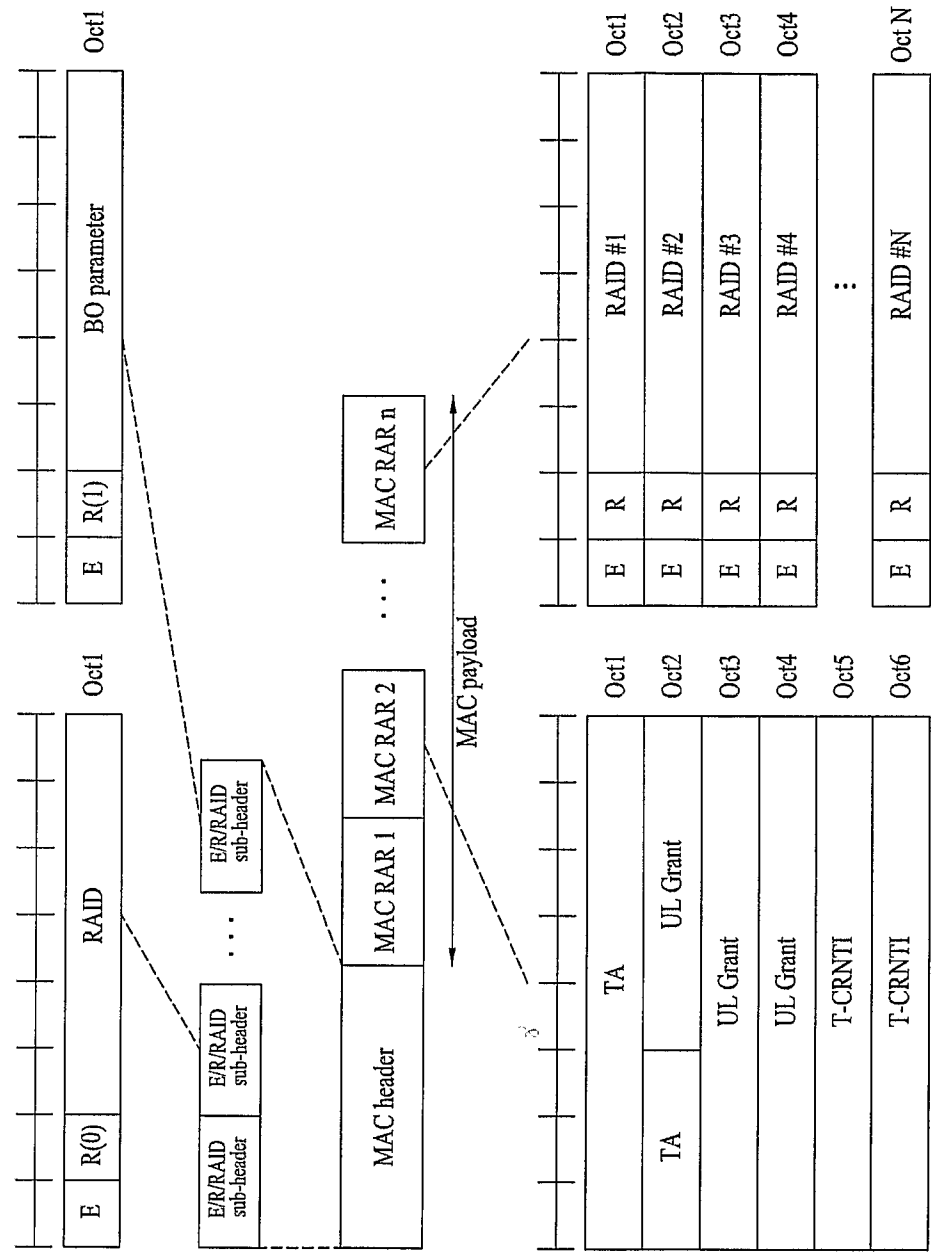


FIG.16

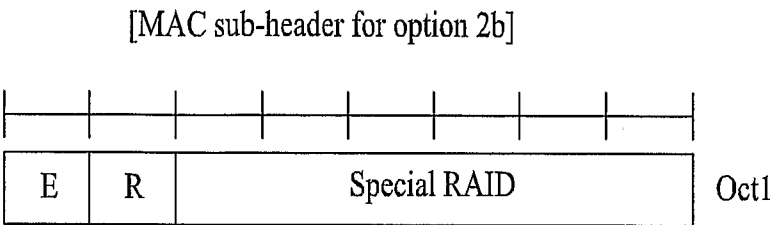


FIG. 17

