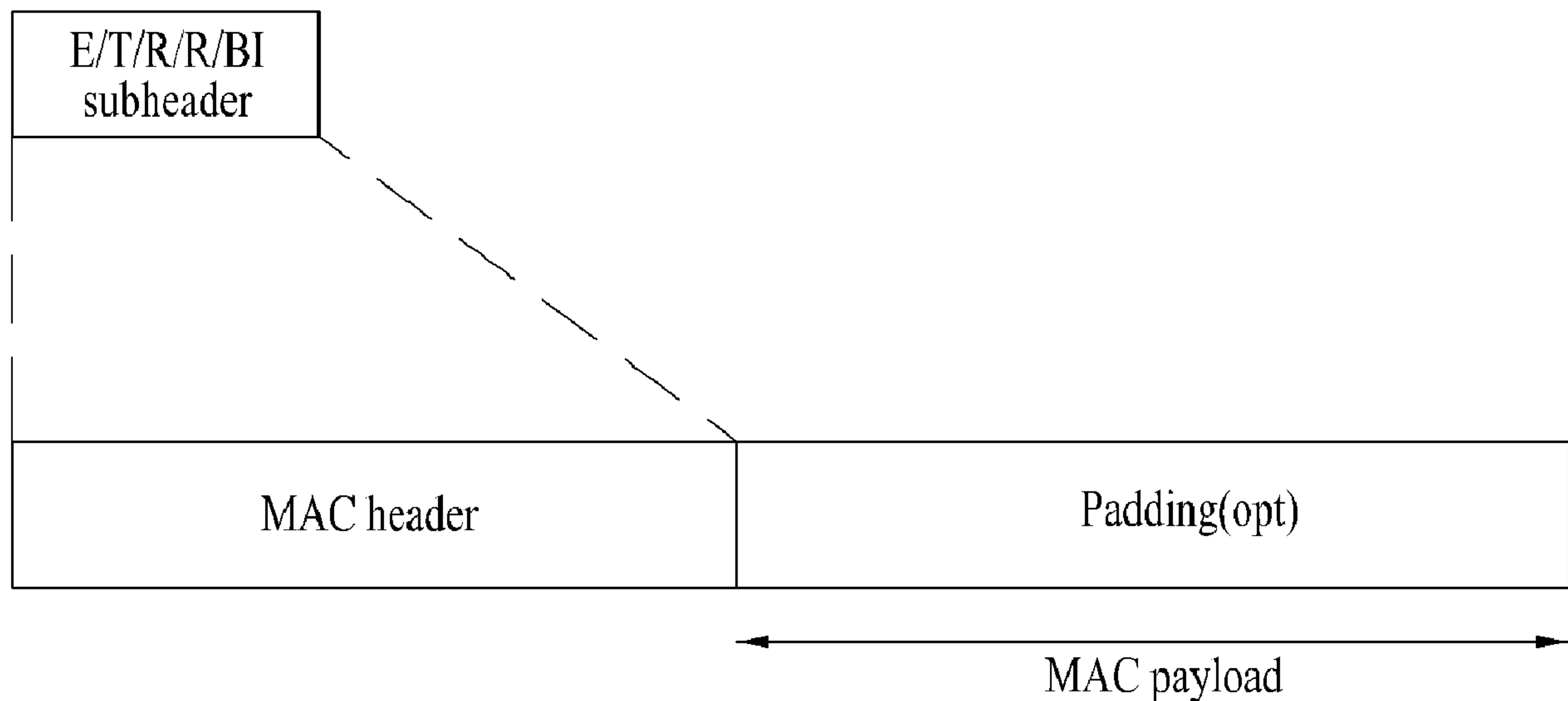




(86) **Date de dépôt PCT/PCT Filing Date:** 2009/12/23
(87) **Date publication PCT/PCT Publication Date:** 2010/07/08
(45) **Date de délivrance/Issue Date:** 2014/04/15
(85) **Entrée phase nationale/National Entry:** 2011/06/29
(86) **N° demande PCT/PCT Application No.:** KR 2009/007700
(87) **N° publication PCT/PCT Publication No.:** 2010/077004
(30) **Priorités/Priorities:** 2009/01/02 (US61/142,257);
2009/01/05 (US61/142,613); 2009/03/13 (US61/159,803);
2009/11/05 (KR10-2009-0106392)

(51) **Cl.Int./Int.Cl. H04W 74/08** (2009.01),
H04W 28/04 (2009.01)
(72) **Inventeurs/Inventors:**
PARK, SUNG JUN, KR;
YI, SEUNG JUNE, KR;
CHUN, SUNG DUCK, KR
(73) **Propriétaire/Owner:**
LG ELECTRONICS INC., KR
(74) **Agent:** SMART & BIGGAR

(54) **Titre : TECHNIQUE D'ACCES ALEATOIRE POUR EQUIPEMENT UTILISATEUR**
(54) **Title: RANDOM ACCESS SCHEME FOR USER EQUIPMENT**



(57) **Abrégé/Abstract:**

A method of efficiently processing a random access response message, when a terminal (or user equipment) performs random access, is disclosed. After the terminal transmits a random access preamble to a base station, the terminal may receive a random access response message having a format of medium access control protocol data unit (MAC PDU) including only a backoff indicator subheader in a MAC header of the MAC PDU, from the base station in response to the random access preamble. As described above, the terminal, which has received the random access response message including only a backoff indicator subheader in the MAC header of the MAC PDU, may consider a random access response reception procedure not successful and may perform a subsequent procedure for a random access responses reception failure.

(12) INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(19) World Intellectual Property Organization
International Bureau(43) International Publication Date
8 July 2010 (08.07.2010)(10) International Publication Number
WO 2010/077004 A3

(51) International Patent Classification:

A61K 9/127 (2006.01) *A61K 8/64* (2006.01)
A61K 8/14 (2006.01)(74) Agents: **KIM, Yong In** et al.; Kbk & Associates, 7th Floor, Hyundai Building, 175-9, Jamsil-dong, Songpa-ku, Seoul 138-861 (KR).

(21) International Application Number:

PCT/KR2009/007700

(81) Designated States (unless otherwise indicated, for every kind of national protection available):

AE, AG, AL, AM, AO, AT, AU, AZ, BA, BB, BG, BH, BR, BW, BY, BZ, CA, CH, CL, CN, CO, CR, CU, CZ, DE, DK, DM, DO, DZ, EC, EE, EG, ES, FI, GB, GD, GE, GH, GM, GT, HN, HR, HU, ID, IL, IN, IS, JP, KE, KG, KM, KN, KP, KZ, LA, LC, LK, LR, LS, LT, LU, LY, MA, MD, ME, MG, MK, MN, MW, MX, MY, MZ, NA, NG, NI, NO, NZ, OM, PE, PG, PH, PL, PT, RO, RS, RU, SC, SD, SE, SG, SK, SL, SM, ST, SV, SY, TJ, TM, TN, TR, TT, TZ, UA, UG, US, UZ, VC, VN, ZA, ZM, ZW.

(22) International Filing Date:

23 December 2009 (23.12.2009)

(25) Filing Language:

English

(26) Publication Language:

English

(30) Priority Data:

61/142,257	2 January 2009 (02.01.2009)	US
61/142,613	5 January 2009 (05.01.2009)	US
61/159,803	13 March 2009 (13.03.2009)	US
10-2009-0106392	5 November 2009 (05.11.2009)	KR

(84) Designated States (unless otherwise indicated, for every kind of regional protection available):

ARIPO (BW, GH, GM, KE, LS, MW, MZ, NA, SD, SL, SZ, TZ, UG, ZM, ZW), Eurasian (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European (AT, BE, BG, CH, CY, CZ, DE, DK, EE, ES, FI, FR, GB, GR, HR, HU, IE, IS, IT, LT, LU, LV, MC, MK, MT, NL, NO, PL, PT, RO, SE, SI, SK, SM, TR), OAPI (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, ML, MR, NE, SN, TD, TG).

(71) Applicant (for all designated States except US): **LG ELECTRONICS INC.** [KR/KR]; 20, Yeouido-dong, Yeongdeungpo-gu, Seoul 150-721 (KR).

(72) Inventors; and

(75) Inventors/Applicants (for US only): **PARK, Sung Jun** [KR/KR]; LG Institute, #533, Hogue 1(il)-dong, Dongan-gu, Anyang-si, Gyeonggi-do 431-080 (KR). **YI, Seung June** [KR/KR]; LG Institute, #533, Hogue 1(il)-dong, Dongan-gu, Anyang-si, Gyeonggi-do 431-080 (KR). **CHUN, Sung Duck** [KR/KR]; LG Institute, #533, Hogue 1(il)-dong, Dongan-gu, Anyang-si, Gyeonggi-do 431-080 (KR).

Published:

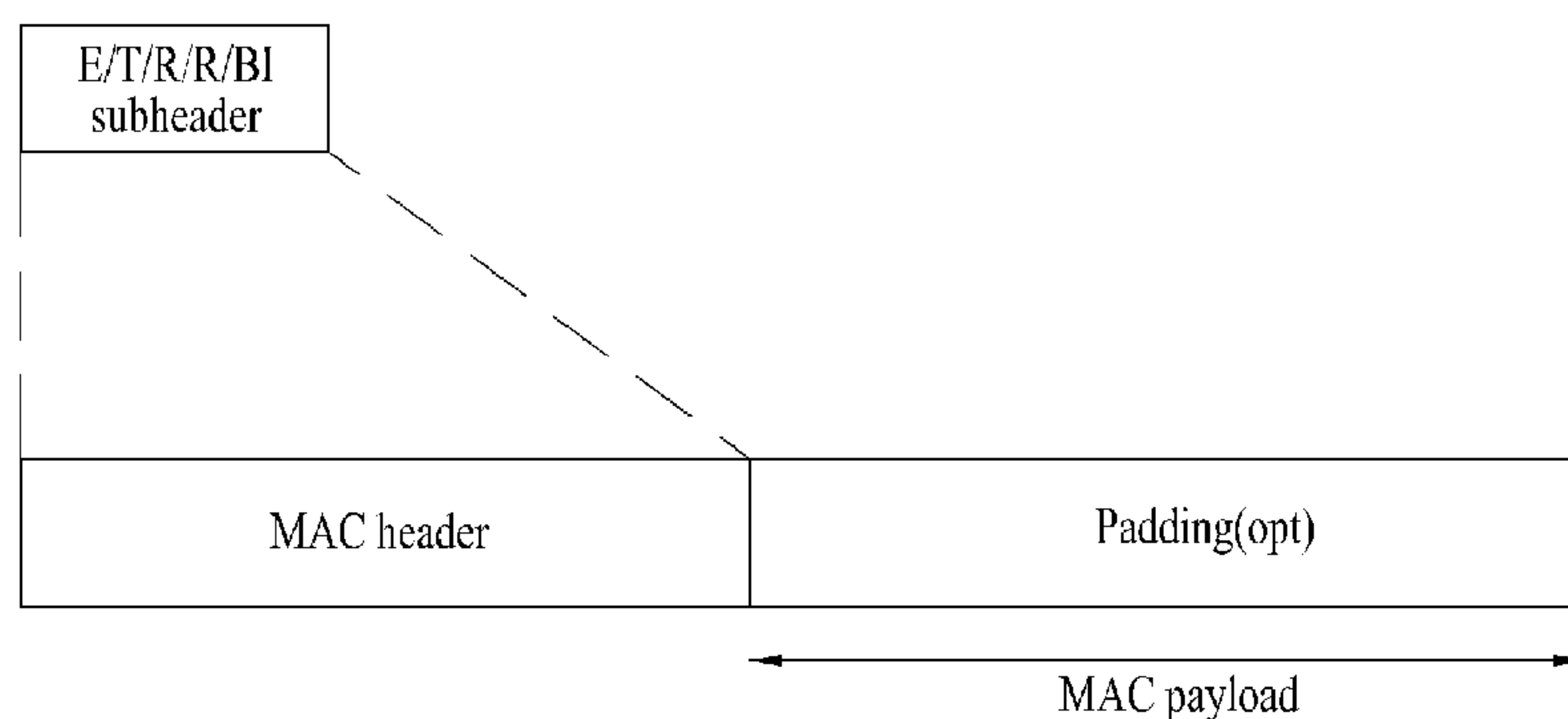
- with international search report (Art. 21(3))
- before the expiration of the time limit for amending the claims and to be republished in the event of receipt of amendments (Rule 48.2(h))

(88) Date of publication of the international search report:

21 October 2010

(54) Title: RANDOM ACCESS SCHEME FOR USER EQUIPMENT

[Fig. 10]



(57) **Abstract:** A method of efficiently processing a random access response message, when a terminal (or user equipment) performs random access, is disclosed. After the terminal transmits a random access preamble to a base station, the terminal may receive a random access response message having a format of medium access control protocol data unit (MAC PDU) including only a backoff indicator subheader in a MAC header of the MAC PDU, from the base station in response to the random access preamble. As described above, the terminal, which has received the random access response message including only a backoff indicator subheader in the MAC header of the MAC PDU, may consider a random access response reception procedure not successful and may perform a subsequent procedure for a random access responses reception failure.

WO 2010/077004 A3

Description

Title of Invention: RANDOM ACCESS SCHEME FOR USER EQUIPMENT

Technical Field

- [1] The present invention relates to a mobile telecommunications technology, and more particularly, to a random access scheme for a user equipment. Although the present invention is suitable for a wide scope of applications, it is particularly suitable for efficiently processing a random access response message when performing random access in a mobile telecommunications terminal.

Background Art

- [2] As an example of a mobile telecommunications system to which the present invention may be applied, a 3rd generation partnership project long term evolution (hereinafter referred to as "LTE" (3GPP LTE) telecommunications system will now be broadly described.
- [3] FIG. 1 illustrates a general view of an E-UMTS network structure as an example of a mobile telecommunications system. Herein, the evolved universal mobile telecommunications system (E-UMTS) corresponds to a system evolved from the conventional universal mobile telecommunications system (UMTS). The 3GPP is presently carrying out a basic standardization process for the E-UMTS. Generally, the E-UMTS may also be referred to as an LTE system.
- [4] The E-UMTS network may be broadly divided into an evolved-UMTS terrestrial radio access network (E-UTRAN) 101 and a core network (CN) 102. The E-UTRAN 101 consists of a user equipment (hereinafter referred to as a UE) 103, a base station (hereinafter referred to as an eNode B or an eNB) 104, and an access gateway (hereinafter referred to as an AG) 105, which is located at an end of the network. The AG 105 may be divided into a portion for processing user traffic and a portion for processing control traffic. At this point, a new interface may be used between a new AG for processing the user traffic and an AG for processing control traffic, thereby enabling the AGs to communicate to and from one another.
- [5] At least one or more cells may exist in a single eNode B. An interface for user traffic or control traffic may be used between each eNode B. CN 102 may be configured of a node used for user registration of AG 105 and other UEs 103. Additionally, an interface for differentiating E-UTRAN 101 from CN 102 may also be used.
- [6] Layers of a radio interface protocol between a user equipment (or terminal) and a network may be divided into an L1 (*i.e.*, a first layer), an L2 (*i.e.*, a second layer), and an L3 (*i.e.*, a third layer), based upon 3 lower layers of an open system interconnection

(OSI) reference model, which is generally and broadly known in a telecommunications system. Herein, a physical layer belonging to the first layer provides an information transfer service using a physical channel. Also, a radio resource control (hereinafter referred to as RRC) layer located in the third layer performs a function of controlling radio source between the terminal and the network. For this, the RRC layer enables the user equipment and the network to exchange RRC messages to and from one another. The RRC layer may be dispersed in network nodes, such as the eNode B 104 and the AG 105, or the RRS layer may be located only in either one of the eNode B 104 and the AG 105.

- [7] FIG. 2 and FIG. 3 respectively illustrate a structure of a radio interface protocol between a user equipment (or terminal), which is configured based upon a 3GPP radio access network standard, and a UTRAN. The radio interface protocol of FIG. 2 and FIG. 3 is horizontally configured of a physical layer, a data link layer, and a network layer, and the radio interface protocol of FIG. 2 and FIG. 3 is vertically divided into a user plane and a control plane. Herein, the user place is used for transmitting data information, and the control plane is used for delivering control signals (or for control signaling). More specifically, FIG. 2 illustrates each layer of the radio protocol control plane, and FIG. 3 illustrates each layer of the radio protocol user plane. As described above, the protocol layers of FIG. 2 and FIG. 3 may be divided into an L1 (*i.e.*, a first layer), an L2 (*i.e.*, a second layer), and an L3 (*i.e.*, a third layer), based upon 3 lower layers of an open system interconnection (OSI) reference model, which is generally and broadly known in a telecommunications system.
- [8] Hereinafter, each layer of the radio protocol control plane shown in FIG. 2 and the radio protocol user plane shown in FIG. 3 will now be described in detail.
- [9] A physical (PHY) layer, which corresponds to the first layer, uses a physical channel to provide an information transfer service to its higher layer (or upper layer). The PHY layer is connected to a medium access control (MAC) layer, which corresponds to the higher layer of the PHY layer, through a transport channel. And, data are transported (or transmitted) to and from the MAC layer and the PHY layer through the transport channel. At this point, depending upon the sharing of the channel, the transport channel may be broadly divided into a dedicated transport channel and a common transport channel. Furthermore, data are transported (or transmitted) to and from different PHY layers, *i.e.*, to and from the PHY layer of a transmitting system and the PHY layer of a receiving system, through a physical channel by using a radio source.
- [10] Multiple layers exist in the second layer. A medium access control (MAC) layer maps various logical channels to various transport channels. And, the MAC layer also performs logical channel multiplexing, wherein multiple logical channel are mapped to a single transport channel. The MAC layer is connected to its higher layer (or upper

74420-516

layer), a radio link control (RLC) layer³, through a logical channel. And, depending upon the type of information that are being transported, the logical channel may be broadly divided into a control channel, which transports information of a control plane, and a traffic channel, which transports information of a user plane.

- [11] The radio link control (RLC) layer of the second layer performs segmentation and concatenation on the data received from its higher layer, thereby adjusting the size of the data so that its lower layer can adequately transport the processed data to a radio section. Also, in order to ensure diverse quality of service (QOS) requested by each radio bearer (RB), the RLC layer provides three different operation modes, a transparent mode (TM), an un-acknowledged mode (UM), and an acknowledged mode (AM). Particularly, the AM RLC performs a re-transport function through an automatic repeat and request (ARQ) function in order to transport (or transmit) reliable data.
- [12] A packet data convergence protocol (PDCP) layer of the second layer performs a header compression function that reduces the size of an IP packet header, which has a relatively large data size and carries (or contains) unnecessary control information, in order to efficiently transport IP packets in a radio section having a small bandwidth, when transporting IP packets, such as IPv4 or IPv6. By allowing only the information absolutely necessary in the header portion of the corresponding data to be transported, the header compression function enhances the transport efficiency of the radio section. Furthermore, in the LTE system, the PDCP layer also performs a security function. Herein, the security function consists of ciphering and integrity protection. More specifically, ciphering prevents data monitoring (or data surveillance) by a third party, and integrity protection prevents data altering by a third party.
- [13] A radio resource control (RRC) layer of the third layer, which corresponds to the uppermost layer in the third layer, is defined only in the control plane. Being associated to the configuration, re-configuration, and release of radio bearers (RBs), the RRC layer controls logical channels, transport channels, and physical channels. Herein, the RB signifies a logical path provided by the first and second layers of a radio protocol, in order to deliver data between the user equipment and the UTRAN. Generally, the configuration of an RS refers to a process of regulating the characteristics of a radio protocol layer and channel, which are required for providing a specific service, and of respectively configuring each specific parameter and operating method. The RB is then divided into signaling RB (SRB) and data RB (DRB). Herein, the SRB is used as a path for transporting an RRC message from the control plane (C-plane), and the DRB is used as a path for transporting user data from the user plane (U-plane).
- [14] Downlink transport channels transporting (or transmitting) data from the network to

the user equipment include a broadcast channel (BCH) and a downlink shared channel (SCH). More specifically, the BCH transports system information, and the downlink SCH transports other user traffic or control messages. A downlink multicast or a traffic or control message may either be transported through the downlink SCH or may be transported through a separate downlink multicast channel (MCH). Meanwhile, uplink transport channels transporting data from the user equipment to the network include a random access channel (RACH) and an uplink shared channel (SCH). More specifically, the RACH transports initial control messages, and the uplink SCH transports other user traffic or control messages.

[15] Additionally, downlink physical channels transporting information, which are transported to the downlink transport channel, to the radio section between the network and the user equipment include a physical broadcast channel (PBCH), a physical multicast channel (PMCH), a physical downlink shared channel (PDSCH), and a physical downlink control channel (PDCCH) (or a downlink (DL) L1/L2 control channel). More specifically, the PBCH transports information of the BCH, and the PMCH transports information of the MCH. The PDSCH transports information of the PCH and the downlink SCH. And, the PDCCH transports control information provided from the first layer and the second layer, such as a downlink or uplink (DL/UL) scheduling grant. Meanwhile, uplink physical channels transporting information, which are transported to the uplink transport channel, to the radio section between the network and the user equipment include a physical uplink shared channel (PUSCH), a physical random access channel (PRACH), and a physical uplink control channel (PUCCH). More specifically, the PUSCH transports information of the uplink SCH, and the PRACH transports information of the RACH. Furthermore, the PUCCH transports control information provided from the first layer and the second layer, such as an HARQ ACK or NACK, a scheduling request (SR), and a channel quality indicator (CQI) report.

[16] Hereinafter, based upon the above description, the method for performing a random access from the user equipment to a base station (or an eNode B or eNB) will be described in detail. Firstly, the user equipment performs a random access process (or procedure) under the following circumstances:

- [17] - when the user equipment performs an initial access, due to an absence of an RRC connection between the user equipment and the eNode B
- [18] - when the user equipment performs a first access to a target cell, during a handover process
- [19] - when a random access process is requested by a command from the eNode B
- [20] - when data that are to be transported through an uplink are generated, in case time synchronization of the uplink does not match, or in case a designated radio source is

74420-516

5

- not allocated, the designated radio source being used for requesting a radio source
- [21] - when performing a recovery process, in case of a radio link failure or a handover failure
- [22] In the LTE system, during the procedure of selecting a random access preamble, a contention based random access procedure, wherein the user equipment randomly selects and uses a preamble from a specific group, and a non-contention based random access procedure, which uses a random access preamble allocated from the base station (or eNode B) only to a specific user equipment, are both provided. However, the non-contention based random access procedure may be used only during the handover procedure (or process) or only upon request from the base station (or eNode B).
- [23] Meanwhile, the process of the user equipment performing a random access with a specific base station (or eNode B) may broadly include the steps of (1) having the user equipment transport an random access preamble to the eNode B (or base station) (also referred as a "Message 1 transporting step", in case there is no confusion hereinafter), (2) receiving a random access response from the eNode B with respect to the transported random access preamble (also referred as a "Message 2 receiving step", in case there is no confusion hereinafter), (3) transporting an uplink message from the random access response message by using the received information (also referred as a "Message 3 transporting step", in case there is no confusion hereinafter), and (4) receiving a message corresponding to the uplink message from the eNode B (also referred as a "Message 4 receiving step", in case there is no confusion hereinafter).
- [24] In the above-described random access procedure, the user equipment stores data that are to be transported through message 3 in a message 3 buffer (or Msg3 buffer). Then, the user equipment transports (or transmits) the data stored in the message 3 buffer with respect to the reception of an uplink grant (or UL grant) signal. The UL grant signal corresponds to a signal notifying information on an uplink radio source, which may be used when the user equipment transports a signal to the base station (or eNode B). Herein, in case of the above-described LTE system, the UL grant signal is received through a random access response (RAR) message, which is received through the physical downlink control channel (PDCCH) or the physical downlink shared channel (PDSCH). Hereinafter, the method for receiving a random access response message of the user equipment will be described in more detail.
- [25] FIG. 4 illustrates a method for receiving and processing a random access response message according to a current LTE standard. After transporting the random access preamble, the user equipment attempts to receive its own random access response from within a random access response reception window, which is designated by the base station (or eNode B) through a system information or handover command. More specifically, the random access response information may be transported in a MAC

74420-516

6

packet data unit (MAC PDU) format. And, the MAC PDU for transporting random access response information includes a MAC payload and a MAC subheader respective of the MAC payload. Herein, the MAC payload corresponds to random access response message information for at least one or more user equipments. The MAC PDU may further include a MAC subheader including a backoff indicator, which may be used when the user equipment reattempts random access. The MAC PDU may be transported through the physical downlink shared channel (PDSCH). Accordingly, in step 601, it is determined whether or not a received random access response message exists within the predetermined random access response reception window. If it is determined that a received random access response message does not exist within the predetermined random access response reception window, it is concluded (or determined) that the reception of the random access response message has failed. Subsequently, the procedure moves on to step 604, so that the operations according to the failure of receiving the random access response message can be performed.

- [26] Alternatively, if it is determined that a received random access response message exists within the predetermined random access response reception window, the system determines, in step 602, whether each of the random access response messages received within the random access response reception window includes a random access identifier (*e.g.*, RA-RNTI), which does not correspond to (or match) the random access preamble already transported from the user equipment. If it is determined that all of the random access response messages received within the random access response reception window include a random access identifier, which does not correspond to (or match) the random access preamble already transported from the user equipment, the user equipment concludes that the reception of the respective random access response message has failed. Thereafter, the procedure moves on to step 604, so that the operations according to the failure of receiving the random access response message can be performed. On the other hand, if it is determined that at least one or more random access response messages received within the random access response reception window include a random access identifier, which corresponds to the random access preamble already transported from the user equipment, the procedure moves on to step 603, so that the corresponding random access response message can be processed.

Summary

- [27] Accordingly, an embodiment of the present invention is directed to a random access scheme for a user equipment that substantially obviates one or more problems due to limitations and disadvantages of the related art.

74420-516

7

- [28] A random access response message reception algorithm of the above-described user equipment has been designed under the assumption that the MAC PDU for transporting random access response information essentially consists of a MAC payload, which functions as a random access response message for at least one or more user equipments, and a MAC subheader respective to the MAC payload. More specifically, in the current LTE standard, it is assumed that the MAC PDU includes a MAC subheader including a backoff indicator, which can be used when the user equipment reattempts random access, as an optional element, and that the MAC PDU includes the MAC payload, functioning as a random access response message for at least one or more user equipments, and the MAC subheader respective to the MAC payload, as its essential elements.
- [29] In a system, wherein a load increases within a cell that provides a service, or in a system which uses a low bandwidth, the following description proposes a base station that allows, when required, the transport (or transmission) of a MAC PDU including only the backoff indicator. An embodiment of the present invention may provide a random access scheme for a user equipment that can provide a random access response message reception algorithm for a user equipment improved according to the proposed system.
- [30] Additional advantages, objects and features will be set forth in part in the description which follows and in part will become apparent to those having ordinary skill in the art upon examination of the following or may be learned from practice of the invention. The objectives and other advantages may be realized and attained by the structure particularly pointed out in the written description and claims hereof as well as the appended drawings.
- [31] To achieve these objects and other advantages and in accordance with embodiments of the invention, as broadly described herein, in a method of performing a random access to a base station by a terminal, the method includes the steps of transmitting a random access preamble to the base station, receiving, from the base station in response to the random access preamble, a random access response message having a format of medium access control protocol data unit (MAC PDU) including only a backoff indicator subheader in a header portion of the MAC PDU, and no medium access control random access response in the MAC PDU, and considering a random access response reception procedure not successful, and performing a subsequent procedure for a random access responses reception failure.
- [32] At this point, the terminal may be adapted to consider the random access response reception procedure not successful, if no random access response is received within a random access response window having a predetermined time length, or if none of all

74420-516

8

received random access responses contains a random access preamble identifier corresponding to the transmitted random access preamble.

- [33] More specifically, a case in which none of all received random access responses contains the random access preamble identifier corresponding to the transmitted random access preamble may include a first case in which all received random access responses contain random access preamble identifiers that do not match the transmitted random access preamble, and a second case in which the random access response message having the format of MAC PDU including only the backoff indicator subheader in the header portion of the MAC PDU and no medium access control random access response in the MAC PDU is received.
- [34] In the embodiment of the present invention, the subsequent procedure for the random access response reception failure may be performed by incrementing a random access preamble transmission counter by 1, and indicating to an higher layer higher than a medium access control (MAC) layer that there is a random access problem, if the random access preamble transmission counter becomes a predetermined maximum number of preamble transmission.
- [35] Additionally, the subsequent procedure for the random access response reception failure may also be performed by incrementing a random access preamble transmission counter by 1, and delaying a subsequent random access preamble transmission by a backoff time selected using a backoff indicator in the backoff indicator subheader, if the random access preamble is selected by a medium access control (MAC) layer.
- [36] Herein, it may be assumed that the random access response message having the format of MAC PDU is received through a physical downlink shared channel (PDSCH).
- [37] In another aspect of the present invention, in a terminal performing a random access to a base station, the terminal includes a transmitting (Tx) module transmitting a random access preamble to the base station, a receiving (Rx) module receiving, from the base station in response to the random access preamble, a random access response message having a format of medium access control protocol data unit (MAC PDU), and, a medium access control (MAC) layer module adapted to consider a random access response reception procedure not successful and to perform a subsequent procedure for a random access response reception failure, when the random access response message received by the receiving module corresponds to a MAC PDU including only a backoff indicator subheader in a header portion of the MAC PDU and no medium access control random access response in the MAC PDU.
- [38] At this point, the MAC layer module may be adapted to consider the random access response reception procedure not successful, if the receiving module does not receive any random access response within a random access response window having a predetermined time length, or if none of all received random access responses containing a random access preamble identifier corresponding to the random access preamble

74420-516

9

- transmitted from the transmitting module is received by the receiving module.
- [39] Additionally, a case in which none of all received random access responses containing the random access preamble identifier corresponding to the random access preamble transmitted from the transmitting module is received by the receiving module, may include a first case in which all random access responses received by the receiving module contain random access preamble identifiers that do not match the random access preamble transmitted from the transmitting module, and a second case in which the random access response message having the format of MAC PDU including only the backoff indicator subheader in the header portion of the MAC PDU and no medium access control random access response in the MAC PDU is received by the receiving module.
- [40] In the embodiment of the present invention, as the subsequent procedure for the random access response reception failure, the MAC layer module may increment a random access preamble transmission counter by 1, and may indicate to a higher layer module corresponding to a higher layer higher than a medium access control (MAC) layer that there is a random access problem, if the random access preamble transmission counter becomes a predetermined maximum number of preamble transmission. And, the MAC layer module may also delay a subsequent random access preamble transmission by a backoff time selected using a backoff indicator in the backoff indicator subheader, if the random access preamble is selected by a medium access control (MAC) layer.
- [41] In a system, wherein a load increases within a cell that provides a service, or in a system which uses a low bandwidth, the following description of the present invention proposes a base station that allows, when required, the transport (or transmission) of a MAC PDU including only the backoff indicator. An advantage of some embodiments of the present invention is to provide a random access scheme for a user equipment that can provide a random access response message reception algorithm for a user equipment improved according to the proposed system.
- [42] It is to be understood that both the foregoing general description and the following detailed description are exemplary and explanatory and are intended to provide further explanation of the invention as claimed.
- Brief Description of Drawings**
- [43] The accompanying drawings, which are included to provide a further understanding of the invention and are incorporated in and constitute a part of this application, illustrate embodiment(s) of the invention and together with the description serve to explain the principle of the invention. In the drawings:
- [44] FIG. 1 illustrates a general view of an E-UMTS network structure as an example of a

mobile telecommunications system;

- [45] FIG. 2 and FIG. 3 respectively illustrate a structure of a radio interface protocol between a user equipment (or terminal), which is configured based upon a 3GPP radio access network standard, and a UTRAN;
- [46] FIG. 4 illustrates a method for receiving and processing a random access response message according to a current LTE standard;
- [47] FIG. 5 illustrates operation process of the user equipment (or terminal) and the base station in a contention based random access procedure;
- [48] FIG. 6 and FIG. 9 illustrate a MAC PDU format for random access response message transmission and structures of MAC subheaders and a random access response configuring the MAC PDU;
- [49] FIG. 10 illustrates a MAC PDU format including only a backoff indicator according to an embodiment of the present invention;
- [50] FIG. 11 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to an embodiment of the present invention;
- [51] FIG. 12 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to another embodiment of the present invention;
- [52] FIG. 13 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to yet another embodiment of the present invention;
- [53] FIG. 14 illustrates a method of performing random access in a user equipment according to an embodiment of the present invention; and
- [54] FIG. 15 illustrates a structure of a user equipment according to an embodiment of the present invention.

Mode for the Invention

- [55] Reference will now be made in detail to the preferred embodiments of the present invention, examples of which are illustrated in the accompanying drawings. It should be understood that the present invention is not limited solely to the following embodiment. The following description includes specific details for providing a full understanding of the present invention. However, it is apparent to anyone skilled in the art that the present invention may also be embodied without such specific details. For example, although it is assumed in the following detailed description of the present invention that the mobile telecommunications system according to the present invention corresponds to a 3GPP LTE system, the present invention may also be applied to other random mobile telecommunications systems, excluding the typical

characteristics of the 3GPP LTE system.

[56] In some cases, to avoid any ambiguity in the concept of the present invention, structures or devices of the disclosure may be omitted, or the embodiment of the present invention may be illustrated in the form of block views focusing on the essential functions of each structure and device. Also, wherever possible, the same reference numbers will be used throughout the drawings to refer to the same or like parts.

[57] Furthermore, in the following description of the present invention, it is assumed that a terminal collectively refers to a mobile or fixed-type user-end device, such as a user equipment (UE), a mobile station (MS), and so on. Additionally, it is also assumed that a base station collectively refers to a random node of a network-end communicating with the terminal, such as Node B, an eNode B, a base station, and so on.

[58] As described above, the following description of the present invention proposes a base station that allows, when required, the transport (or transmission) of a MAC PDU including only the backoff indicator. Hereinafter, in order to understand the necessity of the above-described MAC PDU including only the backoff indicator, the random access procedure and the format of message 2 will be described in detail.

[59] FIG. 5 illustrates operation process of the user equipment (or terminal) and the base station in a contention based random access procedure.

[60] (1) Transmitting Message 1

[61] The user equipment randomly selects a random access preamble from a group of random access preambles designated by a system information or handover command and, then, the user equipment may select and transport a physical RACH (PRACH) source that can transport the selected random access preamble (S501).

[62] (2) Receiving Message 2

[63] After the user equipment transmits a random access preamble, as shown in S501, the base station attempts to receive its random access response from within a random access response reception window designated by a system information or handover command (S502). More specifically, random access response information may be transmitted in a MAC PDU format, and the MAC PDU may be delivered (or transported) through a physical downlink shared channel (PDSCH). Also, in order to allow the user equipment to adequately receive the information being delivered to the PDSCH, it is preferable that the user equipment monitors a physical downlink control channel (PDCCH). In other words, it is preferable that the PDCCH includes information on the user equipment that should receive the PDSCH, frequency and time information of a radio source of the PDSCH, and a transmission (or transport) format of the PDSCH. Once the user equipment succeeds in receiving the PDCCH that is transported (or transmitted) thereto, the user equipment may adequately receive

random access responses transmitted to the PDSCH depending upon the information of the PDCCH. Furthermore, the random access response may include a random access preamble identifier (RAPID), an uplink (UL) grant notifying the uplink radio source, a temporary cell identifier (C-RNTI), and a timing advance command (TAC).

- [64] As described above, a random access preamble identifier (RAPID) is required in the random access response because random access response information for at least one or more user equipments may be included in a single random access response, and also because information specifying (or indicating) the user equipment, in which the UL grant, the temporary C-RNTI, and the TAC are valid, is required. In this step, it is assumed that the user equipment selects a random access preamble identifier (RAPID) matching the random access preamble selected by the user equipment itself in step 502. Through the selected RAPID, the user equipment may receive the UL grant, the temporary cell identifier (C-RNTI), and the timing advance command (TAC).

- [65] (3) Transmitting Message 3

- [66] When the user equipment receives a valid random access response, each set of information included in the random access response is processed. More specifically, the user equipment adopts the TAC and stores the temporary cell identifier (C-RNTI). Also, data that are to be transmitted with respect to the reception of the valid random access response may be stored in the message 3 buffer (or Msg3 buffer).

- [67] Meanwhile, the user equipment uses the received UL grant, so as to transmit the data (*i.e.*, message 3) to the base station (S503). Message 3 shall include an identifier of the user equipment. In the contention based random access procedure, the base station is incapable of knowing which user equipments perform the random access process. Accordingly, the user equipments should be identified in order to resolve collision is a later process.

- [68] Two methods for including identifiers for each user equipment are proposed herein. In the first method, if the user equipment is already provided with a valid cell identifier, which had been allocated from the corresponding cell prior to the random access process, the user equipment transmits its cell identifier through a UL transport (or transmission) signal corresponding to the UL grant. Conversely, if the user equipment has not been assigned with a valid cell identifier prior to the random access process, the user equipment transmits data including its unique identifier (*e.g.*, S-TMSI or random ID). Generally, the unique identifier is longer than the cell identifier. If the user equipment has transmitted the data corresponding to the UL grant, a contention resolution timer (hereinafter referred to as a "CR timer" for resolving contention (or collision) is disclosed.

- [69] (4) Receiving Message 4

- [70] After the user equipment has transmitted the data containing its identifier through the

UL grant included in the random access response, the user equipment waits for instructions from the base station for resolving collision (or contention). More specifically, in order to a specific message, the user equipment attempts the reception of the PDCCH (S504). Two methods for receiving the PDCCH have been proposed. As described above, if the identifier of message 3, which is transmitted with respect to the UL grant, is transmitted by using a cell identifier, the user equipment uses its own cell identifier to attempt PDCCH reception. On the other hand, if the identifier corresponds to a unique identifier, the user equipment may attempt PDCCH reception by using a temporary cell identifier included in the random access response. Thereafter, in the former case, if the user equipment receives the PDCCH through its own cell identifier prior to the expiration of the CR timer, the user equipment determines that the random access process has been successfully performed, thereby ending the random access process. In the latter case, if the user equipment receives the PDCCH through a temporary cell identifier prior to the expiration of the CR timer, the user equipment verifies the data transmitted through the PDSCH, which is designated by the PDCCH. Then, if the data content contains the unique identifier, the user equipment determines that the random access process has been successfully performed, thereby ending the random access process.

- [71] Hereinafter, data configuration format of the random access response provided from the base station to the user equipment, in the above-described contention based random access procedure, will be described in detail.
- [72] FIG. 6 and FIG. 9 illustrate a MAC PDU format for random access response message transmission and structures of MAC subheaders and a random access response configuring the MAC PDU.
- [73] As shown in FIG. 6, the MAC PDU for transmitting the random access response message may include a MAC header and a MAC payload and may also include a padding unit as an optional element. The MAC payload may include N number of random access response information (hereinafter referred to as "MAC RAR"), wherein N is an integer. The MAC header may include MAC subheaders respective to each of the N number of MAC RARs. Herein, the MAC subheaders corresponding to each MAC RAR may be respectively configured of E/T/RAPID fields, as shown in FIG. 7. Furthermore, the MAC header may further include a MAC subheader containing a backoff indicator for notifying the backoff indicator used for determining a time delay when the user equipment reattempts random access. Herein, the MAC subheader containing the backoff indicator (*i.e.*, backoff subheader) may be configured of E/T/R/R/BI fields, as shown in FIG. 8.
- [74] The E field (or extension field) of the MAC subheader corresponding to each MAC RAR and of the MAC subheader containing the backoff indicator, as shown in FIG. 7

and FIG. 8, indicates whether or not MAC subheaders respective to additional MAC RARs exist. For example, when the E field is set to '1', this indicates that MAC subheaders of the E/T/RAPID format exist after the corresponding MAC header. When the E field is set to '0', the user equipment determines that there exist no more MAC subheaders subsequent to the corresponding MAC subheader. And, occasionally, the user equipment may determine that a MAC payload begins at the ending point of the corresponding MAC subheader, thereby being capable of performing a MAC payload reading process.

[75] Meanwhile, the T field (or type field) shown in FIG. 7 and FIG. 8, may indicate whether the corresponding MAC subheader corresponds to MAC subheader respective to the MAC RAR, or whether the corresponding MAC subheader corresponds to the MAC subheader containing the backoff indicator. If the T field is set to '1', this indicates that a random access preamble identifier exists in the corresponding MAC subheader. And, if the T field is set to '0', this indicates that the corresponding MAC subheader corresponds to the MAC subheader containing the backoff indicator.

[76] The RAPID field of the MAC subheader respective to the MAC RAR, shown in FIG. 7, indicates a random access preamble identifier indicating which random access preamble the respective MAC RAR corresponds to. Also, the BI field of the MAC subheader containing the backoff indicator, shown in FIG. 8, indicates the backoff indicator. In FIG. 8, the RS field corresponds to a reserved field.

[77] Meanwhile, FIG. 9 illustrates a structure of the MAC RAR. As shown in FIG. 9, the MAC RAR consists of 4 fields: an R field indicating a reserved bit, a "Timing advance command field" for determining (or setting up) uplink signal synchronization, a UL grant field, and a temporary cell identifier (or temporary C-RNTI) field.

[78] As shown in FIG. 6 to FIG. 9, and more particularly, as shown in FIG. 6, in the current LTE standard, it is defined that the MAC PDU for transmitting a random access response message consists of MAC RARs and MAC subheader respective to each MAC RAR, as essential elements, and also of a MAC subheader containing a backoff indicator, as an optional element. However, in case the random access response containing only the backoff indicator cannot be supported, the following problems may occur.

[79] It is assumed that, in a system of a low bandwidth (*e.g.*, 1.25 MHz or 5 MHz), a plurality of user equipments each transmitting a random access preamble for random access at a specific moment (or time point) to the base station exists. At this point, due to reasons such as lack of sources that are to be assigned (or allocated), the base station may not be able to provide UL grant via random access response to all of the user equipments transmitting a random access preamble. In this case, apart from the user equipments being provided with the UL grant, it is preferable for the base station to use

a backoff indicator on the remaining user equipments that are not provided with the UL grant, so as to delay transmission by a random time value (or a random period of time), thereby instructing retransmission of the random access preamble. This is because, if the retransmission for random access is not adequately dispersed in a random access overload situation, the retransmission of the random access preamble and a new access may be accumulated and generated. Thus, the above-described overload situation may be aggravated.

[80] However, in addition to the above-described scenario, when random access of multiple user equipments occurs, and when the base station is incapable of providing a UL grant to any one of the multiple user equipments, in a system regulating the MAC RAR and the respective MAC subheader as the essential elements configuring the MAC PDU, which is used to transmit random access response messages, backoff indicators cannot be transmitted to the user equipments. Accordingly, in a random access overload situation, backoff indicators cannot be transmitted to the respective user equipments, not to mention that all retransmission of random accesses for the corresponding user equipments will be eventually concentrated in the next (or subsequent) opportunity for random access preamble transmission. Furthermore, new access will also be attempted during the same opportunity for random access preamble transmission. Thus, the above-described overload situation may be even more aggravated.

[81] Therefore, an embodiment of the present invention proposes a MAC PDU format, which includes a MAC header containing only MAC subheaders including backoff indicators that can be used according to the embodiment of the present invention. FIG. 10 illustrates a MAC PDU format including only a backoff indicator according to an embodiment of the present invention.

[82] When comparing the MAC PDU format shown in FIG. 10 with the MAC PDU shown in FIG. 6, FIG. 10 illustrates a MAC PDU format including only backoff indicator MAC subheaders in the MAC header portion, wherein each MAC subheader consists of E/T/R/R/BI fields. As described above, by supporting the MAC PDU including only the backoff indicator MAC subheaders, in case the base station is incapable of performing source assignment (or allocation) to any user equipment via RAR, the base station may deliver (or transport) backoff indicators to each user equipment, thereby adequately distributing retransmission timing of the random access preamble.

[83] Evidently, in this embodiment of the present invention, the MAC RAR and the respective MAC subheader may be included in the MAC header portion, as shown in FIG. 6. More specifically, it is proposed that the MAC PDU for transmitting random access response message according to the embodiment of the present invention

includes one MAC header and zero, one, or multiple MAC RARs. The MAC PDU that does not include any MAC RAR indicates that the corresponding MAC PDU does not include any MAC subheader respective to the absent MAC RAR. Nevertheless, the MAC PDU according to the embodiment of the present invention may further and optionally include a padding unit.

[84] By using the MAC PDU shown in FIG. 10, even when the base station is incapable of transmitting UL grants to the respective user equipments through the MAC RAR, the base station may transmit a backoff indicator so as to delay random access of the corresponding user equipment(s) by a specific period of time. Furthermore, when receiving a MAC subheader having an E field of '0' and a T field of '0' the user equipment does not attempt the reception of the MAC RAR. Instead, the user equipment may interpret the BI field of the corresponding MAC subheader, so as to apply the interpreted content to the retransmission of a subsequent random access preamble.

[85] Meanwhile, in case a MAC PDU including only a backoff indicator MAC subheader is being supported according to an embodiment of the present invention, it will be examined through the following description whether or not one or more problems exist in the random access response reception procedure described with reference to FIG. 4.

[86] In case the user equipment receives a MAC PDU including only a backoff indicator MAC subheader from the base station, since a received random access response message exists within the random access response reception window in step 601 of FIG. 4, the user equipment moves on to step 602. Then, in step 602 of FIG. 4, not all of the RARs received by the user equipment include a random access preamble identifier, which are non-respective to (or do not match) the random access preambles transmitted by each RAR itself. Therefore, the user equipment moves on to step 603. More specifically, in case the user equipment receives only the backoff indicator MAC subheader from the base station, since no random access preamble indicator is included, the user equipment does not determine that the reception of the random access response has failed.

[87] However, in case the user equipment receives only the backoff indicator MAC subheader from the base station, since the MAC PDU does not include any MAC RAR transmitted from the corresponding user equipment, it is impossible to process the respective MAC RAR.

[88] Therefore, in a preferred embodiment of the present invention, the support of a MAC PDU format random access response message, wherein the MAC PDU format includes only the backoff indicator MAC subheader, as shown in FIG. 10, is additionally considered. Thus, in case the user equipment receives the MAC PDU format random access response message, wherein the MAC PDU format includes only the backoff

indicator MAC subheader, from the base station, the user equipment determines that the reception of the corresponding random access response message has failed.

Thereafter, as proposed in the preferred embodiment of the present invention, the user equipment is set (or adapted) to perform operations according to the failure to receive the random access response message.

[89] More specifically, when the user terminal receives a random access response message including only the backoff indicator MAC subheader from the base station according to the embodiment of the present invention, it is proposed in the description of the present invention that the user equipment is set (or adapted) to perform the following operations according to the failure to receive the random access response message.

[90] 1. Increment a PREAMBLE_TRANSMISSION_COUNTER value by 1.

[91] 2. If the condition of PREAMBLE_TRANSMISSION_COUNTER = PREAMBLE_TRANS_MAX + 1 is satisfied, the problem of random access is indicated (or reported) to the higher (or upper) layer.

[92] 3. If the random access procedure currently under process corresponds to the contention based random access procedure, *i.e.*, if a random access preamble is selected by a medium access control (MAC) layer, select a backoff time according to a uniform distribution, so as to delay transmission by the selected backoff time, thereby performing the random access source selection process (*i.e.*, the random access preamble retransmission process).

[93] Hereinafter, improvements of the algorithm for receiving and processing a random access response message of a MAC layer module included in the user equipment according to the embodiments of the present invention will be described in detail.

[94]

[95] Improvement 1

[96] FIG. 11 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to an embodiment of the present invention.

[97] In comparison with FIG. 4, step 1203 and a procedure for processing a subsequent random access response message depending upon the result determined in step 1203 are additionally included in the algorithm for receiving and processing a random access response message of FIG. 11. More specifically, when the user equipment receives a random access response message including only the backoff indicator MAC subheader from the base station, since the conditions of step 1201 and step 1202 are not satisfied, as described above, the procedure moves on to step 1203. In step 1203, when the user equipment receives a random access response message including only the backoff indicator MAC subheader, it is determined that the reception of the random access

response message has failed. And, accordingly, in step 1203, the random access transmission counter is incremented by '1', thereby reporting the random access problem to the higher layer according to the random access transmission counter value, or attempting retransmission of the random access preamble by applying the received backoff indicator.

[98]

[99] Improvement 2

[100] FIG. 12 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to another embodiment of the present invention.

[101] In comparison with FIG. 4, the conditions of step 1301 are improved in the algorithm for receiving and processing a random access response message of FIG. 12. More specifically, in step 1301, the user equipment determines whether a random access response including a random access preamble identifier is received in a predetermined random access response reception window. If the user equipment receives a random access response including only a backoff indicator MAC subheader from the base station, this indicates that the random access response including a random access preamble identifier has not been received in the predetermined random access response reception window. Thus, the procedure moves on to step 1304, so as to perform operations according to the failure to receive the random access response. And, accordingly, in step 1304, the user equipment increments the random access transmission counter by '1', thereby reporting the random access problem to the higher layer according to the random access transmission counter value, or attempting retransmission of the random access preamble by applying the received backoff indicator.

[102] The remaining parts and respective operations of the algorithm for receiving and processing a random access response message of FIG. 12 are identical to those described in FIG. 4.

[103]

[104] Improvement 3

[105] FIG. 13 illustrates an algorithm for receiving and processing a random access response message of a user equipment improved according to yet another embodiment of the present invention.

[106] In comparison with FIG. 4, the conditions of step 1402 are improved in the algorithm for receiving and processing a random access response message of FIG. 13. More specifically, when the user equipment receives a random access response including only a backoff indicator MAC subheader, as shown in FIG. 10, in the random access response reception window, from the base station, this indicates that the random access

response has been received in the corresponding window in step 1401. Therefore, the procedure moves on to step 1402.

- [107] In step 1402, the user equipment determines the existence (or presence) of a random access response including a random access preamble identifier respective to the random access preamble transmitted by the corresponding user equipment. As shown in the above-described example, in the user equipment that received the random access response including only a backoff indicator MAC subheader, the corresponding random access response does not correspond to (or match) the random access response including the random access preamble identifier respective to the random access preamble transmitted from the corresponding user equipment. Therefore, the procedure moves on to step 1404, thereby performing operations according to the failure to receive the random access response.

[108]

- [109] Among the above-described improvements 1 to 3, improvement 1 represents an algorithm directly connecting the case of receiving a random access response including only a backoff indicator MAC subheader, as shown in FIG. 10, to the operations according to the failure to receive the random access response. And, improvement 2 and improvement 3 respectively represent algorithms modifying (or changing) one specific condition of the conventional determination steps, shown in FIG. 4, in order to connect the case of receiving a random access response including only a backoff indicator MAC subheader, as shown in FIG. 10, to the operations according to the failure to receive the random access response.

- [110] Meanwhile, exemplary operations of the user equipment according to the above-described embodiments of the present invention will now be described in detail with reference to FIG. 14. FIG. 14 illustrates a method of performing random access in a user equipment according to an embodiment of the present invention.

- [111] First of all, in step 1501, the user equipment may transmit a random access preamble to the base station in a contention or non-contention based random access procedure. According to an embodiment of the present invention, the base station may transmit a MAC PDU format random access response message including only the backoff indicator MAC subheader, as shown in FIG. 10, to the user equipment (S1502). Also, according to a preferred embodiment of the present invention, the user equipment, which has received MAC PDU format random access response message including only the backoff indicator MAC subheader, is set (or adapted) to determine that the random access response reception has failed (S1503). For this, the above-described algorithm for receiving a random access response message of the user equipment with reference to FIG. 11 to FIG. 13 may be used. Such algorithm for receiving a random access response message of the user equipment may be set and used in a MAC layer module

of a user equipment processor in the form of hardware or software.

[112] Meanwhile, the user equipment, which has determined a random access response failure in step 1503, may perform the following operations according to the failure to receive the random access response message (S1504).

[113] A. Increment a PREAMBLE_TRANSMISSION_COUNTER value by 1.

[114] B. If the condition of PREAMBLE_TRANSMISSION_COUNTER = PREAMBLE_TRANS_MAX + 1 is satisfied, the problem of random access is indicated (or reported) to the higher (or upper) layer.

[115] C. If the random access procedure currently under process corresponds to the contention based random access procedure, select a backoff parameter value according to a uniform distribution, so as to delay transmission by the selected backoff parameter value, thereby performing the random access source selection process.

[116]

[117] After performing the above-described random access source selection process, the user equipment may transmit a random access preamble to the base station (S1505).

[118]

[119] Hereinafter, the device structure of the above-described user equipment performing random access to the base station will be described in detail. FIG. 15 illustrates a structure of a user equipment according to an embodiment of the present invention.

[120] The user equipment according to the embodiment of the present invention and, more particularly, a processor of the user equipment may include a physical (PHY) layer module 1600, a medium access control (MAC) layer module 1700, and a higher layer module. Herein, the PHY layer module 1600 includes a transmitting module (Tx module) 1610 and a receiving module (Rx module) 1620. The MAC layer 1700 performs the actual processing operations of the random access procedure. And, the higher layer module corresponds to a layer higher than the above-described MAC layer, with reference to FIG. 2 and FIG. 3.

[121] First of all, the user equipment according to the embodiment of the present invention may be configured so as to be capable of transmitting a random access preamble through the Tx module 1610 of the PHY layer module 1600. Additionally, the user equipment according to the embodiment of the present invention may be configured so as to be capable of receiving a medium access control protocol data unit (MAC PDU) format random access response message through the Rx module 1620 of the PHY layer module 1600 from the base station. Herein, the MAC PDU format random access response message corresponds to the response for the random access preamble transmitted by the user equipment. More specifically, it is proposed that the Rx module 1620 of the user equipment is configured so as to receive the MAC PDU format random access response message including zero, one, or multiple MAC RARs.

- [122] Meanwhile, the MAC layer module 1700 of the user equipment according to the embodiment of the present invention may include a random access response (RAR) reception control module 1710 and a processing module for RAR reception failure 1720. The RAR reception control module 1710 of the MAC layer module 1700 according to the embodiment of the present invention corresponds to a module for embodying (or realizing) the algorithm for receiving and processing random access response messages, as described in FIG. 11 to FIG. 13, in the form of hardware or software. More specifically, if the random access response message received by the Rx module 1610 corresponds to the MAC PDU including only the backoff indicator subheader in the MAC header, the reception of the random access response message is determined to have failed. Accordingly, the user terminal according to the present invention is set (or adapted) to induce the operations of the processing module for RAR reception failure 1720, so as to perform subsequent operations according to the RAR reception failure.
- [123] More specifically, the RAR reception control module 1710 is set, so as to determine the case where the Rx module 1620 fails to receive a random access response message within the random access response message reception window, which has a predetermined time length, or to determine the case where the Rx module 1620 fails to receive a random access response message including a random access preamble identifier respective to the random access preamble transmitted by the Tx module 1610, as a case where the RAR message reception has failed. Herein, if the Rx module 1620 has failed to receive the random access response message including the random access preamble identifier respective to the random access preamble transmitted by the Tx module 1610, two different cases are assumed. In a first assumed case, all RAR messages received by the Rx module 1620 within the RAR message reception window include a random access preamble identifier that does not correspond with (or match) the random access preamble identifier respective to the random access preamble transmitted from the Tx module 1610. And, in the second assumed case, the Rx module 1620 receives a MAC PDU format RAR message including only the backoff indicator subheader within the MAC header.
- [124] Accordingly, the processing module 1720 performing operations according to the random access reception failure of the MAC layer module 1700 increments the random access preamble transmission counter by '1'. And, in case, the random access preamble transmission counter has reached a maximum transmission counter value, it is reported to a higher layer module that a problem in random access exists. Thereafter, if the current random access procedure corresponds to the contention based random access procedure, the backoff indicator is used so as to select a backoff parameter value based upon a uniform distribution and to delay transmission by the selected backoff

parameter value, thereby performing the random access source selection process. After performing the above-described random access source selection process, the user equipment may retransmit the random access preamble through the Tx module 1610.

- [125] The MAC PDU format RAR message received by the Rx module 1620 of the user equipment may be received through a physical downlink shared channel (PDSCH). And, the Rx module 1620 of the user equipment may monitor the PDCCH within the predetermined RAR reception window, so as to receive the MAC PDU format RAR message.

Industrial Applicability

- [126] The random access scheme for a user equipment according to the present invention has the following advantages. By supporting MAC PDU format RAR messages, which include only a backoff indicator subheader in the MAC header according to the embodiments of the present invention, a concentration in random access from multiple user equipments, in case of a lack in system source, may be efficiently distributed. Furthermore, when algorithm for receiving and processing RAR messages according to the embodiments of the present invention is improved, even when supporting MAC PDU format RAR messages including only a backoff indicator subheader in the MAC header, as described above, the user equipment may perform subsequent operations without any error.
- [127] It will be apparent to those skilled in the art that various modifications and variations can be made in the present invention without departing from the scope of the inventions. Thus, it is intended that the present invention covers the modifications and variations of this invention provided they come within the scope of the appended claims and their equivalents. The scope of the claims should not be limited by the embodiments set forth in the examples herein, but should be given the broadest interpretation consistent with the description as a whole.

74420-516

23

CLAIMS:

1. A method of performing a random access to a base station by a terminal, the method comprises:

transmitting a random access preamble to the base station;

5 receiving, from the base station in response to the random access preamble, a random access response message having a format of medium access control protocol data unit (MAC PDU) including only a backoff indicator subheader in a header portion of the MAC PDU and no medium access control random access response in the MAC PDU; and

10 considering a random access response reception procedure not successful, and performing a subsequent procedure for a random access responses reception failure.

2. The method of claim 1, wherein the terminal is adapted to consider the random access response reception procedure not successful, if no random access response is received within a random access response window having a predetermined time length, or if none of all received random access responses contains a random access preamble identifier
15 corresponding to the transmitted random access preamble.

3. The method of claim 2, wherein a case in which none of all received random access responses contains the random access preamble identifier corresponding to the transmitted random access preamble comprises:

20 a first case in which all received random access responses contain random access preamble identifiers that do not match the transmitted random access preamble; and

a second case in which the random access response message having the format of MAC PDU including only the backoff indicator subheader in the header portion of the MAC PDU and no medium access control random access response in the MAC PDU is received.

74420-516

24

4. The method of claim 1, wherein the subsequent procedure for the random access response reception failure is performed by:

incrementing a random access preamble transmission counter by 1; and

indicating to a higher layer higher than a medium access control (MAC) layer

5 that there is a random access problem, if the random access preamble transmission counter becomes a predetermined maximum number of preamble transmission.

5. The method of claim 1, wherein the subsequent procedure for the random access response reception failure is performed by:

delaying a subsequent random access preamble transmission by a backoff time

10 selected using a backoff indicator in the backoff indicator subheader, if the random access preamble is selected by a medium access control (MAC) layer.

6. The method of claim 1, wherein the random access response message having the format of MAC PDU is received through a physical downlink shared channel (PDSCH).

7. A terminal performing a random access to a base station, the terminal
15 comprises:

a transmitting (Tx) module for transmitting a random access preamble to the base station;

a receiving (Rx) module for receiving, from the base station in response to the random access preamble, a random access response message having a format of medium
20 access control protocol data unit (MAC PDU); and

a medium access control (MAC) layer module adapted to consider a random access response reception procedure not successful and to perform a subsequent procedure for a random access response reception failure, when the random access response message received by the receiving module is a MAC PDU including only a backoff indicator

74420-516

25

subheader in a header portion of the MAC PDU and no medium access control random access response in the MAC PDU.

8. The terminal of claim 7, wherein the MAC layer module is adapted to consider the random access response reception procedure not successful, if the receiving module does
5 not receive any random access response within a random access response window having a predetermined time length, or if none of all received random access responses containing a random access preamble identifier corresponding to the random access preamble transmitted from the transmitting module is received by the receiving module.

9. The terminal of claim 8, wherein a case in which none of all received random
10 access responses containing the random access preamble identifier corresponding to the random access preamble transmitted from the transmitting module is received by the receiving module, comprises:

a first case in which all random access responses received by the receiving module contain random access preamble identifiers that do not match the random access
15 preamble transmitted from the transmitting module; and

a second case in which the random access response message having the format of MAC PDU including only the backoff indicator subheader in the header portion of the MAC PDU and no medium access control random access response in the MAC PDU is received by the receiving module.

20 10. The terminal of claim 7, wherein, as the subsequent procedure for the random access response reception failure, the MAC layer module increments a random access preamble transmission counter by 1, and

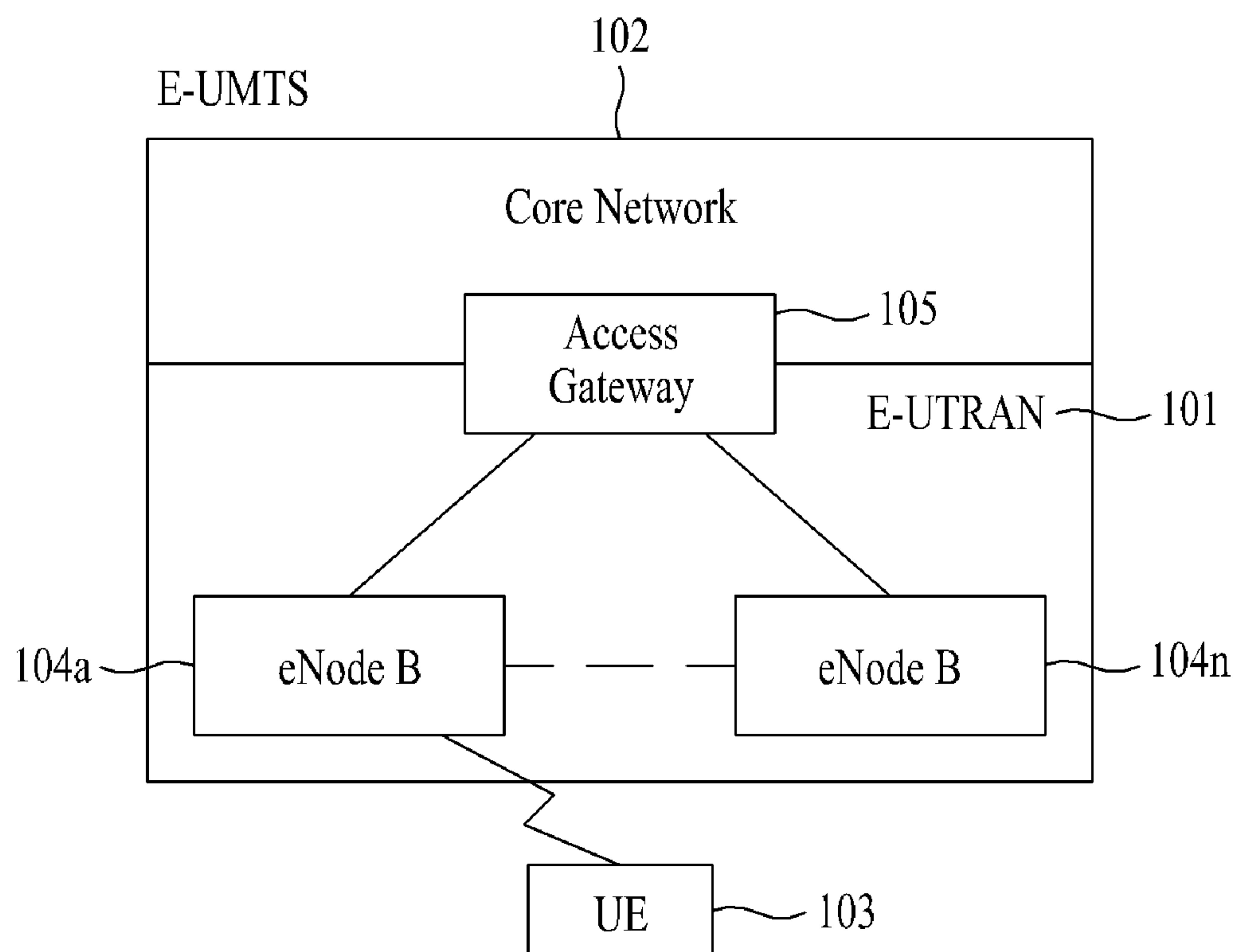
wherein the MAC layer module indicates to a higher layer module corresponding to a higher layer higher than a medium access control (MAC) layer that there is
25 a random access problem, if the random access preamble transmission counter becomes a predetermined maximum number of preamble transmission.

74420-516

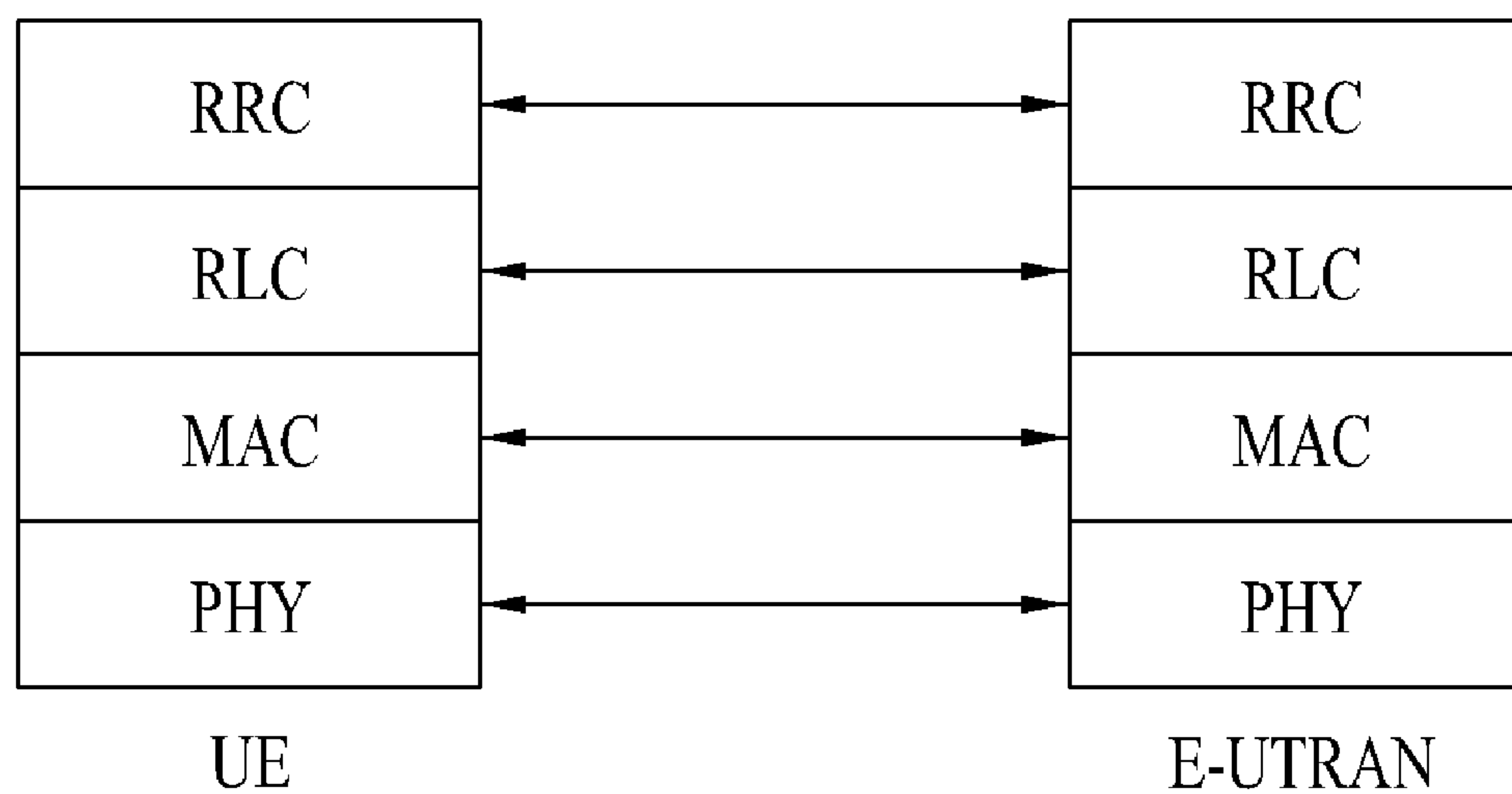
26

11. The terminal of claim 7, wherein, as the subsequent procedure for the random access response reception failure, the MAC layer module delays a subsequent random access preamble transmission by a backoff time selected using a backoff indicator in the backoff indicator subheader, if the random access preamble is selected by a medium access control
- 5 (MAC) layer.
12. The terminal of claim 7, wherein the receiving module receives the random access response message having the format of MAC PDU through a physical downlink shared channel (PDSCH).

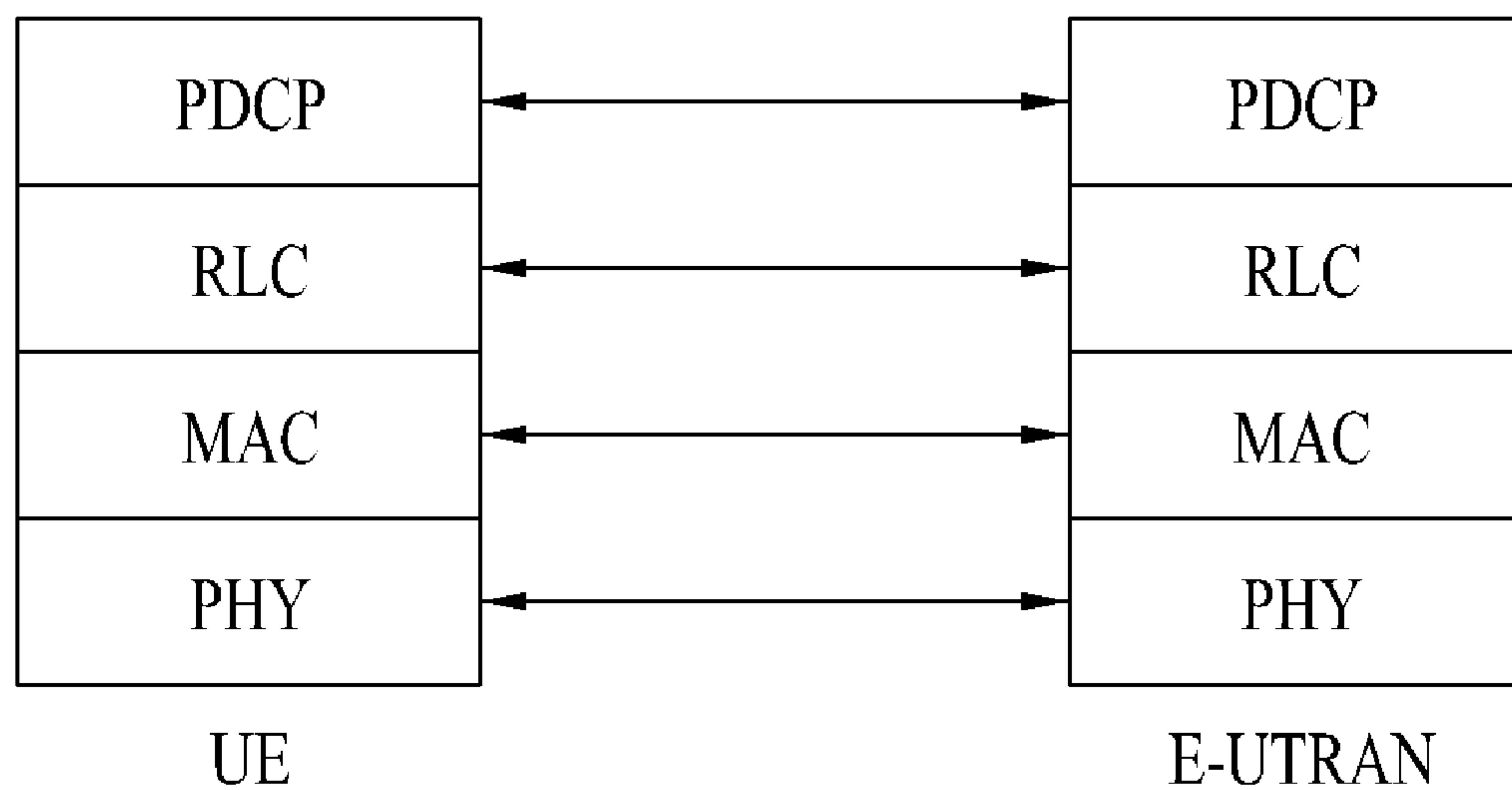
[Fig. 1]



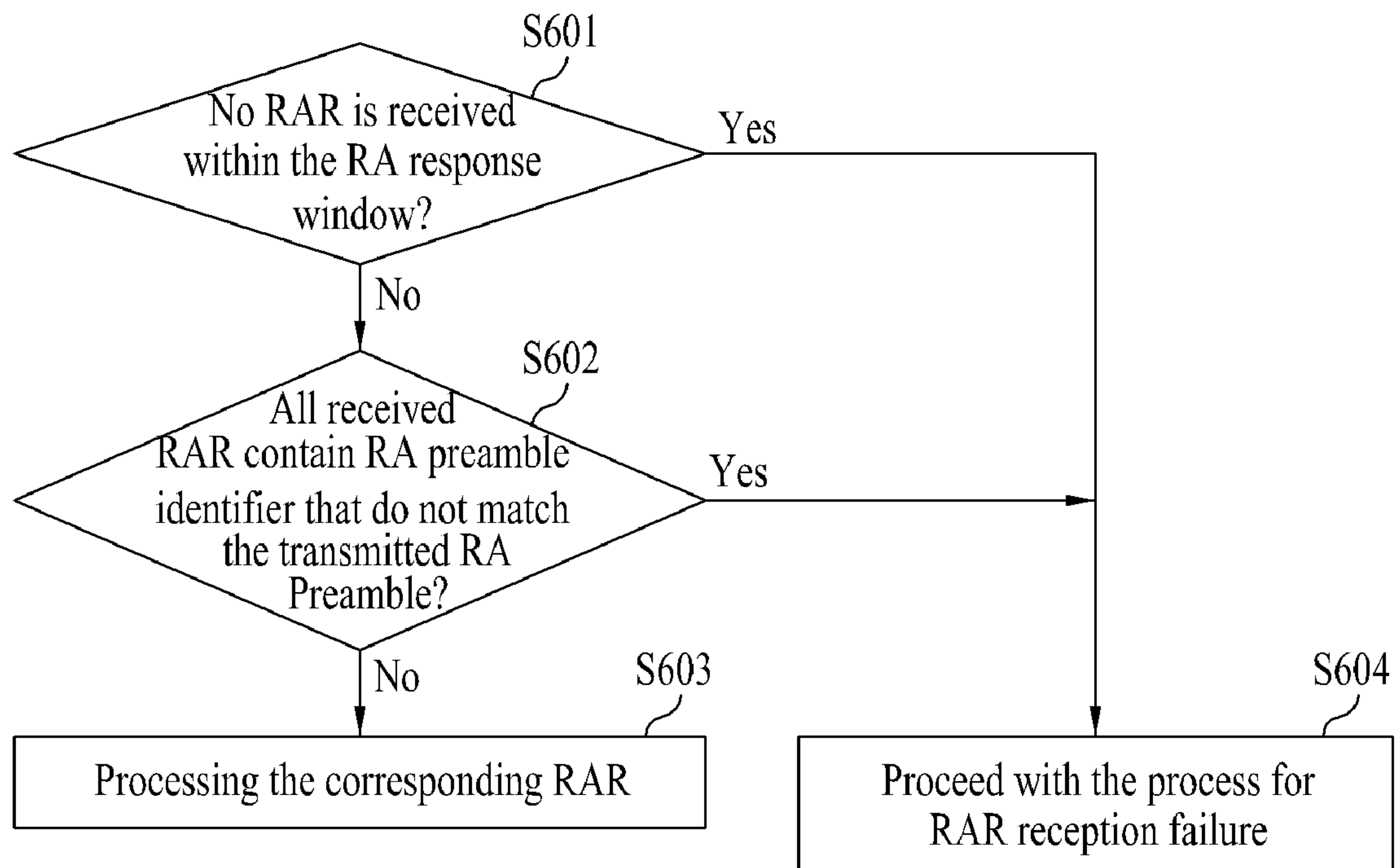
[Fig. 2]



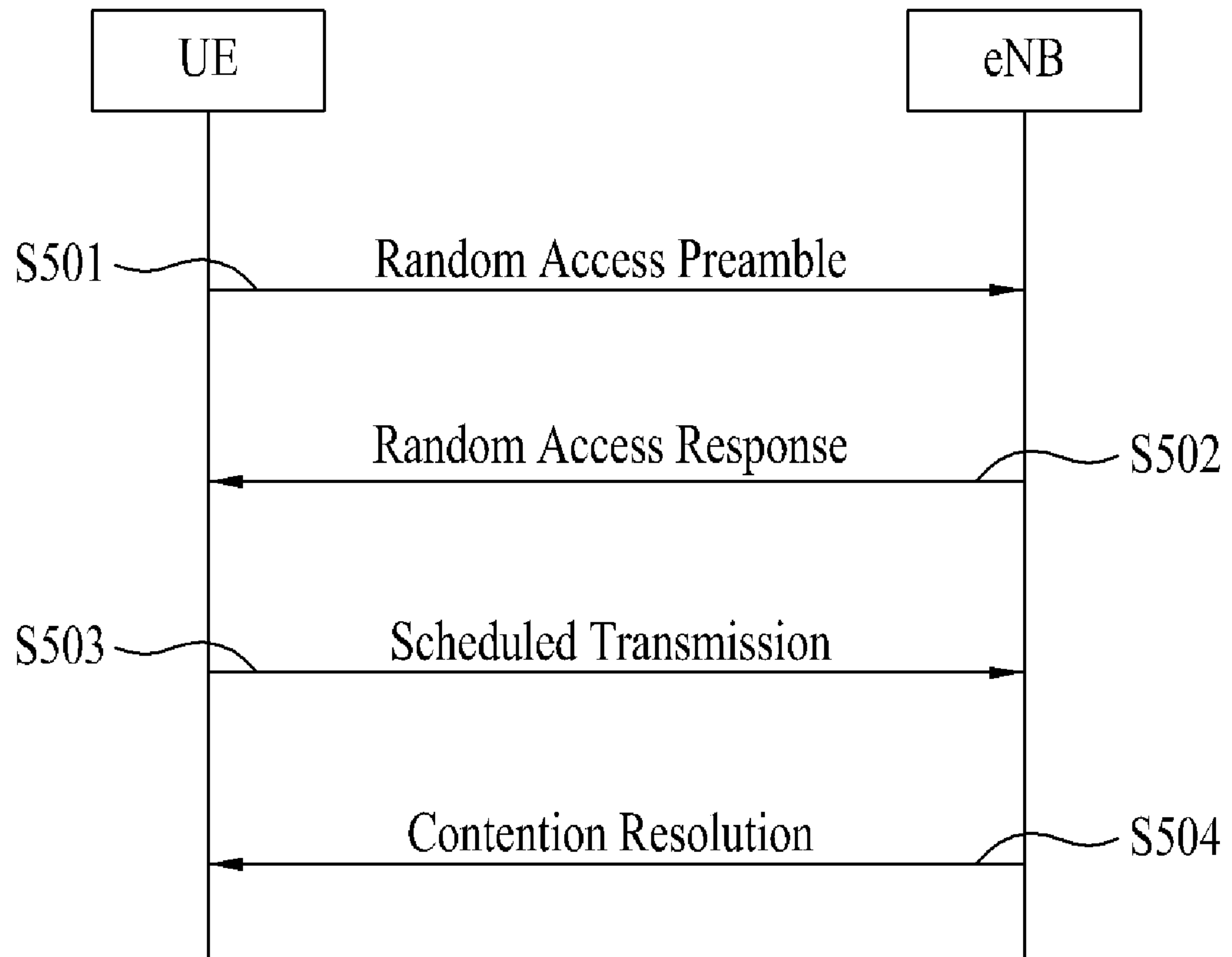
[Fig. 3]



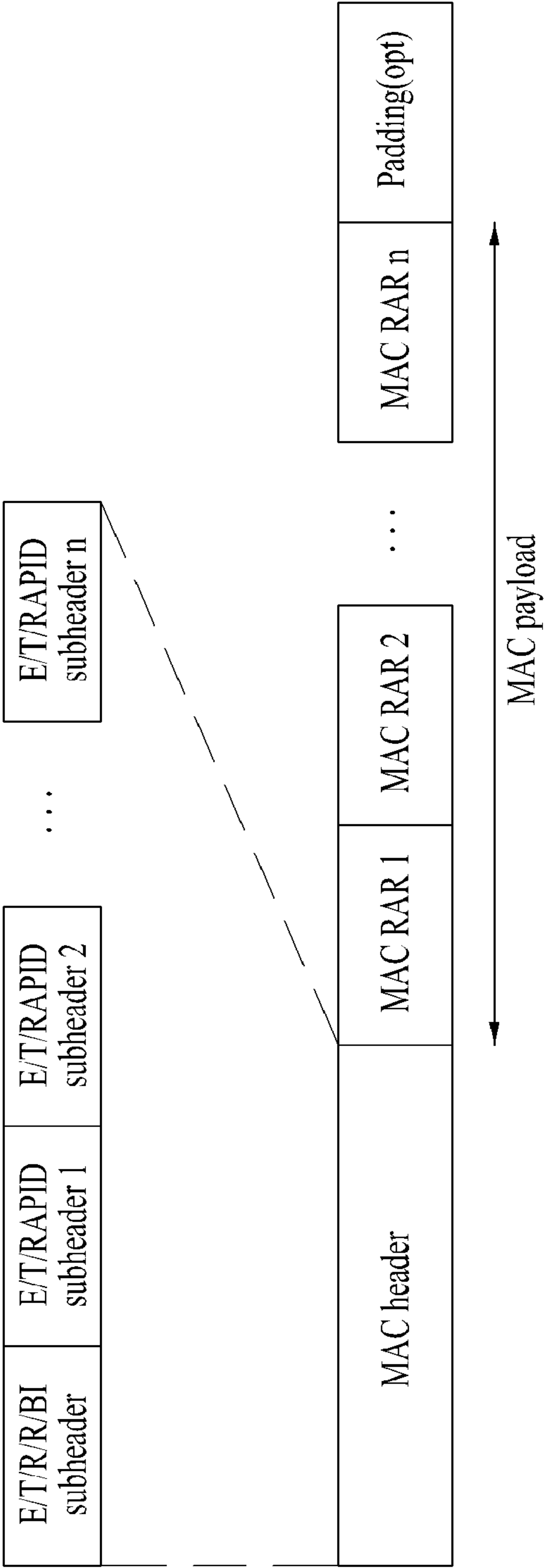
[Fig. 4]



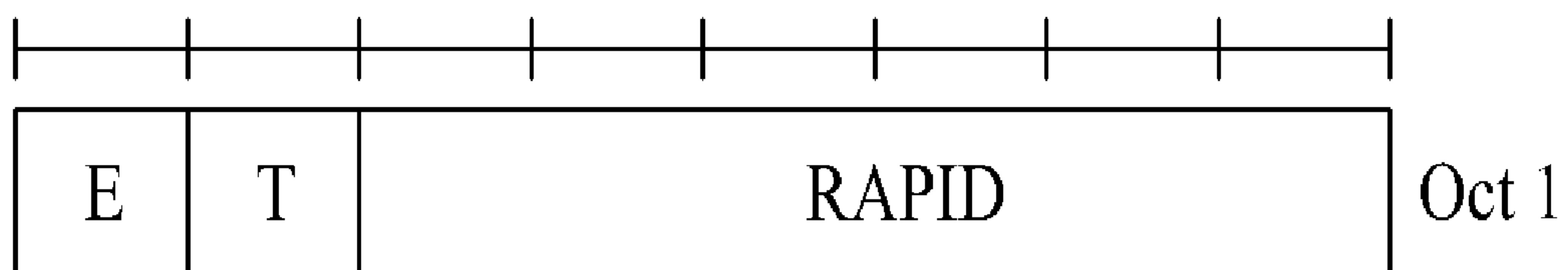
[Fig. 5]



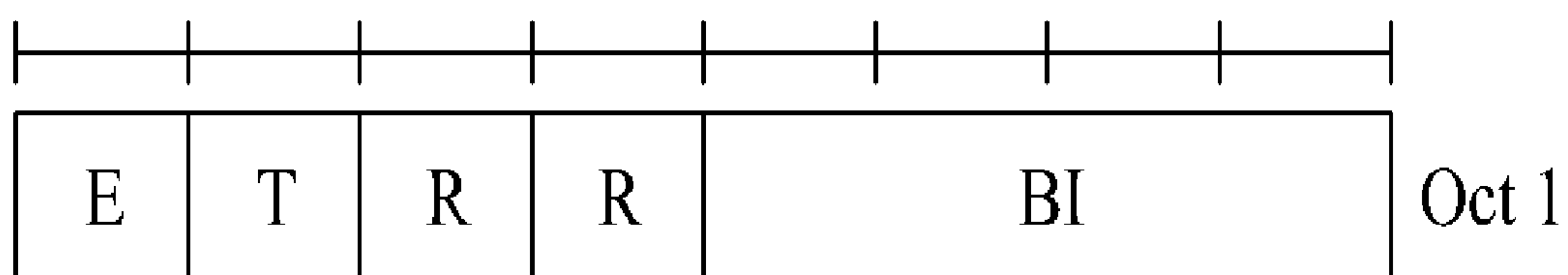
[Fig. 6]



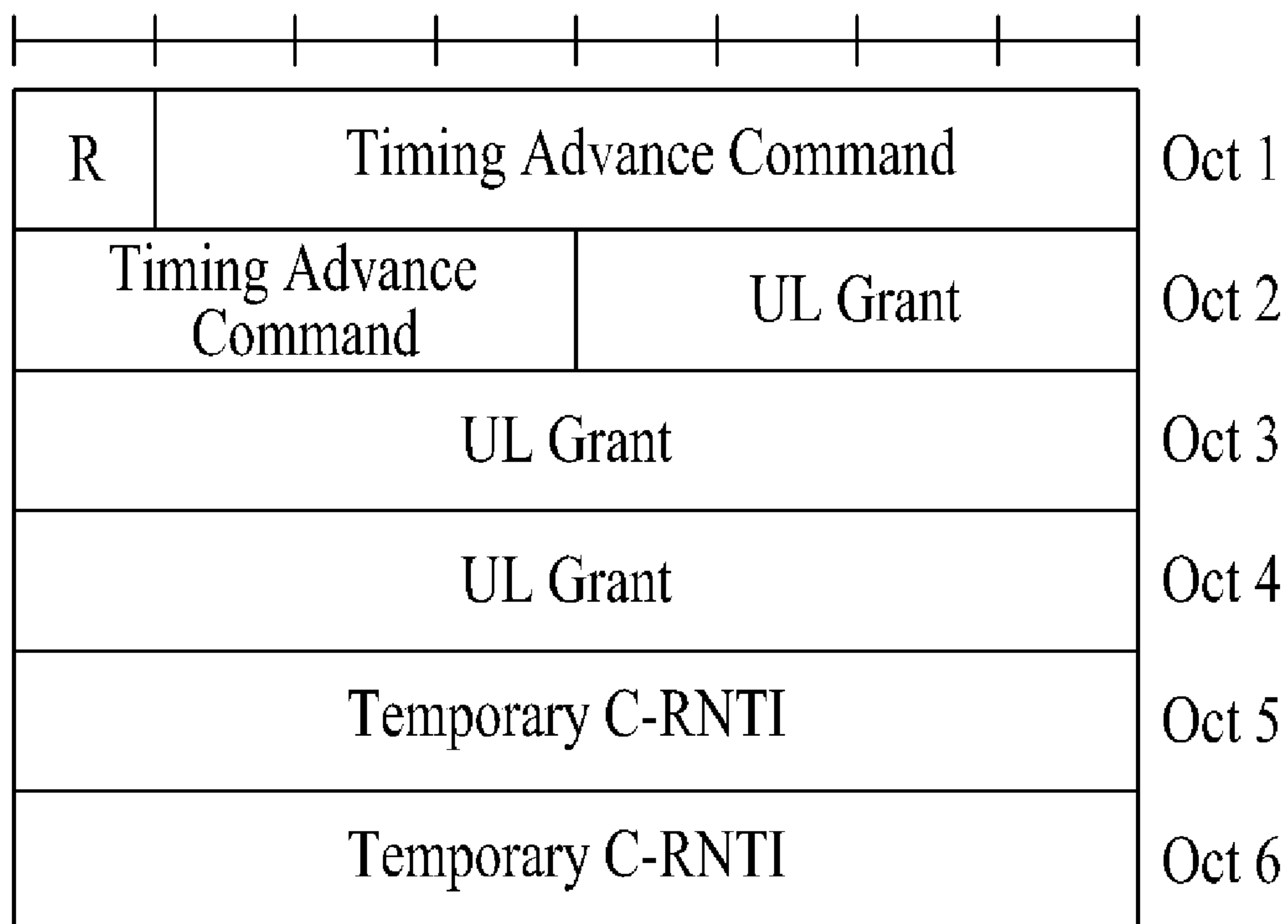
[Fig. 7]



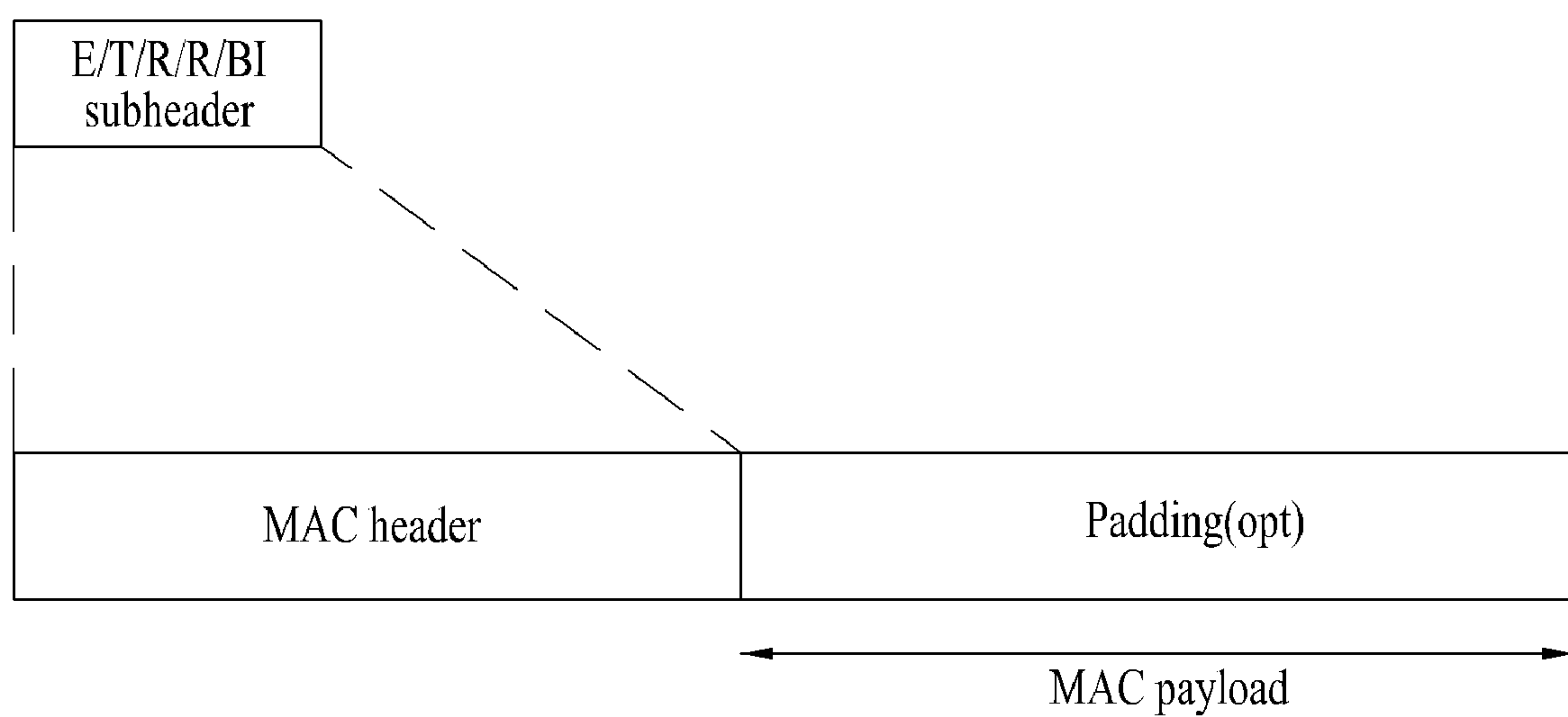
[Fig. 8]



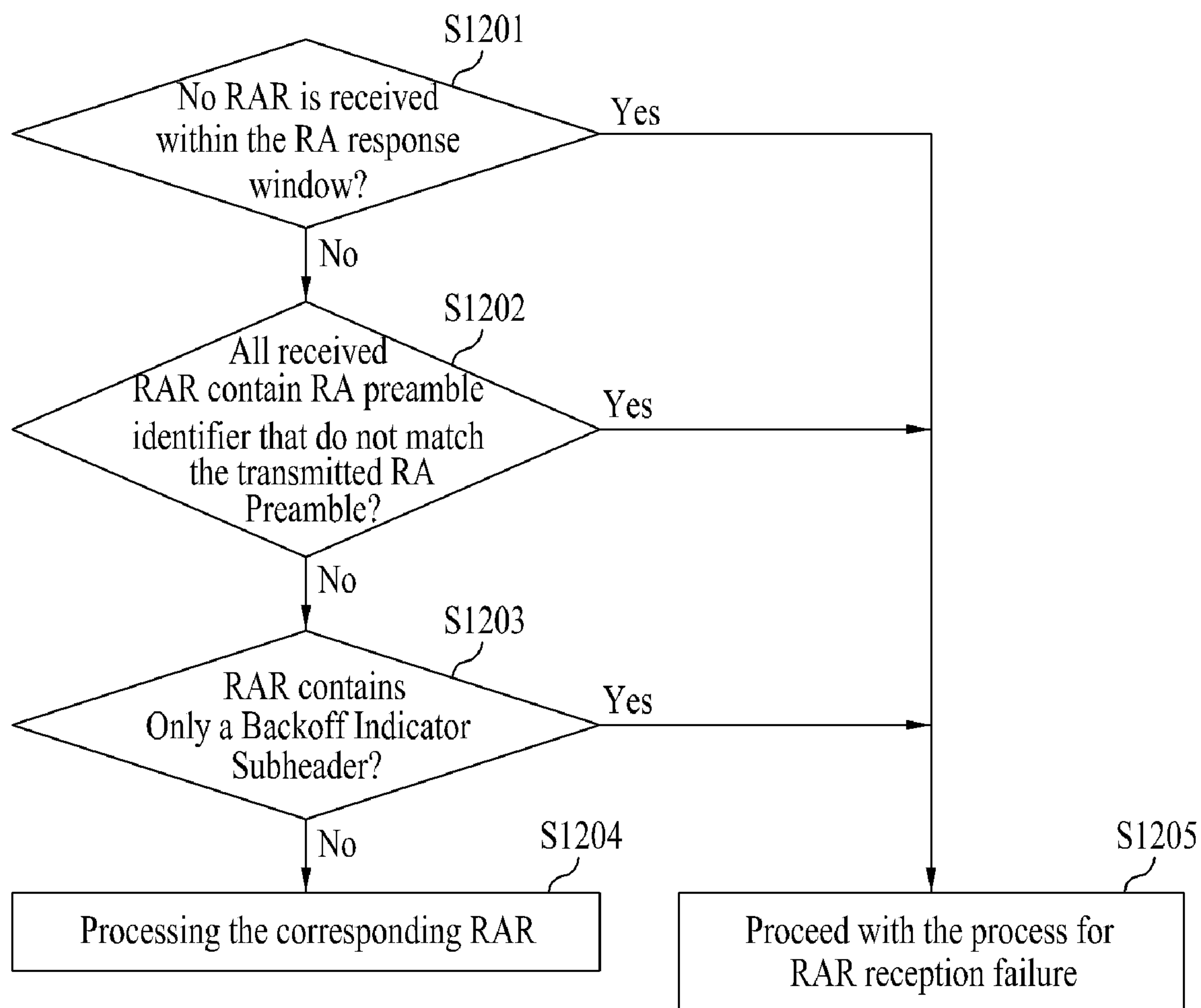
[Fig. 9]



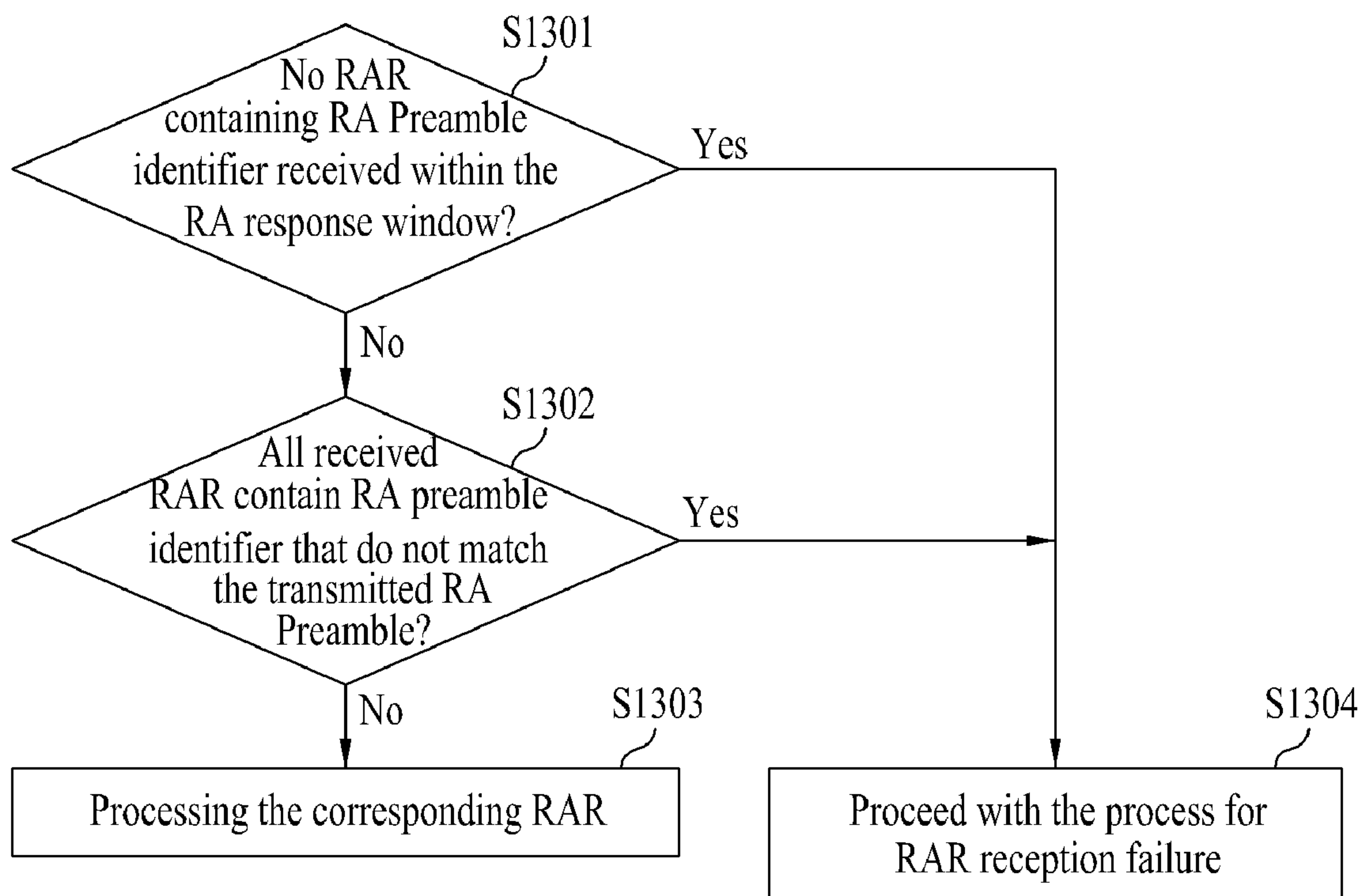
[Fig. 10]



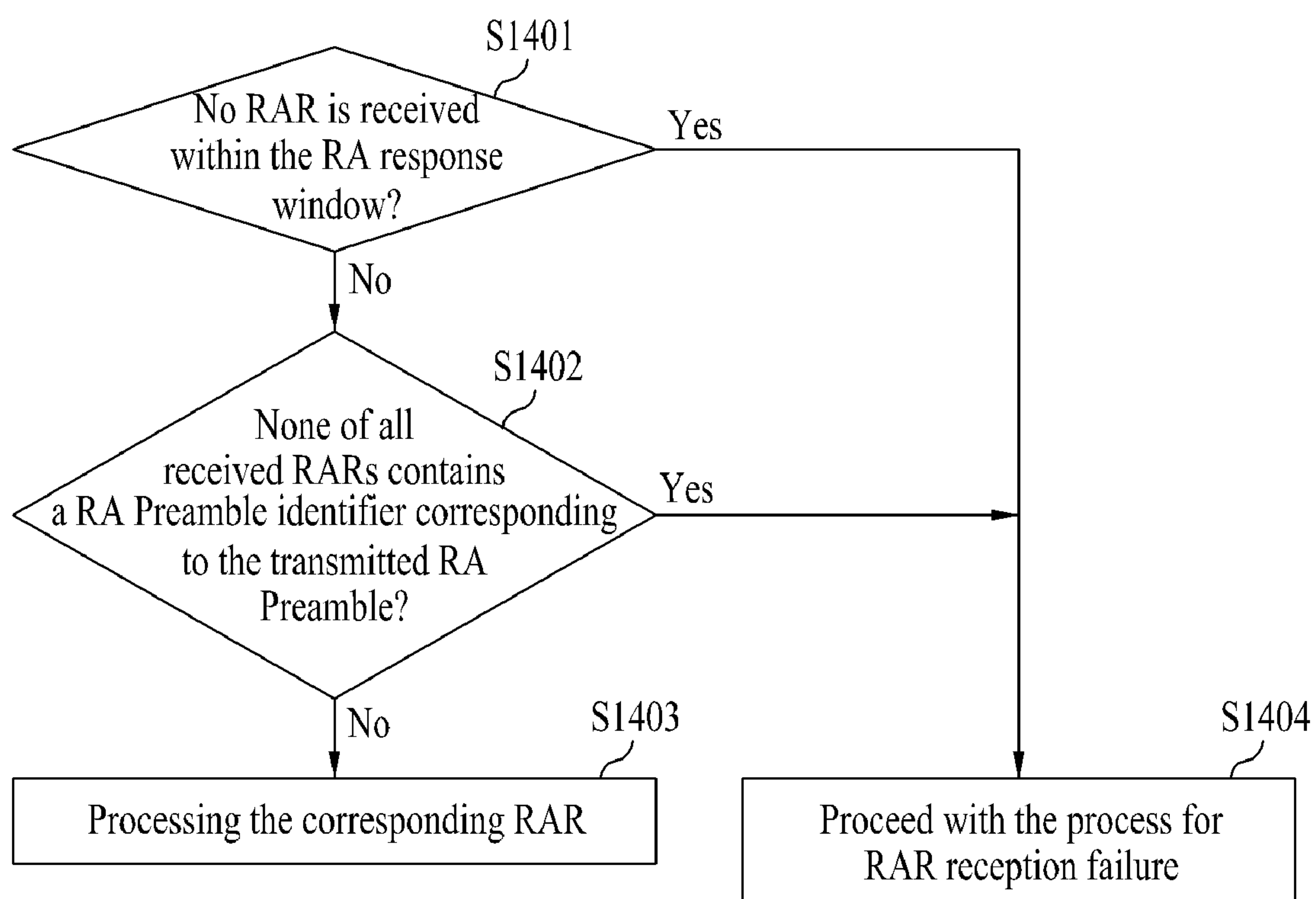
[Fig. 11]



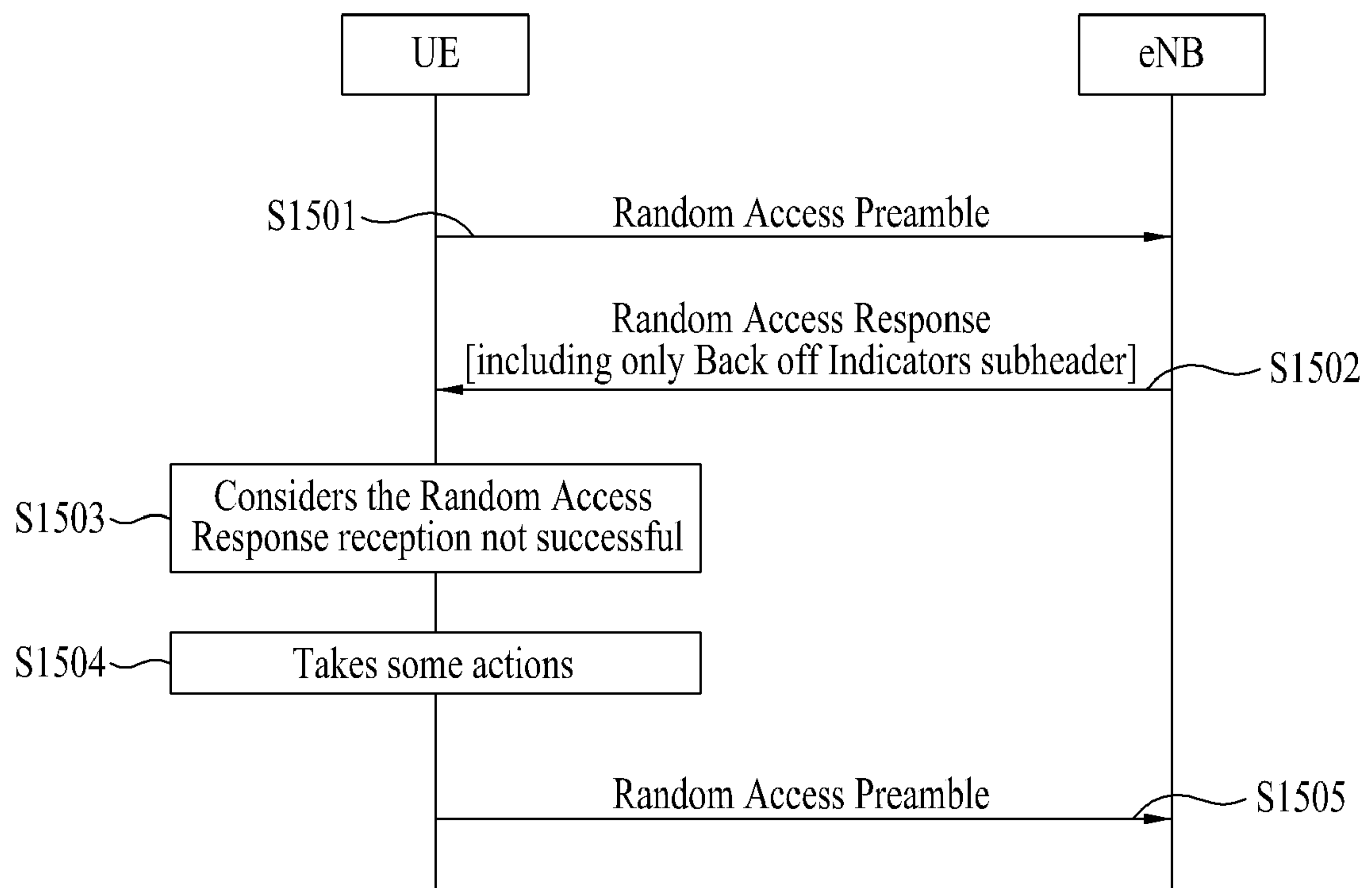
[Fig. 12]



[Fig. 13]



[Fig. 14]



[Fig. 15]

