

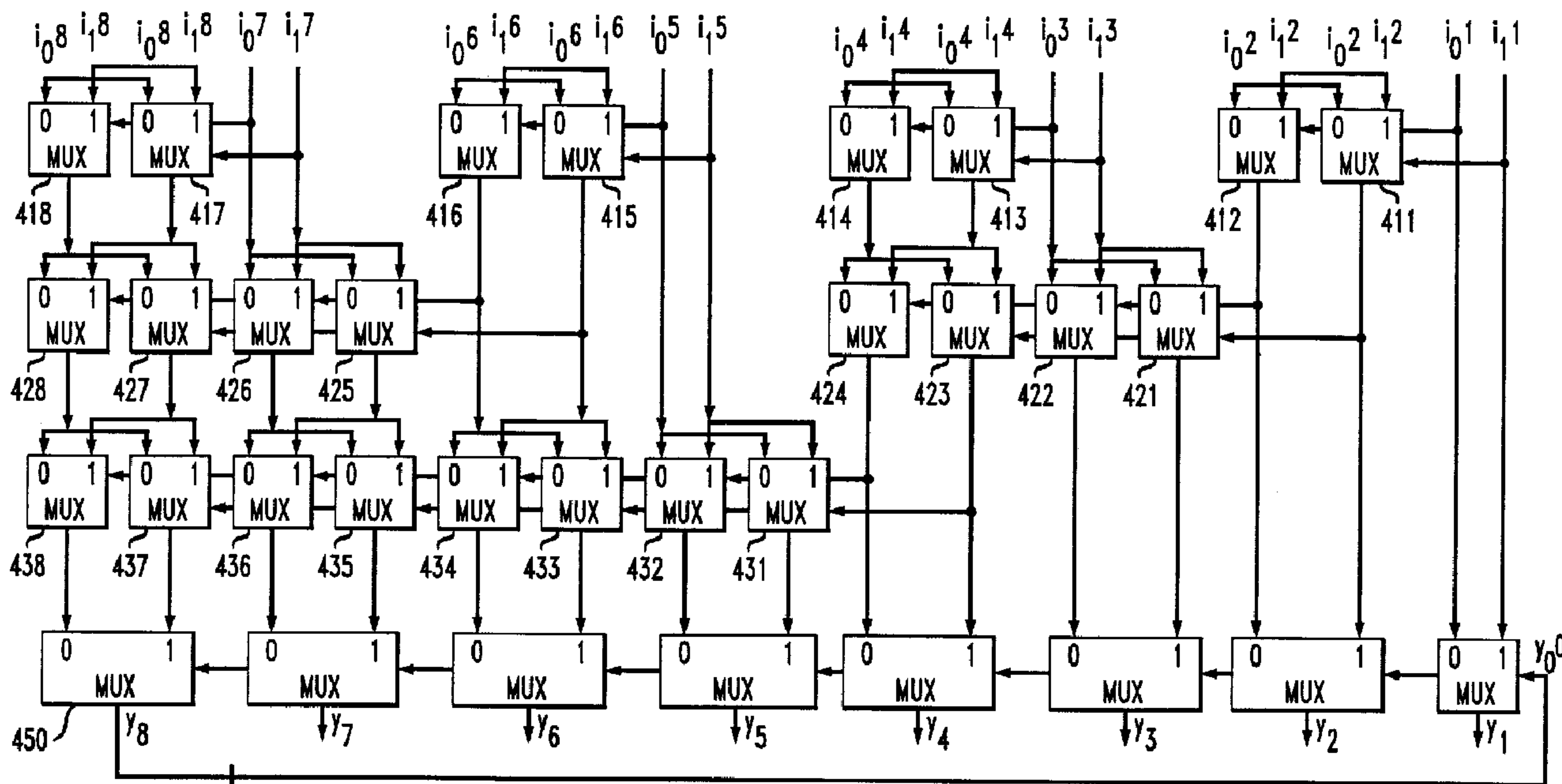


(22) Date de dépôt/Filing Date: 2000/05/29  
(41) Mise à la disp. pub./Open to Public Insp.: 2000/12/04  
(45) Date de délivrance/Issue Date: 2004/04/06  
(30) Priorité/Priority: 1999/06/04 (09/326,781) US

(51) Cl.Int.<sup>7</sup>/Int.Cl.<sup>7</sup> H04L 27/01, H04B 3/04  
(72) Inventeurs/Inventors:  
AZADET, KAMERAN, US;  
YU, MENG-LIN, US  
(73) Propriétaire/Owner:  
LUCENT TECHNOLOGIES INC., US  
(74) Agent: KIRBY EADES GALE BAKER

(54) Titre : EGALISATEUR A DECISION RETROACTIVE DE TRAITEMENT PARALLELE AVEC TRAITEMENT ANTICIPE  
(54) Title: A PARALLEL PROCESSING DECISION-FEEDBACK EQUALIZER (DFE) WITH LOOK-AHEAD PROCESSING

400



(57) Abrégé/Abstract:

A method and apparatus are disclosed for increasing the effective processing speed of a parallel decision-feedback equalizer by combining block processing and look-ahead techniques in the selection (multiplexing) stage. The present invention extends a parallel decision-feedback equalization by using look-ahead techniques in the selection stage to precompute the effect of previous blocks on each subsequent block, and to thereby remove the serial output dependency. The parallel decision-feedback equalization includes a multiplexor tree structure that selects an appropriate output value for each block and precomputes the effect of previous blocks on each subsequent block. A multiplexing delay algorithm on the order of logN is

**(57) Abrégé(suite)/Abstract(continued):**

employed to resolve the output dependency and thus speeds up parallel block processing decision-feedback equalizations. The disclosed decision-feedback equalization architecture can be combined with pipelining to completely eliminate the critical path problem. Pipelining reduces the required critical path timing to one multiplexing time. The disclosed multiplexor tree circuitry for the parallel decision-feedback equalizer groups multiplexor blocks into groups of two, referred to as block pairs, and provides at least one multiplexor for each block,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values. The output of each parallel block depends on the possible precomputed values generated by the look-ahead processors for the block, as well as the actual values that are ultimately selected for each previous block. In order to reduce the delay in obtaining each actual output value, the present invention assumes that each block contains each possible value, and carries the assumption through to all subsequent blocks. Thus, the number of multiplexors required to select from among the possible values grows according to  $N\text{-log}N$ , where  $N$  is the block number.

**ABSTRACT**

A method and apparatus are disclosed for increasing the effective processing speed of a parallel decision-feedback equalizer by combining block processing and look-ahead techniques in the selection (multiplexing) stage. The present invention extends a parallel decision-feedback equalization by using look-ahead techniques in the selection stage to precompute the effect of previous blocks on each subsequent block, and to thereby remove the serial output dependency. The parallel decision-feedback equalization includes a multiplexor tree structure that selects an appropriate output value for each block and precomputes the effect of previous blocks on each subsequent block. A multiplexing delay algorithm on the order of  $\log N$  is employed to resolve the output dependency and thus speeds up parallel block processing decision-feedback equalizations. The disclosed decision-feedback equalization architecture can be combined with pipelining to completely eliminate the critical path problem. Pipelining reduces the required critical path timing to one multiplexing time. The disclosed multiplexor tree circuitry for the parallel decision-feedback equalizer groups multiplexor blocks into groups of two, referred to as block pairs, and provides at least one multiplexor for each block,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values. The output of each parallel block depends on the possible precomputed values generated by the look-ahead processors for the block, as well as the actual values that are ultimately selected for each previous block. In order to reduce the delay in obtaining each actual output value, the present invention assumes that each block contains each possible value, and carries the assumption through to all subsequent blocks. Thus, the number of multiplexors required to select from among the possible values grows according to  $N \cdot \log N$ , where  $N$  is the block number.

## **A PARALLEL PROCESSING DECISION-FEEDBACK EQUALIZER WITH LOOK-AHEAD PROCESSING**

### **Field of the Invention**

The present invention relates generally to channel equalization techniques, and  
5 more particularly, to decision-feedback equalizers (DFEs) employing parallel (block)  
processing.

### **Background of the Invention**

Signals arriving at a receiver are typically corrupted by intersymbol  
interference (ISI), crosstalk, echo, and other noise. Thus, receivers must jointly equalize the  
10 channel, to compensate for such intersymbol interference and other distortions, and decode  
the encoded signals at increasingly high clock rates. Decision-feedback equalization is a  
widely-used technique for removing intersymbol interference where noise enhancement  
caused by a linear feed-forward equalizer (FFE) may introduce performance problems. For a  
detailed discussion of decision feedback equalizers, see, for example, R. Gitlin et al., *Digital*  
15 *Communication Principles*, (Plenum Press 1992) and E.A. Lee and D.G. Messerschmitt,  
*Digital Communications*, (Kluwer Academic Press, 1988). Generally, decision-feedback  
equalization utilizes a nonlinear equalizer to equalize the channel using a feedback loop based  
on previously decided symbols. Since decision-feedback equalization techniques are  
nonlinear, the decision-feedback equalization process does not enhance the channel noise.

20 In many high-speed applications, such as Fast Ethernet 100BASE-TX, mass  
storage, Gigabit Ethernet, or Synchronous Optical Networks (SONET), the symbol rates are  
high. For example, the Fast Ethernet 100BASE-TX standard uses a clock rate of 125MHz.  
Thus, the equalization and decoding computations performed by the decision-feedback  
equalization must be performed at a clock period of 8 nanoseconds. Other advanced data  
25 networking applications, such as the SONET standard, may require even shorter clock  
periods. In many cases, such a short clock period makes the implementation of a decision-  
feedback equalizer challenging or infeasible and forms the critical path in their  
implementation.

A number of methods have been proposed or suggested for speeding up the  
30 decision-feedback equalization processing. For example, K. Parhi, "Pipelining in Algorithm  
with Quantizer Loops". *IEEE Transactions on Circuits and Systems*, Vol. 38, No. 7, 745-54

(July 1991), proposes a look-ahead architecture. Generally, a look-ahead decision-feedback equalization implementation transforms the original decision-feedback equalization feedback loop into a look-ahead structure (with duplicate decision-feedback equalizations for each possible value) and a simple loop with only a large selection multiplexor and one memory device. The look-ahead technique precomputes the symbol value for each possible variation, then utilizes the multiplexor to select the appropriate symbol when the actual value is determined. The complexity of the look-ahead implementation is  $M^L$ , where  $M$  is the size of symbol alphabet, and  $L$  is the number of coefficient taps in the decision-feedback equalization. The speed of the look-ahead method is limited by a select-and-latch operation in the transformed simple loop.

Block (parallel) processing techniques have also been proposed for speeding up adaptive filters for high-speed communication applications. For a discussion of block processing techniques, see, for example, G.A. Clark et al., "Block Implementation of Adaptive Digital Filters", Trans. on ASSP, Vol. ASSP-29, No. 3, (June 1981). Generally, block processing increases the throughput of the system by processing several inputs (made available through proper buffering) in one clock cycle using duplicated hardware. In return, the clock speed can be set at a lower and more feasible speed. For example, if ten inputs are processed in the same clock cycle, the processing clock speed can be lowered to ten percent (10%) of the clock speed of the received signal, while maintaining the same throughput.

While block processing techniques effectively reduce the required processing speed for general adaptive filter applications, block processing cannot easily be directly used with decision-feedback equalizers because of an output dependency in the feedback loop of decision-feedback equalizations. In the feedback loop of a decision-feedback equalizer, each output depends on previous output values that may not be available from the previous cycle. As apparent from the above-described deficiencies with conventional decision-feedback equalizers, a need exists for a technique that reduces the output dependency and further speeds up decision-feedback equalization processing. United States Patent No. 6,363,112, filed December 7, 1998, entitled "A Parallel Processing Decision-Feedback Equalizer," assigned to the assignee of the present invention (hereinafter referred to as "the Azadet System"), discloses a parallel implementation of a decision-feedback equalization that speeds up decision-feedback equalization processing. While the Azadet System employs look-ahead techniques in each of a plurality of parallel blocks to achieve performance gains, the output

dependencies in the disclosed parallel decision-feedback equalization nonetheless produce a delay on the order of  $N$ , where  $N$  is the number of blocks.

### **Summary of the Invention**

Generally, a method and apparatus are disclosed for increasing the effective  
5 processing speed of a parallel decision-feedback equalizer by combining block processing and look-ahead techniques in the selection (multiplexing) stage. The parallel decision-feedback equalizer of the Azadet System receives a plurality of symbol blocks in parallel using a plurality of corresponding input branches and utilizes look-ahead techniques within a given block to precompute all possible output values for the given block. The present invention  
10 extends the Azadet System by using look-ahead techniques in the selection stage to precompute the effect of previous blocks on each subsequent block, and to thereby remove the serial output dependency.

The present invention reduces the delay in the critical path for parallel decision-feedback equalizations by employing block processing and look-ahead techniques in  
15 the selection (multiplexing) stage to select the actual output values from among the generated possible values. According to one aspect of the invention, the parallel decision-feedback equalization includes a multiplexor tree structure that selects an appropriate output value for each block and precomputes the effect of previous blocks on each subsequent block. A multiplexing delay algorithm on the order of  $\log N$  is employed to resolve the output  
20 dependency and thus speeds up parallel block processing decision-feedback equalizations.

According to another aspect of the invention, the disclosed decision-feedback equalization architecture can be combined with pipelining to completely eliminate the critical path problem. The present invention permits a pipeline implementation of the disclosed multiplexor array circuit, because there are no dependencies from one row of the multiplexor  
25 array to another row and the select signal,  $y_0$ , is needed only at the bottom row of multiplexors. Pipeline latches can be added after any row of multiplexors in the multiplexor array circuit and each pipelined segment of the multiplexor array can be processed simultaneously. The number of pipeline segments that can be formed is between 2 and  $\log N$ , where  $N$  is the number of parallel blocks.

30 The disclosed multiplexor tree circuitry for the parallel decision-feedback equalizer groups multiplexor blocks into groups of two, referred to as block pairs, and

provides at least one multiplexor for each block,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values. In addition, block pairs are also progressively grouped into block groups, such that the first block group has one block pair, the second block group has two block pairs, and so on. The output of each parallel block depends on the possible precomputed values generated by the look-ahead processors for the block, as well as the actual values that are ultimately selected for each previous block. In order to reduce the delay in obtaining each actual output value, the present invention assumes that each block contains each possible value, and carries the assumption through to all subsequent blocks. Thus, the number of multiplexors required to select from among the possible values grows according to  $N \cdot \log N$ , where  $N$  is the block number.

For example, the first block is not influenced by previous blocks and requires only one multiplexor to select from among the possible precomputed values. The output of the second parallel block depends on the possible precomputed values generated by the look-ahead processors for the second block, as well as the actual values that are ultimately selected for the first block. Thus, the second block includes a multiplexor for processing each of the assumed possible values of the first block (selection signal), with each multiplexor receiving the possible precomputed values generated by the look-ahead processors for the second block as inputs. In addition, the second block includes a final multiplexor for selecting the final output value,  $y_2$ . The inputs to the final multiplexor for the second block are also applied as inputs to multiplexors in the second block group (blocks three and four). The second block group is processed in a similar manner, and the inputs to the final multiplexor for the fourth block are also applied as inputs to multiplexors in the next block group comprised of the third and fourth block pair (blocks five through eight).

A more complete understanding of the present invention, as well as further features and advantages of the present invention, will be obtained by reference to the following detailed description and drawings.

### **Brief Description of the Drawings**

FIG. 1 illustrates a conventional decision-feedback equalizer structure;

FIG. 2A illustrates two neighboring multiplexors  $m_1$  and  $m_2$ , where the output of  $m_1$  is connected to the select input of  $m_2$ ;

FIG. 2B illustrates the transformation of the multiplexor m2 of FIG. 2A in accordance with the present invention, such that all possible values are precomputed and applied to the multiplexor m2' for selection by the select signal of multiplexor m1;

FIG. 3 illustrates a conventional block processing implementation of a decision-feedback equalizer, where the number of parallel blocks is eight;

FIG. 4 illustrates a multiplexor tree array in accordance with the present invention for the parallel decision-feedback equalizer of FIG. 3 to produce a delay on the order of  $\log N$ ; and

FIG. 5 illustrates an alternate implementation of the multiplexor tree array of FIG. 4 in accordance with the present invention having a delay on the order of  $\log N$  and reduced complexity.

### **Detailed Description**

The present invention speeds up decision-feedback equalization processing by combining block processing and look-ahead techniques in the selection stage to produce a delay on the order of  $\log N$ . A block processing decision-feedback equalization with a block factor of  $N$  takes  $N$  inputs and produces  $N$  outputs in parallel. Since each output computation depends on previous decisions, the computation of the  $N$  outputs forms a dependency chain in block processing decision-feedback equalizations and generally requires a multiplexing delay time of  $N-1$  when block processing is done sequentially. The present invention employs a multiplexing delay algorithm on the order of  $\log N$  to resolve the output dependency and thus speeds up parallel block processing decision-feedback equalizations.

The output of each parallel block depends on the possible precomputed values generated by the look-ahead processors for the block, as well as the actual values that are ultimately selected for each previous block. In order to reduce the delay in obtaining each actual output value, the present invention assumes that each block contains each possible value, and carries the assumption through to all subsequent blocks. Thus, the number of multiplexors required in each block to select from among the possible values grows according to  $N \cdot \log N$ , where  $N$  is the block number. A novel multiplexor tree architecture selects the actual output values with a significantly reduced delay. As soon as the possible values are computed, the output is selected through multiplexing.

The disclosed decision-feedback equalizer architecture groups multiplexor blocks into groups of two, and provides one or more multiplexors for each possible precomputed value to ultimately compute the output values,  $y$ , independently and concurrently. The decision-feedback equalization is comprised of a tree of multiplexors, and a transform operation, discussed below, is applied to non-overlapping neighboring multiplexor pairs, in parallel. A transformed pair is treated as a single multiplexor with multiple outputs and is again grouped into neighboring pairs. This transform and group-in-pairs operation can then be repeated, and the select signals closest to the beginning of the dependency chain will double its direct control distance each time. By properly grouping the original  $N$  multiplexors in a tree fashion and repeatedly using the “transform” operation, the entire dependency chain can be computed in  $D$  plus one steps, where  $D$  is the tree depth and is equal to  $\log N$  (because the select signal propagates a distance of  $2^{\log N}$  equal to  $N$ ).

### DECISION-FEEDBACK EQUALIZERS

FIG. 1 illustrates a conventional decision-feedback equalization structure 100.

The decision-feedback equalization filter output is  $z[n] = \sum_{k=1}^L w_{-k} y[n-k]$ , where  $L$  is the length of the filter and  $w_{-k}$  is the  $k$ th tap weight. Thus, the output of the slicer 110 can be expressed as:

$$y[n] = Q\left(x[n] + \sum_{k=1}^L w_{-k} y[n-k]\right), \quad (1)$$

where the function,  $Q$ , is a nonlinear function employed by the slicer 110, mapping signals to symbol alphabets  $\{\alpha_i\}$  for  $0 \leq i < M$ . Equation 1 leads to a straightforward implementation, namely, at clock cycle  $n$ , when  $x\{n\}$  is available and all previous  $y\{n-k\}$ , for all  $k > 0$  are already computed and known, the equation can be evaluated during the given clock period.

## MULTIPLEXOR TERMINOLOGY

Generally, a multiplexor function is represented as  $m(i_0, i_1, \dots, i_{N-1}; s_0, s_1, \dots, s_{n-1})$ , where  $s_k$  indicates the select signals of the multiplexor,  $i_k$  indicates the data inputs, and  $N=2^n$ . If the unsigned binary number  $s_{n-1}s_{n-2}\dots s_0$  represents a number  $j$ , then the output of the multiplexor function,  $m(i_0, i_1, \dots, i_{N-1}; s_0, s_1, \dots, s_{n-1})$ , is  $i_j$ . The two neighboring multiplexors,  $m1$  and  $m2$ , shown in FIG. 2A, can be represented as  $m1(i_0^1, i_1^1, \dots, i_{N-1}^1; s_0^1, s_1^1, \dots, s_{n-1}^1)$  and  $m2(i_0^2, i_1^2, \dots, i_{N-1}^2; s_0^2, s_1^2, \dots, s_{n-1}^2)$ , where  $s_i^2 = s_{i+1}^1$ , for  $0 \leq i \leq N-2$  and  $s_{n-1}^2 = m1()$ .

FIG. 2A illustrates two neighboring multiplexors  $m1$  and  $m2$ , where the output of  $m1$  is connected to the select input of  $m2$ . According to a feature of the present invention, the structure of  $m2$  is transformed by creating multiple copies of  $m2$ , with one copy for each of the possible output values of  $m1$ . The correct output of  $m2$  is then selected using the select signal of  $m1$ . As shown in FIG. 2B, the transform operation changes the multiplexor function of  $m2$  to  $m2'$  ( $m2(i_0^2, i_1^2, \dots, i_{N-1}^2; s_1^2, \dots, s_{n-2}^2, i_0^1)$ ,  $m2(i_0^2, i_1^2, \dots, i_{N-1}^2; s_1^2, \dots, s_{n-2}^2, i_1^1)$ , ...,  $m2(i_0^2, i_1^2, \dots, i_{N-1}^2; s_1^2, \dots, s_{n-2}^2, i_{n-1}^1)$ ;  $s_0^1, s_1^1, \dots, s_{n-1}^1$ ). The multiplexor  $m2'$  is consistent with original functionality of the multiplexor  $m2$ . The transformation from  $m2$  to  $m2'$  can be viewed as propagating the control/select signals of  $m1$  to  $m2$  through one multiplexor delay. In addition, the transformation from  $m2$  to  $m2'$  can also be viewed as doubling the distance under the direct control of the select signals of  $m1$ . In this manner, all the possible values are precomputed and applied to the multiplexor  $m2'$  and the appropriate value is selected (as opposed to computed) using the select signal of multiplexor  $m1$ .

As discussed further below in conjunction with FIG. 4, the transform operation can be applied to non-overlapping neighboring multiplexor pairs in parallel. A transformed pair can then be treated as a single multiplexor with multiple outputs and again be grouped into neighboring pairs. This transform and group-in-pairs operation can then be repeated, and the select signals closest to the beginning of the dependency chain will double its direct control distance each time. By properly grouping the original  $N$  multiplexors in a tree fashion and repeatedly using the "transform" operation discussed above in conjunction with FIGS. 2A and 2B, the entire dependency chain can be computed in  $D$  steps, where  $D$  is the tree depth and is equal to  $\log N$  (because the select signal propagates a distance of  $2^{\log N} = N$ ).

FIG. 3 provides a conceptual representation for a block processing implementation of a decision-feedback equalizer 300, such as the implementation of the Azadet System, where the number of parallel blocks 311-318 is eight. Thus, if the clock rate

of the received signal is  $C$ , the processing clock rate of the decision-feedback equalizer 300 can be  $C/8$ . The illustrative decision-feedback equalization 300 shown in FIG. 3 is an implementation for  $k=1$  tap, with  $M=2$  possible values (levels) for each symbol or bit.

As shown in FIG. 3, each parallel block 311-318 of the decision-feedback equalizer 300 includes a single tap decision-feedback equalizer, such as the decision-feedback equalizations 321-a, 321-b for the first block 311, for precomputing each of the possible output values,  $y_i$ . Thereafter, the possible precomputed values are applied to the corresponding  $i_N$  data inputs of a multiplexor 331-338 for each block 311-318. In the illustrative implementation, a two-level (binary) signaling scheme is employed. Thus, as shown in FIG. 3, the two possible values from the look-ahead decision-feedback equalizers at each block, such as the decision-feedback equalizations 321-a, 321-b for the first block 310, corresponding to the two possible values (0/1) of  $y_i$ , are applied to the corresponding  $i_0$  and  $i_1$  inputs of each multiplexor 311-318. Once the actual value,  $y_i$ , of a given block 311-318 is determined, the actual value is applied to the select signal of the multiplexor  $i + 1$  (for the next block), to select the appropriate next symbol or bit for the next block  $y_{i+1}$ .

FIG. 4 illustrates a multiplexor array circuit 400 in accordance with the present invention that utilizes block processing and look-ahead techniques in the selection (multiplexing) stage to produce a delay on the order of  $\log N$ . The illustrative multiplexor array circuit 400 shown in FIG. 4 is an implementation for  $k=1$  tap, with  $M=2$  possible values (levels) for each symbol or bit. The multiplexor array circuit 400 of FIG. 4 groups the multiplexors 331-338 of FIG. 3 into groups of two, and provides an array of multiplexors, in a manner described further below, for selecting the appropriate precomputed value for each block pair independently and concurrently. Thus, the  $i_0^1$  and  $i_1^1$  inputs of block 311 are grouped with the  $i_0^2$  and  $i_1^2$  inputs of block 312, as shown in FIG. 4. Likewise, the  $i_0^3$  and  $i_1^3$  inputs of block 313 are grouped with the  $i_0^4$  and  $i_1^4$  inputs of block 314 ( $i_0^5$  and  $i_1^5$  inputs are grouped with  $i_0^6$  and  $i_1^6$  inputs, and  $i_0^7$  and  $i_1^7$  inputs are grouped with  $i_0^8$  and  $i_1^8$  inputs in a similar manner).

As shown in FIG. 4, the transform operation of the present invention is applied to non-overlapping neighboring multiplexor pairs, in parallel. A transformed pair is treated as a single multiplexor with multiple outputs and is again grouped into neighboring pairs. This transform and group-in-pairs operation can then be repeated, and the select signals closest to the beginning of the dependency chain, such as  $y_0$  in FIG. 4, will double its direct control

distance each time. By properly grouping the original  $N$  multiplexors 311-318 in a tree fashion and repeatedly using the "transform" operation discussed above in conjunction with FIGS. 2A and 2B, the entire dependency chain can be computed in  $D$  steps, where  $D$  is the tree depth and is equal to  $\log N$  (because the select signal propagates a distance of  $2^{\log N}$  equal to  $N$ ).

The output of block 311 (FIG. 3) is determined by the two possible values  $i_0^1$  and  $i_1^1$  and the select signal,  $y_0^0$ . The output of each subsequent block 312-318 (FIG. 3) is determined by the two possible input values  $i_0^j$  and  $i_1^j$  and the output of the previous block 311-317. The present invention utilizes look-ahead techniques in the selection stage to produce a delay on the order of  $\log N$ .

In the following discussion of FIG. 4, it is assumed that all multiplexing operations take an equal amount of time. At a time,  $t$ , equal to 0, when all  $i_0$ 's and  $i_1$ 's are available it is not known whether  $i_0^k$  or  $i_1^k$  will be chosen, for  $k=1$  through 8. The present invention, however, utilizes the fact that either  $i_0^k$  or  $i_1^k$  will be the final correct value. Initially, for each independent block pair, the correct value is assumed to be  $i_1^k$  ( $i_1^1, i_1^3, i_1^5, i_1^7$ ), and the  $i_1^k$  value is applied to the corresponding multiplexors 411, 413, 415, 417 in the first row of each block pair, as shown in FIG. 4. The  $i_1^k$  value selects  $i_0^2$  or  $i_1^2$  for multiplexor 411,  $i_0^4$  or  $i_1^4$  for multiplexor 413,  $i_0^6$  or  $i_1^6$  for multiplexor 415, and  $i_0^8$  or  $i_1^8$  for multiplexor 417.

In addition, for each independent block pair, the alternate correct value (in the two-level illustration) is also assumed to be  $i_0^k$  ( $i_0^1, i_0^3, i_0^5, i_0^7$ ), and the  $i_0^k$  value is applied to the corresponding multiplexors 412, 414, 416, 418 in the first row of each block pair, as shown in FIG. 4. The  $i_0^k$  value selects  $i_0^2$  or  $i_1^2$  for multiplexor 412,  $i_0^4$  or  $i_1^4$  for multiplexor 414,  $i_0^6$  or  $i_1^6$  for multiplexor 416, and  $i_0^8$  or  $i_1^8$  for multiplexor 418.

Therefore, at a time,  $t$ , equal to one multiplexor delay time, for each block pair (i) the output of the multiplexors 411, 413, 415, 417 contain the correct output value if the corresponding  $i_1^1, i_1^3, i_1^5, i_1^7$  value is the correct value, or (ii) the output of the multiplexors 412, 414, 416, 418 contain the correct output value if the corresponding  $i_0^1, i_0^3, i_0^5, i_0^7$  value is the correct value. The outputs of the multiplexors 411, 412 represent the possible values of output  $y_2$  of the block 312 (FIG. 3). It is also noted that even though which assumption is correct remains unknown, the choice of which multiplexor for each block pair in the first row contains the correct value no longer depends on the output value of the preceding block  $y_1$ ,

$y_3, y_5, y_7$ . The duplicated multiplexing for each possible precomputed value provided by the present invention allows the choice of which multiplexor for each block pair in the first row containing the correct value to depend solely on  $y_0$ . In other words, the distance (selections) under the direct control of  $y_0$  is doubled from 1 to 2 blocks (for outputs  $y_1$  and  $y_2$ ).

5                    Similarly, in processing the second row of multiplexors 421-428 in FIG. 4, which of multiplexors 411 or 412 for the first block pair or multiplexors 415 or 416 for the third block pair contains the final correct value remains unknown. Initially, it is assumed that multiplexors 411 and 415 contain the correct value and the corresponding value is applied to multiplexors 421 and 423 in the second block pair and multiplexors 425 and 427 in the fourth  
10 block pair. At the same time, it is assumed that multiplexors 412 and 416 contain the correct value and the corresponding value is applied to multiplexors 422 and 424 in the second block pair and multiplexors 426 and 428 in the fourth block pair.

                  Therefore, at a time,  $t$ , equal to two multiplexor delay times, if the multiplexor 411 contains the correct value (which in turn means the  $i_1^1$  value is the correct value), the  
15 output of the multiplexor 411 selects the outputs of the multiplexors 421 and 423 as the  $y_3$  and  $y_4$  outputs. Likewise, if the multiplexor 412 contains the correct value (which in turn means the  $i_0^1$  value is the correct value), the output of the multiplexor 412 selects the outputs of the multiplexors 422 and 424 as the  $y_3$  and  $y_4$  outputs. The output of multiplexor 421 and  
20 multiplexor 423 and 424 represent the possible values of output  $y_3$  of the block 313 (FIG. 3). The output of multiplexor 422 and 424 represent the possible values of output  $y_4$  of the block 314 (FIG. 3). The choices for  $y_1, y_2, y_3, y_4$  at this point are solely determined by  $y_0$ . In other words, the distance under the direct control of  $y_0$  is doubled to 4 blocks. Similarly, at a time,  $t$ , equal to three multiplexor delay times, all eight output values are directly controlled (selected) by  $y_0$ , with the selection being performed by the last row of multiplexors.

25                    FIG. 5 illustrates a decision-feedback equalizer 500 in accordance with the present invention having a delay on the order of  $\log N$ . The multiplexor array circuit 500 of FIG. 5 is a simplified version of the multiplexor array circuit 400 of FIG. 4. The first multiplexor 511 of the multiplexor array circuit 500 selects the correct value for the first stage. The simplified architecture is achieved by a utilizing a multiplexor 511 at the first  
30 stage, controlled by the select signal  $y_0^0$ .

The complexity of the multiplexor array circuit 500 (FIG. 5) is  $\log N * N - N + 2$  multiplexors, compared with a complexity of  $N * \log N + N$  for the multiplexor array circuit 400 of FIG. 4. Savings are more significant for small N.

The architecture of the multiplexor array circuit 400 of FIG. 4 lends itself to a pipeline implementation, because the select signals,  $y_0$ , are needed only at the bottom row of multiplexors 450. In addition to a traditional pipeline implementation of the speculative finite impulse response (FIR) filters, pipeline latches can be added after any row of multiplexors in the multiplexor array circuit 400 of FIG. 4. Specifically, since there are no dependencies from one row of the multiplexor array 400 to another row, pipeline latches can be added after any row of multiplexors in the multiplexor array circuit 400 of FIG. 4. Thus, each segment of the multiplexor array 400 can be processed simultaneously. The number of pipeline segments that can be formed is between 2 and  $\log N$ . In one implementation, the multiplexor array 400 is divided into two segments after the second row of multiplexors, utilizing pipelining techniques. When  $\log N$  segments are used, the architecture has the most relaxed critical path timing constraint, with just a single multiplexor delay plus one register latch delay. It is noted that while the critical path seems to be equal to that of a traditional look-ahead implementation, such as those described in K. Parhi, "Pipelining in Algorithm with Quantizer Loops". *IEEE Transactions on Circuits and Systems*, Vol. 38, No. 7, 745-54 (July 1991), the clock is operated at a reduced frequency of  $f/N$ . The decision-feedback equalization architecture shown in FIGS. 4 and 5 thus allows decision-feedback equalizations to operate at a very high speed. In theory, the architecture of the present invention removes the limit on how fast a decision-feedback equalization can operate.

It is to be understood that the embodiments and variations shown and described herein are merely illustrative of the principles of this invention and that various modifications may be implemented by those skilled in the art without departing from the scope and spirit of the invention. For example, while the invention has been illustrated with a binary implementation using 2-level signals, the present invention can be easily generalized to multi-level signals.

**Claims:**

1. A method for equalizing a signal received from a dispersive channel, said method comprising the steps of:

receiving a plurality of symbol blocks in parallel;

5 computing all possible values of a decision feedback signal for each of said corresponding blocks;

providing a multiplexor array for processing said possible values for each of said blocks and the possible values for any previous blocks on each subsequent block; and

selecting an output value for each of said blocks using said multiplexor array.

10 2. The method of claim 1, wherein said multiplexor array selects an appropriate output value for each block and processes the effect of previous blocks on each subsequent block.

3. The method of claim 1, wherein pairs of multiplexors in said multiplexor array are grouped into block pairs.

15 4. The method of claim 3, wherein each of said block pairs provides at least one multiplexor for each of said blocks,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values.

5. The method of claim 1, wherein said multiplexor array is produced by (i) transforming non-overlapping neighboring multiplexor pairs in said multiplexor array, in  
20 parallel, (ii) treating each transformed multiplexor pair as a single multiplexor with multiple outputs, and (iii) again grouping said effective single multiplexors into neighboring pairs.

6. The method of claim 1, wherein said multiplexor array assumes that each block contains each possible value, and carries said assumed possible values through to all subsequent blocks.

25 7. The method of claim 1, wherein the number of multiplexors in each block of said multiplexor array grows according to  $N \cdot \log N$ , where  $N$  is the block number.

8. The method of claim 1, further comprising the step of utilizing pipeline techniques to simultaneously process various segments of said multiplexor array.

9. A decision-feedback equalizer for equalizing a signal received on a plurality of parallel blocks from a dispersive channel, comprising:

at least one look-ahead processor for computing possible output values of a decision feedback signal for each of said corresponding blocks; and

5 a multiplexor array for selecting an output value for each of said blocks from said possible values for each of said blocks and wherein said multiplexor array employs look-ahead techniques to precompute the effect of previous blocks on each subsequent block.

10 10. The decision-feedback equalizer of claim 9, wherein said multiplexor array selects an appropriate output value for each block and processes the effect of previous blocks on each subsequent block.

11. The decision-feedback equalizer of claim 9, wherein pairs of multiplexors in said multiplexor array are grouped into block pairs.

15 12. The decision-feedback equalizer of claim 11, wherein each of said block pairs provides at least one multiplexor for each of said blocks,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values.

20 13. The decision-feedback equalizer of claim 9, wherein said multiplexor array is produced by (i) transforming non-overlapping neighboring multiplexor pairs in said multiplexor array, in parallel, (ii) treating each transformed multiplexor pair as a single multiplexor with multiple outputs, and (iii) again grouping said effective single multiplexors into neighboring pairs.

14. The decision-feedback equalizer of claim 9, wherein said multiplexor array assumes that each block contains each possible value, and carries said assumed possible values through to all subsequent blocks.

25 15. The decision-feedback equalizer of claim 9, wherein the number of multiplexors in each block of said multiplexor array grows according to  $N \cdot \log N$ , where  $N$  is the block number.

16. The decision-feedback equalizer of claim 9, wherein pipeline techniques are employed to simultaneously process various segments of said multiplexor array.

17. The decision-feedback equalizer of claim 16, wherein pipeline latches are provided after at least one row of said multiplexor array and each pipelined segment of the multiplexor array can be processed simultaneously.

18. A decision-feedback equalizer for equalizing a signal received on a plurality of parallel blocks from a dispersive channel, comprising:

at least one look-ahead processor for computing possible output values of a decision feedback signal for each of said corresponding blocks; and

a multiplexor array for selecting an output value for each of said blocks, wherein said multiplexor array employs look-ahead techniques to precompute the effect of previous blocks on each subsequent block.

19. The decision-feedback equalizer of claim 18, wherein said multiplexor array selects an appropriate output value for each block and processes the effect of previous blocks on each subsequent block.

20. The decision-feedback equalizer of claim 18, wherein pairs of multiplexors in said multiplexor array are grouped into block pairs.

21. The decision-feedback equalizer of claim 20, wherein each of said block pairs provides at least one multiplexor for each of said blocks,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values.

22. The decision-feedback equalizer of claim 18, wherein said multiplexor array is produced by (i) transforming non-overlapping neighboring multiplexor pairs in said multiplexor array, in parallel, (ii) treating each transformed multiplexor pair as a single multiplexor with multiple outputs, and (iii) again grouping said effective single multiplexors into neighboring pairs.

23. The decision-feedback equalizer of claim 18, wherein said multiplexor array assumes that each block contains each possible value, and carries said assumed possible values through to all subsequent blocks.

24. The decision-feedback equalizer of claim 18, wherein the number of multiplexors in each block of said multiplexor array grows according to  $N \cdot \log N$ , where  $N$  is the block number.

25. The decision-feedback equalizer of claim 18, wherein pipeline techniques are employed to simultaneously process various segments of said multiplexor array.

26. The decision-feedback equalizer of claim 25, wherein pipeline latches are provided after at least one row of said multiplexor array and each pipelined segment of the multiplexor array can be processed simultaneously.

27. A method for equalizing a signal received from a dispersive channel, said method comprising the steps of:

receiving a plurality of symbol blocks in parallel;

10 computing all possible values of a decision feedback signal for each of said corresponding blocks; and

selecting an output value for each of said blocks using look-ahead techniques to precompute the effect of previous blocks on each subsequent block.

28. The method of claim 27, wherein said selecting step selects an appropriate output value for each block and processes the effect of previous blocks on each subsequent block.

29. The method of claim 27, wherein said selecting step is performed using a multiplexor array.

30. The method of claim 29, further comprising the step of providing at least one multiplexor for each of said blocks,  $i$ , to select an output value,  $y_i$ , from among the possible precomputed values.

31. The method of claim 27, further comprising the step of performing said selecting step using a multiplexor array produced by (i) transforming non-overlapping neighboring multiplexor pairs in said multiplexor array, in parallel, (ii) treating each transformed multiplexor pair as a single multiplexor with multiple outputs, and (iii) again grouping said effective single multiplexors into neighboring pairs.

32. The method of claim 27, wherein said selecting step assumes that each block contains each possible value, and carries said assumed possible values through to all subsequent blocks.

33. The method of claim 27, further comprising the step of utilizing pipeline techniques to simultaneously process various segments of said multiplexor array.

34. A decision-feedback equalizer for equalizing a signal received on a plurality of parallel blocks from a dispersive channel, comprising:

means for receiving a plurality of symbol blocks in parallel;

5 means for computing all possible values of a decision feedback signal for each of said corresponding blocks;

means for providing a multiplexor array for processing said possible values for each of said blocks and wherein said multiplexor array employs look-ahead techniques to precompute the effect of previous blocks on each subsequent block; and

means for selecting an output value for each of said blocks using said multiplexor array.

35. A decision-feedback equalizer for equalizing a signal received on a plurality of parallel blocks from a dispersive channel, comprising:

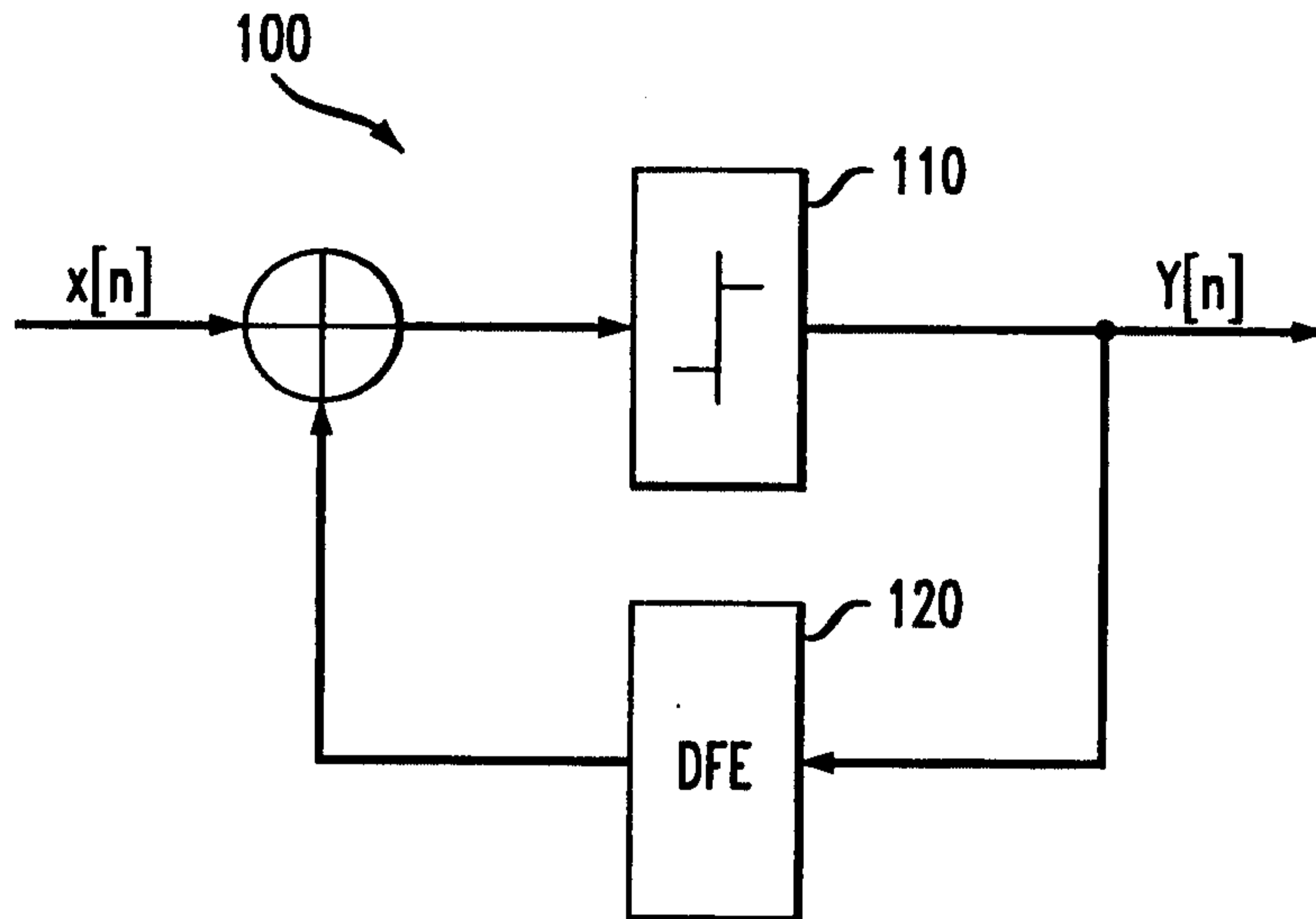
means for receiving a plurality of symbol blocks in parallel;

10 means for computing all possible values of a decision feedback signal for each of said corresponding blocks; and

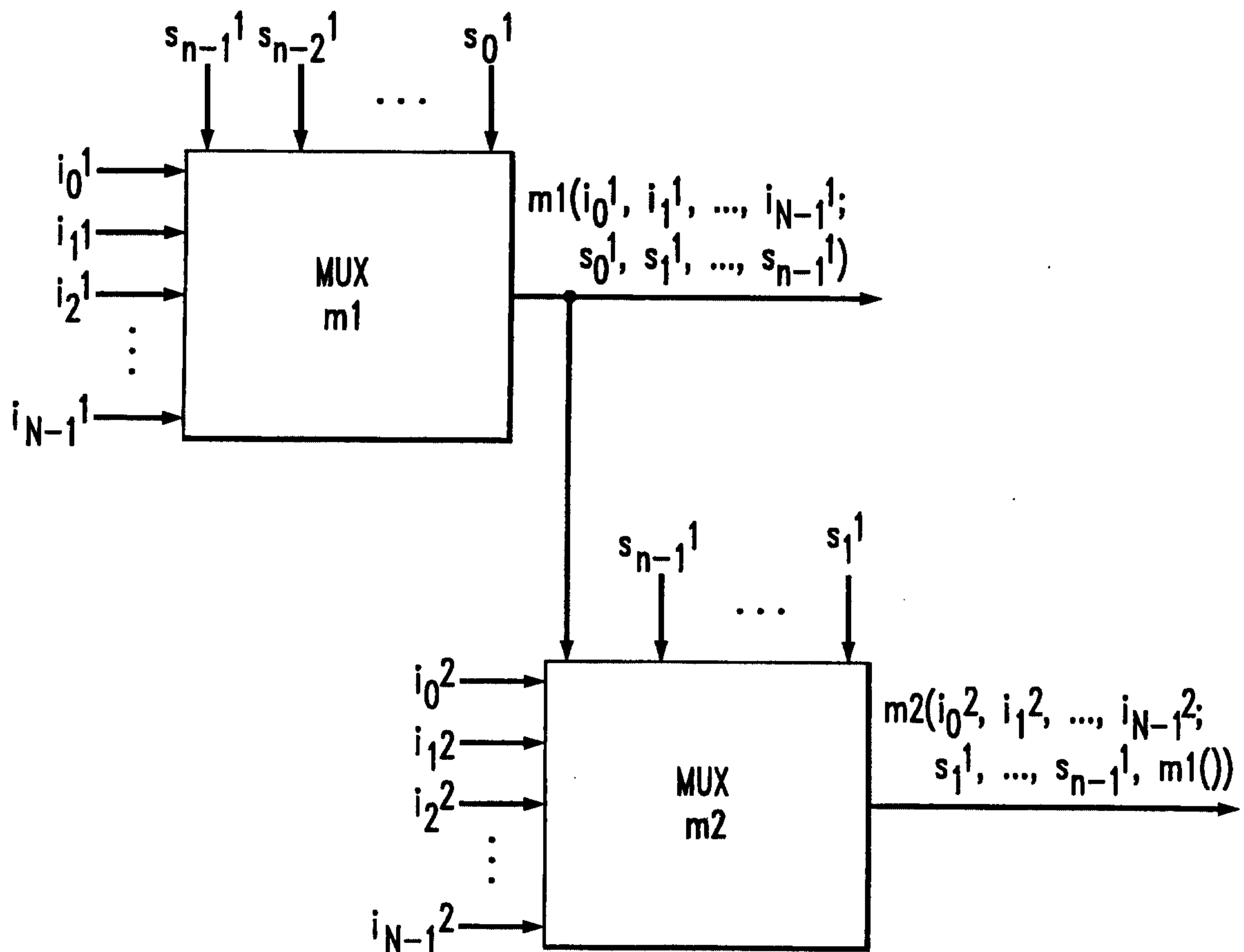
means for selecting an output value for each of said blocks using look-ahead techniques to precompute the effect of previous blocks on each subsequent block.

1/5

**FIG. 1**  
PRIOR ART



**FIG. 2A**



2/5

FIG. 2B

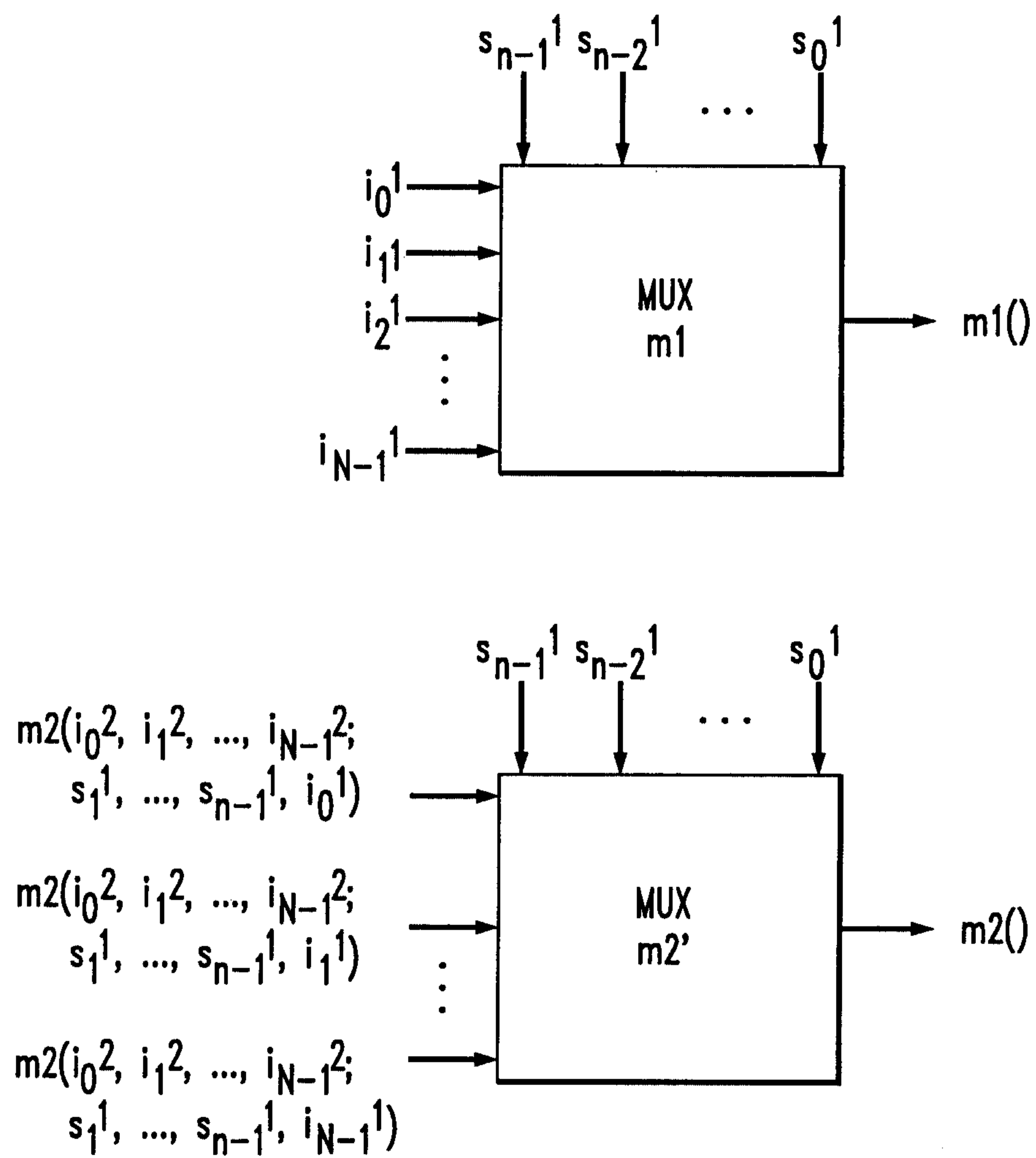
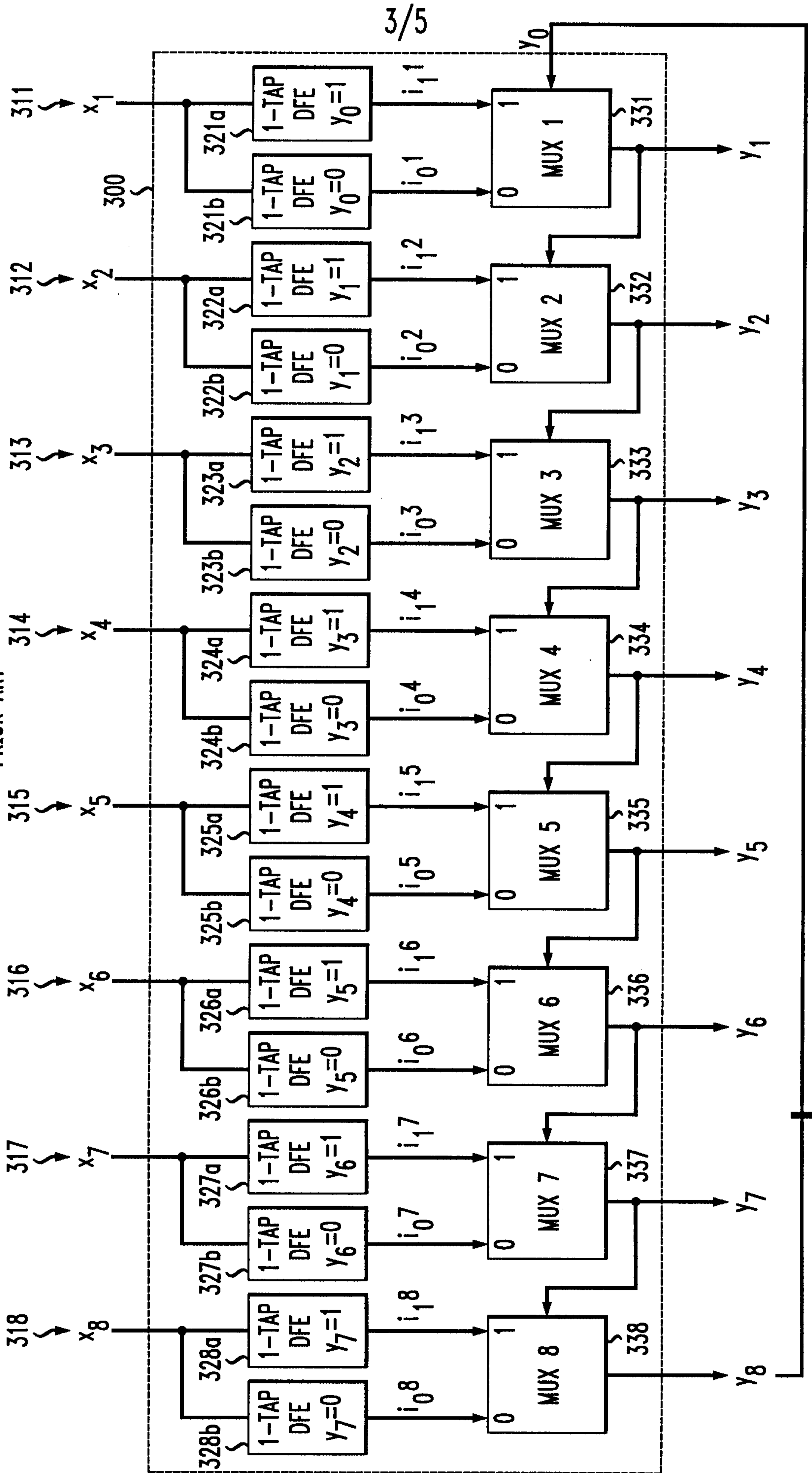


FIG. 3

PRIOR ART



4/5

FIG. 4  
400

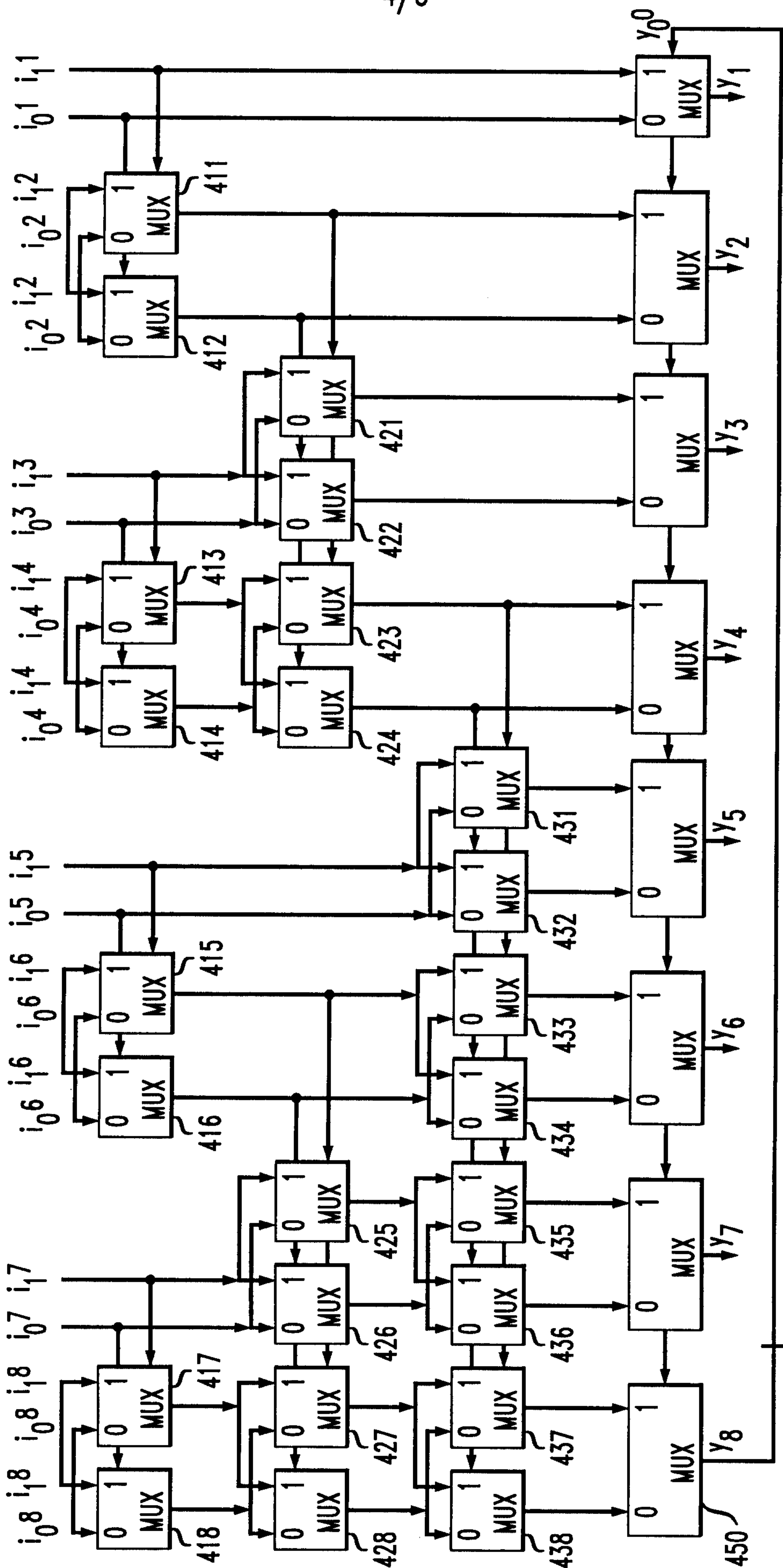


FIG. 5

