

(19) World Intellectual Property Organization
International Bureau



(43) International Publication Date
11 December 2003 (11.12.2003)

PCT

(10) International Publication Number
WO 03/102778 A2

(51) International Patent Classification⁷: **G06F 12/00**

(21) International Application Number: PCT/EP03/03529

(22) International Filing Date: 4 April 2003 (04.04.2003)

(25) Filing Language: English

(26) Publication Language: English

(30) Priority Data:
02012080.4 31 May 2002 (31.05.2002) EP

(71) Applicant (for all designated States except LU, US):
INTERNATIONAL BUSINESS MACHINES CORPORATION [US/US]; New Orchard Road, Armonk, NY 10504 (US).

(71) Applicant (for LU only): **IBM DEUTSCHLAND GMBH** [DE/DE]; Pascalstrasse 100, 70569 Stuttgart (DE).

(72) Inventors; and

(75) Inventors/Applicants (for US only): **HANSMANN, Uwe** [DE/DE]; Birkenstrasse 30/1, 71155 Altdorf (DE).

HEPPER, Stefan [DE/DE]; Dorfackerstrasse 24, 72074 Tuebingen (DE). **MERK, Lothar** [DE/DE]; Kirchwiesenstrasse 5, 71093 Weil im Schönbuch (DE). **STOBER, Thomas** [DE/DE]; Schubartweg 8, 71032 Boeblingen (DE).

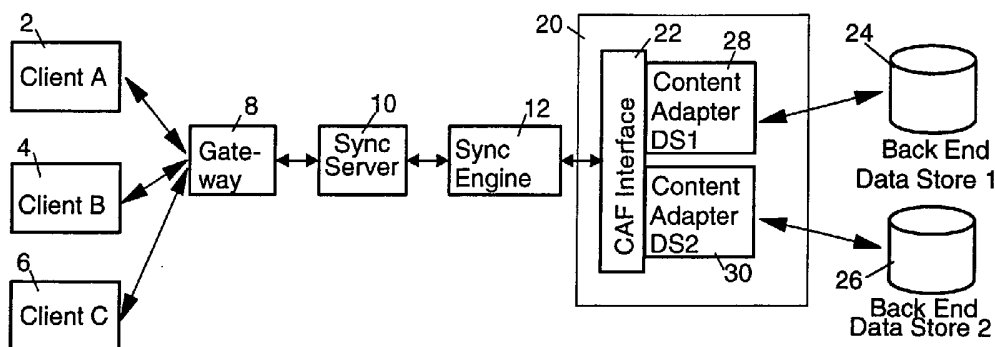
(74) Agent: **KLEIN, Hans-Jörg**; IBM Deutschland GmbH, Intellectual Property, 70548 Stuttgart (DE).

(81) Designated States (*national*): AE, AG, AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, BZ, CA, CH, CN, CO, CR, CU, CZ, DE, DK, DM, DZ, EC, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LV, MA, MD, MG, MK, MN, MW, MX, MZ, NI, NO, NZ, OM, PH, PL, PT, RO, RU, SC, SD, SE, SG, SK, SL, TJ, TM, TN, TR, TT, TZ, UA, UG, US, UZ, VC, VN, YU, ZA, ZM, ZW.

(84) Designated States (*regional*): ARIPO patent (GH, GM, KE, LS, MW, MZ, SD, SL, SZ, TZ, UG, ZM, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, BG, CH, CY, CZ, DE, DK, EE, ES, FI, FR, GB, GR, HU, IE, IT, LU, MC, NL, PT, RO, SE, SI, SK, TR), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, ML, MR, NE, SN, TD, TG).

[Continued on next page]

(54) Title: SYSTEM AND METHOD FOR ACCESSING DIFFERENT TYPES OF BACK END DATA STORES



(57) Abstract: The present discloses a framework that allows a synchronization engine to synchronize data between a mobile device and Back End data stores independently from the architecture and data formats of that Back End data store. The framework introduces content adapters, which access synchronization data from backend data systems. These adapters convert the data into a Back End data store independent representation, which can be used by all applications or modules which need to access different back-ends in a generic manner. A generic synchronization engine for the purpose of conflict detection and resolution is one example for a module of this kind. Other applications that could use the content adapter are Notification Frameworks or Portals and all other applications aggregating data.



Published:

— without international search report and to be republished
upon receipt of that report

For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

D E S C R I P T I O N

System and method for accessing different types of Back End data stores

The present invention relates to a method and system for exchange or synchronization of data between different clients, and in particular to synchronization of data between clients by using a central synchronization server linked with different types of Backend data store.

Synchronization can be defined as keeping data consistent between different clients, e.g. a Notebook calendar and Personal Digital Assistant (PDA) calendar, and storing at these clients data into different data formats with different identifiers. The synchronization data can also consist of complete computer programs or parts of these programs.

Frequently, a central synchronization server is used through which data can be exchanged between different clients (see FIG.1). The central synchronization server can either store all data locally (i.e. when using a proprietary data format), which is normally used by carriers (e.g. Yahoo) with high loads, or can directly access Backend data stores like DB2, Domino, or Exchange (see FIG.2).

Back End data stores have their own protocols and mechanisms to access and store information. Although relational databases use ODBC as a common interface, for instance databases storing Personal Information Management (PIM) data are usually accessed in a proprietary way and rely on very specific data structures of the content. This results in a dependency of Web-Server based applications on the specific databases and

- 2 -

the particular types of the accessed content. It requires big efforts when adopting new databases and new types of content for an existing software.

For instance, developing a big scale synchronization solution, which involves many different database manufacturers and multiple multimedia content is quite difficult under these circumstances: The synchronization engine, which includes the logic for synchronizing multiple client devices, needs to be adapted for each supported type of content and each connected database. In fact, today's synchronization engines depend strongly on the backend store, which contains the data. Significant investments for implementing a sophisticated sync engine is for the benefit of one or few backend systems only (see FIG.3)

The same problem appears if a Notification System wants to inform a user by sending data which is collected and stored in different backend systems. Also portals which aggregate data from different systems need a system independent method to access the information.

US Patent 5,974,238 describes an apparatus for performing dynamic synchronization between data stored in a handheld computer and a host computer, each having a plurality of data sets including at least one common data set, each computer having a copy of the common data set. The handheld computer has a processor, a communication port, and a data synchronization engine. The data synchronization engine has a pseudo-cache and one or more tags connected to the pseudo cache. Data is synchronized whenever data is written to main memory and/or when the associated pseudo-cache tag is invalidated. By strict adherence to a set of protocols, data

- 3 -

coherency is achieved because the system always knows who owns the data, who has a copy of the data, and who has modified the data. The data synchronization engine resolves any differences in the copies and allows the storage of identical copies of the common data set in the host computer and in the handheld computer.

This prior art patent is not directed to a synchronization architecture using a Back End data store. Therefore, the above mentioned problems related to the Back End data store are neither dealt nor solved by that patent.

It is therefore object of the present invention to provide a new method and system for exchange or synchronization of data in an architecture using a central synchronization server linked to different Back End data store types however avoiding the disadvantages of the prior art solutions.

That object is solved by the features of the independent claims. Further preferred embodiments are laid down in the dependent claims.

The present invention discloses a framework that allows a synchronization engine to synchronize data between a mobile device and Back End data stores independently from the architecture and data formats of that back-end store. The framework introduces content adapters, which access synchronization data from Back End data stores. These adapters convert the data into a Back End data store independent representation, which can be used by all applications or modules which need to access different back-ends in a generic manner. A generic synchronization engine for the purpose of conflict detection and resolution is one example for a module

- 4 -

of this kind. Other applications that could use the content adapter are Notification Frameworks or Portals and all other applications aggregating data. Any Back End data store specific issues are handled by the Back End dependent part of the content adapters, which can easily be provided by third parties and plugged into the framework.

The present invention will be described in more detail with the accompanying drawings in which:

FIG.1 shows a simplified synchronization architecture on which the present invention may be based,

FIG.2 shows the prior art synchronization architecture with direct access to the Back End data store,

FIG.3 shows the prior art synchronization architecture with different Back End data stores,

FIG.4 shows the SyncML communication protocol which may be preferably implemented by the present invention,

FIG.5 shows the basic architecture of the inventive content adapter framework (CAF),

FIG.6 shows a preferred implementation of the CAF,

FIG.7 shows the CAF interfaces,

FIG.8 shows a communication flow between client, sync engine, CAF and backend-system with session authentication,

FIG.9 shows the inheritance model as used by the CAF, and

- 5 -

FIG.10 shows the CAF specific process flow.

Synchronization between different clients using a central synchronization server is based on a synchronization protocol which typically consist of the following steps:

Pre-Synchronization: To prepare the actual synchronization some action must be taken before this can happen. These actions fall into the following groups: authentication, authorization, and determine client capabilities. Authentication ensures that the server is who it claims to be, and that the client is, who it claims to be. Authorization checks, whether the client is allowed to perform the requested action (e.g. delete, update, or only create new entries). Finally the server determines the device capabilities (e.g. maximum buffer size) to optimize the data flow to the client.

Synchronization: This is the part, where the synchronization data is exchanged. Between two synchronization partners all local IDs of data entries are mapped to global IDs known to both partners. Every partner therefore has a mapping table, to map local to global IDs. Then only the updated, new or deleted entries are exchanged. If both partners update the same data entry there will be a conflict. This update conflict can be resolved in different ways: try to merge the updates, duplicate the entries, let one entry win over the other, or simply do nothing and report the conflict so that the user can solve it.

Post-synchronization: At post-synchronization all the cleanup tasks are performed, like updating the mapping tables, report unresolved conflicts, and so on.

- 6 -

A widely used synchronization protocol is SyncML. SyncML provides an XML-based transport protocol for synchronization that is independent of the transport protocol. Each synchronization Message is a XML-document. A typical SyncML system is shown in FIG.4 where Application B (e.g. Personal Digital Assistant) is sending synchronization data via its Client Sync Engine to the SyncML Framework. The SyncML Framework translates the API calls (e.g. Update, Create) and the data into a valid SyncML document and sends it to the Server. On the Server side the SyncML Framework receives the document, parses it, and then sends the command and data to the Server Sync Engine, which then talks to the Application A (e.g. Lotus Notes Calendar of a notebook).

FIG.5 shows the basic architecture of the inventive content adapter framework (CAF) used in a communication architecture between mobile clients and different Back End data store types.

The different mobile Clients 2,4,6 access the Sync Engine 12 via a wireless or wired gateway 8 and through a Web Server 10 and the Sync Engine 12 talks via CAF 20 to the different Back End data store types 24,26. The CAF 20 provides the infrastructure to access data of different Back End data stores 24,26 through a single backend neutral interface (CAF-interface 22) and to easily add new Back End data stores. The CAF 20 consists of at least a single CAF-interface 22 and one or more content adapters 28,30.

The CAF-interface 22 represents a single interface for the Sync Server 10 to access Back End data and therefore separates

- 7 -

the content retrieval from the Sync Server. Through the CAF-interface 22 a Sync Engine 12 is able to access content independently from a specific Back End data store 24,26. For the data exchange between CAF-interface 22 and Sync Server 10 preferably Data Objects are used as data format.

Basically the content adapter provides all data store specific dependencies.

In preferred embodiment of the present invention each content adapter 28,30 includes an abstract Back End independent part and a Back End dependent part. The Back End dependent part contains all data store specific dependencies. It implements access to the synchronized Back End data and creates a data store independent representation of that data, which is provided to the Sync Engine 12 or application layer using the CAF interface 22. The CAF specific process flow is managed by the Back End independent part of the content adapter 28,30. Back End independent part provides the functionality common to all Back End data stores, e.g. queuing mechanism, communication handling.

In order to include semantic information, which can be used by the application (e.g. sync engine), a class hierarchy of common data objects is defined: Special subclasses of data objects describe the typical properties of the supported types of data, e.g. address, calendar, multimedia information, relational data bases, etc. The properties describing a data object for a particular kind of information can be taken from common standards, such as vCard (standard format for exchanging business card information) for address book information or vCal (standard format for exchanging calendar

- 8 -

information) for calendar entries. Applying XML allows even representing customer specific databases independent from a particular database.

Finally the framework 20 provides the infrastructure to easily integrate a caching mechanism between the Sync Engine 12 and the Back End data stores 24,26 for high volume systems or slow back-end systems (see FIG. 6)

The content adapter provide fast read/write access, adaptable to different backend systems (e.g. Domino, DB2, Exchange), support multiple SyncML messages, always have consistent data, and adaptable to different content formats.

The method carried out by the basic architecture of CAF may be briefly summarized as follows: Client requests sync session with Sync Server. Server authenticates client and accepts sync session. Client sends update to server. Sync server creates data objects and fills in the update received from the client. Sync server calls then CAF interface and hands over data objects. CAF selects the appropriate Back End specific part of the content adapter.

CAF calls the Back End specific part and passes the data objects to it. Back End specific part of the content adapter transforms data objects in a Back End specific format and calls Back End specific API (application programming interface).

One of the main advantages of the present invention compared with the prior art is making the access specific back-end databases independent from the calling application. CAF

- 9 -

integrates components from different database providers and offers access to their database functionality through a high-level interface.

By using common interfaces, the content adapters ensure interoperability of application with multiple Back End systems. A Sync Engine, for example, does not depend on proprietary commands of a particular database. Additionally, the components hide the complexity of the content, which is exchanged between a database and an application. This significantly reduces the programming efforts and the complexity of solutions. Also CAF allows the backend end system provider to just develop one interface for different applications accessing that backend data store. This saves both parties (the application provider and the backend supplier) a huge amount of time and money.

CAF allows:

- sync engines to talk to different back-end systems using the same protocol and API
- back-end system providers to create their own content adapter tailored towards their back-end system
- provides a session between sync engine, CAF and the back-end system handling the authentication to re-use back-end connections for efficiency
- allows the usage of a caching system to achieve low latency time for the communication sync engine - CAF
- allows for load-balancing and fail-save distribution of components

FIG.6 shows a preferred implementation of the CAF. The CAF 100 comprises a CAF-interface 22, a Content Manager 30, and a caching mechanism.

- 10 -

The CAF interface 22 provides a single interface for the Sync Server to access Back End data stores 24 and therefore separates the content retrieval from the Sync Server 10.

The Content Manager 30 forwards authentication and backend management requests (e.g. get a sync anchor) to the Back End Manager 80, writes new data to the cache 50 using the Persistent Store 40, and gets updates from the cache 50 through the Persistent Store 40. Search and execute commands are performed on the Back End.

The caching mechanism provides a permanent cache 50 and a mechanism for buffering of updates into the cache 50 and synchronizing buffered updates with the respective clients. The permanent cache 50 may be a relational data base like Oracle or IBM DB2 and may be accessed for example via JDBC calls. The caching mechanism preferably consists of a Cache Monitor 70, a Backend Monitor 60, a Back End Manager 80, and a persistent store 40. The Backend End Manger 80 includes an abstract Backend End Manager 80'' with its Back End specific parts 80' (Content Adapter), the Cache Monitor 70 includes an abstract Cache Monitor 70'' with its Back End specific Cache Monitor parts 70' (Content Adapter), and the Backend Monitor 60 includes an abstract Backend Monitor 60'' with its Back End specific Backend Monitor parts 60' (Content Adapter).

The Cache Monitor 70 is primarily used to replicate all new data from the cache to the back end data store. Depending on the Back End requirements different replication strategies, such as batch or trickle may be adopted. If the primary objective is to better support sync clients instead of regular back end clients the batch mode is preferred. The Back End dependent part 70' of the Cache Monitor 70 is specific for

- 11 -

each Back End data store and must exploit the features of the Back End data store (e.g. DB2, Domino). It also translates Sync Objects into a content storage specific format (e.g. Lotus Domino or MS Exchange).

The Back End Monitor 60 trickles updates that occur in the Back End data store 24 from outside the sync server (e.g. a regular Lotus Notes client updating a database) into the cache 50. This allows sync Clients to synchronize always with the latest back end data without requiring the overhead of a full replication for each sync session. The Back End specific part 60' of the Back End Monitor 60 is specific for each Back End data store and translates the content storage specific format (e.g. Lotus Domino or MS Exchange) into CAF Sync Objects.

The Back End Monitor can have different update policies, including aggressive or lazy updates, to optimize the overall system performance.

The Back End Manager 80 provides access to administrative functionality of the Back End data store. The following functionality preferably offered for supporting the CAF: validation of user authentication, retrieval of access permissions for authenticated users, retrieval of the current back end specific timestamp (current "sync anchor"), and adding/removing the URIs CAF wants to monitor for changes. The Back End specific part 80' of the Back End Manager 80 is specific for each Back End data store and translates the content storage specific format (e.g. Lotus Domino or MS Exchange) into CAF Sync Objects.

- 12 -

The data to be synchronized can either be stored directly in the remote content store (backend) or can be cached persistently on the server (locally) for performance reasons. The Persistent Store 40 uses a persistent storage medium as a cache to optimize read/write access to the Back End data store, however, the architecture does not prevent that the Persistent Store directly connects to the backend data store via the Cache 50 and Back End Monitor.

The communication flow within above CAF implementation may be briefly summarized as follows:

The Content Manager 30 receives the requests from the Sync Engine 12 and forwards them either to the persistent store 40, if data needs to be retrieved or stored, or to the Back End Manager 80, if a timestamp is needed or authentication is requested. The Back End Registry, if available, contains all available Back End Managers 80 and Monitors 60,70 and is accessed from Content Manager 30 and Persistent Store 40. The Cache Monitor 70 gets updated data from the Persistent Store 40, translates these to the Back End format, and forwards the data to the Back End data store 24 by using the Back End dependent part 70' of the Cache Monitor. The Back End dependent part 60' of the Back End Monitor 60 receives the updates from the Back End data store 24, translates them to Data Objects, and forwards them to the Persistent Store 40.

In case the cache is not available or for a given database a direct access to the Back End 24 is specified, the Content Manager 30 forwards the getUpdates call to the Back End Monitor 60 and the items to be updated to the Cache Monitor

- 13 -

70. Both monitors will use the provided user ID and password to access the backend.

Authentication Sync Engine with CAF

The sync engine has two options for authentication its requests to CAF

1. provide for each command the required backend user ID and password
2. request an authentication token from CAF for one backend by providing a user ID and a password. Each subsequent command to CAF for this sync session can be authenticate with this token (similar to a LTPA token). The following token types may be supported: read-only access, read/write access, or unrestricted read access, but write must be authenticated each time at the Back End.

Authentication CAF to Back End data store

The required authentication level to the Back End data store is stored in Access Control Lists for each Back End data store and checked by the Sync Adapter. Dependent on this list, CAF authenticates itself to the Back End data store either with a group user and group password valid for all user that update their data through the sync server or on a per user basis.

FIG.7 shows the interfaces (I/F) of the CAF in the preferred implementation of FIG.6.

The Sync Engine 12 uses the CAF interface 22 to access the Back End data stores in a generic way. To efficiently exchange data between generic Sync Engine 12 and CAF interface 22, CAF

- 14 -

interface 22 preferably uses the raw value binary encoding scheme for data exchange. These raw data are embedded as ActionData object in the SyncObj together with the CAF Meta data.

CAF Meta data are:

- Timestamp represents the current sync anchor of the SyncObj. It is stored as object of the class `java.sql.Timestamp`.
- ActionType represents the data action type of the SyncObj. Possible values are defined in class constants (CREATE, UPDATE, DELETE). It is stored as short value.
- GUID represents the backend specific global id of the SyncObj. It is stored as `java.lang.String` value.
- CUID represents the CAF specific cache id of the SyncObj. It is stored as `java.lang.String` value.
- databaseURI represents the URI of the backend database to which the SyncObj belongs to. It is stored as object of the class `com.ibm.caf.URI`.
- userID represents the user id of the client that initiated the synchronization to which the SyncObj belongs to. It is stored as `java.lang.String` value.

The interfaces of the CAF specific process flow

The Abstract Monitor Class offers basic functionality for the integration of Back End specific Monitors into the CAF architecture. Two abstract classes called Abstract Cache Monitor and Abstract Backend Monitor are provided that can be used via inheritance for the creation of new Monitor classes (see FIG.8).

The abstract classes provide the following functionality:

- 15 -

- Queuing mechanisms for different update policies (trickle / batch updates)
- Registration / Deregistration at the Backend Registry
- Reading und registration for updates at the System Preferences Component
- Handling of communication with CAF Persistent Store (Cache)

The CAF specific process flow is managed by the Abstract Backend Monitor 60 and the Abstract Cache Monitor 70 through the interfaces defined in this patent application. The direct communication with the Back End data store 24 is implemented in Back End specific components 70', 60' that are inherited from the Abstract Monitor classes (see FIG. 9)

Each Monitor implements an internal queue that enables different update and data propagation policies. Dependent on the usage scenario, it may be necessary to send updated data in groups to the backend, or in large time intervals.

The Abstract Monitor component implements a configurable queuing mechanism that offers the following update policies:

Amount Trigger:

The Amount Trigger monitors the size of the internal queue. It will propagate the collected items when a certain configurable threshold is reached. This update policy can be used for configuring a batch (threshold > 1 item) or trickle (threshold = 1 item) update mechanism in the specific monitor.

Interval Trigger:

The Interval Trigger monitors the time that has been passed since the last time of propagation. It will send the collected

- 16 -

items when a certain configurable time interval is reached.

Combined Trigger:

The Combined Trigger utilizes both above-mentioned policies: It will propagate the collected items whenever one of the triggers is activated.

Session Handling

The topic of session handling applies primarily to the Cache Monitor, which is part of the overall CAF processing. Session handling in the Backend Monitor is completely dependent on the backend-specific implementation.

There are two different usage patterns for the Cache Monitor: with and without the CAF Cache. When the Cache is used, no session handling is supported at all. The updates are replicated asynchronously from the cache to the backend. Since access to the backend database is performed through a admin-like user account, pooling of connections is possible for all occurring updates.

In the case that no cache is used, session handling is performed with the methods `beginConnect()` and `endConnect()`. All updates are grouped on a "per user" basis, which allows connection pooling for each synchronizing user with the given credentials.

FIG.8 shows a synchronization flow for a two-way synchronization with session authentication preferably applied by the inventive CAF.

A two-way synchronization between Client and Server is performed where the Client as updated item A,B and F, deleted C and created a new item D. Via an external client (e.g. Notes client) E was created on the Back End and B, C and F were

- 17 -

updated.

Package 1 from client sends the credential for the Back End data store. The Sync Engine forwards these credentials to the CAF for verification. CAF asks the responsible Back End to verify the credentials and returns an authentication token valid for a synchronization session to the Sync Engine. This token needs to be included into any future request to the CAF for this synchronization session.

In package 3 the clients sends its updates to the Sync server. The Sync Engine requests the Back End updates from the CAF by presenting the authentication token. Now the Sync Engine compares the lists, resolves the conflicts, and populates the updated entries to the CAF and client. CAF stores these updates in its cache and replicate these changes later to the Back End system.

The Sync Engine without interaction of CAF does the handling of package 5 with mapping table information.

- 18 -

C L A I M S

1. System for exchange of data between different clients (2,4,6) by using a central synchronization server (10) having a connection to said clients (2,4,6) and a connection to a Back End data store (24,26), wherein said clients having a program for creation of data to be synchronized, and a Sync Engine (12) for performing synchronization with said central synchronization server (10), wherein said system is characterized by the further components:

a single Back End neutral interface (CAF- interface; 22)) with said Sync Engine (12), and

at least one component (content adapter; 28,30) comprising a Back End dependent part having an interface with said single Back End neutral interface (22) and said assigned Back End data store (24, 26).

2. System according to claim 1, wherein said component(28,30) further comprises an abstract Back End independent part, wherein said abstract Back End independent part provides common functionalities for use by all the Back End dependent parts.

3. System according to claim 2, wherein each Back End data store type (24,26) is assigned an own component(28,30).

4. Server according to claim 1, wherein said exchange of data is synchronization of data.

5. System according to claim 2, further comprises a cache

- 19 -

(50) for permanently buffering of updates of said Backend data store (24) and said Clients, and said component (28,30) comprises a caching mechanism for controlling and executing buffering updates into said cache and replicating buffered updates to said respective clients and said Backend data store (24).

6. System according to claim 5, wherein said caching mechanism having a Back End Monitor (60).

7. System according to claim 5, wherein said caching mechanism further having a Cache Monitor (70).

8. System according to claim 6, wherein said caching mechanism further having a Back End Manager (80).

9. System according to claim 6, wherein said caching mechanism provides for each Back End data store type (24) an own Back End Monitor, Cache Monitor, and Back End Manager with its Back End dependent (60',70',80') and its abstract Back End independent part (60'',70'',80'').

10. System according to claim 5, wherein said caching mechanism further comprises a persistent store (40).

11. System according to claim 7, said Cache Monitor (70) replicates updates from said cache to the Back End data store (24) in a batch or a continuous trickle mode.

12. System according to claim 6, wherein said Back End Monitor (60) replicates updates between said cache (50) and the Backend data store (24) in a batch or a continuous trickle mode.

- 20 -

13. Systems according to claim 5, wherein said cache (50) and said Backend data store (24) are databases.

14. System according to claim 1, wherein said clients are mobile clients.

15. System according to claim 4, wherein the synchronization protocol is SyncML.

16. Method for synchronization of data by using a system according to claim 1-15, comprising the steps of:

Receiving sync session request from said Client,

Authenticating said Client against said Sync Server,

Receiving update from said client,

Authenticating said client against Back End data store via said CAF interface using Back End Monitor,

Creating of data objects and filling in said update received from said client by said Sync Server,

Calling said CAF interface and handing over said data objects,

Selecting the appropriate Back End specific part of said component assigned to said Back End data store,

Transforming said data objects CAF into a Back End specific format,

- 21 -

Executing said updates by calling Back End specific part and passing the data objects to it.

17. Method according to claim 16, wherein said Back End specific part is inherited from said abstract Back End independent part assigned to said Back End data store.

18. Method according to claim 16, wherein said data objects are used to pass said client request to Back End specific parts.

19. Method according to claim 18, wherein said data objects contain meta data.

20. Method according to claim 16, wherein synchronization protocol used between said clients and said central synchronization server is SyncML and said updates received by said central synchronization server are presented as XML documents.

21. Computer program product stored in the internal memory of a digital computer, containing parts of software code to execute the method in accordance with claims 16 to 20 if the product is run on the computer.

1 / 6

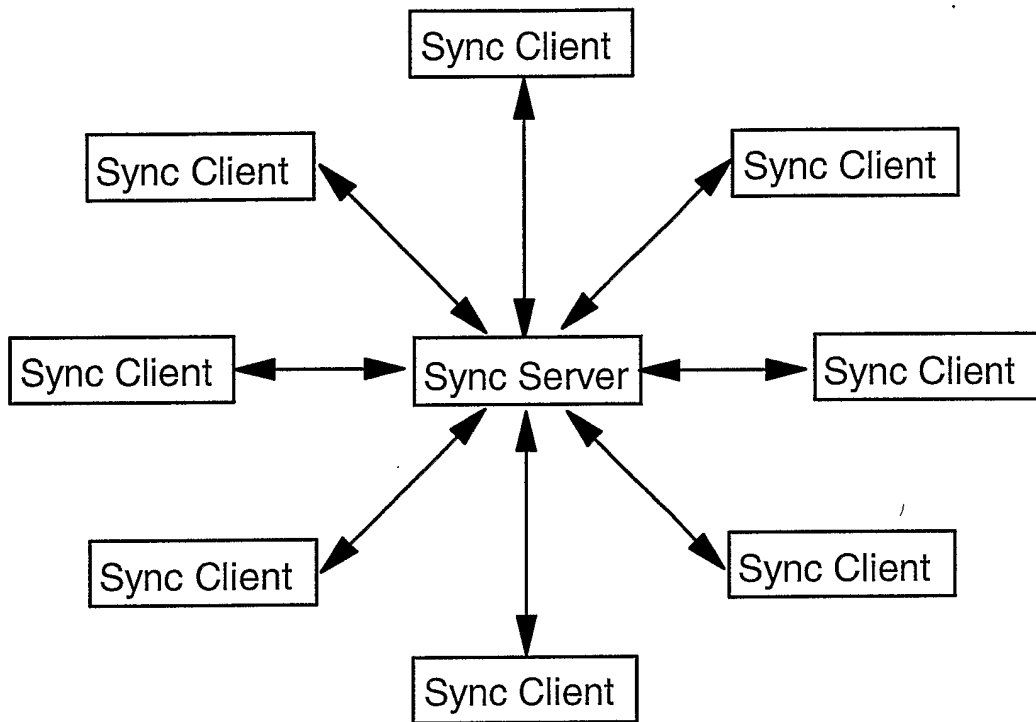


FIG. 1

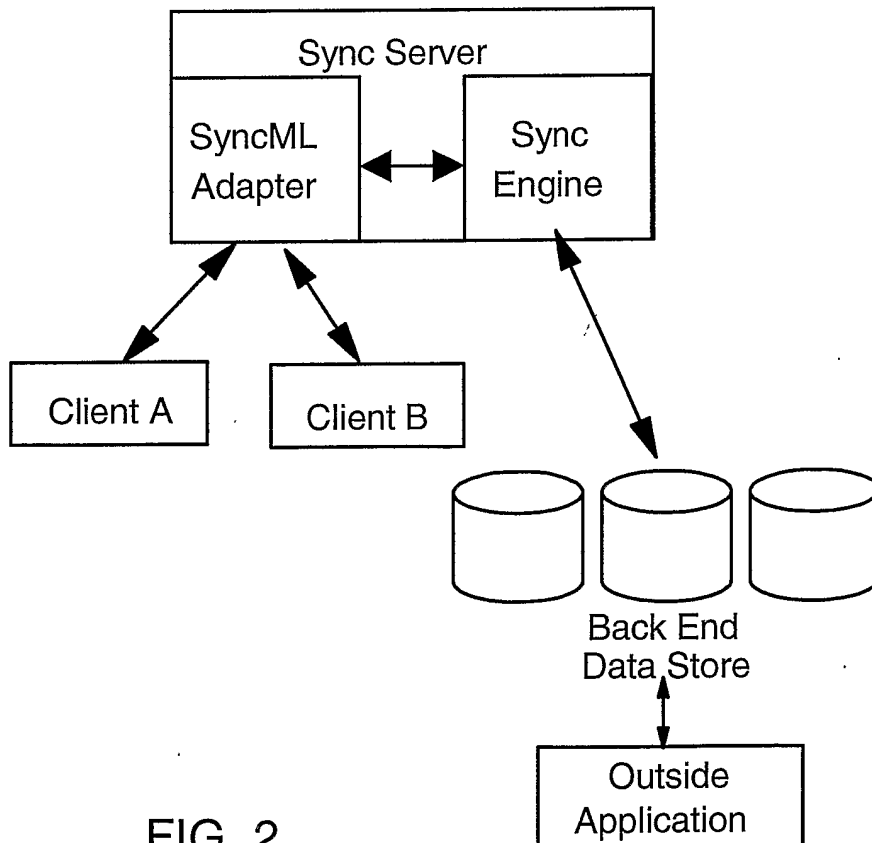


FIG. 2

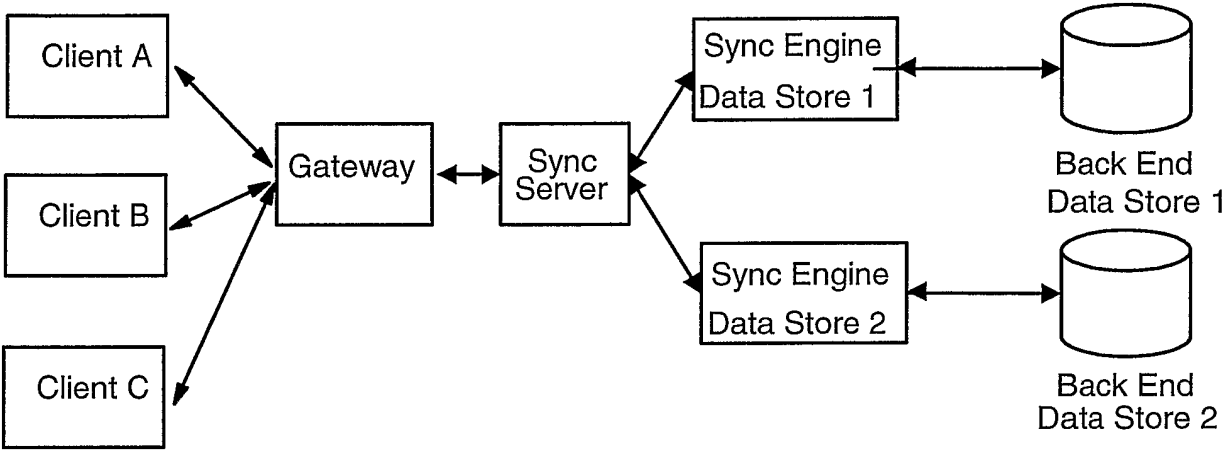


FIG. 3

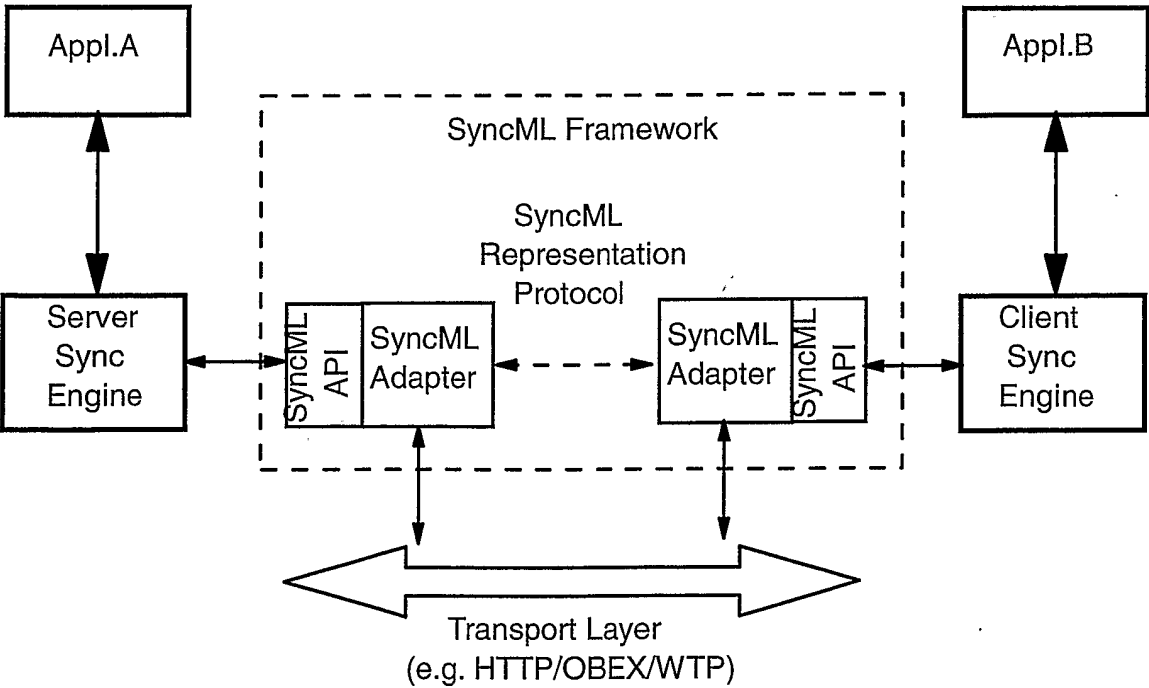


FIG. 4

3 / 6

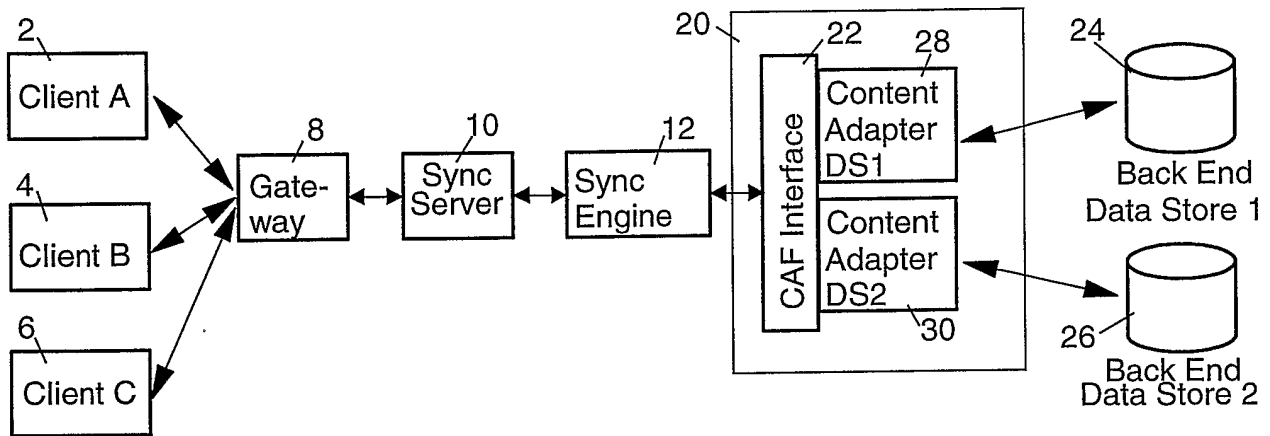


FIG. 5

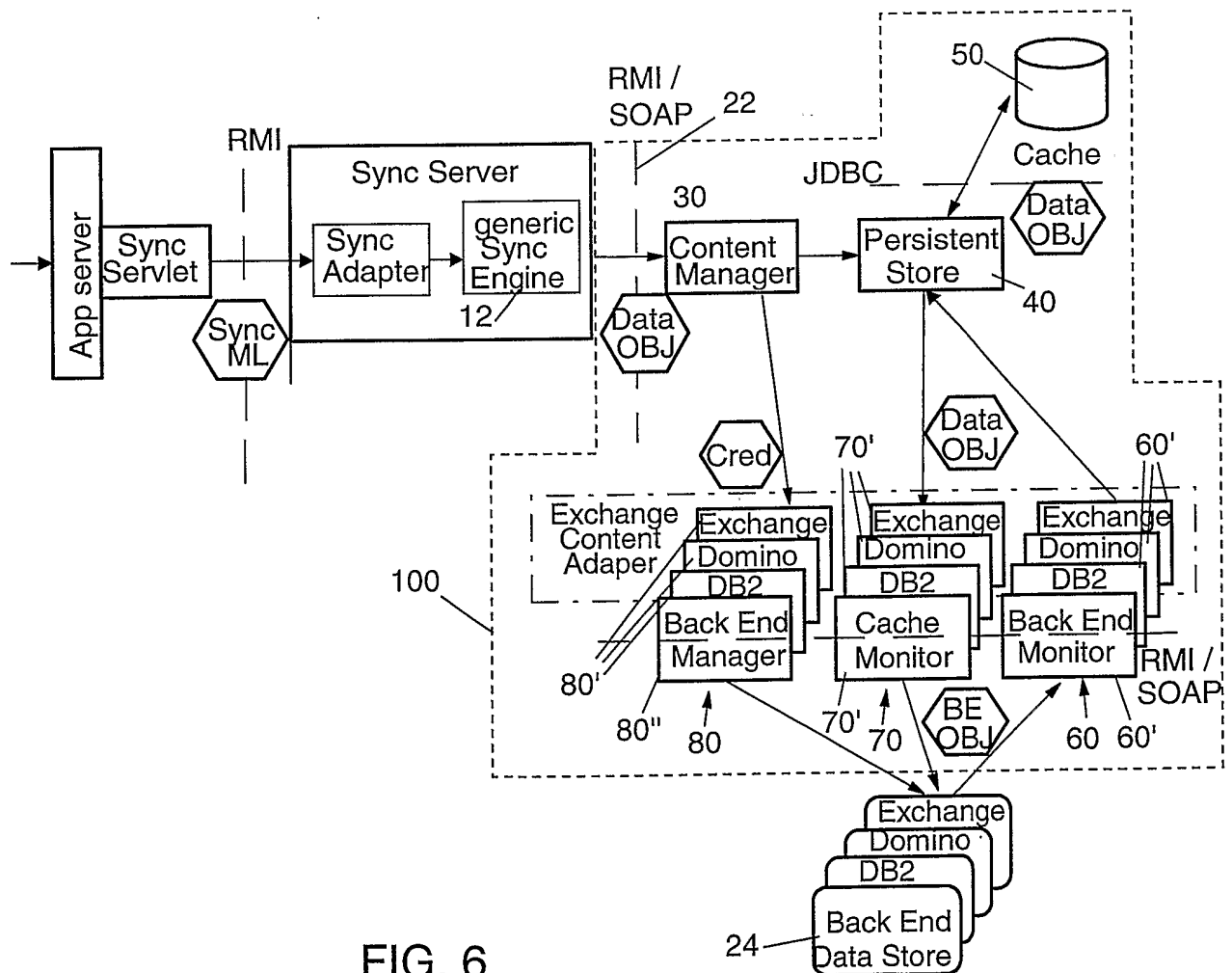


FIG. 6

4 / 6

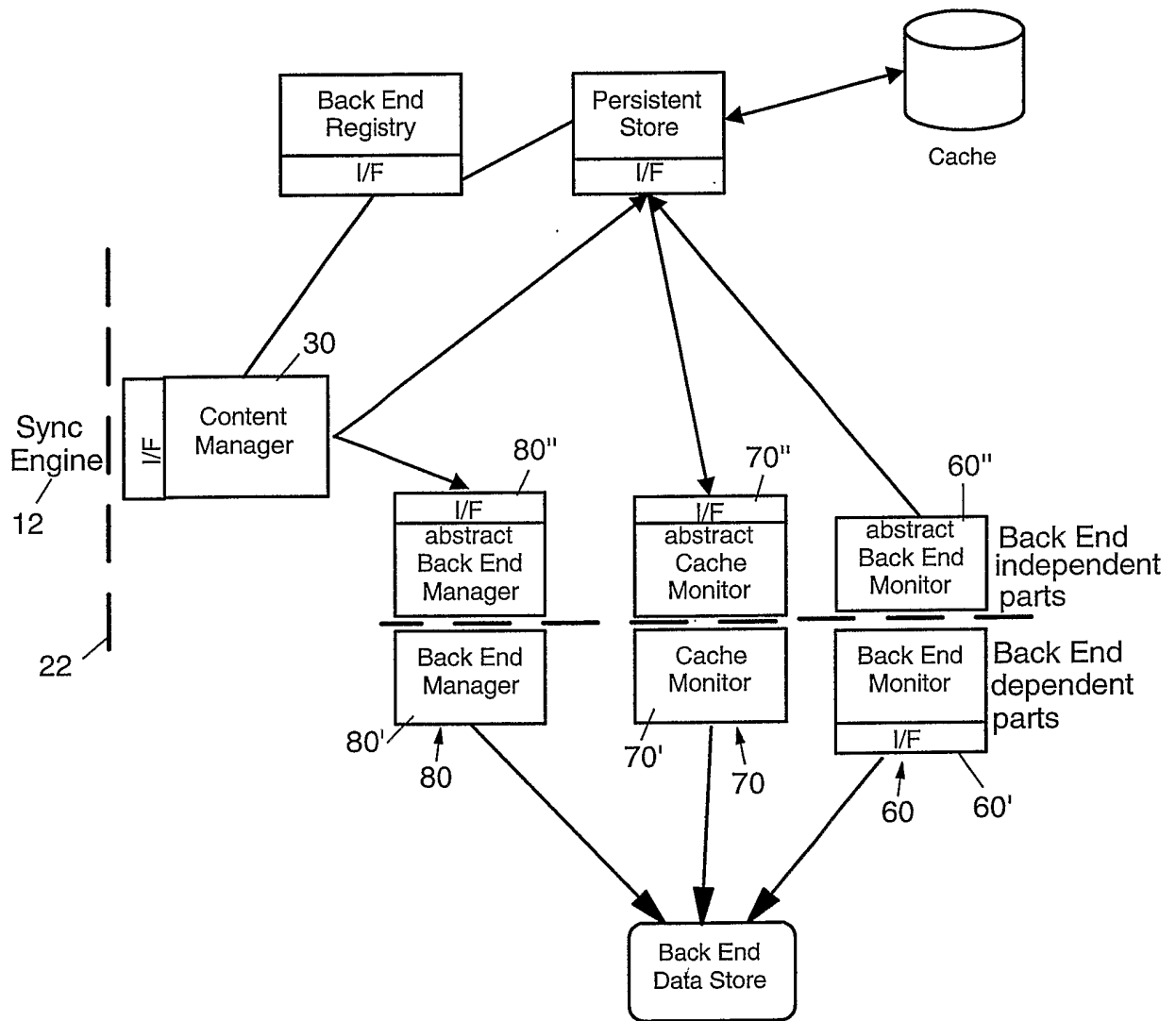


FIG. 7

5 / 6

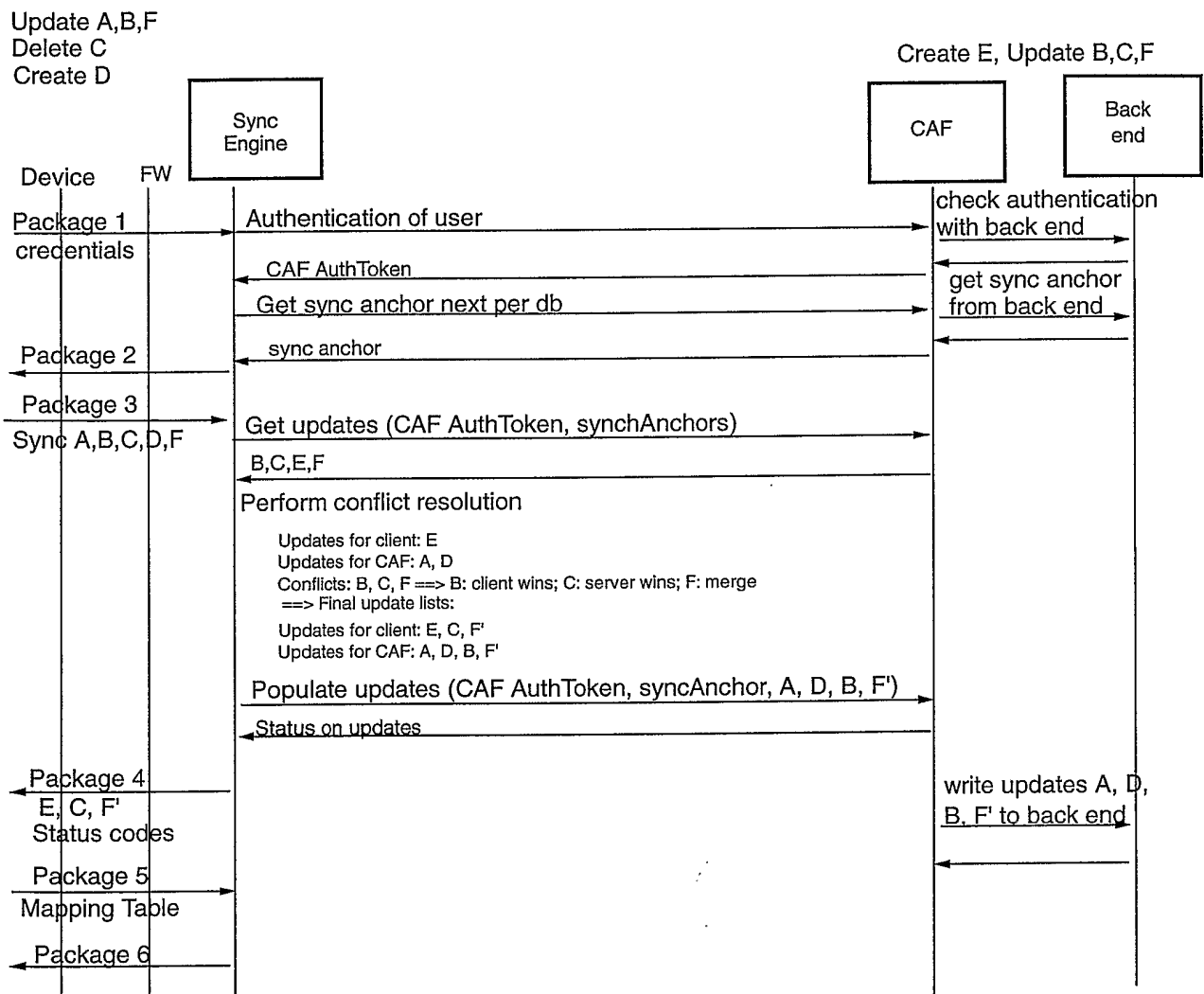


FIG. 8

6 / 6

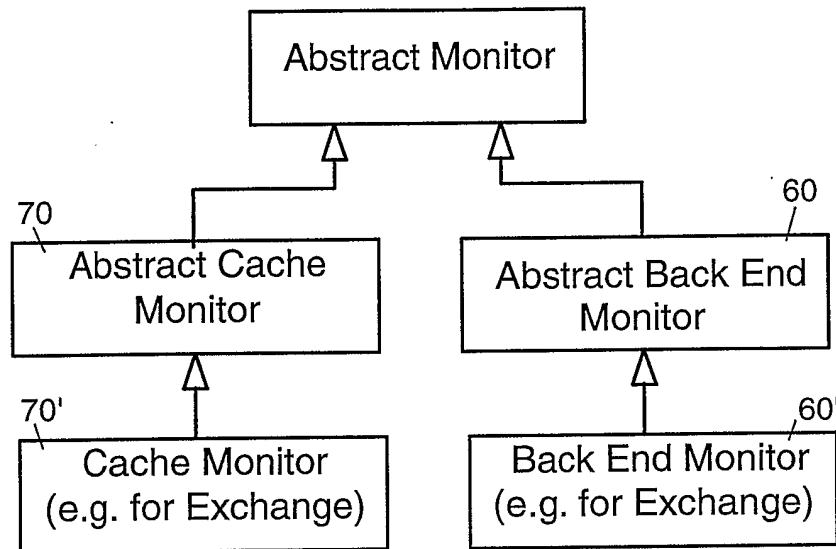


FIG. 9

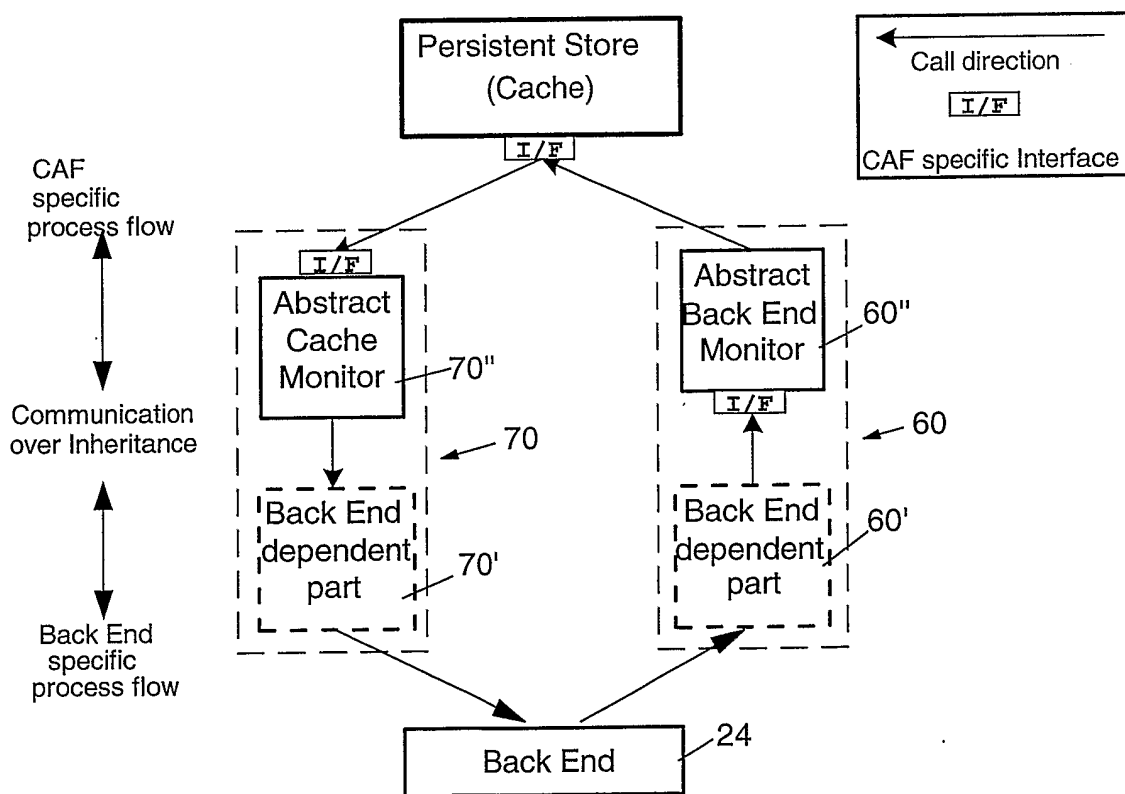


FIG. 10