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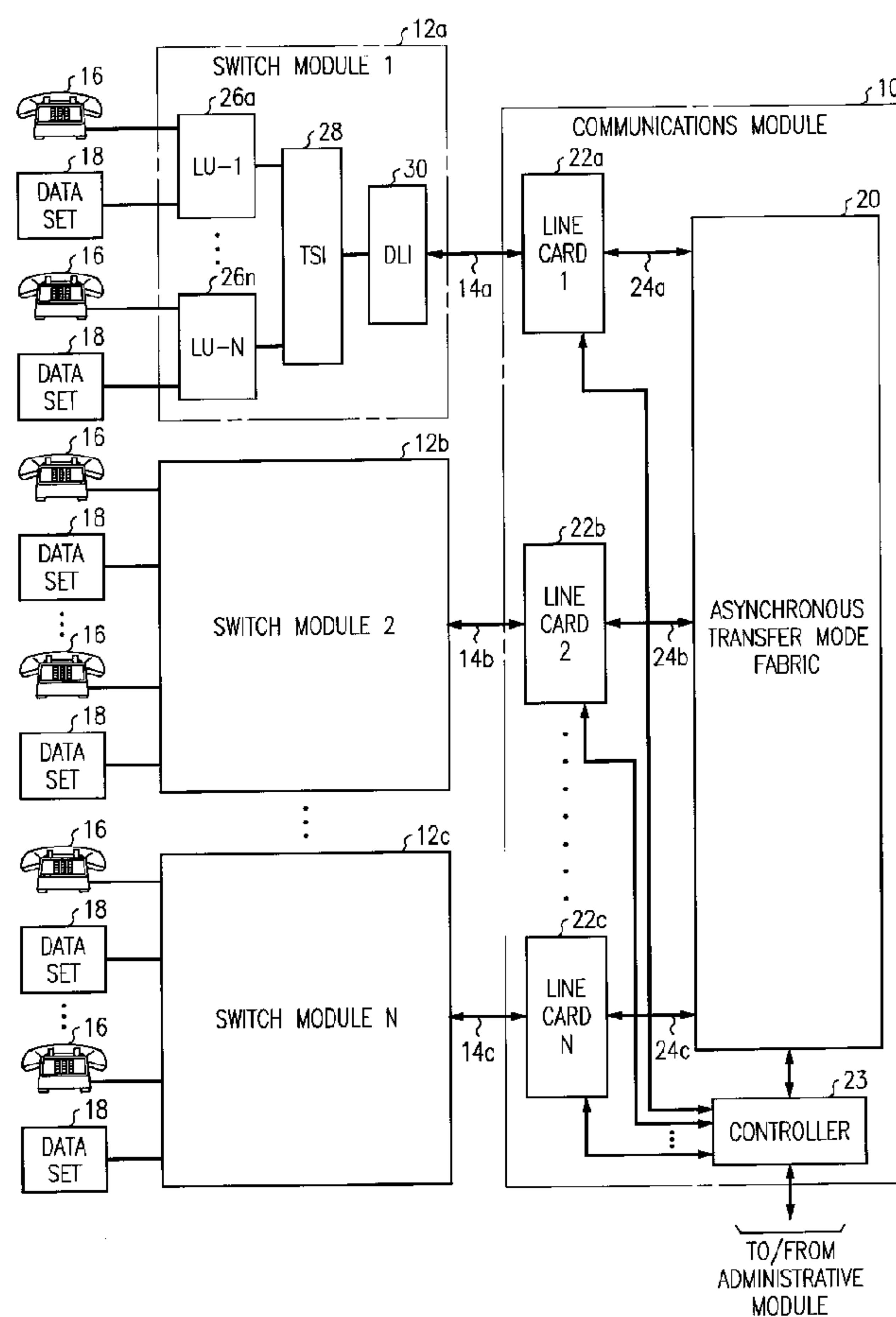
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(30) 1994/12/07 (350,926) US

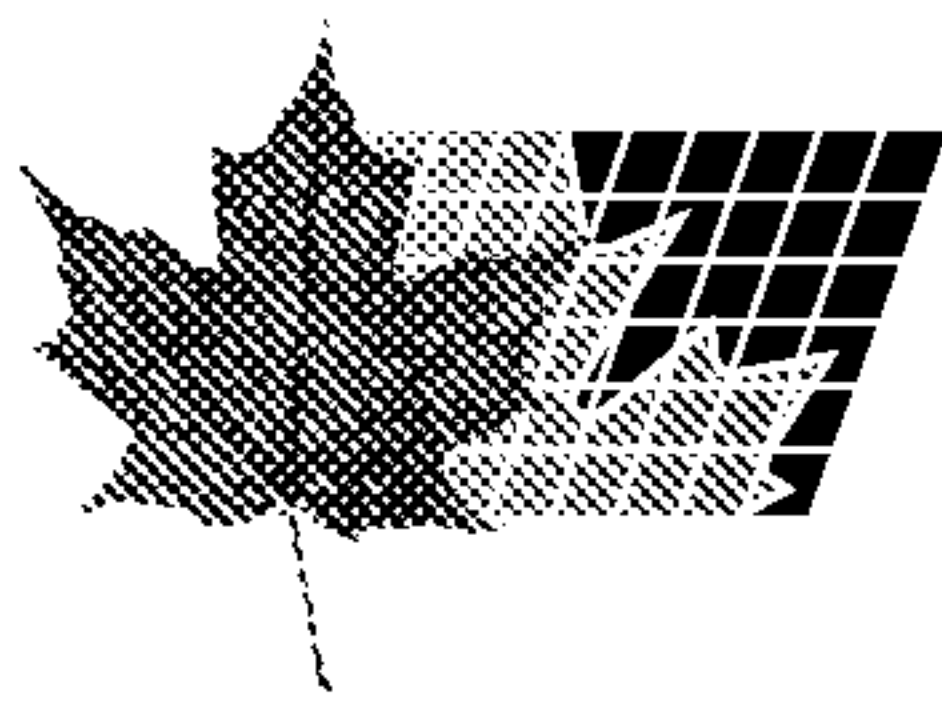
(54) **SYSTEME DE TELECOMMUNICATION A DISPOSITIF DE
DETECTION ET DE CONTROLE DES COLLISIONS DE
PAQUETS**

(54) **TELECOMMUNICATION SYSTEM WITH DETECTION AND
CONTROL OF PACKET COLLISIONS**



(57) Line cards provide an input/output interface between synchronous communication channels and an asynchronous network. The input line card translates information received from the synchronous communication channels into packets for transmission to the asynchronous network. Each line card generates and stores packets corresponding to a time frame associated with the synchronous communication channels. A sequence is selected for the sequential





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(45) 2000/04/04

transmission of packets by each input line card to the asynchronous network. The cumulative number of packet collisions that would occur at each output line card during the predetermined time frame is determined. If the cumulative number of collisions exceeds a predetermined delay threshold, the number of packets actually transmitted to the system is reduced preferably by denying service to new call origination requests. This technique keeps the maximum cumulative collisions for each output line card within the predetermined delay threshold.

Abstract of the Disclosure

Line cards provide an input/output interface between synchronous communication channels and an asynchronous network. The input line card translates information received from the synchronous communication channels into packets for transmission to the asynchronous network. Each line card generates and stores packets corresponding to a time frame associated with the synchronous communication channels. A sequence is selected for the sequential transmission of packets by each input line card to the asynchronous network. The cumulative number of packet collisions that would occur at each output line card during the predetermined time frame is determined. If the cumulative number of collisions exceeds a predetermined delay threshold, the number of packets actually transmitted to the system is reduced preferably by denying service to new call origination requests. This technique keeps the maximum cumulative collisions for each output line card within the predetermined delay threshold.

TELECOMMUNICATION SYSTEM WITH DETECTION AND CONTROL
OF PACKET COLLISIONS

Background of the Invention

5 This invention is directed to a telecommunication system in
which a plurality of channels (calls) are transported using an
asynchronous node as at least part of the system. More
specifically, this invention addresses the monitoring and
controlling of the network in the presence of packet collisions
10 which result in delays arising from the need to store the colliding
packets in queues.

A conventional telecommunication system may utilize a switch
such as an AT&T 5ESS[®] switch consisting of an administrative module,
a communications module, and switching modules which support a
15 plurality of telephone subscribers. The switch modules and the
communications module each provide synchronous communication
channels in which an originating and a terminating subscriber are
assigned respective channels or time slots by which information is
transmitted between the subscribers. Since the communications
20 module provides a synchronous communication link with the switch
modules, delays or queuing problems are not encountered because each
communication path is assigned a time slot which is maintained for
the duration of the call.

Asynchronous data communication systems such as those using
25 packets to carry information provide an alternative to the
synchronous communication systems typically utilized for voice
communications. Since an asynchronous communication system will not
have a unique time slot for each packet, it is possible that a
plurality of packets destined for the same node may result in the

need to buffer or store packets for delivery to the node. The depth or number of packets which must be stored in a queue provides an increased timing delay for packets traversing such a system. Such timing delays may become critical in certain applications such as when voice communications are also carried by other equipment in the system having time delay limits such as synchronous communications equipment. Thus, there exists a need in such a network to monitor and limit time delays due to collisions and queuing delays to maintain the quality of communications.

10 **Summary of the Invention**

It is an object of the present invention to provide a method and a corresponding apparatus which monitors and limits delays in a telecommunication system due to queuing delays associated with the collision of packets.

15 In accordance with an embodiment of the present invention, a telecommunication system includes input and output line cards which provide an interface between synchronous communication channels interconnected by an asynchronous network. The input line card translates information received from the synchronous communication channels into packets for transmission to the asynchronous network. The input line cards concurrently transmit packets to the asynchronous network. Each line card generates and stores packets corresponding to a time frame associated with the synchronous communication channels. A sequence is selected for the sequential transmission of packets by each input line card to the asynchronous network. The cumulative number of collisions that would occur for each output line card during the predetermined time frame is determined. If the cumulative number of collisions exceeds a predetermined delay threshold, the number of packets actually

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transmitted to the system is reduced from the number of calculated packets by denying service to a new call origination request. This technique keeps the maximum cumulative collisions for each output line card within the predetermined delay threshold. Thus, a
5 predetermined maximum time delay or skew can be maintained even though an asynchronous network is utilized.

Brief Description of the Drawings

FIG. 1 is a block diagram of a portion of a telecommunication system in accordance with the present invention.

10 FIG. 2 illustrates an exemplary format of a packet in accordance with the present invention.

FIG. 3 is a functional block diagram of an exemplary line card in accordance with an embodiment of the present invention.

15 FIG. 4 is a block diagram of an asynchronous transmission fabric or network which may be utilized in accordance with the present invention.

FIG. 5 illustrates a block diagram of an embodiment of an implementation of a line card in accordance with the present invention.

20 FIG. 6 illustrates a block diagram of an embodiment of a controller as shown in FIG. 1.

Detailed Description

FIG. 1 illustrates a portion of a telecommunication system which incorporates an embodiment of the present invention. A
25 communications module 10 is shown illustrated connected to switch modules 12a, 12b, and 12c by corresponding transmission channels 14a, 14b and 14c. Each of the switch modules may support a plurality of customer premises equipment such as telephone sets 16 and data sets 18. The conventional telephone sets 16 may comprise

conventional analog telephones or digital telephones. The data sets
18 represent data communications equipment such as a personal
computer which communicates over an analog line utilizing a modem,
or a multimedia device which could include digitized video as well
5 as audio or data information.

The communications module 10 includes an asynchronous network
20 which may consist of an asynchronous transfer mode (ATM) fabric.
A plurality of line cards 22a, 22b, and 22c are coupled by
respective transmission channels 24a, 24b, and 24c to switch 20.
10 The ATM fabric 20 may comprise a known asynchronous ATM system which
transports packets of data between ports of the asynchronous system.
However, different asynchronous systems could be utilized to
transfer packets between the line cards. A controller 23 is coupled
to the ATM fabric 20 and each of the line cards 1-N. It controls
15 the routing of calls between the switch module just as controllers
used by synchronous communications modules such as used in an AT&T
5ESS switch.

Since the switch modules are substantially identical, only
switch module 12a (switch module 1) will be described. A plurality
20 of line units 26a-26N provide an interface to the connected
subscriber equipment 16, 18. A time slot interchange (TSI) unit 28
translates digital information from the line units into
corresponding time slots which are transmitted to a digital link
interface (DLI) unit 30. The unit 30 provides a transmission
25 interface between the TSI unit 28 and a line card in the
communications module 10. In accordance with the present invention
each DLI unit 30 is supported by a corresponding line card. The
switch modules provide the known function of assembling digital
information representing a plurality of calls, transmitting it to

the communications module 10, and translating the digital information received from the communications module into a format suited for delivery to the customer premises equipment. Switch modules in accordance with the illustrative embodiment are available from AT&T Corp. as part of the 5ESS switch environment.

FIG. 2 illustrates a format for a packet 32 such as utilized in the ATM network 20 of the embodiment of the present invention. The packet includes a 5-byte header 34 and a payload 36 consisting of 48 bytes of information. Each payload byte carries a time slice of information corresponding to one communication channel. Of course, more than one of the payload bytes can be associated with a single information source to provide a larger effective bandwidth. The packets 32 are carried by the ATM network 20 and are received from and transmitted to the line cards. The header of each packet contains destination information sufficient to permit the ATM fabric 20 to determine which of the line cards 1-N is the destination line card. There is a probability that a plurality of packets destined for the same line card will arrive for delivery at substantially the same time since packets are launched in synchronism from the line cards. Such a condition is referred herein as a collision which means that two or more packets are to be delivered to the same destination at substantially the same time. Minimizing the probability of collisions is explained below in accordance with an embodiment of the present invention. Further, a technique associated with managing collisions is explained below in accordance with another aspect of the present invention.

FIG. 3 is a functional block diagram illustrating an embodiment of line card 22A in accordance with the present invention. Each of the line cards are identical in structure and

function. However, it will be appreciated by those skilled in the art that each of the line cards has a separate and unique address enabling it to be individually identified as a destination for packets. Digitized representations of the communication channels generated by a TSI and sent by a DLI of a switch module is received on transmission channel 14A and stored in a shift register 38. In the illustrative embodiment it is assumed that each switch module operates on a synchronous time frame of 125 microseconds and that during each frame, six thousand separate time slots, which could correspond to six thousand different channels (calls), are generated and transmitted to shift register 38. FIG. 3 illustrates the storage of time slots TS1-TS6000 in register 38 representing the time slices received during the frame.

A cell assembly unit 40 assembles the six thousand time slices in the illustrative embodiment into 180 cells which are stored in buffer 42. Each of cells 1-180 as shown in FIG. 3 consists of a packet payload such as shown in FIG. 2. Thus, each cell could carry up to 48 time slices (bytes) resulting in sufficient transmission capability to carry more than the required six thousand time slices in a given 125 microsecond frame. In accordance with the exemplary design, the number of cells to be transmitted during a 125 microsecond frame was selected since the ATM network had a capacity for transferring 180 packets every 125 microseconds for each line card.

A cell organization unit 44 operates on the cells stored in buffer 42 to reorder the sequence in which the cells will be transmitted to the ATM network. The cell organization unit 44 stores the cells in buffer 48 in the order as determined by unit 44 and adds destination information as the packet header to define a

corresponding packet. Thus, buffer 48 contains a predetermined sequence of 180 packets to be transmitted sequentially (1, 2, ..., 180) to the ATM network 20 via channel 24a.

An important aspect in accordance with the present invention resides in the recognition of the problem that an excessive number of packet collisions results in queuing delays which can adversely impact the performance of the synchronous portion of the system. In accordance with the present invention the cell organization unit 44 orders the packets for sequential transmission to the asynchronous network to minimize the probability that more than one of the input line cards will concurrently transmit a packet to the same destination output card.

Referring to FIG. 3, the cell assembly unit 40 assembles the time slices contained in register 38 into one of 180 cells contained in buffer 42. Each of the time slices selected by the assembly unit 40 for assembly into a particular given cell will have the same output line card destination since each of the cells in buffer 42 are routed to a single destination output line card. The cell organization unit 44 provides a non-random cell organization technique in which the cells in buffer 42 are ordered by sequence of transmission and placed into buffer 48 for transmission to the ATM network 20. That is, packet 1 is transmitted first, packet 2 second, etc. The line cards are preferably synchronized so that each concurrently transmits a packet to the ATM network. The input line cards sequentially transmit packets 1-180 to the ATM network. The illustrated line card functions both as a packet transmitting and packet receiving mechanism.

Referring to FIG. 1, the line cards 1-N are preferably synchronized in the transmitting of packets to the ATM network so

that each concurrently transmits a packet to the ATM network 20. The line cards operate independently in ordering the packets for transmission to the ATM network. The ATM network receives the packets and transfers the packets to the respective output line card destination as deduced from the information contained in the header of the packet. As will be explained in more detail below, a technique is provided in accordance with the present invention which minimizes the probability that two or more input line cards will concurrently transmit a packet having the same output line card destination.

FIG. 4 illustrates a block diagram of an exemplary ATM fabric 20 in accordance with the present invention. The input portion of the bi-directional transmission channels 24A, 24B, and 24C are coupled, respectively, to input registers 50A, 50B, and 50C. Each of these registers stores a packet transmitted by the corresponding line card. Each register transmits the stored packet to the ATM packet routing unit 52 which routes the received packets to the output destination line card. A plurality of output queues 54a, 54b, and 54c receive packets from the packet routing unit 52 and provide a buffer or storage for storing up to M packets waiting to be delivered. It will be apparent to those skilled in the art that it is desirable to minimize the number of packets stored in the output queues in order to control the associated time delays. Further, if the number of packets held in an output queue approaches a predetermined maximum threshold level, the delivery of the packets to the synchronous portion of the communication system may jeopardize the required time synchronization, thereby resulting in the loss of packets or degradation of communication quality to end users.

FIG. 5 is a block diagram illustrating an exemplary architecture of a line card in accordance with the present invention. It will be apparent to those skilled in the art that the functional operations illustrated in FIG. 3 may be implemented by the architecture shown in FIG. 5. The time slot data is received on channel 14a by register 60 that stores the data for further processing by logic unit 62 which may comprise a very large scale integrated (VLSI) circuit having predetermined data handling functions including read, write, and index capabilities. The random access memory (RAM) 64 contains a map comprised of addresses that identify locations in RAM 66; the map is written to RAM 64 by the CM controller 23. Logic unit 62 reads the map in RAM 64 and stores each time slice in register 60 at a location in RAM 66 as defined by a corresponding address stored in RAM 64. Predetermined areas in RAM 66 correspond to the cells 1-180 and function as the buffer 42 in FIG. 3. Each of the time slots received by register 60 are written by logic unit 62 to a memory location in RAM 66 corresponding to a cell in accordance with the destination map contained in RAM 64.

Logic unit 68 which may also comprise a VLSI circuit reads sequence map stored in RAM 70 which defines the transmission sequence of each of the cells stored in RAM 66. CM controller 23 writes the map to RAM 70 to control the transmission order of the packets. Logic unit 68 sequentially reads each map element which identifies one cell stored in RAM 66 and contains a destination address for each cell. Logic unit 68 reads the identified cell, appends an appropriate packet header based on the destination address, and transmits the packet on channel 24a to the ATM fabric.

Packets on channel 24a from the ATM fabric are received and temporarily stored by register 72 of the line card shown in FIG. 5. Logic unit 74 which may also comprise a VLSI circuit reads each received packet stored in register 72 and, based upon disassembly map stored in RAM 76, writes each time slice (byte) carried by each packet to a location in RAM 76. The CM controller 23 writes a disassembly map to RAM 76 for every 125 microsecond frame in preparation for the disassembly of the packets. After logic unit 74 receives all packets transmitted during a frame and provides the corresponding disassembly, RAM 78 will contain each time slice to be delivered during the frame to the switching module connected to the particular line card.

Following the translation of all received packets into corresponding time slice data stored in RAM 78, logic unit 80 which may comprise a VLSI circuit transmits each of the stored time slice data in RAM 78 to the associated switch module by channel 14a based upon a time slice map stored in RAM 82. This map identifies the location in RAM 78 of the time slot data that corresponds to each of the 1-6000 time slots. The time slice organizational map stored in RAM 82 is written by the controller 23 for each frame. Logic unit 80 reads time slice data stored in RAM 78 as defined by the map stored in RAM 82 and transmits it to the associated switch module by channel 14a.

In view of the required transmission speeds, it is preferable that the line cards employ hardware logic circuits which can operate on clock cycles of very short duration in order to provide maximum throughput. It will be apparent to those skilled in the art that the intelligence required to create the maps contained in RAM's 64, 70, 76, and 82 could be generated by a single intelligent node or

could be generated by distributed intelligence including intelligence contained upon each line card.

FIG. 6 is an illustrative block diagram of an exemplary controller 23 as shown in FIG. 1. A microprocessor 90 operates in cooperation with program instructions and data stored in read-only memory (ROM) 92, random access memory (RAM) 94, and nonvolatile data storage device 96 which may comprise one of various types of nonvolatile data storage including a hard disk drive. An input/output interface circuit 98 is coupled to microprocessor 90 and facilitates the reception and transmission of digital data to the line cards, the ATM fabric, and an administrative module. Digital data written to the Ram's 64, 70, 76, and 82 of each line card control the handling of time slots and packets routed by each line card as described above with regard to FIG. 5. The microprocessor 90 operates under the control of program instructions to accomplish the processing of information required for the general operation of the communications module including information for the operation of each line card. Controllers such as employed by AT&T 5ESS communication modules are known.

In accordance with the present invention, the controller 23 determines the transmission order sequence for the outbound packets by the sequence map stored in RAM 70. This sequence map is based on the following calculation:

$$\text{slot} = (\text{input LC \#} + \text{output LC \#}) \text{ modulo } N$$

Where slot is a number upon which the transmission sequence order of a packet is based, the input LC# is the identification number of the input line card, output LC# is the identification number of the output line card, and N is the number of total output line cards. In accordance with an example of an embodiment of the present

invention, assume that 16 input and output line cards are utilized with both the input and output line cards numbered 1-16, respectively. Since the above calculation is a function of modulo 16, all combinations of input and output line card numbers will result in the slot being a number between 1 and 16. Each of the cells in RAM 66 (buffer 42) is initially calculated to be one of 16 (1-16) slots. Each cell is then assigned a packet order number in the map of RAM 70 (position in buffer 48) in accordance with the following calculation:

$$\text{packet \#} = \text{slot} + (\text{N} * \text{I})$$

where packet # is the packet order number, slot is the number assigned in accordance with the preceding calculation, N is the number of the modulo function of the preceding calculation, and I is a positive integer beginning at zero which represents the previous number of cells which have been assigned the same slot number.

Examples of the ordering of packets to minimize collisions follows in Table 1.

Table 1

| input LC # | output LC # | slot | I | packet # |
|------------|-------------|------|-----|------------------|
| 1 | 7 | 8 | 0 | 8 |
| 2 | 7 | 9 | 0 | 9 |
| 2 | 6 | 8 | 1 | 24 (8+16) |
| ... | ... | ... | ... | ... |
| 15 | 9 | 8 | 10 | 168 (8+160) |
| 16 | 8 | 8 | 11 | >180 (8+176=184) |

The first two rows in Table 1 illustrate the benefit of the slot calculation by illustrating that the same output line card destination (7) for packets to be transmitted by the first and second line cards are assigned to packets 8 and 9, respectively.

Thus, a collision conflict at output line card 7 is avoided. The third row illustrates the assignment of packets where a prior slot has been previously utilized, i.e. $I > 0$. The penultimate row illustrates the assignment of slot 8 where $I = 10$ resulting in a packet assignment of 168. The last row illustrates an example in which the calculation for assigning the packet number would result in the assignment of a number greater than 180. The calculation for assigning packet numbers is followed where the resulting number is less than or equal to 180. Where the packet number calculation results in a number greater than 180, a remaining available packet slot is selected based on the first non-assigned packet slot starting at packet slot 1. Alternatively, any non-assigned packet slot could be randomly selected.

Where the packet number calculation results in a number of 180 or less, such packet transmissions will not collide with other packets. However, when the packet calculation exceeds 180, it will be apparent that it is possible that the selected packet may collide with another packet transmitted from another line card having the same output line card destination. In normal traffic loading, the embodiment in accordance with the present invention will assign a majority of the output packet numbers for each of the line cards in accordance with the calculation. This minimizes the statistical likelihood of collisions as contrasted with a system in which only random selection is utilized.

Another aspect in accordance with an embodiment of the present invention resides in the detection and control of cells which would result in excessive collisions. Table 2 below illustrates the storage of packet assignment information on a per-packet basis.

Table 2

| packet # | input LC 1 | input LC 2 | ... | input LC 16 |
|----------|------------|------------|-----|-------------|
| 1 | 7 | | | |
| 2 | | 7 | | |
| 3 | | | | |
| ... | ... | ... | ... | ... |
| 180 | | | | |

5 The information stored in Table 2 shows, for each packet number and each input line card, the corresponding output destination line card assignment. For example, in packet 1, the input line card 1 has assigned a packet having a destination output line card 7. For packet number 2, for input line card 2, an assignment has also been
 10 made to transmit a packet to output line card 7. This table is based upon the assignment of packets as determined in accordance with Table 1 in the above-described calculations.

Table 3 below illustrates for each packet number and for each output destination line card, the number of colliding packets which
 15 will be received by the respective output line card after each of the 180 packet transmissions. The information contained in Table 3 can be derived from Table 2 by summing across the rows in Table 2 to determine where there are more than one entry in the row for the same output line card number. For example, if the number 2 occurred
 20 twice in packet number 1 in Table 2, this would indicate that at least two input line cards would be concurrently transmitting a packet having a destination of output line card 2. This would correspond to a collision of two packets for delivery during the same transmission time. This result is shown in Table 3 under the
 25 output line card 2 by indicating that two collisions will occur during the transmission of the first packet in the sequence.

Table 3

| packet # | output LC 1 | output LC 2 | ... | output LC 16 |
|----------|-------------|-------------|-----|--------------|
| 1 | | 2 | | |
| 2 | 3 | | | |
| 3 | 2 | | | |
| ... | ... | ... | ... | ... |
| 180 | | | | 3 |

5 Table 4 below represents the cumulative number of packets queued for output to each output line card. The information contained in this table can be derived from Table 3.

Table 4

10

| packet # | output LC 1 | output LC 2 | ... | output LC 16 |
|----------|-------------|-------------|-----|--------------|
| 1 | | 1 | | |
| 2 | 2 | | | |
| 3 | 3 | | | |
| ... | ... | ... | ... | ... |
| 180 | | | | 2 |
| total | 3 | 1 | | 2 |

15 For example, consider packet number 1, output line card 2 in Table 4 has an entry of "1" which equals the corresponding matrix position in Table 3 minus one. One is subtracted on each occurrence of the total number of colliding cells since one of the colliding cells will be delivered during each of the packet number slots. In Table 4, packet number 3, output line card 1 contains a value of 3 which is calculated from Table 3. For packet number 2, output line card 1 in Table 3 has three collisions and for packet number 3, has 20 two collisions. In Table 4, for output line card 1, packet number 2 shows a value of 2 (3 minus 1) and in packet number 3, shows a value of 3 (2 plus (2 minus 1)). For each of the respective output line cards and for each packet number, a cumulative value is calculated

representing the number of packets to be held in storage or queued awaiting delivery to the corresponding output line card based on each packet transmission. The total in Table 4 represents the final number of remaining packets, if any, which would be held in a queue still to be delivered to the corresponding output line card after all 180 packets in a frame had been sent to the ATM network. This number can be visualized by referring to FIG. 4. Each of the output queues in would contain the number indicated in the last row of FIG. 4 for each corresponding line card.

The largest number of cells contained in any one of the output queues following the transmission of 180 packets represents the maximum time delay or skew which would be required in order to complete the transmission of all packets. A finite time can be associated with each packet stored in an output queue and hence, depending upon the time constraints of the synchronous portion of the system, a predetermined threshold can be calculated of the maximum number of packets which can be held in an output queue in order to still satisfy the timing requirements as constrained by the synchronous network. Thus, a number greater than the maximum number of permitted cells to be held in the output queue would represent a time delay not permitted in order to maintain proper operation of the synchronous network.

Not all of the time slices associated with a given packet may be occupied, i.e. less than 48 telephone calls may be carried by a packet so that some of the 48 bytes of the packet are not used. Additional telephone calls are preferably handled without creating a new packet (cell) by using unallocated bytes in an existing cell having the same destination as needed for the new call, thereby only increasing its payload to carry the additional traffic. This

assignment technique can be implemented by maintaining a table where each existing packet's destination and bytes used are stored. Such a table is consulted when new calls are requested to determine if an existing packet is available to handle a new call. A calculation is
5 made in accordance with the exemplary embodiment of the present invention to generate a new packet (a packet number assignment is tentatively made for the new packet) only if an additional byte in an existing cell with the proper destination is unavailable to accommodate the new call.

10 A maximum payload of time slices is carried by 180 packets, however, a fewer number of required time slices may result in the transmission of a fewer number of packets. Where a new packet is required to be originated to handle a new call request, and where the tentative packet number assignment results in a collision that
15 increases the number of packets which will be held in an output queue, the request to establish the new call can be denied in accordance with a preferred embodiment of the present invention to prevent an excessive timing skew or degradation of performance. Thus, call requests which would result in the generation of an
20 additional packet that would in turn lead to an excessive number of packets held in an output queue, will be denied. The denial of this request is communicated by the controller 23 back to the originating central office which in turn provides the requesting subscriber with an indication such as an "all trunks busy" audible signal.

25 Various modifications and alternative embodiments can be utilized within the scope of the present invention. For example, the functions of the line cards of the illustrative embodiment could be incorporated in other apparatus, e.g. either the switch modules or the ATM fabric could include such functions.

Alternative packet formats may contain a varying number of payloads corresponding to one or a plurality of calls.

Although an embodiment of the present invention has been described above and shown in the drawings, the scope of the
5 invention is defined by the claims which follow.

I Claim:

1. A method in a telecommunication system in which input and output line cards provide an interface between synchronous communication channels that are interconnected by an asynchronous network, the method comprising the steps of:

generating and storing packets by each input line card during a predetermined time frame, each packet carrying information received from the synchronous communication channels and having an output line card destination;

selecting a sequence for the sequential transmission of the packets by each input line card to the asynchronous network;

determining, based on the sequence of packets selected by the selecting step, a cumulative number of packets that will be stored for delivery to each output line card due to collisions that would occur at each output line card during the predetermined time frame, a collision being the substantially concurrent arrival at an output line card of more than one packet;

comparing the cumulative number for each output line card with a predetermined delay threshold;

reducing the number of packets that will be transmitted to an output line card if the cumulative number exceeds the predetermined delay threshold in order to maintain time delay within a predetermined limit.

2. The method according to claim 1 wherein the reducing step includes the step of denying new call requests associated with a packet stored for delivery to an output line card in which the corresponding cumulative number exceeds the predetermined delay threshold.

3. The method according to claim 1 further comprising the step of accumulating, for each output line card for each transmission sequence, a first number of packets, wherein the first number represents the number of collisions for each output line card
5 on a per transmission sequence basis.

4. The method according to claim 3 further comprising the steps of storing the first numbers for each output line card for each transmission sequence, and updating the stored first numbers based on new first numbers calculated for each subsequent time
10 frame, whereby the stored first numbers provide a running account of the number of packets stored for delivery to each output line card for each transmission sequence.

5. The method according to claim 1 wherein each packet carries telephone calls, the method of claim 1 further comprising
15 the steps of determining the output line card that would be a first destination associated with a request for a new call, determining if an existing packet having the first destination has capacity to carry another call, and if the existing packet with a first destination has said capacity, assigning the new call to said
20 existing packet.

6. The method according to claim 5 further comprising the step of creating a new packet to carry a new call associated with said first destination only if no existing packet having the first destination has capacity to carry the new call.

25 7. In a telecommunication system in which communication channels are interconnected by an asynchronous network, a method comprising the steps of:

storing packets during a predetermined time frame, each packet carrying information received from the communication channels, the asynchronous network having inputs and outputs;

selecting a sequence for the sequential transmission of the
5 packets on each input of the asynchronous network;

determining, based on the sequence of packets selected by the selecting step, a number of packets that will be stored for delivery to each output due to collisions that would occur at each output during the predetermined time frame, a collision being the arrival
10 at an output of more than one packet during one time frame;

comparing said number for each output with a predetermined delay threshold;

reducing the number of packets that will be transmitted to an output if the number exceeds the predetermined delay threshold in
15 order to maintain time delay within a predetermined limit.

8. The method according to claim 7 wherein the reducing step includes the step of denying new call requests associated with a packet stored for delivery to an output in which the corresponding number exceeds the predetermined delay threshold.

20 9. The method according to claim 7 further comprising the step of accumulating, for each output for each transmission sequence, a first number of packets, wherein the first number represents the number of collisions for each output on a per transmission sequence basis.

25 10. The method according to claim 9 further comprising the steps storing the first numbers for each output for each transmission sequence, and updating the stored first numbers based on new first numbers calculated for each subsequent time frame, whereby the stored first numbers provide a running account of the

number of packets stored for delivery to each output for each transmission sequence.

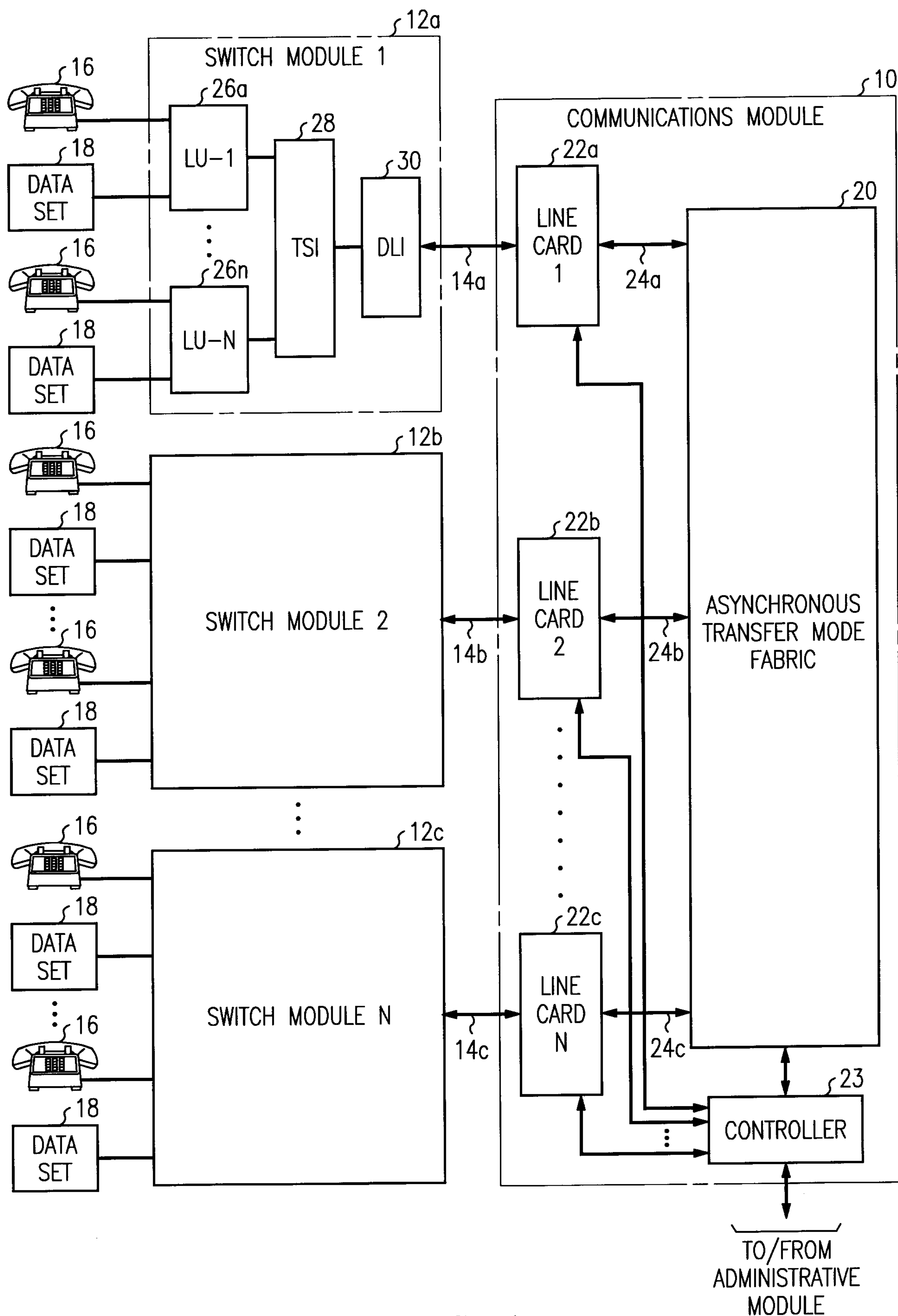


FIG. 1

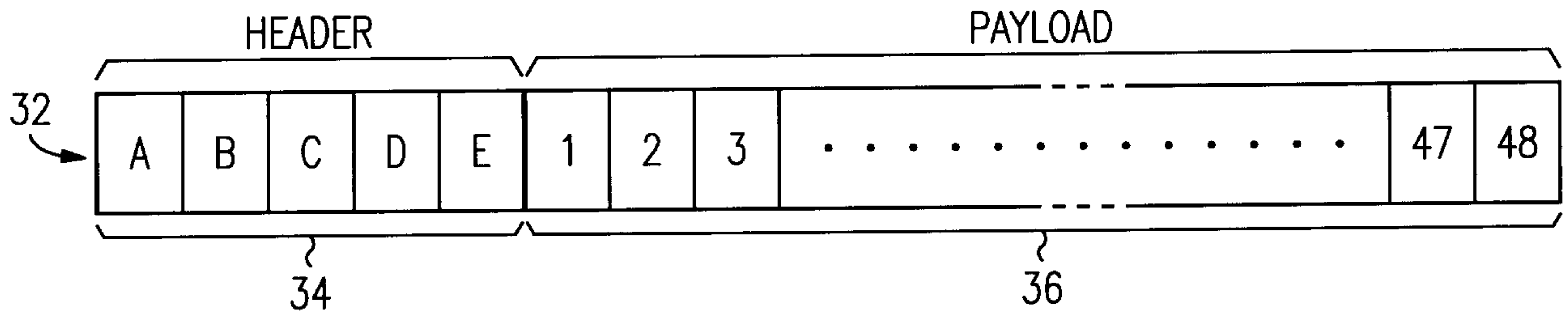


FIG. 2

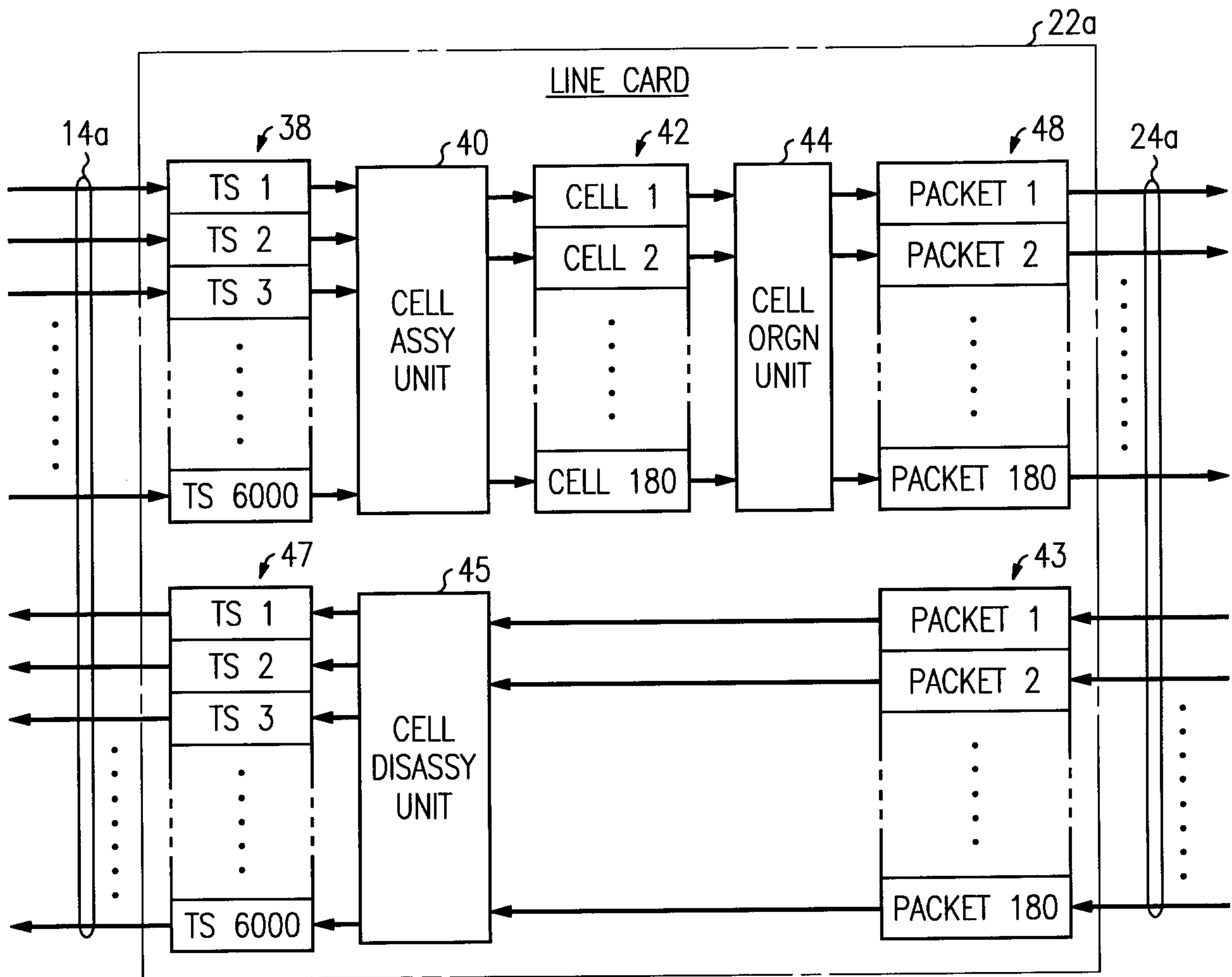


FIG. 3

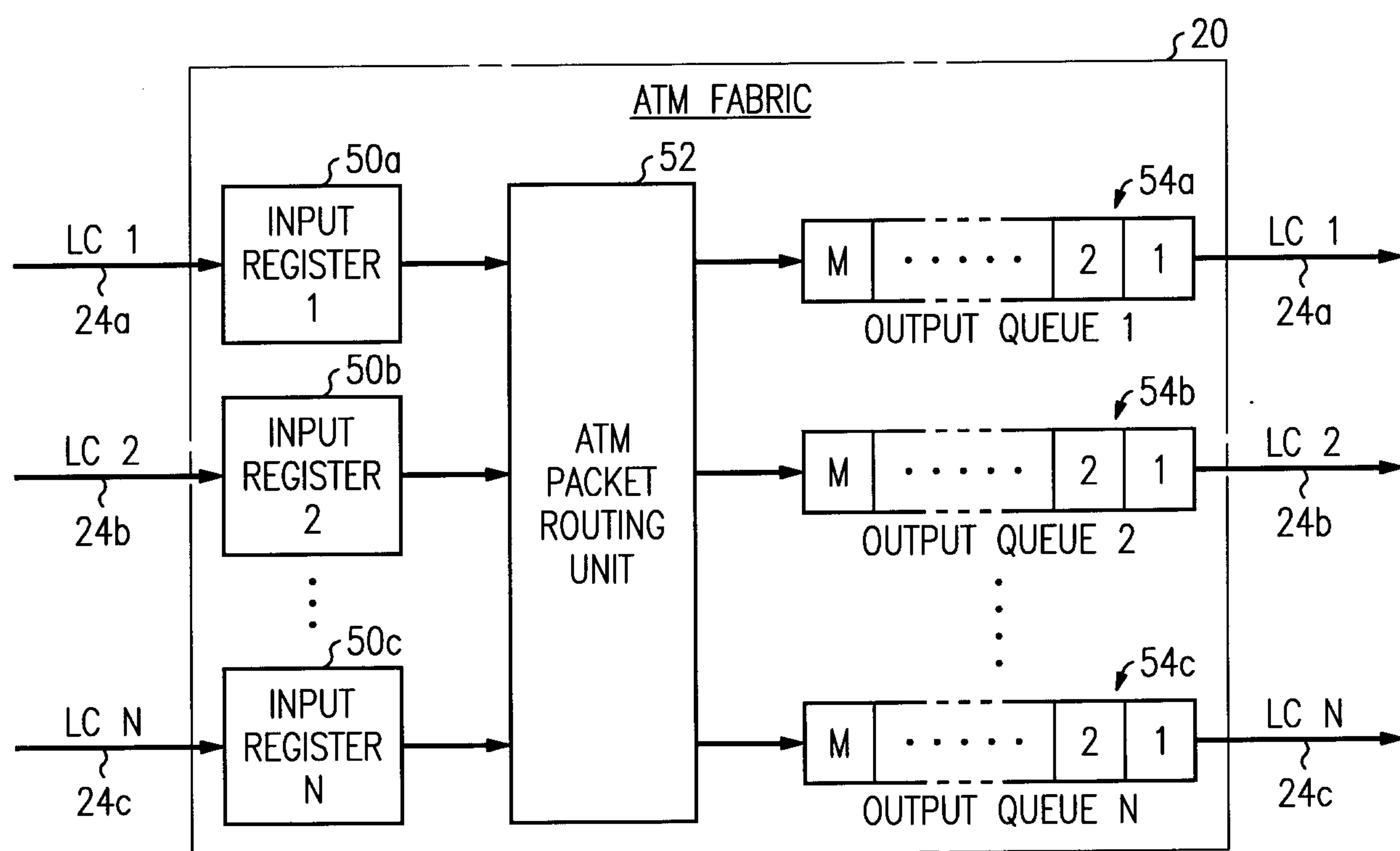


FIG. 4

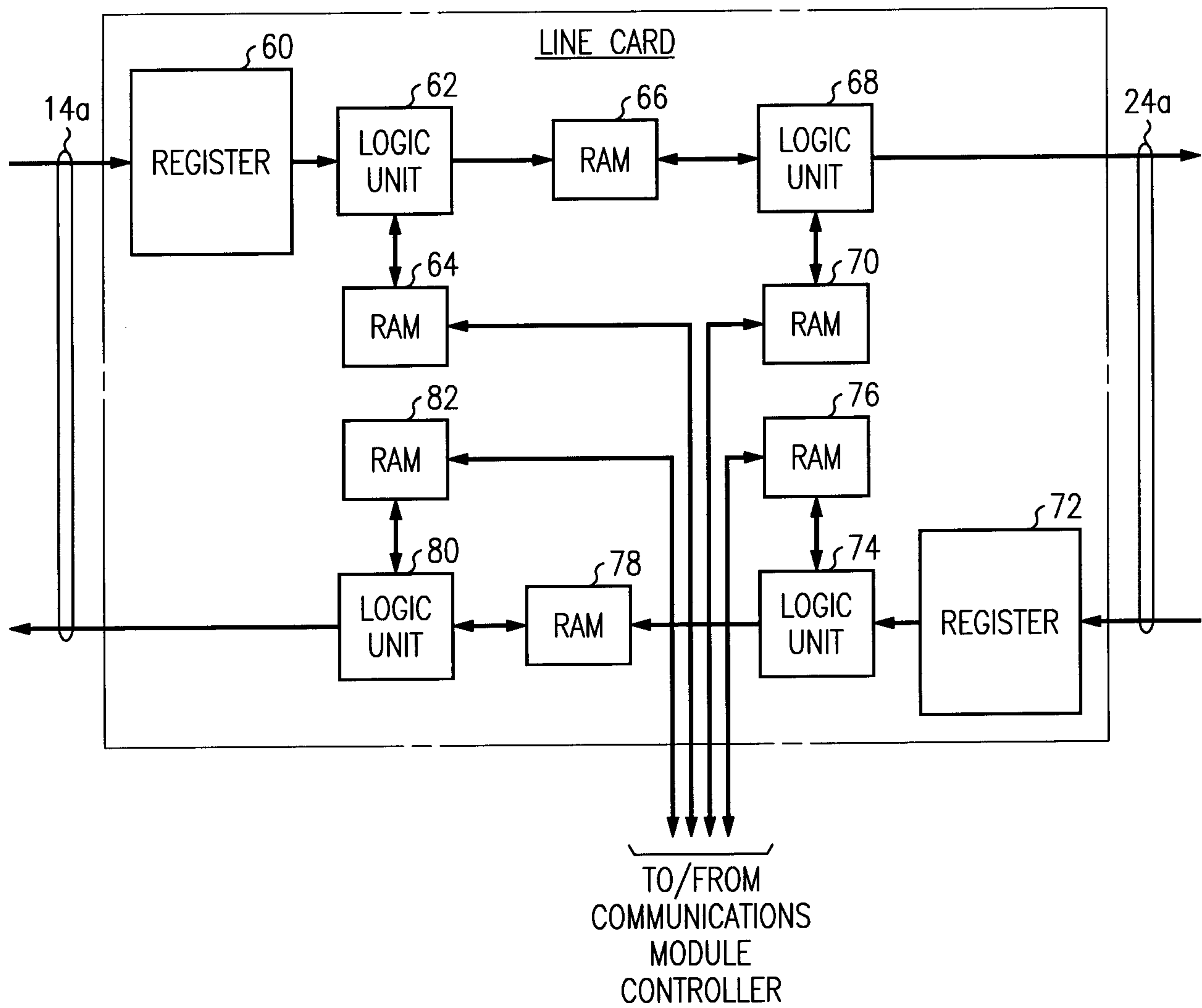
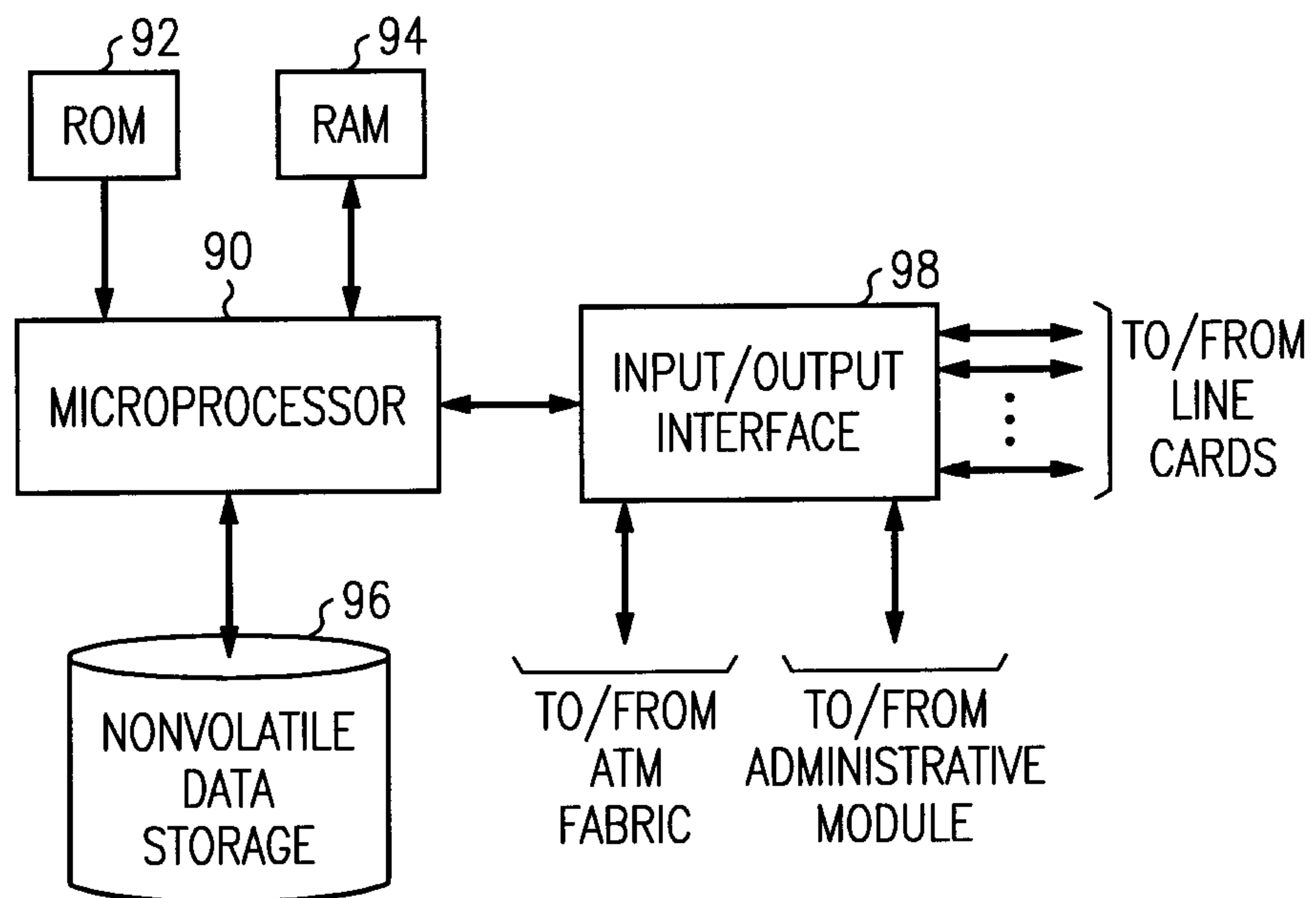


FIG. 5

*FIG. 6*