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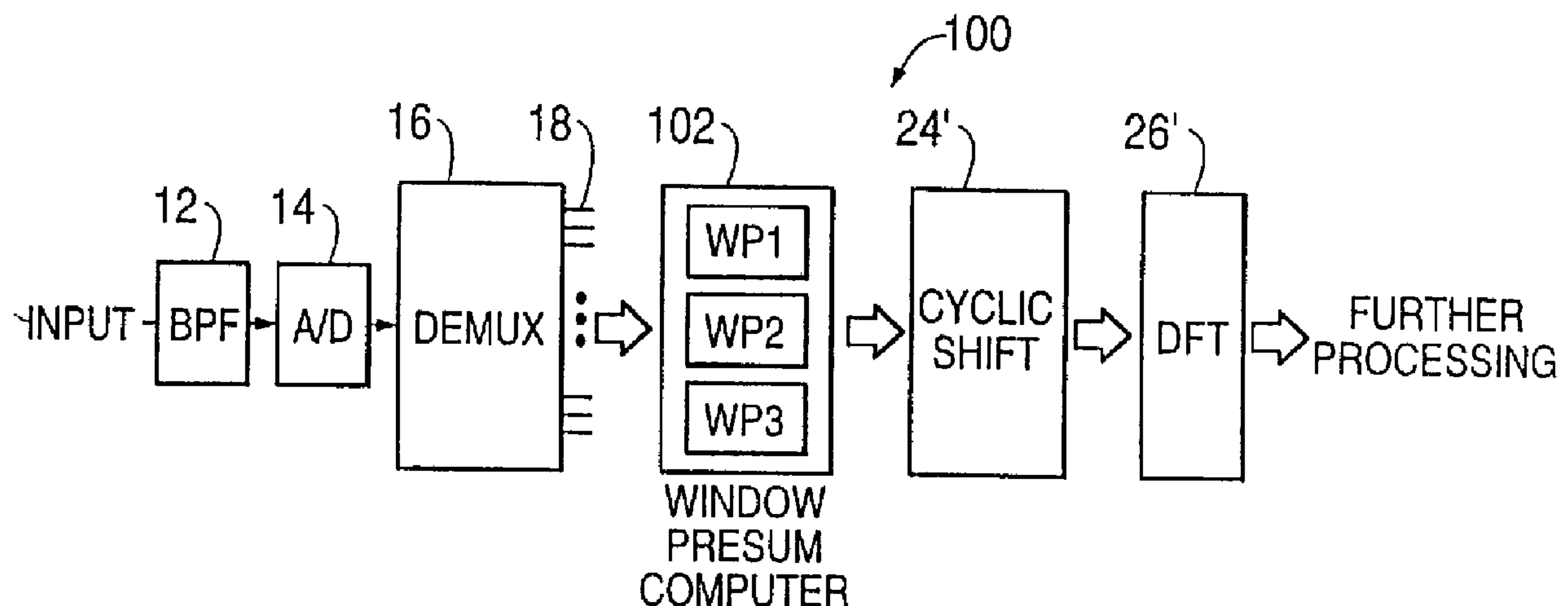
(72) Inventeurs/Inventors:
CASO, GREGORY S., US;
MORETTI, VINCENT C., US

(73) Propriétaire/Owner:
NORTHROP GRUMMAN CORPORATION, US

(74) Agent: GOWLING LAFLEUR HENDERSON LLP

(54) Titre : DISPOSITIF NUMERIQUE DE DECOUPAGE EN CANAUX AYANT UNE ARCHITECTURE EFFICIENTE POUR EFFECTUER LA PRESOMMATION EN FENETRE ET UTILISANT L'ARITHMETIQUE REPARTIE AFIN D'EFFECTUER LA PRESOMMATION EN FENETRE EN UN CYCLE D'HORLOGE

(54) Title: DIGITAL CHANNELIZER HAVING EFFICIENT ARCHITECTURE FOR WINDOW PRESUM USING DISTRIBUTED ARITHMETIC FOR PROVIDING WINDOW PRESUM CALCULATIONS IN ONE CLOCK CYCLE



(57) Abrégé/Abstract:

The invention is a digital channelizer and a process for dividing input bandwidth into at least some of N channels. A digital channelizer which divides an input bandwidth into at least some of N channels in accordance with the invention includes an analog to digital converter (14) which encodes the input bandwidth into a serial digital data stream of words; a demultiplexer (16), coupled to the analog to digital converter, which divides the serial digital data stream into parallel data streams of data words; a window presum (102), coupled to the parallel streams of data words, responsive to a window presum function, storing L data words in a window containing R subparts and outputting N outputs which are each respectively capable of being processed into a different one of the N channels, each of the N outputs being a product of R words each respectively from a different one of the R subparts and a coefficient of the window presum function preassigned to each of the R words from a different one of the R subparts with $R=L/N$ and each of the N outputs being calculated using distributed arithmetic during a single clock cycle; and a discrete Fourier transform apparatus (26'), coupled to the N outputs of the window presum, having N inputs and which performs a discrete Fourier transform on the N inputs to output at least some of the N channels.

DIGITAL CHANNELIZER HAVING EFFICIENT
ARCHITECTURE FOR WINDOW PRESUM USING
DISTRIBUTED ARITHMETIC FOR PROVIDING WINDOW
PRESUM CALCULATIONS IN ONE CLOCK CYCLE

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Abstract of the Disclosure

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10 N channels. A digital channelizer which divides an input bandwidth into at least some of N channels in accordance with the invention includes an analog to digital converter (14) which encodes the input bandwidth into a serial digital data stream of words;
15 a demultiplexer (16), coupled to the analog to digital converter, which divides the serial digital data stream into parallel data streams of data words; a window presum (102), coupled to the parallel streams of data words, responsive to a window presum function, storing L data words in a window
20 containing R subparts and outputting N outputs which are each respectively capable of being processed into a different one of the N channels, each of the N outputs being a product of R words each respectively from a different one of the R subparts and a coefficient of the
25 window presum function preassigned to each of the R words from a different one of the R subparts with $R=L/N$ and each of the N outputs being calculated using distributed arithmetic during a single clock cycle; and a discrete

Fourier transform apparatus (26'), coupled to the N outputs of the window presum, having N inputs and which performs a discrete Fourier transform on the N inputs to output at least some of the N channels.

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CROSS-REFERENCE TO RELATED APPLICATIONS

10 Reference is made to the following patent applications
which are filed on even date herewith:

(1) U.S. Patent No. 6,351,759

(Attorney Docket: TRW 22-0037/ATS&K 199.36554X00), entitled

15 "Digital Channelizer Having Efficient Architecture For
Discrete Fourier Transformation and Operation Thereof";

(2) U.S. Patent No. 6,349,118

(Attorney Docket: TRW 22-0038/ATS&K 199.36555X00), entitled "Digital Channelizer Having Efficient Architecture For Cyclic Shifting and Method of Operation Thereof";

5 (3) U.S. Patent No. 6,393,066

(Attorney Docket: TRW 22-0045/ATS&K 199.36557X00), entitled "Efficient Digital Channelizer System and Method of Operation Thereof"; and

(4) U.S. Patent No. 6,327,603

10 (Attorney Docket: TRW 22-0058/ATS&K 199.36574X00), entitled "Digital Channelizer Having Efficient Architecture For Presum Discrete Fourier Transformation Selectively of Real or Complex Data and Method of Operation Thereof".

15

BACKGROUND OF THE INVENTION

Field of the Invention

The present invention relates to filters for dividing an input bandwidth into a plurality of channels and more particularly, to a digital channelizer for satellite
20 communication applications using discrete Fourier transformation (DFT) to divide the input bandwidth into channels.

Description of the Prior Art

25 Digital channelizers in satellite communication systems have several design constraints. High computation

complexity is required which requires highly complex integrated circuit logic function and interconnections. High power consumption by integrated circuits can lead to high operating temperatures which could contribute to channelizer malfunction or failure. The system clock rate is required to be sufficiently high to support a high data throughput but should be as low as possible to lessen power consumption which contributes to the aforementioned possible high operating temperatures. The power consumption of a digital channelizer is proportional to the clock rate and the type of integrated circuits which implement the required high computational complexity. Furthermore, excess hardware can interfere with processing efficiency and be a source of potential malfunction.

Fig. 1 illustrates a block diagram of a prior art digital channelizer 10 which functions as a down converter and filter which divides a wideband input bandwidth into a plurality of equally spaced channels. The channelizer 10 is representative of channelizers using DFT which have been described in the literature. See Multirate Digital Signal Processing, published in 1983 by Prentice Hall, Englewood Cliffs, New Jersey, written by Crochiere and Rabiner. Such systems have applications in wideband satellite communication systems.

The INPUT signal is applied to a bandpass filter 12 which passes a selected wide bandwidth for division into N equally spaced channels each of a narrower bandwidth. For example, a wideband signal of 320 MHz. may be passed by the bandpass filter 12 for division into sixteen 20 MHz. wide channels. The bandpass filtered signal is applied to analog to digital converter 14 which samples the bandpass filtered signal. A representative frequency spectrum resultant from sampling is described further below in conjunction with Figs. 3A and 3B. Each sample is comprised of a multiple bit word. A serial stream of multiple bit words is outputted by the analog to digital converter 14 as an input to demultiplexer 16 which produces D outputs 18. The variable D may equal the variable M known as the decimation rate. The demultiplexer 16 functions as a multiple tapped delay line with each parallel output being outputted from a different tap of the delay line. The D outputs 18 of the demultiplexer are applied to window presum computer 20. The window presum computer 20 functions in a well-known manner to process the sequence of words within a window of L words by multiplying each corresponding word in a plurality of equal subparts R of the window containing N words by a window presum function coefficient and summing the resultant multiplication products to produce the sum of the multiplication products where $R=L/N$. The individual summed multiplication products, after further

processing including DFT as described below, are outputted as individual ones of the N channels. The number of outputted channels may be selected to be less than N.

For example, a sequence of ninety-six words, outputted
5 by the demultiplexer 16, is broken up into four subparts each containing twenty-four words. Each of R corresponding words, e.g. words 0, 24, 48 and 72, from a different subpart, are multiplied by their preassigned window presum function coefficient and summed to produce an output
10 summation which is subsequently processed into one of the N output channels. The window presum computer 20 has D inputs and N outputs. The relationship between M, D and N, which is the DFT size and the number of possible channels, affects the architecture of the window presum
15 computer 20. Words stored in a number of registers (not illustrated), e.g. words 0, 24, 48 and 72, equal to the R subparts in the window processed by the window presum computer 20, are summed after multiplication by their preassigned window presum function coefficient to produce
20 the output summation.

The window presum computer 20 has been implemented by the Assignee with parallel data processing paths using integrated circuits with M not being equal to N. The number of processing paths I used by the Assignee to perform
25 parallel data processing satisfies the relationship I equals

the greatest common divisor of N and M which is expressed hereafter as $GCD(N,M)$.

The N outputs from the window presum computer 20 are applied to a cyclic shift 24, which provides phase
5 adjustment, for processing into each channel by DFT. The phase produced by the cyclic shift 24 is applied to the resultant N word outputs from the window presum computer by a calculated number of shifts. The operation of cyclic shifting is well known and is, for example, described in the
10 aforementioned publication on pp. 320-323. The number of shifts of the output words of the window presum computer 20 by the cyclic shift 24 is determined by computing the value of the relationship $mM \bmod N$ or $-mM \bmod N$. The variable m is an output index variable which ranges from
15 zero upward to positive integers. The output from the cyclic shift 24, which has N channels, is applied to a discrete Fourier transform apparatus 26 having N inputs which transforms the output from the cyclic shift 24 into the N output channels.

20 Fig. 2 illustrates a conceptual block diagram of the window presum algorithm which represents the window presum processing performed by the system of Fig. 1. The sampled output of L individual words is shifted into a shift register which stores the sequential words outputted by the
25 analog to digital converter 14. The input data are shifted into the shift register, which has a number of subparts R,

e.g. 4 in the above example. The number of words per subpart (the DFT size) is equal to the number N of output channels. The shift register has an analysis window L words long which is R times the size N of the discrete Fourier transform. The sum of the individual R subparts contains the words which are further processed to individual channels by DFT. The data in the shift register are weighed with a time reverse window according to equation 7.70 on page 317 of the aforementioned publication to produce a windowed sequence as illustrated. The sequence is processed as blocks of samples starting at $r=0$ which are time aliased. The resultant summation is processed by a cyclic shift 24 through a number of shifts equal to $mM \bmod N$ or $-mM \bmod N$ and is applied to discrete Fourier transform 26.

The analog to digital converter 14 of Fig. 1, in accordance with digital sampling theory, produces a spectrum of frequency domain signals centered about zero frequency as illustrated in Figs. 3A and 3B which respectively illustrate groups of twelve and twenty four frequency domain signals. The sampling frequency of f_s for real signals produces counterpart frequency domain signals centered about the zero frequency extending to $f_s/2$ in both the positive (real) and negative (conjugate) frequencies. The positive frequencies may be expressed mathematically as $a=x+iy$ and the negative frequencies may be expressed as $a=x-iy$ with corresponding

positive and negative frequencies being conjugates of each other. Also, in accordance with digital sampling theory, the frequency domain signals of Figs. 3A and 3B repeat periodically with a period f_s for successively higher positive frequencies and successively lower negative frequencies. These upper repeating frequency domain signals have been omitted from the illustration. The counterpart frequency domain signals of Fig. 3A are 1 and 11, 2 and 10, 3 and 9, 4 and 8, and 5 and 7, and the counterpart channels of Fig. 3B are 1 and 23, 2 and 22, 3 and 21, 4 and 20, 5 and 19, 6 and 18, 7 and 17, 8 and 16, 9 and 15, 10 and 14 and 11 and 13. Frequency domain signals 0 and 6 in Fig. 3A and 0 and 12 in Fig. 3B do not have counterparts. The information of each frequency domain signal is transformed to its counterpart conjugate by a sign reversal of the imaginary term iy .

Fig. 4A illustrates a diagram of the window presum function of a window containing forty eight (L) real words having four (R) subparts each containing 12(N) real words and Fig. 4B illustrates a window presum function of a window containing twenty four complex words containing an imaginary part identified by the letter "i" following a number and a real part identified by the letter "r" following a number.

The window presums are identical except that the window presum function of Fig. 4B has half as many words in view of each word having a real and a imaginary part. Complex data

in Fig. 4B is, for example, obtained when conversion of a spread spectrum transmission is down converted in the tuner of the receiver. The individual words of Figs. 4A and 4B are multiplied by their preassigned window presum
5 function coefficients and subsequently summed with other products of corresponding words from other subparts to produce the output of the window presum computation which is subsequently processed into N channels by DFT.

Corresponding words in each subpart R are summed to
10 produce a number of sums equal to the number of words per window, e.g. P0-P11 or P0_r-P5_i. The summations P0-P11 and P0_r-P5_i; are processed with the window multiplication process by the use of stored coefficients to compute a value of y for each of the N channels which represents the
15 summation of products which is applied to the cyclic shift 24.

The summation process, when the rate of decimation M is equal to the number of channels N, may be implemented efficiently with an array of registers storing the
20 individual words identified in the vertical columns of Fig. 4A to produce the outputs P0-P11.

The summation of the products of corresponding individual words (e.g. 0, 12, 24 and 36 in Fig. 4A or word parts 0r, 6r, 12r and 18r in Fig. 4B) times their
25 preassigned window presum function coefficients may be implemented in a number of ways. One method is illustrated

in Fig. 5 which has the disadvantage of using substantial hardware requiring multipliers 40 and summation calculation 42. The number of multipliers 40 is equal to R and the number of adders in summation calculation 42 is equal to (R-1) in the worst case. This method computes the output summation y in one clock cycle (pipelined). The word values x_0 , x_1 , x_2 and x_3 represent corresponding word values from each of the R subparts of the window which is L words in length, e.g. words 0, 12, 24 and 36 in Fig. 4A or real word parts 0r, 6r, 12r and 18r in Fig. 4B. While this implementation for computing the summation y is computationally fast, it has the disadvantage of requiring a substantial number of gates, other hardware and interconnections which have the disadvantages described above especially in an environment involving satellites.

Fig. 6 illustrates a block diagram of a finite impulse response filter proposed in "Applications of Distributed Arithmetic to Digital Signal Processing: A Tutorial Review", by Stanley A. White, in IEEE ASSP Magazine, July 1989, pp. 1-19. The illustrated filter computes the summation y with an input of four eight bit words x_0 , x_1 , x_2 , x_3 in a serial fashion requiring eight clock cycles to process eight bit words. The overall operation is to compute individual products of input words x_0 , x_1 , x_2 and x_3 and their multiplying window coefficients W_0 , W_1 , W_2 , and W_3 to generate a sum $\sum_i x_i W_i$. The individual products are not

computed in isolation and then added. Each bit of every word determines whether to add or subtract a multiple of their respective window coefficients, and all the bits at the same position within the words are processed at the same time.

5 The result is that the overall sum of products is generated not by simply summing up products, but by summing multiples of different sum combinations of window coefficients and their negations. The bit select 50 selects bit slices from the eight bit words, e.g. the bits of the least significant
10 bit LSB in ascending order to the most significant bit MSB.

The logic circuit 52 exploits symmetry in the DA ROM 54 to eliminate half of the values that need to be stored for distributed arithmetic to operate correctly as a process. The process of reducing the number of bits is described on
15 pages 5 et. seq. in the aforementioned paper. The DA (distributed arithmetic) ROM 54 stores all the possible sum combinations of window coefficients and their negations.

The bit slice from the input words act as the address into the DA ROM 54 to choose the proper sum of coefficients. The
20 shifter 56 outputs the proper multiple by a power of two of the DA ROM output to the adder 58. The output of shifter 56 is applied to a summation calculation 58 which sums multiples of different sum combinations of window coefficients and their negations. Feedback from register 60
25 provides the current sum which is summed with the new sum by summation calculation 58 for each successive bit slice.

The serial implementation of Fig. 6 suffers from the disadvantage of requiring a high clock rate to compute the summation y for large data words. A high clock rate in satellite applications requires high energy consumption which can cause heating in integrated circuits and for systems requiring high data rates, represents a potential speed processing barrier. Processing one bit at a time has undesirable latency.

Fig. 7 illustrates an implementation of distributed arithmetic used by the Assignee in the window presum computer 20 to sum the products of the corresponding words and their preassigned window presum function coefficient. This system performs processing similar to Fig. 6, except that three clock cycles are used to respectively process four bit nibbles, which are inputted from twelve bit words x_0, x_1, x_2, x_3 . The processing of the four bit nibbles is in parallel but otherwise is analogous to Fig. 5. The summation calculation 62 sums for each of the three clock cycles the outputs from the shifters 56. Register 64 stores the resultant summation outputted by summation calculation 62 and feeds the summation back to the summation calculation to sum the current summation with the summation of the next clock cycle processing.

This implementation of distributed arithmetic has disadvantages for high word processing throughputs. It requires a higher clock rate to process the four bit nibbles

in three cycles for each word which increases energy consumption when compared to processing all the bits of the word in one cycle. The clock rate required to perform three processing cycles per word could, for certain satellite
5 processing applications, limit the word processing throughput below that which is required for a desired system performance.

Window presum computers 20 are well known which utilize an array of registers to store words x_0 , x_1 , x_2 , x_3 processed
10 in accordance with the aforementioned processes for computing the summation y . Some applications have their decimation rate M equal to the DFT size N and the number of channels which are outputted. However, the Assignee has implemented a window presum calculator 20 having the
15 decimation rate M not equal to the DFT size N . These systems provide the corresponding input words from each of the aforementioned subparts of the window, e.g. words 0, 12, 24 and 36 from Fig. 4A or corresponding parts of words from Fig. 4B from storage in registers for multiplication and
20 summation to produce the output y of the window presum computer for each of the N channels.

Discrete Fourier transforms are well known. A stand alone discrete Fourier transform apparatus of N inputs provides a frequency response at N outputs at specific
25 equidistant frequencies. The N inputs are time domain signals and the N outputs are frequency domain signals

determined at singular frequencies. The discrete Fourier transform apparatus itself can extract channel information, although it samples at only a singular frequency that represents the channel information.

5 In a channelizer, an output does provide a single extracted channel. The N outputs correspond to N time domain signals that each contain information from one of N frequency bands (equal in bandwidth) that divide the input signal frequency spectrum. This does not exclude the input
10 signal from having more or fewer channels than N. The spectrum is simply divided into N frequency bands. All N inputs are required in the computation for every one of the N outputs and each of the N outputs represents the time-domain signal of one of the frequency bands dividing the
15 frequency spectrum.

Fig. 8 is a diagram representing a prior art discrete Fourier transform device 118 which converts the output of twelve cyclically shifted summations y , produced by the cyclic shift 24, into twelve frequency domain outputs.
20 Various algorithms are known for computing a DFT. The Winograd algorithm used for non-power of two discrete Fourier transforms is used in Fig. 8. The DFT apparatus 118 has twelve time domain inputs "in 0-in 11" each representing multiple bit words and twelve frequency domain outputs
25 "out 0-out 11" each representing multiple bit channel outputs 0-11. Some of the outputs, which are represented in

Fig. 8 as real numbers, in fact are complex numbers containing a real term and an imaginary term. The illustration of complex input words has been omitted in order to simplify illustration.

5 The discrete Fourier transform apparatus 118 includes an input discrete Fourier transform computation stage 120 comprised of six two point DFT's 122 of known construction each having a pair of time domain inputs and a pair of frequency domain outputs, an intermediate discrete Fourier
10 transform computation stage 124 comprised of four three point DFTs 126 of known construction each having three inputs and three outputs and an output discrete Fourier transform computation stage 128 comprised of six two point
15 DFTs 130 of known construction each having two inputs and producing two frequency domain outputs. The outputs of the two point DFTs 122 are inputs to the individual DFTs 126 of the intermediate discrete Fourier transform computation stage 124 and the outputs of the three point DFTs 126 of the intermediate discrete Fourier transform computation stage
20 are inputs to the individual discrete Fourier transforms 130 of the output discrete Fourier transform computation stage 128.

Fig. 9 illustrates a prior art pruned discrete Fourier transform apparatus 140 which is representative of
25 modifications performed by the Assignee to eliminate unnecessary DFTs 130 in the output discrete Fourier

transform computation stage 128 when all of the N possible frequency domain outputs, equal to the number of time domain inputs, are not needed for further processing. The discrete Fourier transform 118 of Fig. 8 has been modified in Fig. 9 to eliminate two output stages 130 in view of only frequency domain signals 2, 3, 4, 5 and 8, 9, 10 and 11 representative of the frequency domain signals of Fig. 3A being required for further processing. The two point discrete Fourier transforms 130 which produce frequency domain signals 0 and 6 and 1 and 7 have been eliminated. The pairs of frequency domain inputs 0 and 6 and 1 and 7 are not represented. Output 5 is a counterpart and conjugate of output 7. However, output 5 in Fig. 9 is not further processed downstream in place of output 7.

Fig. 10 illustrates another form of prior art discrete Fourier apparatus 300 having all possible frequency domain signals as outputs. The frequency domain input discrete Fourier computation stage 302 has three four point discrete Fourier transforms 304 of known construction which each have four time domain inputs and four outputs which respectively are coupled to an output discrete Fourier transform computation stage 304 having four three point DFTs 306. This structure does not have an intermediate discrete Fourier computation stage like Figs. 8 and 9.

Fig. 11 is a diagram representing a prior art discrete Fourier transform apparatus 400 which is twenty four time

domain inputs and all of the possible twenty four frequency domain outputs as outputs. The discrete Fourier transform has an input discrete Fourier computation stage 402 comprised of three eight point preweaves 404, three
5 intermediate discrete Fourier computation stages, 410, 412 and 414, respectively comprised of eight three point preweaves 416, a multiply stage and eight three point postweaves 418 and an output discrete Fourier computation stage 419 comprised of three eight point postweaves 420.

10

SUMMARY OF THE INVENTION

The present invention is a digital channelizer and a method which divides an input bandwidth, such as the wideband signal which is received by a satellite, into at
15 least some of N possible channels.

A digital channelizer in accordance with the invention has an efficient architecture, organization and movement of data from the window presum through the discrete Fourier transform device. The decimation rate M is not limited to
20 being equal to N. The use of distributed arithmetic reduces hardware required for the window presum operation compared to the prior art of Fig. 5. An efficient layout of window presum computations permits efficient cyclic shifts, which map directly into the discrete Fourier transforms. The
25 output discrete Fourier computation stage may be simplified when not all channels are required as outputs.

The digital channelizer includes a window presum computer having a modular processing architecture which transfers words stored within memory elements, which in the preferred form of the invention is a connected array of registers, in a systematic and periodic pattern to complete window presum computations during a single clock cycle. Minimizing the number of operations which must be performed to complete the window presum operation reduces power consumption and permits the system to operate at higher throughputs. Furthermore, a modular implementation of a window presum computer as parallel window presum circuits simplifies the memory structure of the registers in the window presum by permitting the same register array within an integrated circuit to be used for each of the parallel processing paths.

The window presum operation is performed in parallel in modular window presum circuits which efficiently map into the cyclic shift and discrete Fourier transform device which are also implemented as a modular architecture. The number of parallel paths, which is equal to the number of window presum circuits, may be determined by the value of the $\text{GCD}(N,M)$. When the number of window presum circuits is determined by the value of $\text{GCD}(N,M)$, hardware use is reduced in view of the processings from the window presum computations, cyclic shifting and DFT being efficiently mapped into a minimum amount of hardware. The processing

speed of each of the integrated circuits within the parallel processing paths is at a lower clock rate. Parallel processing within the window presum computer, cyclic shift and discrete Fourier transform apparatus permits slower, but
5 more power efficient integrated circuit technologies, such as CMOS, to be used to perform the required operations. Slower parallel operations lessen the generation of heat caused by high clock rates.

The window presum computer processes a block of words
10 produced by analog to digital conversion which have a data length L . The data sequence is windowed by multiplying each word by the preassigned window presum function coefficient which is chosen to provide the filter requirement as, for example, illustrated in Figs. 4A and 4B. Each of
15 R individual subparts of the window, which are N words long, are processed word by word to provide products of the word value times the preassigned window presum function coefficient. While the invention is not limited thereto, the preferred form of summing the individual products of the
20 words times the window presum function coefficients is with distributed arithmetic which calculates the summation of the products of words and their preassigned window presum coefficients within a single clock cycle.

A window presum in accordance with the invention is
25 responsive to parallel data streams of words which are used to produce N outputs for every M input words, which are

subsequently processed into N channels such that each channel data rate has been decimated by a factor of M from the original. The N outputs each are a function of a window function and a function of a plurality of inputs to the
5 window presum.

A digital channelizer which divides an input bandwidth into at least some of N channels in accordance with the invention includes an analog to digital converter which encodes the input bandwidth into a serial digital data
10 stream of words; a demultiplexer, coupled to the analog to digital converter, which divides the serial digital data stream into parallel data streams of data words; a window presum, coupled to the parallel streams of data words, responsive to a window presum function, storing L data words
15 in a window containing R subparts and outputting N outputs which are each respectively capable of being processed into a different one of the N channels, each of the N outputs being a product of R words each respectively from a different one of the R subparts and a coefficient of the
20 window presum function preassigned to each of the R words from a different one of the R subparts with $R=L/N$ and each of the N outputs being calculated using distributed arithmetic during a single clock cycle; and a discrete Fourier transform apparatus, coupled to the N outputs of the
25 window presum, having N inputs and which performs a discrete

Fourier transform on the N inputs to output at least some of the N channels.

A process for dividing an input bandwidth into at least some of N channels in accordance with the invention includes
5 encoding the input bandwidth into a serial data stream of words; dividing the serial data stream into parallel data streams of data words; processing the data words of the parallel data streams to produce products of data words and a window presum function which are summed to produce
10 N outputs, the window presum function having a window containing L data words divided into R subparts, each of the N outputs being a product of R samples each respectively from a different one of the R subparts and a coefficient of the window presum function preassigned to each of the R
15 words from a different one of the R subparts with $R = L/N$ and each of the N outputs being calculated using distributed arithmetic during a single clock cycle; and performing a discrete Fourier transform on N inputs, which are coupled to N outputs produced by processing the data words of the
20 parallel data streams to produce at least some of the N channels.

It should be understood that the invention is not limited to its elements as summarized above.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 illustrates a block diagram of a prior art digital channelizer.

Fig. 2 illustrates a processing diagram representative
5 of the processing of the prior art of Fig. 1.

Figs. 3A and 3B respectively illustrate frequency domain signals produced by digital sampling of wideband signals in accordance with the prior art.

Figs. 4A and 4B illustrate a prior art window presum
10 function containing respectively real and complex data words.

Fig. 5 illustrates a prior art hardware implementation for calculating a window presum for individual corresponding words in accordance with the window presum function of
15 Fig. 4A.

Fig. 6 illustrates a prior art hardware implementation for providing a finite impulse response filter function.

Fig. 7 illustrates a prior art hardware implementation for calculating a window presum of individual corresponding
20 data words in accordance with the window presum function of Fig. 4A.

Fig. 8 illustrates a first type of prior art discrete Fourier transform apparatus.

Fig. 9 illustrates a prior art pruning of the discrete
25 Fourier transform apparatus of Fig. 8.

Fig. 10 illustrates a second type of prior art discrete Fourier transform apparatus.

Fig. 11 illustrates a third type of prior art discrete Fourier transform apparatus.

5 Fig. 12 illustrates a block diagram of the present invention.

Fig. 13 illustrates a block diagram of a window presum computer in accordance with the present invention.

10 Figs. 14-16 illustrate an embodiment of the individual window presum circuits of the window presum computer of Fig. 13.

Fig. 17 illustrates a block diagram of a preferred form of distributed arithmetic used in the window presum computer of Figs. 12-16 with the practice of the present invention.

15 Fig. 18 illustrates a block diagram of a preferred embodiment of a cyclic shift in accordance with the present invention.

Fig. 19 is a table illustrating cyclic shifts of the cyclic shift of Fig. 18 required for the outputs of the
20 window presum circuits of Figs. 14-16 of the present invention.

Fig. 20 is a first embodiment of a pruned discrete Fourier transform apparatus in accordance with the present invention.

Fig. 21 is a second embodiment of a pruned discrete Fourier transform apparatus in accordance with the present invention.

Fig. 22 is a block diagram of an embodiment of
5 a discrete Fourier transform apparatus in accordance with the present invention which calculates the transformation of real or complex data inputs with a processing architecture which is configured to transform real or complex data in response to command.

10 Fig. 23 is a block diagram of a two-point discrete Fourier transform/recombine for processing real data.

Fig. 24 is a block diagram of a two-point discrete Fourier transform/recombine for recombining complex data.

Fig. 25 is a block diagram of multiplication by $-j$
15 (a twiddle factor) of an output of one of the three point discrete Fourier transforms of Fig. 22.

Fig. 26 is a block diagram of an embodiment of a two point discrete Fourier transform/recombine which performs the operations of Figs. 23-25.

20 Like reference numerals identify like parts throughout the drawings.

DESCRIPTION OF THE PREFERRED EMBODIMENTS

Fig. 12 illustrates a block diagram of a digital
25 channelizer 100 in accordance with the present invention. Differences between the embodiment 100 of Fig. 4 and the

prior art of Fig. 1 are that the window presum computer 102 is implemented as a plurality of parallel window presum circuits WP1, WP2 and WP3, which may vary in number in accordance with window presum processing requirements and the cyclic shift 24' and discrete Fourier transform apparatus 26' have been modified to a similar parallel architecture to provide high speed calculation with lower power consumption and simplified hardware requirements. The window presum circuits WP1, WP2 and WP3 each perform processing in parallel to implement the functions of the window presum computer 20. The number of window presum circuits I is preferably equal to $\text{GCD}(N,M)$, which for the example of $M=15$ and $N=24$ makes $I=3$, which results in an efficient deployment of hardware operating at a lower clock speed in the individual parallel processing paths of data in the window presum circuits WP1, WP2 and WP3, cyclic shift 24' and discrete Fourier transform apparatus 26' as described below which is important in satellite applications.

The window presum computer 102 implemented as a plurality of window presum circuits operating in parallel permits a lower clock rate in each window presum circuit. As a result, slower integrated circuit technologies may be used to implement all of the processing in the I parallel data processing paths extending from the window presum circuits WP1, WP2 and WP3 through the discrete Fourier

transform apparatus 26' as described below which consume less power, such as CMOS circuits.

Parallel processing achieves a high throughput data rate of window presum processing of words required to perform the complete window presum computer function in a wide band satellite application. In view of a preferred embodiment of the present invention being in a satellite, it is highly desirable to lessen heat generating energy consumption resultant from a high clock rate in the window presum computer 102, the cyclic shift 24' and discrete Fourier transform apparatus 26' which is accomplished by implementation of the window presum computer, cyclic shift and discrete Fourier transform apparatus in a parallel architecture in a low-power consuming technology (such as CMOS).

Fig. 13 illustrates a block diagram of the three window presum circuits WP1, WP2 and WP3 which comprise the window presum computer 102. The use of I parallel processing paths within the window presum computer 102, the cyclic shift 24' of Fig. 12 as described below in conjunction with Figs. 18 and 19 and the I parallel processing paths within the discrete Fourier transform apparatus 26' as described below in conjunction with Figs. 20-22 is a highly efficient hardware implementation. The number D of parallel paths of input data from the demultiplexer 16 (also equal to M the decimation rate in this instance) is 15 with input data

words "in_wp_0" thru "in_wp_14" each comprised of multiple bits being shifted into the window presum circuits 102 during one clock cycle. The decimation rate M is not necessarily equal to a number of branches in the

5 demultiplexer 16. Each of the window presum circuits 102 performs window presumming, described below in detail in conjunction with Figs. 14-17, to output for each clock cycle N output words "out_wp_0" thru "out_wp_23" which represent the N channels which are produced by the digital

10 channelizer or with a reduced number of channels produced by pruning the DFT structure as described below in conjunction with Fig. 21 which is used within a channelizer. Each of the output words $y(i)$, which are inputs to the cyclic shift $24'$, is a function of the window presum function

15 calculated by the sum of products of stored coefficients for each of the corresponding R input words with the corresponding input words respectively from the parallel inputs of one of the I groups of input words and summing those R words together.

20 Figs. 14-16 respectively illustrate an implementation of each of the individual window presum circuits WP1, WP2 and WP3 of Fig. 13. The window presum circuits 200 are identical, permitting their use in a modular architecture. Each window presum circuit has a number L' of registers

25 equal to the number of data points in each of the I groups wherein:

$$L' = \frac{L}{GCD(N, M)}$$

In the window presum circuits 200 in Figs. 14-16, input data words are shifted through registers which are connected
5 in a group of R registers to individual distributed arithmetic functions 202 (DA#1-DA#24) which process the input data words to compute the individual products of each data word and its stored window presum function coefficient followed by a summation of the products which is outputted
10 as the data words $y(i)$. The calculation of the summation $y(i)$ does not require the use of distributed arithmetic. However, preferably the calculation of the summation $y(i)$ is performed with distributed arithmetic. A preferred implementation of the calculation of $y(i)$ with
15 distributed arithmetic is described below in conjunction with Fig. 17 which requires only a single clock cycle to complete the calculation of each value of $y(i)$. The parallel processing of each of the window presum circuits lowers the clock rate permitting energy efficient integrated
20 circuit implementations such as CMOS to be used.

The individual numbers appearing inside of the rectangular boxes, representing registers, identify particular stored data words in a data sequence in a window which is L words long, e.g. 96. Each of the words may have
25 any number of bits. Additionally, buffering registers 204

are provided where time shifting for a clock cycle is necessary. The input words are identified by the legend "in_wp_" followed by a number and output words, which are outputted for each clock cycle, are identified by the legend
5 "out_wp_" followed by a number.

Each of the outputs $y(i)$ of the window presum circuits of Figs. 14-16 is a function of a window presum function and data words from a plurality of the parallel data streams which are inputs "in-wp" thereto that are the outputs from
10 the demultiplexer 18. The calculation of each output $y(i)$ is described in further detail below in conjunction with Fig. 17. For example, with reference to Fig. 14, input word "in_wp_12" is coupled to register 87 which is directly coupled to a distributed arithmetic calculator DA#2 and is
15 an input to register 72 which is directly coupled to DA#1. Furthermore, input word "in_wp_3" is directly coupled to register 93 which is coupled to register 78 which is coupled to register 63 which is coupled to register 48 which is directly coupled to DA#1. Therefore, the word output $y(i)$
20 "out_wp_0" is a function of "in_wp_12" and "in_wp_3", as well as being a function of input words "in_wp_9" and "in_wp_0". Furthermore, the number of R inputs respectively from R different registers coupled to each distributed arithmetic function or circuit 202 is equal to the number of
25 subparts of the window of L words processed by the window presum function which each contain N words. Similarly, the

outputs of the other distributed arithmetic functions or circuits 202 are a function of a plurality of input words respectively from R different registers.

As may be seen, corresponding words stored in each of
5 the R subparts are inputs to the distributed arithmetic function 202, e.g. DA#1 has R inputs. Each distributed arithmetic function 202 produces a sum $y(i)$ of four weighted products. The different sums $y(i)$ are also identified as output words "out_wp_0-out_wp_23". Each output $y(i)$
10 mathematically represents the sum value of each of the input words multiplied by that word's preassigned individual window presum function coefficient of the filter, such as the filter functions illustrated in Figs. 4A and 4B.

Certain relationships define a preferred window presum
15 computer 102 implementation in accordance with the present invention. The window presum has N outputs which are equal in number to the number N of possible output channels produced by the digital channelizer and the number of points in the discrete Fourier transform apparatus 26'. The window
20 presum computer 102 as described has a total of D inputs which is equal to M, the decimation rate, and represents the number of data words which are inputted to the window presum computer during each clock cycle. However, the invention is not limited to embodiments in which D equals M.

25 It should be understood that the clock circuit utilized in the register implementation of Figs. 14-16 has been

omitted for purposes of simplicity and operates at a slower speed than the serial input data rate because of the parallel operation of the window presum circuits WP1, WP2 and WP3.

5 The window presum computer 102 comprises I window presum circuits with I equal to $\text{GCD}(N,M)$. Each window presum circuit is responsive to a different one of the groups of data words. As may be seen, the data words, which are inputs to Fig. 14, are data words 0, 3, 6, 9 and 12; the
10 data words which are inputs to Fig. 15 are data words 1, 4, 7, 10 and 13; the data words which are inputs to Fig. 16 are data words 2, 5, 8, 11 and 14. The data words $y(i)$ which are outputs from Fig. 14 are data words 0, 3, 6, 9, 12, 15, 18 and 21; the data words $y(i)$ which are outputs from
15 Fig. 15 are data words 1, 4, 7, 10, 13, 16, 19 and 22; and the data words $y(i)$ which are outputs from Fig. 16 are data words 2, 5, 8, 11, 14, 17, 20 and 23. Each window presum circuit WP1, WP2 and WP3 is responsive to a different one of the groups of the aforementioned data input words and
20 produces N' outputs with each of the N' outputs being a function of the window presum function, as illustrated, for example, in the prior art of Figs. 4A and 4B, and a plurality of the data input words with N' not being equal to M and the ratio of $N/N'=I$.

25 The window presum circuits 200 are useful in a modular architecture. Consumption of power is reduced. Undesirable

heating of integrated circuits, which is disadvantageous for applications such as satellites, is lessened which is a primary field of application of the present invention. Data processing in each of the window presum circuits is
5 identical, systematic and in a periodic pattern which allows operations performed by each of the individual window presum circuits WP1, WP2 and WP3, which are preferably integrated circuits, to perform identical processing with each clock cycle. Identical standard window presum circuits
10 implemented in a modular parallel architecture allow the clock rate of the window presum computer 102 to be lowered which lowers power consumption. The number of operations performed to complete the window presum operation is reduced. Interconnect and logic complexity required to
15 perform presum operations is reduced.

Fig. 17 illustrates a preferred embodiment 202 of distributed arithmetic which may be used with the practice of the present invention. However, it should be understood that the present invention is not limited to the use of
20 distributed arithmetic in the practice of the window presum computer 102 of the present invention nor to the preferred embodiment of distributed arithmetic of Fig. 17. The prior art of Fig. 5 is a possible implementation for calculating $y(i)$ without distributed arithmetic.

25 The inputs x_0 , x_1 , x_2 , and x_3 represent eight bit words which are inputted to the individual distributed arithmetic

functions or circuits 202 of Figs. 14-16. A number of parallel processing paths 203, equal to the number of bits within each word x_0 , x_1 , x_2 , and x_3 , perform the calculations which are summed as the quantity $y(i)$ which is the output
5 from each of the distributed arithmetic functions or circuits 202. Each processing path 203 is comprised of logic 52, distributed arithmetic ROM 54 and shifter 56 which perform functions analogous to the prior art of Fig. 6 but accomplish all processing within one clock cycle instead of
10 with multiple clock cycles as in Fig. 6. Each processing path 203 processes a single different bit slice from the R input data words outputted from the registers of Figs. 14-16. As illustrated in Fig. 17, eight different bit slices are processed respectively during one clock cycle by
15 the different processing channels 203.

The logic functions 52 reduce the number of bits by one which are inputted as addressing bits to the distributed arithmetic ROMs 54 which output a precalculated sum of the coefficients which is addressed by the bits within the bit
20 slice of each of the bits of the bit slice. The shifter 56 multiplies the output by the weighting factor 2^n or 2^{-n} to output the correct total for the position of the particular bit slice, e.g. paths 0-7 to the carry save adder 65 which outputs the total y , representing the individual
25 outputs $y(i)$ of each of the distributed arithmetic functions or circuits 202 of Figs. 14-16.

The distributed arithmetic architecture 202 of Fig. 17 has substantial advantages in an environment which has processing constraints such as a satellite application. The use of only parallel processing for each bit slice (e.g. 0-7) reduces the system clock rate which permits circuits such as CMOS to be used, which is a preferred integrated circuit implementation for reducing power consumption in integrated circuits used in accordance with the present invention. The use of CMOS and other integrated circuit technologies which lower power consumption can provide increased processing capability by lowering operating temperature of the integrated circuits. Furthermore, fewer gates are required to implement the distributed arithmetic 202 of Fig. 17 in comparison to the prior art of Fig. 5.

Figs. 18 and 19 illustrate the implementation of a preferred embodiment of the cyclic shift 24' used in a digital channelizer in accordance with the invention in which divides an input bandwidth into at least some of equally spaced N channels. The cyclic shifting function of the present invention performs the same function as in the prior art but is implemented in a modular architecture of integrated circuits, such as CMOS, in a parallel data processing architecture. The cyclic shift 24', as illustrated in Fig. 18, includes I parallel cyclic shifting paths which are identical in number to the I window presum

circuits utilized in a preferred embodiment of a window
presum computer 102 in accordance with the present
invention. The use of same number of parallel processing
paths in the window presum computer 102 and in the cyclic
5 shift 24' permits a highly efficient interconnection between
the window presum computer 102 and the cyclic shift which
increases the efficiency of the cyclic shift 24' and
substantially reduces power consumption in a digital
channelizer in a satellite application, such as a multibeam
10 regenerative wideband satellite system. Each cyclic shift
path 500 is responsive to a different input group of data
words which correspond to the data words outputted by the
window presum computer in I (three) groups as illustrated in
Fig. 13. The I cyclic shift paths 500 produce I output
15 groups of data words 504 as illustrated on the right hand
side of Fig. 18. The outputs of data words 504 from the
cyclic shift paths 500 are inputs to the discrete Fourier
transform apparatus used in practice of the present
invention which may be in accordance with the prior art of
20 Fig. 11 or in accordance with the discrete Fourier transform
apparatus of Fig. 21 which prunes certain outputs from the
output Fourier transform computation stage 419 where all
frequency domain outputs are not necessary for subsequent
processing into channels by the channelizer. Each cyclic
25 shift path 500 has a number of barrel shifters 502 which is
equal to the number of inputs thereto, from the

corresponding window presum circuit 200 of the window presum computer 102. The inputs to each of the barrel shifters 502 to the right of the barrel shifters having the outputs wp-0, wp-1 and wp-2 in each of the I cyclic shift paths 500 are
5 rotated one position so that the identified output represents the top input but the order of the inputs is unchanged. For example, the barrel shifter 502 having the output wp-15 has its inputs out-wp from top to bottom in the order 15, 6, 21, 12, 3, 18, 9 and 0 and the barrel shifter
10 immediately to the right has its inputs out-wp in the order 6, 21, 12, 3, 18, 9, 0, 15.

The number of cyclic shift paths is equal to $\text{GCD}(N, M)$.

Each output group 504 of data words is shifted through a number of cycles Q equal to $N/\text{GCD}(N, M)$. The number of
15 cycles used by the cyclic shift 24', e.g. 8, determines the number of outputs of the window presum circuits 202 which are grouped together and permits direct mapping of interconnections from the window presum circuits 200 into the parallel cyclic shift paths 500. Each cycle has a shift
20 value defined by $mM \bmod N$ or $-mM \bmod N$ with m varying from zero to Q-1. The foregoing number of cycles and shift values produce an efficient cyclic shift 24' of a parallel architecture which has inputs which efficiently map to the outputs of the window presum circuits 102 and into the
25 inputs of the discrete Fourier transform apparatus 26' without complex interconnections.

Fig. 19 illustrates a table containing the various cycle values for the eight cycles through which each of the word outputs of the barrel shifters 502 are shifted in each of the cyclic shift paths 500. The headings 0-23 respectively indicate for each of the eight cycles the output value of the barrel shifters 502 identified by the headings 0-23 which are the input words to individual cyclic shift paths 500 to the left in Fig. 18. As may be seen, the parallel cyclic shift paths 500 from top to bottom respectively have input words in a first group of 0, 15, 6, 21, 12, 3, 18 and 9, in a second group of 1, 16, 7, 22, 13, 4, 19 and 10 and in a third group of 2, 17, 8, 23, 14, 5, 20 and 11. The first group of 0, 15, 6, 21, 12, 3, 18 and 9 represents the calculated result of mM modulo N for M equalling 15, N equalling 24 and m varying from 0 to 7. When the outputs of the window presum computer 102 produced by the window presum circuits 200 are chosen in this order, the interconnection between the window presum computer 102 and the cyclic shift 24' is simplified. The choice of the order of the words in the second and third groups is in accordance with the same sequence as used in the first group except that each word value is augmented by one in group two relative to group one and is augmented by two in group three relative to group one. Therefore, the required relative shifting of the output between subsequent words in each group is identical and is maintained which permits the

shifting of the parallel cyclic paths 500 to be driven by a modulo eight (3 bit) counter 506 toggled at the clock rate.

As illustrated with the outputs 504 from the cyclic shift 24' it is possible to enable the cyclic shift outputs to efficiently map into the inputs of the discrete Fourier transform apparatus 26' such as in Fig. 21. The cyclic shift paths 500 produce outputs that are independent of a first stage discrete Fourier transform which allows hardware design replication in a parallel architecture. This efficient mapping occurs when the first discrete independent Fourier transform computation stage contains $GCD(N,M)$ discrete Fourier transforms each having Q input points and the prime factor algorithm is used which causes the ordering of the discrete Fourier transform to be such that the outputs 504 may be mapped with simple interconnection into the inputs of the discrete Fourier transform such as that illustrated in the prior art of Fig. 11 and the pruned discrete Fourier transform described below in conjunction with Fig. 21. Other efficient mappings may offer similar simple interconnection characteristics.

The actual output values in Fig. 19 for each of the barrel shifters identified by the headings 0-23 of the cyclic shift paths 500 are obtained by identifying the particular shift cycle of the repeating cycles 1-8.

Therefore, it is seen that the barrel shifter identified

with the heading 15 sequentially for cycles 1-8 outputs inputs identified by out_wp 15, 6, 21, 12, 3, 18, 9 and 0, respectively.

The individual cyclic shift paths 500 contain Q barrel
5 shifters 502 which function as word shifting elements. Each barrel shifter 502 is responsive to an identical group of input words 503 identified by the legend "out_wp" followed by a number and outputs only a single word at a time. The present invention is not limited to barrel shifters to
10 implement the cyclic shift 24' with any circuit or process which can output the word values for the cycle inputs of Fig. 19 being usable in the practice of the invention.

Fig. 20 illustrates a first embodiment 600 of a pruned discrete Fourier transform apparatus in accordance with the
15 present invention which represents a modification of the prior art discrete Fourier transform of Fig. 10. Pruning results in the number of actual outputs P of the discrete Fourier transform being less than a possible number of outputs. Pruning of a discrete Fourier transform reduces
20 computational complexity both in terms of logic complexity and interconnection complexity. Furthermore, pruning also reduces power consumption. Any operations used to generate inputs to the pruned portion of the discrete Fourier transform apparatus can be eliminated from the input
25 discrete Fourier transform stage.

The discrete Fourier transform apparatus 600 represents the circumstance where only output frequency domain signals 2, 3, 4, and 5 are required for subsequent processing from inputs from the cyclic shift 24'. The three point DFT 306' which contains an "X" is unnecessary in view of output frequency domain signal 7 being used in place of frequency domain signal 5 for processing by processing device 602 which represents any downstream processing of the discrete Fourier transform including, but not limited to, applications in a digital channelizer. For real input signals frequency domain signal 7 is a conjugate of frequency domain signal 5, which is not provided in an output of the three remaining three point discrete Fourier transforms 306. However, because frequency domain signal 7, as illustrated in Fig. 3A, is a conjugate of frequency domain signal 5, which is not present in the output from the discrete Fourier transform apparatus 600, a sign negation of the imaginary part of the frequency domain signal 7 permits the information of frequency domain signal 5 to be recovered and available for downstream processing while permitting the discrete Fourier transform apparatus output(s) to be simplified which has the above-stated advantages.

A discrete Fourier transform apparatus in accordance with the invention has at least one discrete Fourier transform computation stage. However, in a preferred embodiment the discrete Fourier transform includes multiple

computation stages 302 and 305', as illustrated in Fig. 20, and, 402, 410, 412, 414 and 419 as illustrated in Fig. 21, as discussed below. While the example illustrated in Fig. 20 represents only the use of frequency domain signal 7 in place of frequency domain signal 5, it should be understood that each of the N inputs which have conjugates (e.g. frequency domain signals 0 and N/2 (for even N) do not have conjugates) may be represented in the output by its conjugate which is subsequently processed in place of the frequency domain signal which is not an actual output. Furthermore, at least one processing device 602 may be used to process at least one conjugate as representative of one of the input signals.

Fig. 21 illustrates another embodiment of a discrete Fourier transform apparatus 700 in accordance with the present invention in which two conjugates of the frequency domain signals illustrated in Fig. 3B are processed as representative of inputs. The difference between the prior art of Fig. 11 and the discrete Fourier transform apparatus 700 resides in the elimination of one of the eight point postweaves 420 which permits channels X[16] and X[19] to be used in place of their conjugate frequency domain signals X[8] and X[5] which are not outputs in view of the elimination of the third eight point postweave 420 contained in Fig. 11. The processing devices 702 process the transformation of the plurality of conjugates of the twenty-

four inputs, which are not transformed into an actual output, in place of the plurality of the twenty-four inputs which are not transformed into an actual output. As a result, the overall structure of the discrete Fourier transform apparatus 700 is simplified in that the output discrete Fourier transform computation stage has an eight point postweave 420 pruned therefrom which lessens power consumption and further reduces the amount of hardware in an application where all of the twenty four possible output channels in a channelizer are not required.

Fig. 22 illustrates a modification of the circular shift 24' and discrete Fourier transform apparatus 26' as illustrated in Fig. 12 under the control of an input command specifying the processing of real or complex data. The command permits real word inputs or complex word inputs to be transformed into channel outputs of a channelizer without external modification of the cyclic shifts 24' and the discrete Fourier transform 26'. In the left-hand portion of Fig. 22, the headings "complex" and "real" respectively identify the selective inputting of twelve data points. The data listed under the heading "complex", signify the real and imaginary parts of six words and, the data listed under the heading "real", signify the real data of twelve input words.

The cyclic shift 24' is comprised of twelve barrel shifters (multiplexers) 700 which receive inputs that

correspond to the data under the "real" heading. However, it should be understood that if complex data is being inputted, that that data is identified by the words "w0_r" through "w5_i". The actual input values are either

5 represented by the imaginary or real component of complex data or real data. The barrel shifters 700 have outputs which, depending upon the particular barrel shifter, are either coupled directly to a two point discrete Fourier transform 702 or an input to switch 704. The switches 704

10 each have at least one input receiving real data as identified by the words W4, W10, W8, W2, W9, W3, W1, and W7 and at least one input receiving a component of complex data as identified by inputs W4_r, W1_r, W2_r, W5_r, W0_i, W3_i, W4_i and W1_i. The input to each of the switches 704, REAL/COMPLEX

15 SELECT is a command which specifies processing of the cyclically shifted real data or the cyclically shifted complex data outputted to the barrel shifters 700 by the cyclic shift 24'. The discrete Fourier transform apparatus 26' is comprised of an input discrete Fourier

20 computation stage 720 comprised of six two-point discrete Fourier transforms 702 and four three-point discrete Fourier transforms 706 and an output stage 728 comprised of six two-point discrete Fourier transforms/recombines 708. In real processing mode, the two point discrete Fourier

25 transforms/recombines 708 function as two point discrete Fourier transforms and as two point recombines in complex

mode. The individual outputs from the output stage 728 are always complex data but, depending upon whether the REAL/COMPLEX SELECT command specifies the processing of real or complex data, the number and order of complex outputs is
5 changed.

The architecture of the cyclic shift 24' and the discrete Fourier transform apparatus 26' of Fig. 22 permits modular construction in diverse applications in which it is desirable to selectively input to discrete Fourier transform
10 apparatus 26' either real or complex data outputs from a window presum computer 102 without having to modify the cyclic shift 24' and discrete Fourier transform apparatus 26'. The only variable in the operation of Fig. 22 is the command REAL/COMPLEX SELECT to transform the inputs to the
15 cyclic shift 24' as either real data or complex data.

A universal architecture for discrete Fourier transforming of real or complex data may be implemented in integrated circuits to process data in parallel processing architectures such as that of the present invention. While
20 Fig. 22 is illustrated as processing as inputs the output from a twelve point window presum calculation, it should be understood that the configuration of Fig. 22 may be modified to parallel architectures, such as the eight point processing of the parallel channels described above.

25 The following relationships exist for the discrete Fourier transform apparatus 26' of Fig. 22. N is a multiple

of two. The N point discrete Fourier transform apparatus 26' includes an input discrete Fourier computation stage 724 which is required to have two N/2 point Fourier transforms 702 and an output discrete Fourier computation stage 728 which is required to have N/2 two-point discrete Fourier transforms/recombines 708. The two point discrete Fourier transforms/recombines 708 are two point discrete Fourier transforms in real mode and two point recombines in complex mode as described in more detail below. The architecture of the discrete Fourier transforms/recombines 708 as two point structures enables them to function in real and complex mode.

In order to permit operations in real and complex mode, the discrete Fourier transforms/recombines 708 contain identical hardware. The discrete Fourier transform/recombine 708 functions in a configurable manner which configuration depends upon the type of data being processed and the location in the output discrete Fourier transforms stage 728 in Fig. 22.

Real processing is illustrated in Fig. 23 during which only real data is subjected to discrete Fourier transformation.

In a complex data processing mode, as illustrated in Fig. 24, one complex input is combined from the N/2 point discrete Fourier transforms 702 that process the "real" inputs, while the other complex input will come from the N/2

point discrete Fourier transforms 702 that process the
"imaginary" inputs. The "imaginary" value must be
multiplied by j (equivalent to i mathematically). Then, the
individual components of the "real" and "imaginary" values
5 are added to achieve the final output.

In the implementation of the twelve point discrete
Fourier transform apparatus 26' (in the 2-3-2 architecture),
the outputs of one of the three point discrete Fourier
transforms 706 requires a multiplication by $-j$ (a twiddle
10 factor) prior to the final two point discrete Fourier
transform/recombine 708. This multiplication operation is
performed by the two point discrete Fourier
transform/recombine 708 as shown in Fig. 25 as (real case-
special).

15 The bottom three two-point discrete Fourier
transform/recombines 708 in Fig. 22 (which have their lower
input from the bottom-three point discrete Fourier
transforms 706), provide a multiplication by $-j$ in the case
of real data processing.

20 The design of the two point discrete Fourier
transforms/recombines 708 must incorporate all possible
operations, including two point discrete Fourier transform,
two point discrete Fourier transform with $-j$ multiplication,
and two point complex recombine function. A high level
25 diagram of a discrete Fourier transform 708 performing all

of these functions is illustrated in Fig. 26 and shows that all output combinations are attainable.

While the invention has been defined in terms of preferred embodiments, it should be understood that numerous
5 modifications may be made thereto without departing from the spirit and scope of the present invention. It is intended that all such modifications fall within the scope of the appended claims.

CLAIMS

1. A digital channelizer which divides an input bandwidth into at least some of N channels comprising:
- 5 an analog to digital converter which encodes the input bandwidth into a serial digital data stream of words;
- a demultiplexer, coupled to the analog to digital converter, which divides the serial digital data stream into parallel data streams of data words;
- 10 a window presum, coupled to the parallel streams of data words, responsive to a window presum function, storing L data words in a window containing R subparts and outputting N outputs which are each respectively capable of being processed into a different one of the N channels, each
- 15 of the N outputs being a product of R words each respectively from a different one of the R subparts and a coefficient of the window presum function preassigned to each of the R words from a different one of the R subparts with $R=L/N$ and each of the N outputs being calculated using
- 20 distributed arithmetic during a single clock cycle; and
- a discrete Fourier transform apparatus, coupled to the N outputs of the window presum, having N inputs and which performs a discrete Fourier transform on the N inputs to output at least some of the N channels.

25

2. A digital channelizer in accordance with claim 1 wherein the window presum further comprises:

N distributed arithmetic calculators, each distributed arithmetic calculator comprising a number of
5 processing paths equal to a number of bits of each data word and an adder which sums products of the R words and the preassigned coefficient of each of the R words from each of the processing paths in one clock cycle.

10 3. A digital channelizer in accordance with claim 2 wherein each processing channel further comprises:

a memory which stores sums of different combinations of preassigned coefficients and negations thereof at addressable locations which are addressed by at
15 least a part of a bit slice of the R words and each addressed sum being coupled to the adder.

4. A digital channelizer in accordance with claim 3 wherein each processing channel further comprises:

20 a shifter, coupled to the addressed sums, which multiplies the addressed sums by a power of two corresponding to a position of the bit slice in the R words.

5. A digital channelizer in accordance with claim 4
25 wherein the adder comprises:

a carry save adder.

6. A digital channelizer in accordance with claim 4 wherein each processing channel further comprises:

5 a logic circuit, responsive to a bit slice of the R words, which reduces by a factor of two a number of addressable sums stored in the memory.

7. A digital channelizer in accordance with claim 1 wherein:

10 each of N outputs is also a function of words from a plurality of the parallel data streams.

8. A digital channelizer in accordance with claim 2 wherein:

15 each of N outputs is also a function of words from a plurality of the parallel data streams.

9. A digital channelizer in accordance with claim 3 wherein:

20 each of N outputs is also a function of words from a plurality of the parallel data streams.

10. A digital channelizer in accordance with claim 4 wherein:

25 each of N outputs is also a function of words from a plurality of the parallel data streams.

11. A digital channelizer in accordance with claim 5
wherein:

each of N outputs is also a function of words from
5 a plurality of the parallel data streams.

12. A digital channelizer in accordance with claim 6
wherein:

each of N outputs is also a function of words from
10 a plurality of the parallel data streams.

13. A process for dividing an input bandwidth into at
least some of N channels comprising:

encoding the input bandwidth into a serial data
15 stream of words;

dividing the serial data stream into parallel data
streams of data words;

processing the data words of the parallel data
streams to produce products of data words and a window
20 presum function which are summed to produce N outputs, the
window presum function having a window containing L data
words divided into R subparts, each of the N outputs being
a product of R samples each respectively from a different
one of the R subparts and a coefficient of the window presum
25 function preassigned to each of the R words from a different
one of the R subparts with $R = L/N$ and each of the N outputs

being calculated using distributed arithmetic during a single clock cycle; and

performing a discrete Fourier transform on N inputs, which are coupled to N outputs produced by processing the
5 data words of the parallel data streams to produce at least some of the N channels.

14. A process in accordance with claim 13 further comprising:

10 providing N distributed arithmetic calculators with each distributed arithmetic calculator calculating one of the N outputs;

each distributed arithmetic calculator having a number of processing paths equal to a number of bits of
15 each data word with each processing path processing a different bit slice of the R samples; and

each of the N outputs being produced by summing a processing result of each of the processing paths.

20 15. A process in accordance with claim 14 wherein:

each distributed arithmetic calculator is provided with a memory which stores sums of different combinations of preassigned coefficients and negations thereof at
addressable locations which are addressed by at least part
25 of a bit slice of the R words;

coupling each addressed sum to an adder which calculates a summation representing one of the N outputs.

16. A process in accordance with claim 15 further
5 comprising:

multiplying each addressed sum from the memory by a power of two corresponding to a position of a bit slice in the R words.

10 17. A process in accordance with claim 16 wherein:
reducing by a factor of two a number of
addressable sums stored in the memory.

18. A process in accordance with claim 13 wherein:
15 each of the N outputs is also a function of words
from a plurality of the parallel data streams.

19. A process in accordance with claim 14 wherein:
each of the N outputs is also a function of words
20 from a plurality of the parallel data streams.

20. A process in accordance with claim 15 wherein:
each of the N outputs is also a function of words
from a plurality of the parallel data streams.

25

21. A process in accordance with claim 16 wherein:
each of the N outputs is also a function of words
from a plurality of the parallel data streams.

5 22. A process in accordance with claim 17 wherein:
each of the N outputs is also a function of words
from a plurality of the parallel data streams.

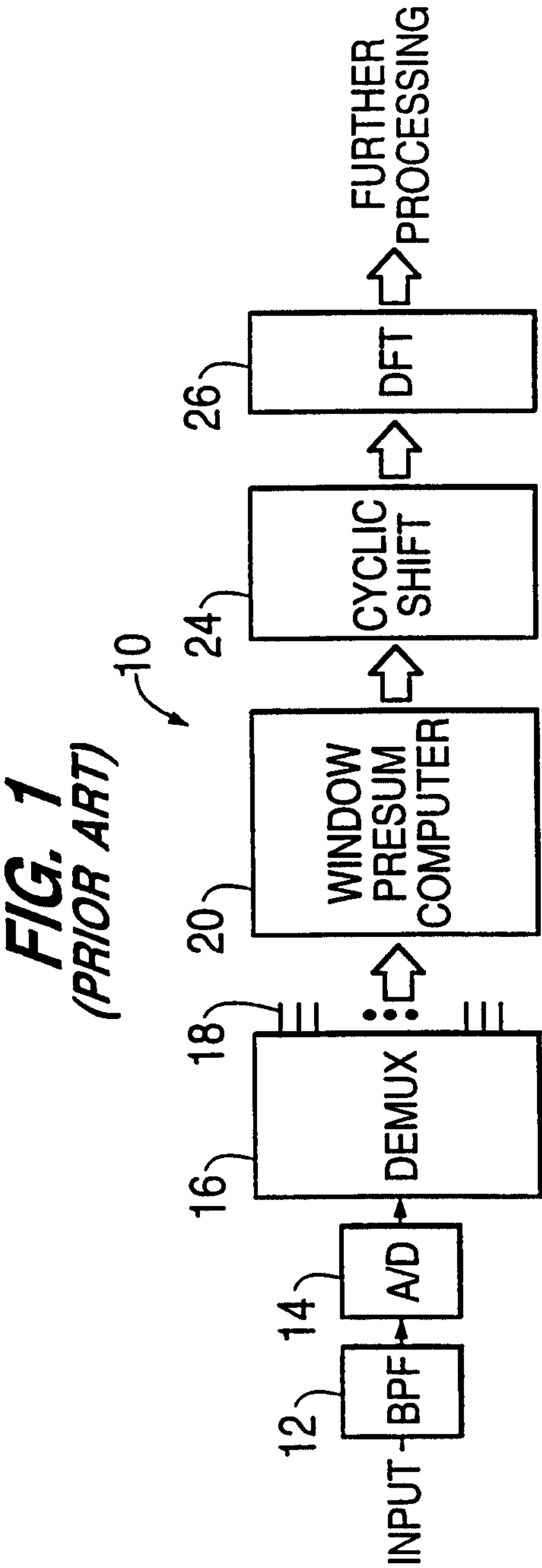


FIG. 2
(PRIOR ART)

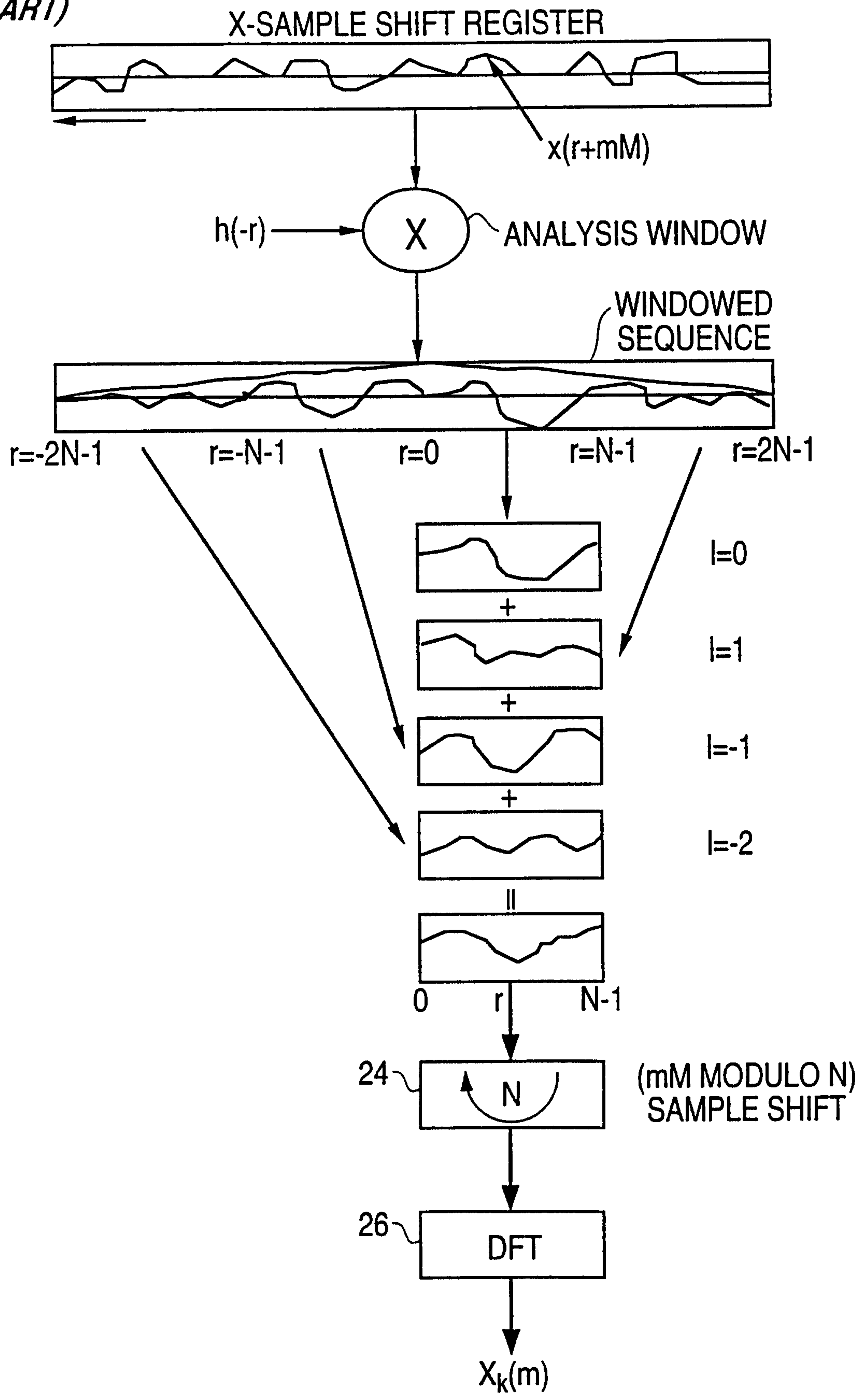


FIG. 3A
(PRIOR ART)

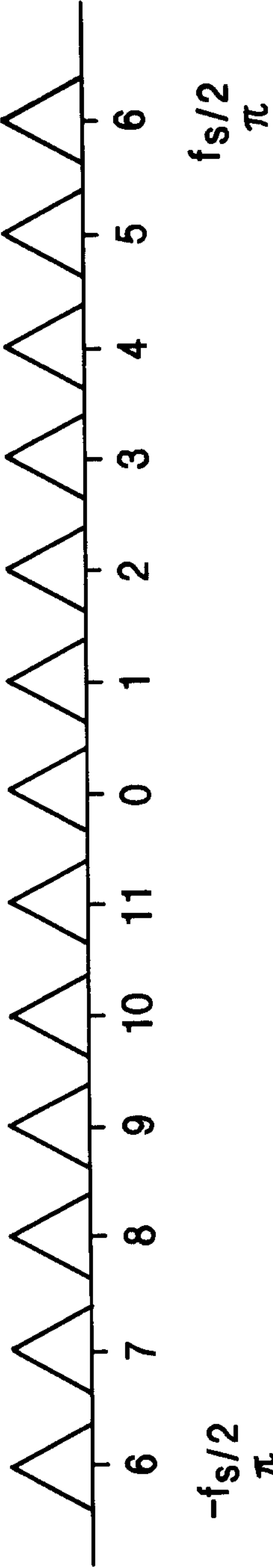


FIG. 3B
(PRIOR ART)

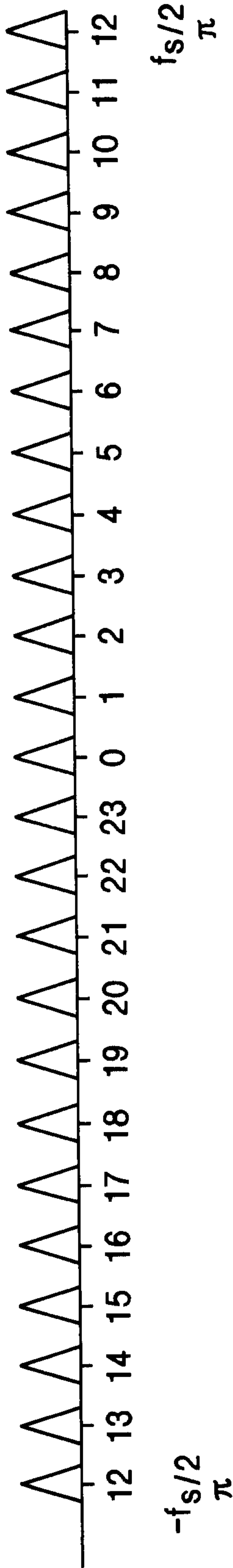


FIG. 4A
(PRIOR ART)

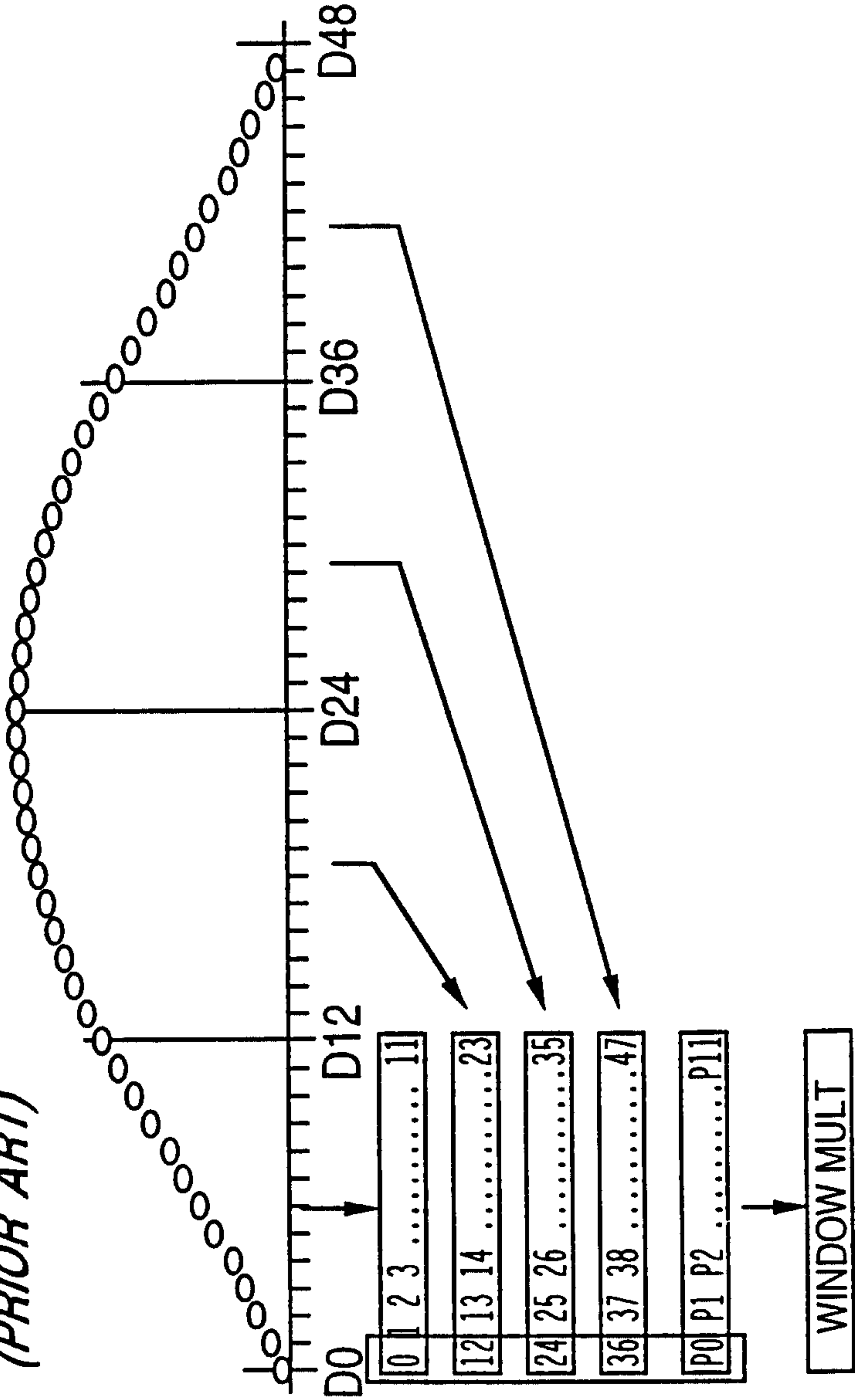


FIG. 4B
(PRIOR ART)

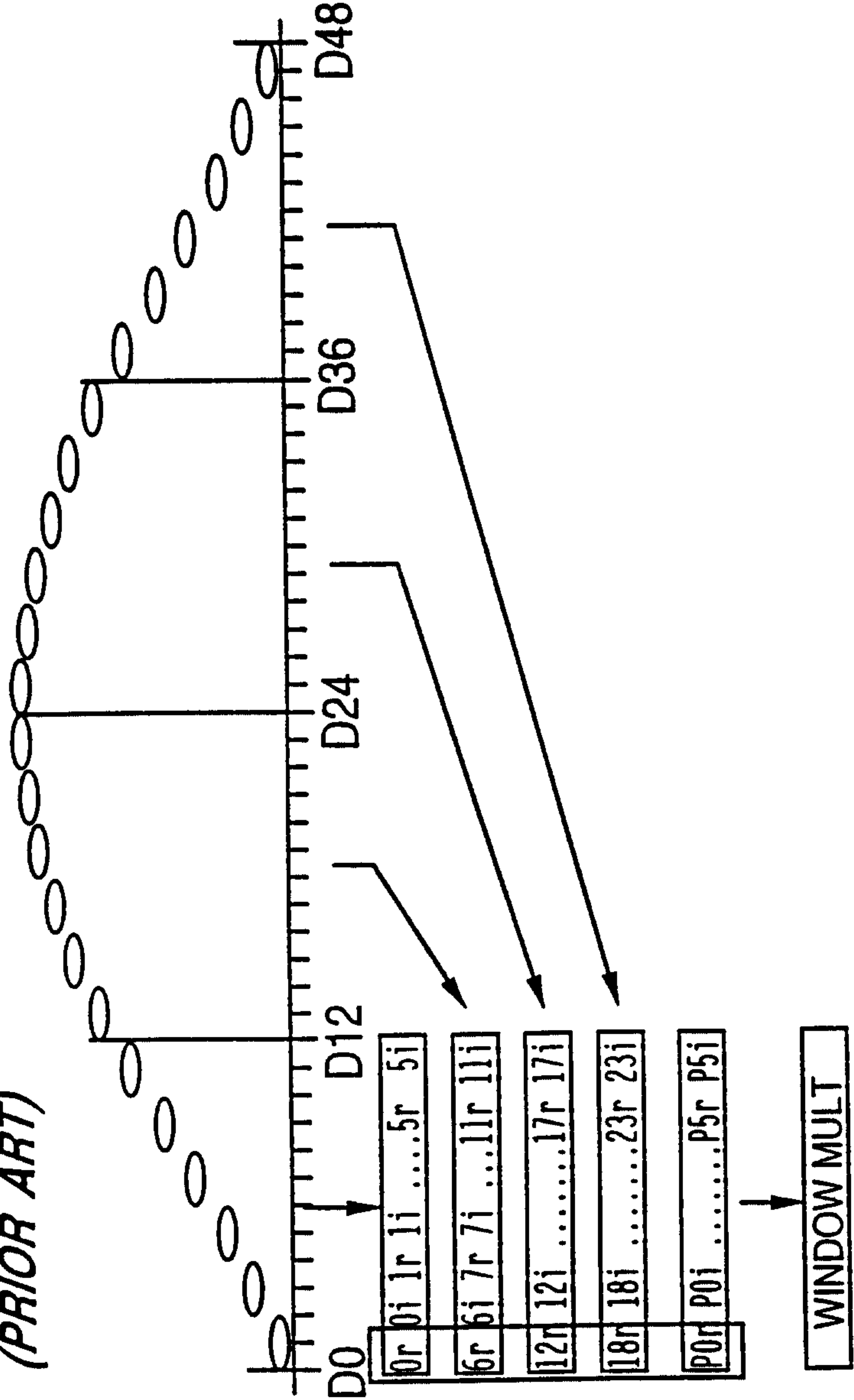


FIG. 5
(PRIOR ART)

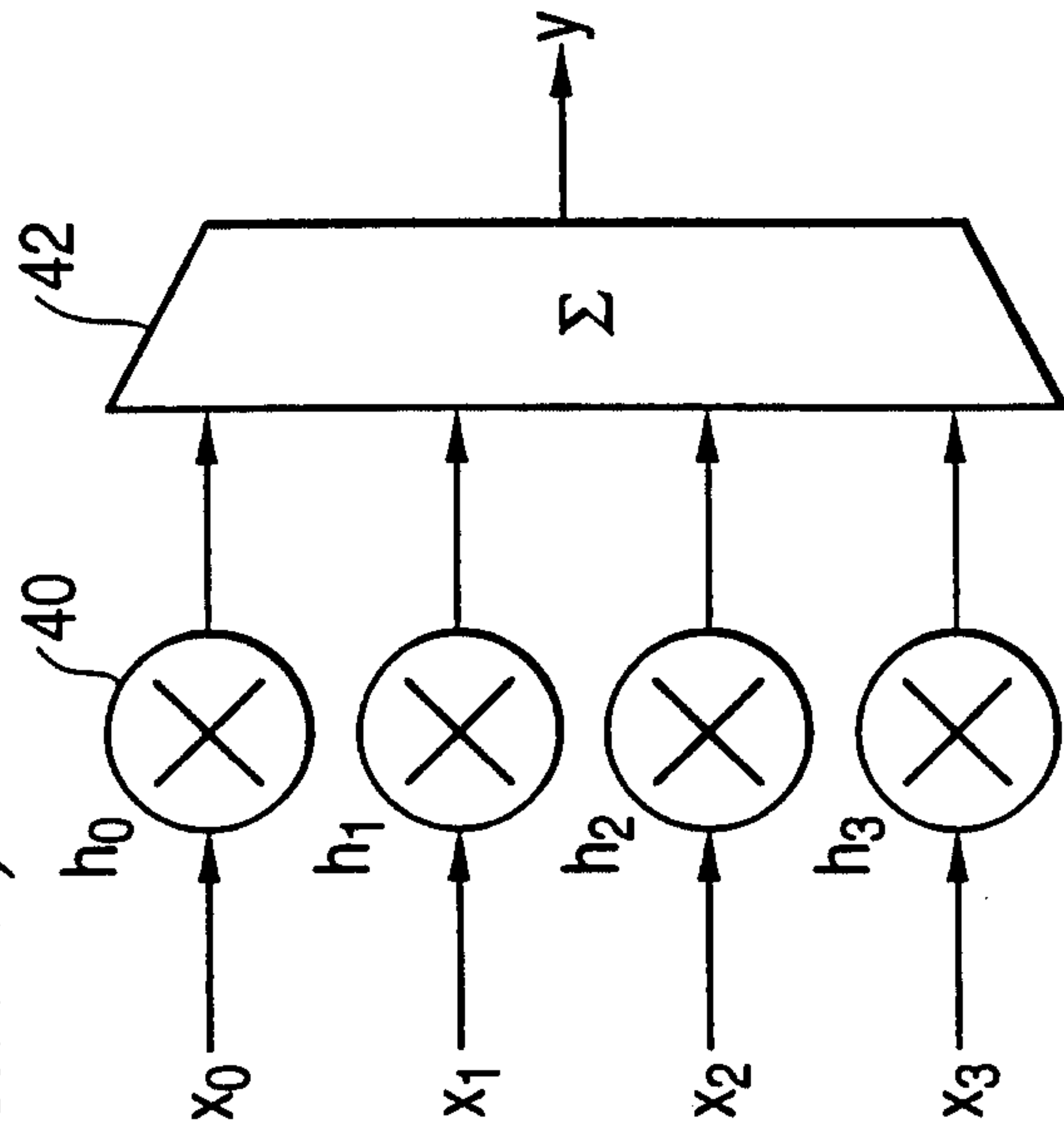


FIG. 17

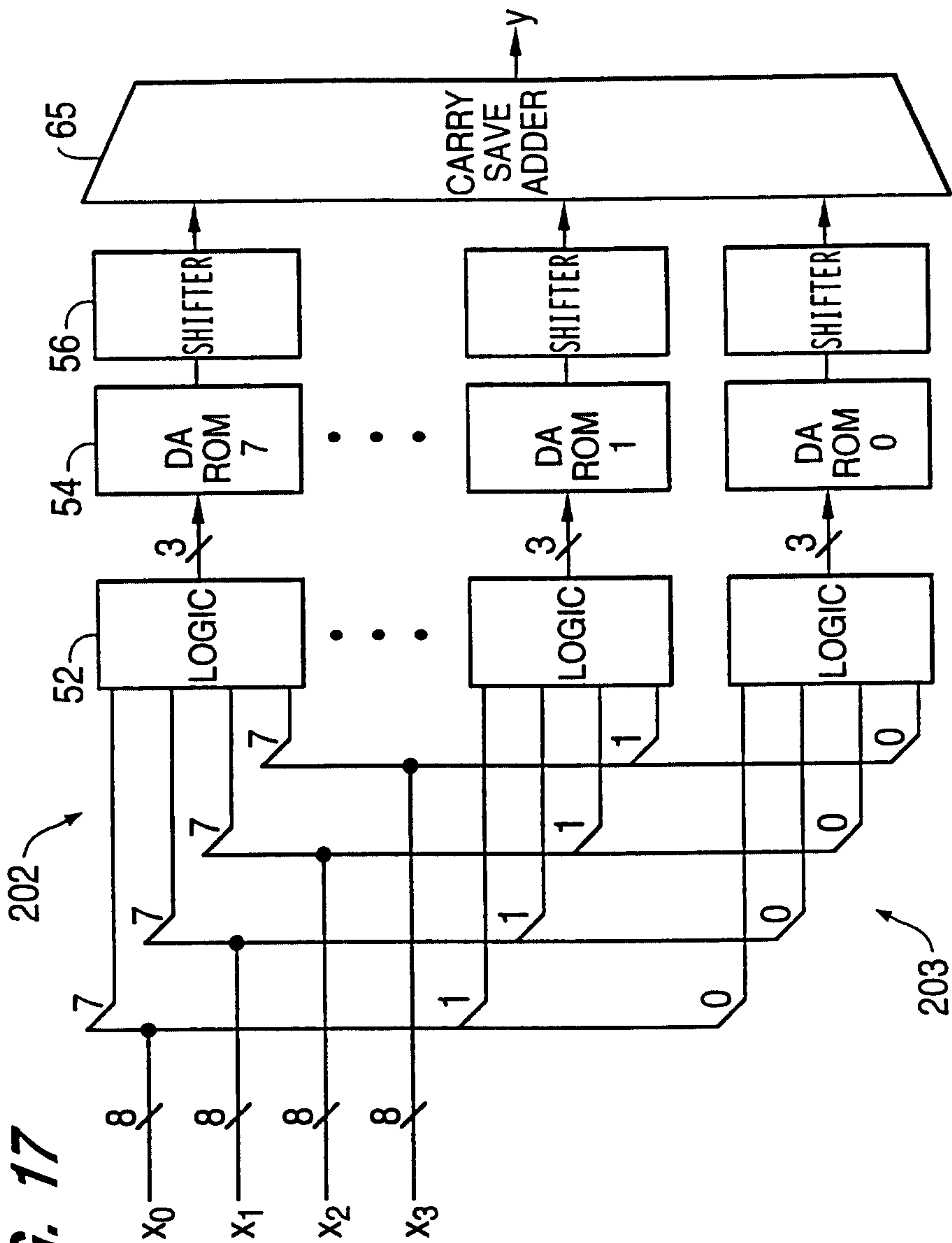


FIG. 6
(PRIOR ART)

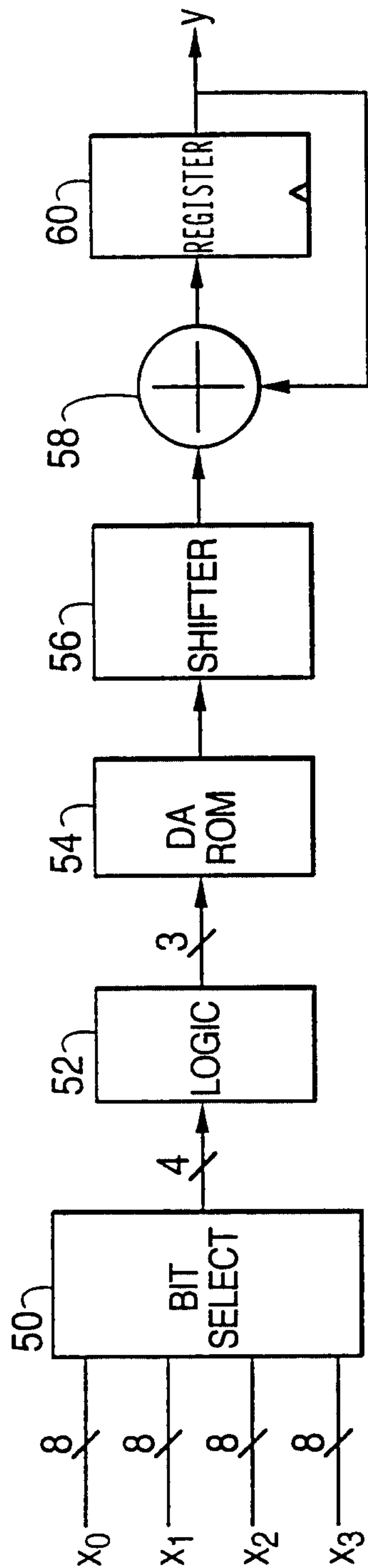


FIG. 7
(PRIOR ART)

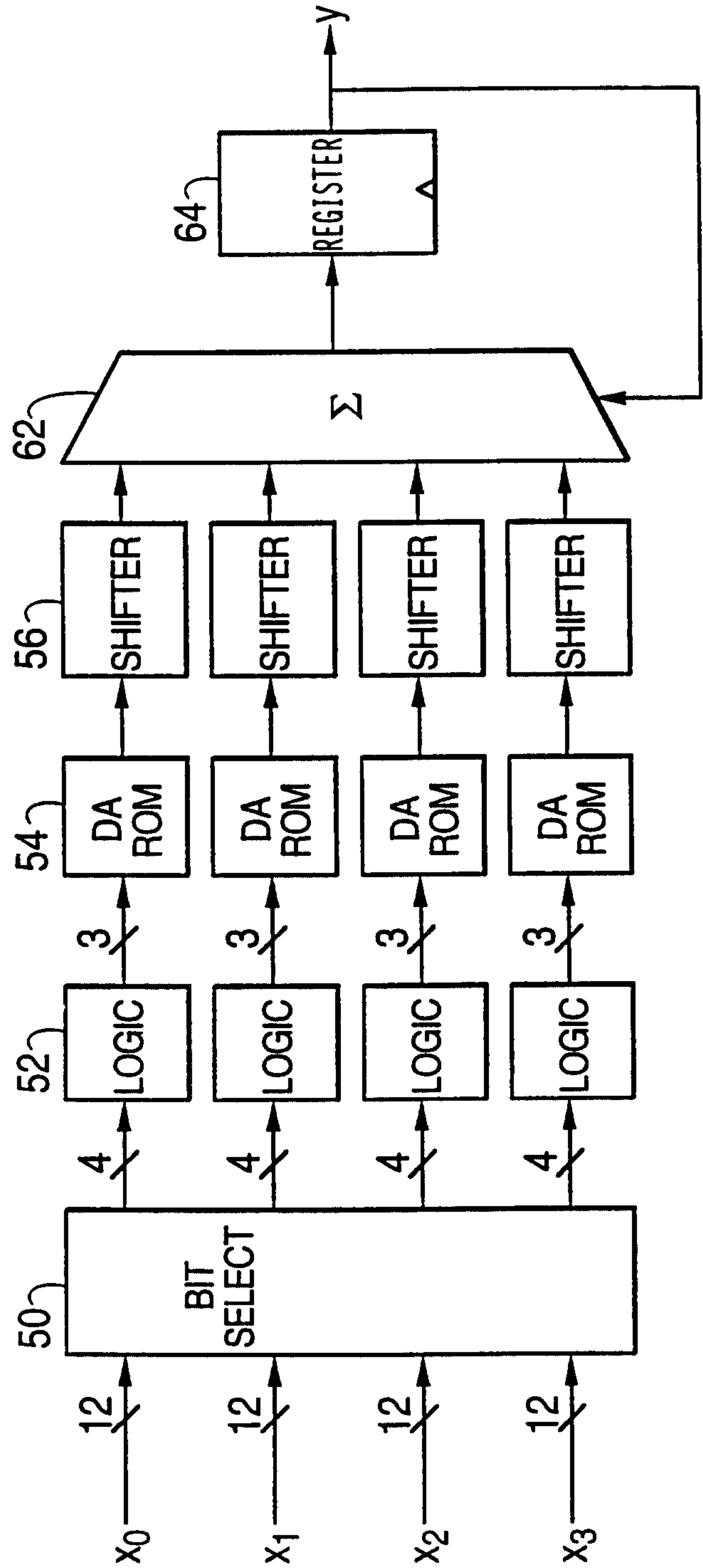


FIG. 8
(PRIOR ART)

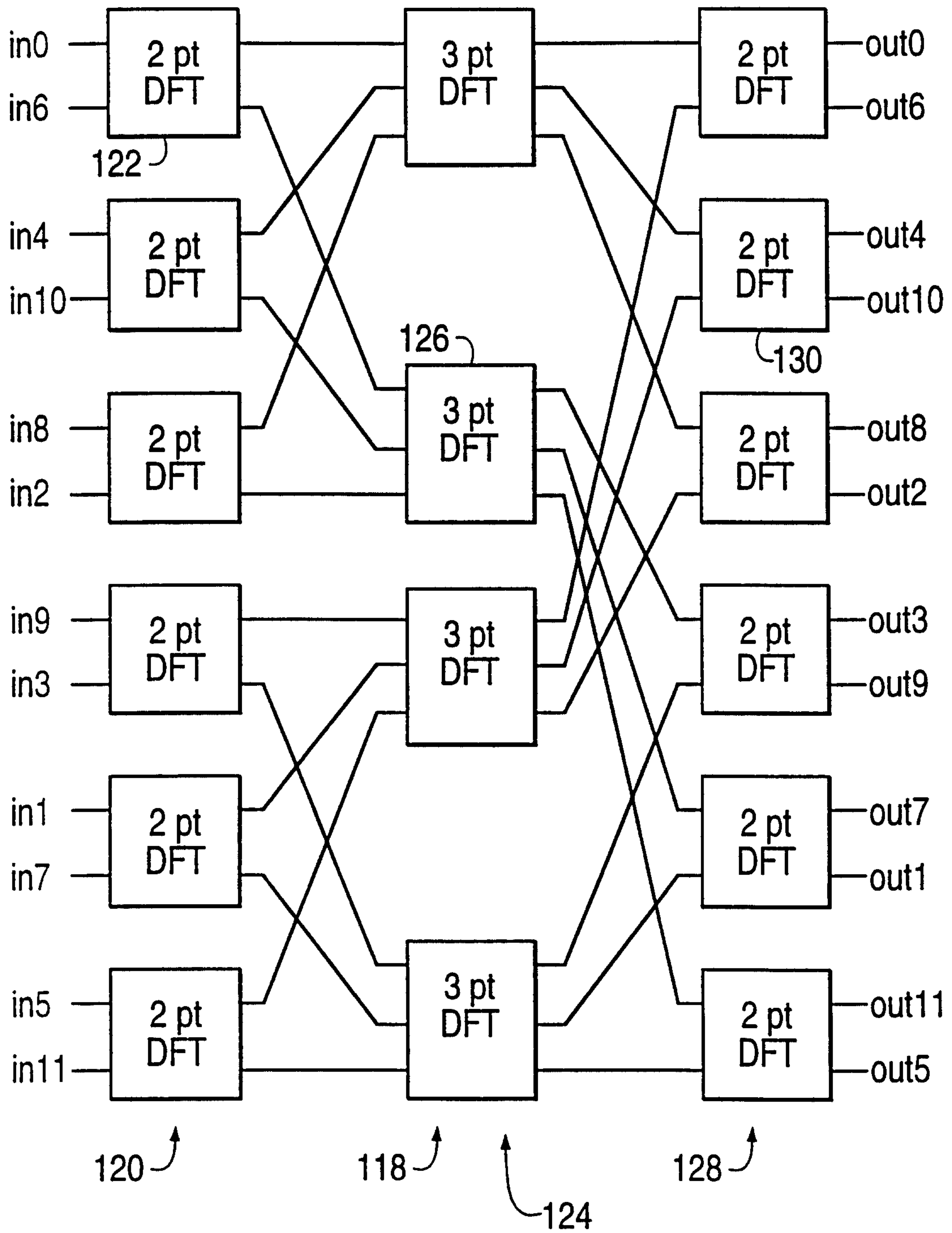


FIG. 9
(PRIOR ART)

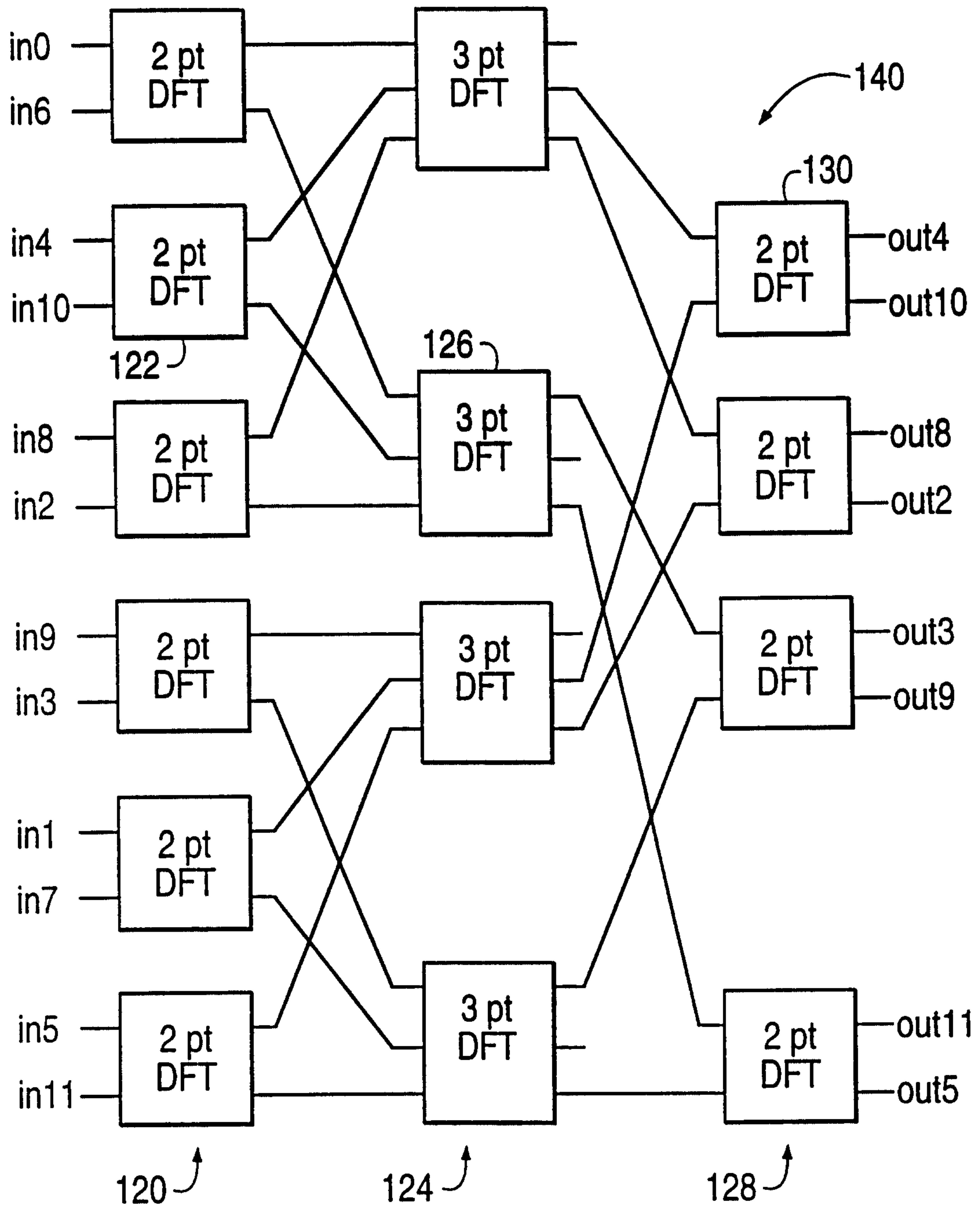


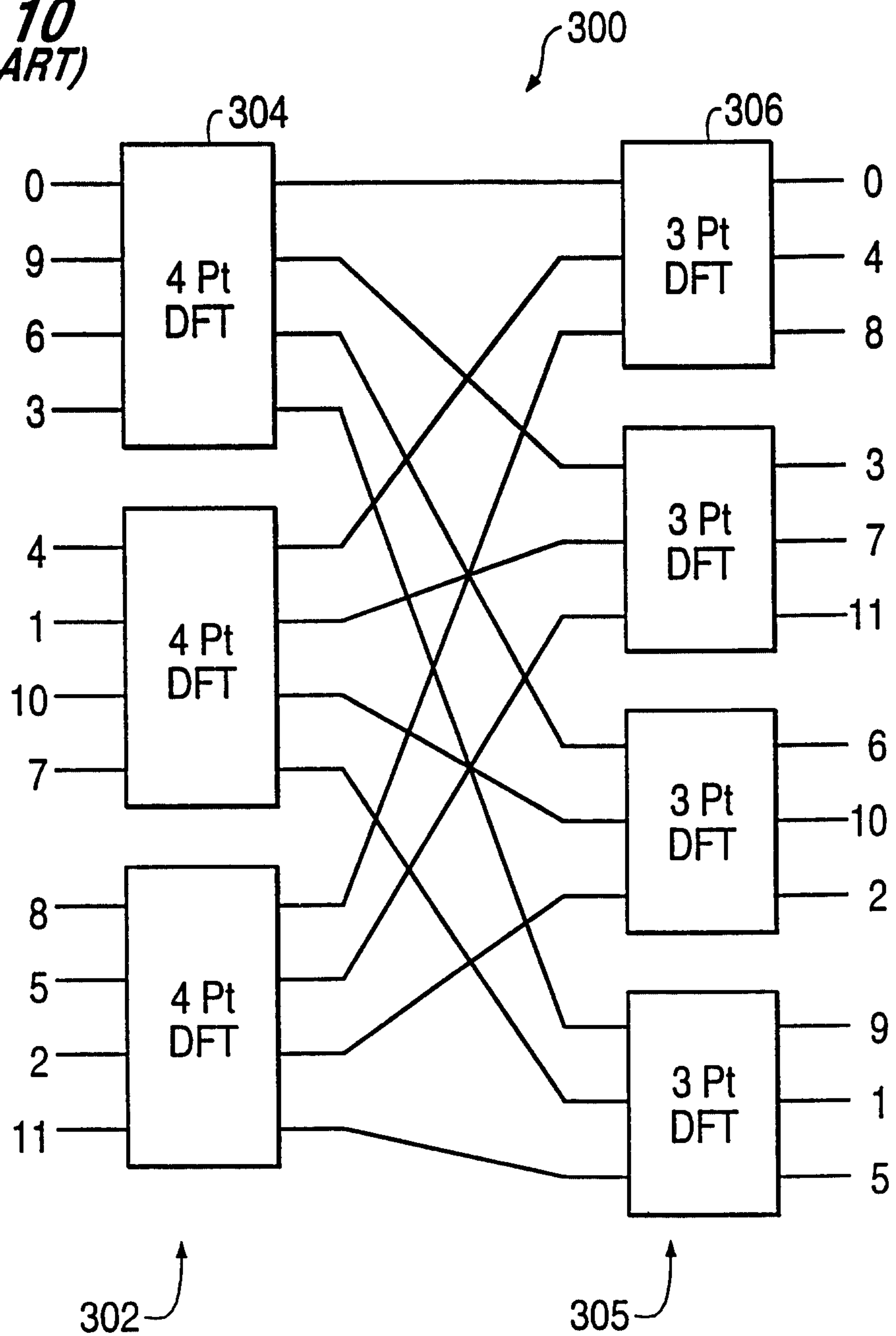
FIG. 10
(PRIOR ART)

FIG. 11
(PRIOR ART)

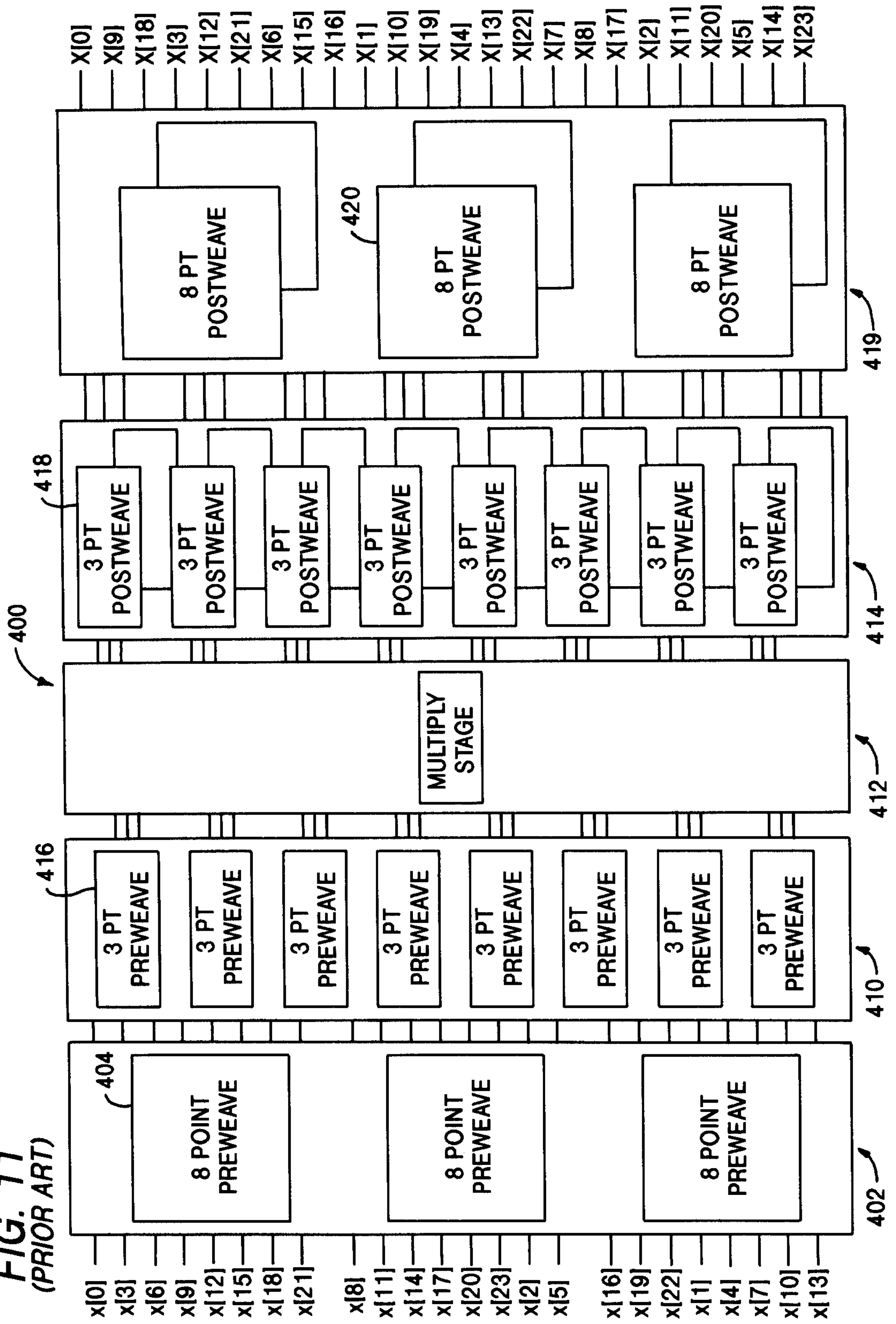


FIG. 12

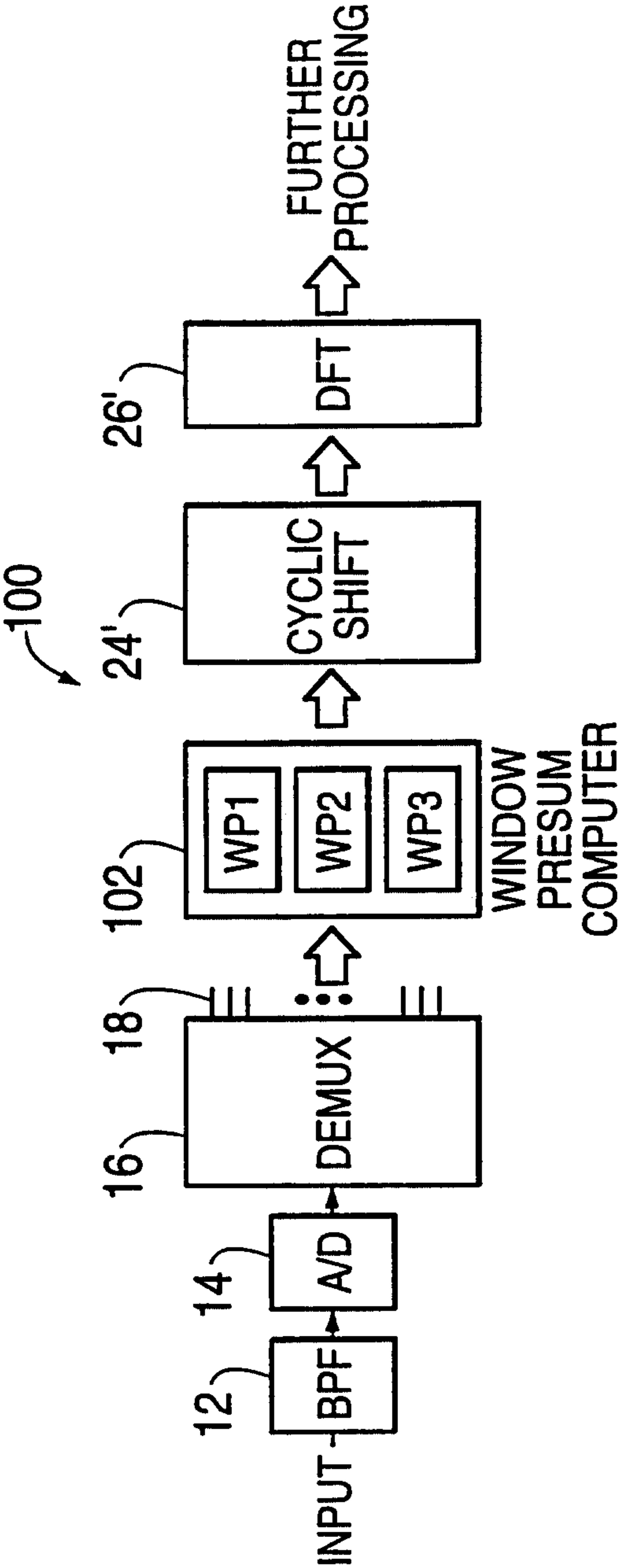


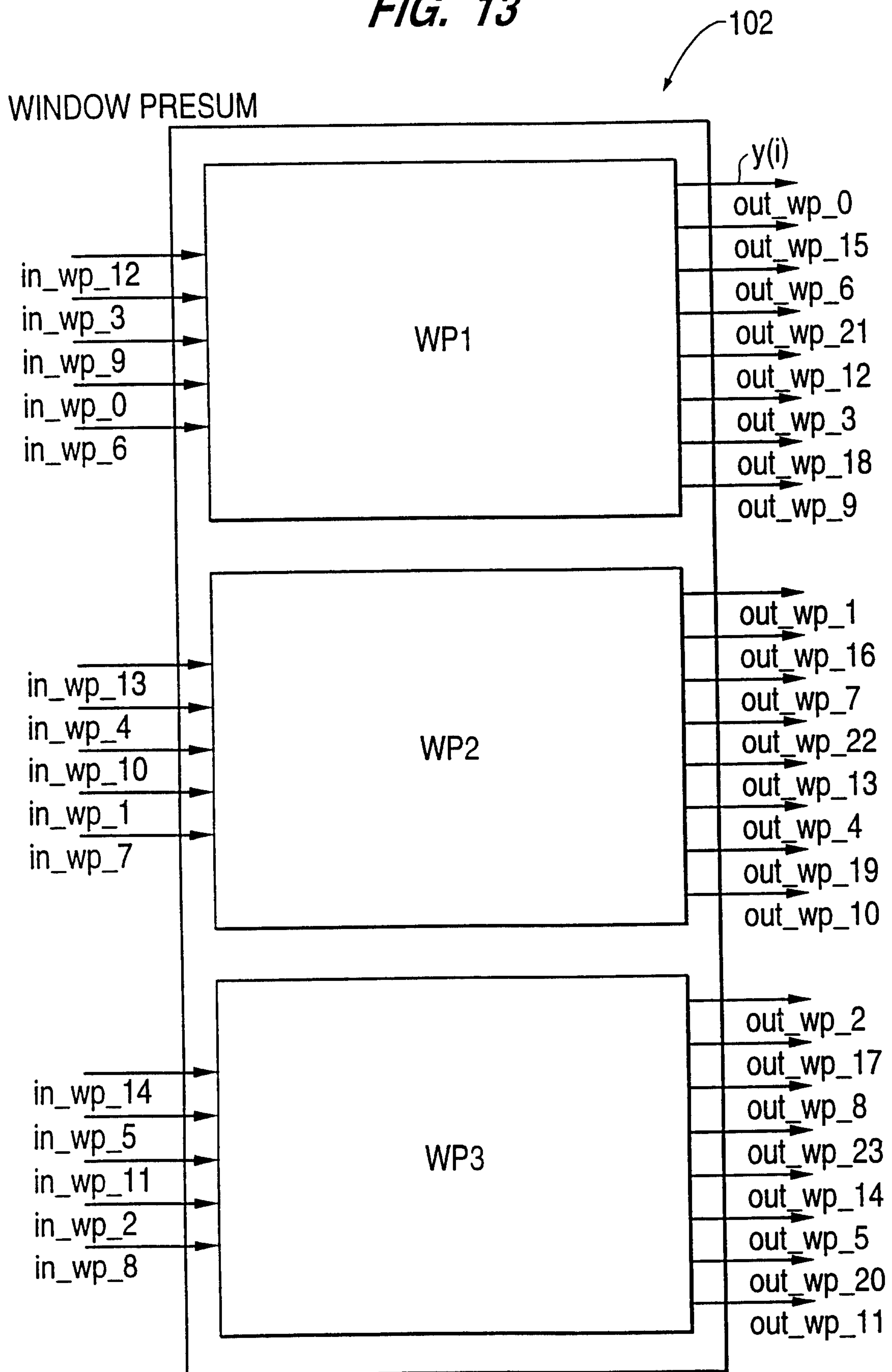
FIG. 13

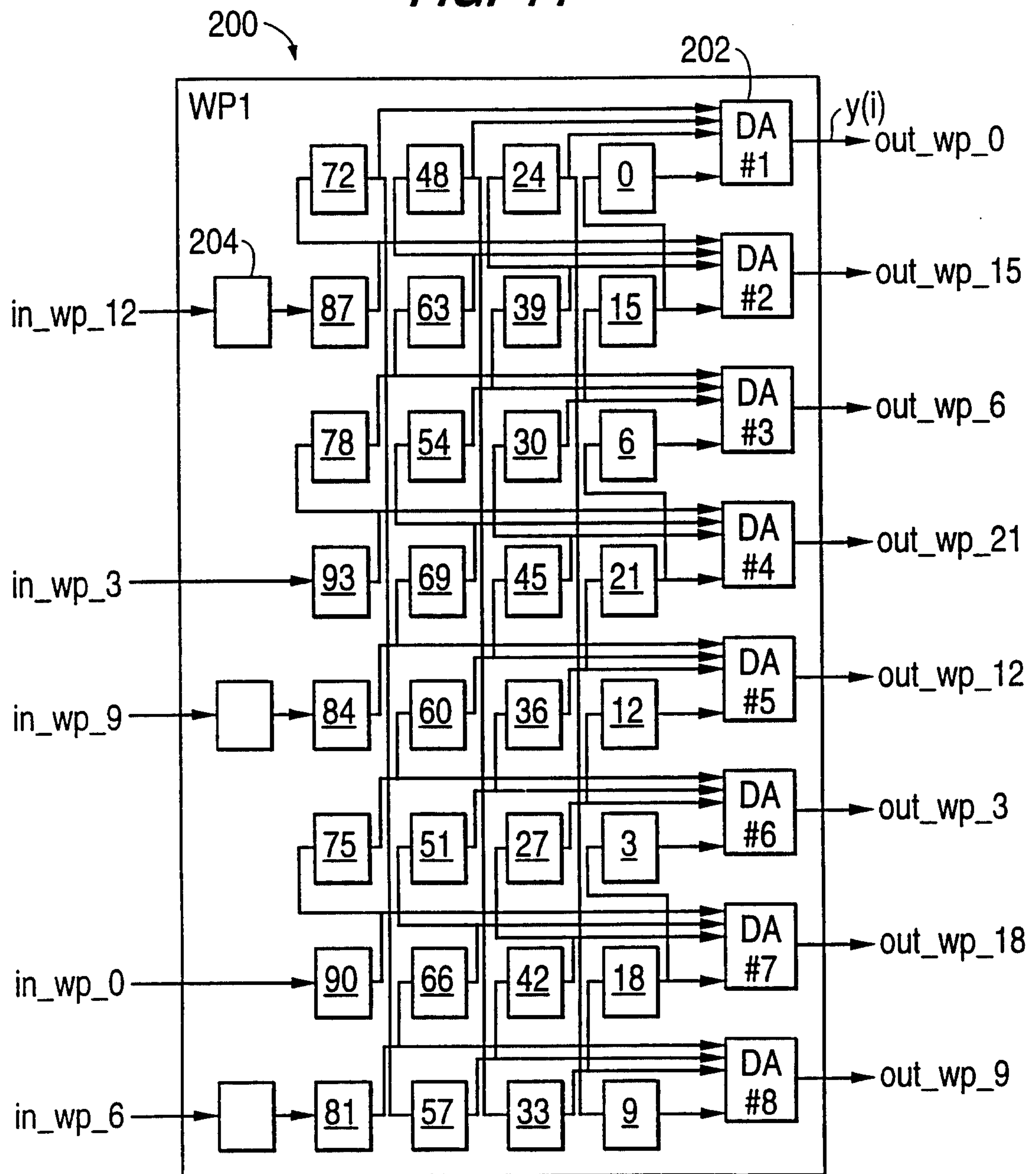
FIG. 14

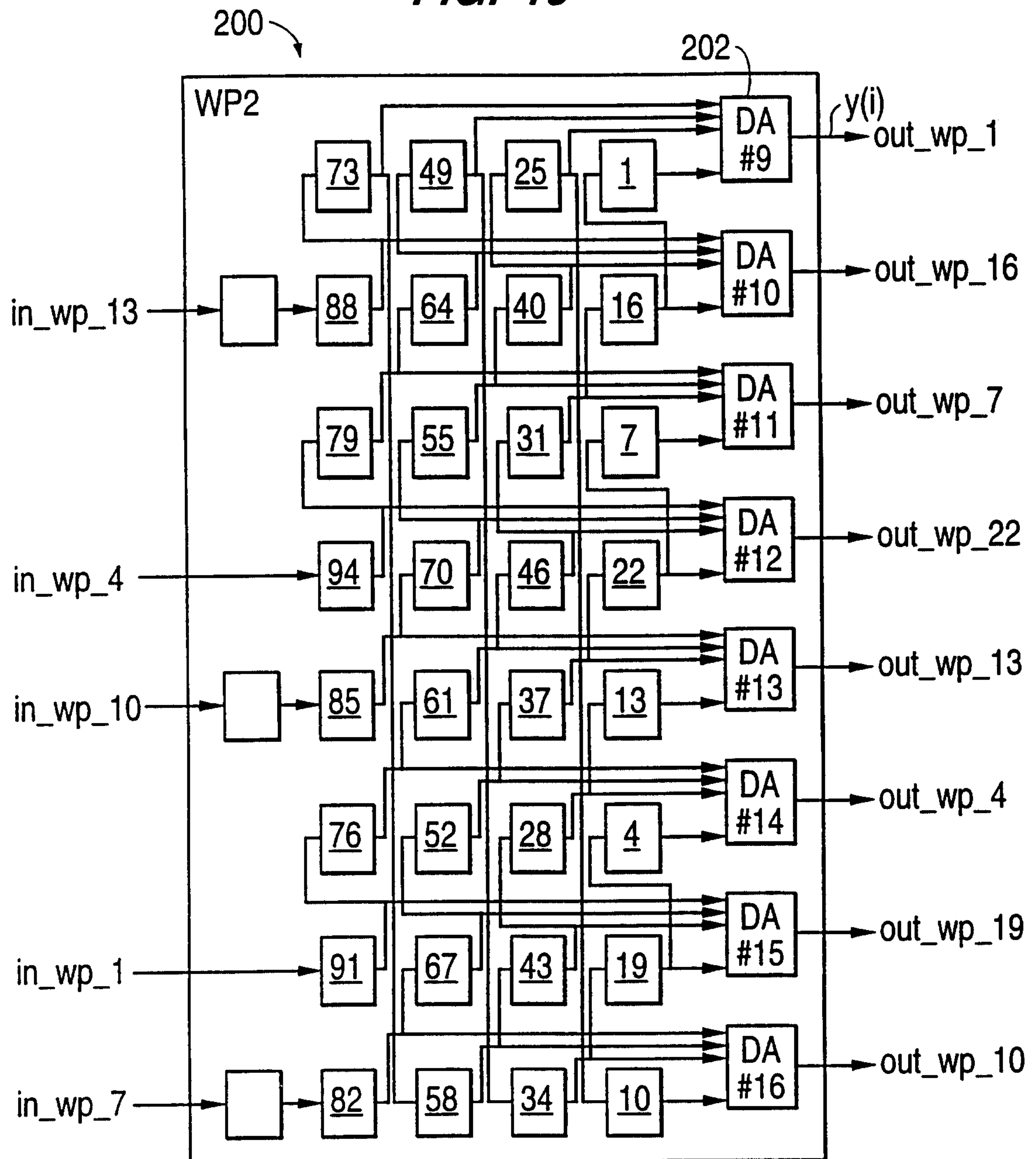
FIG. 15

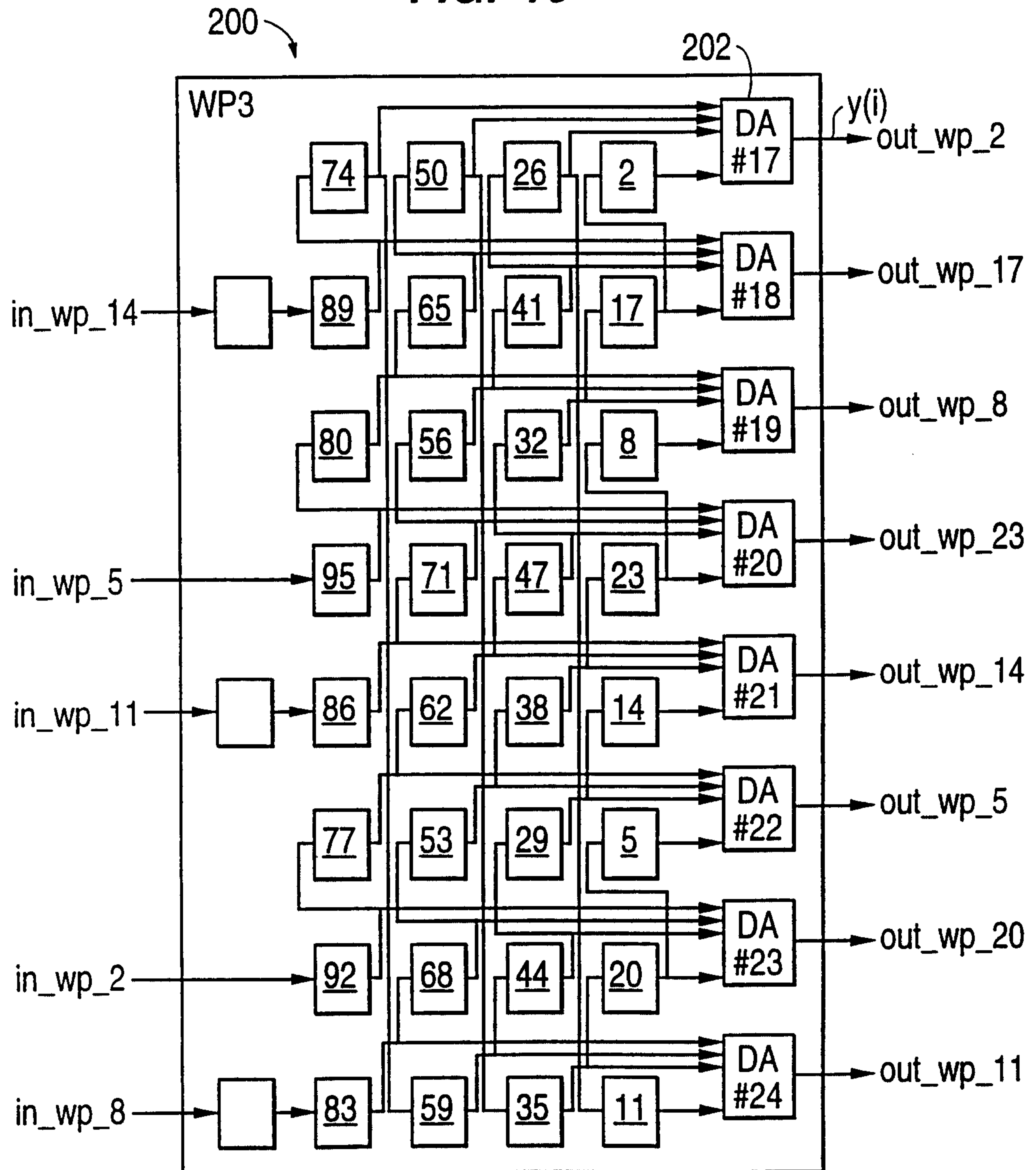
FIG. 16

FIG. 18

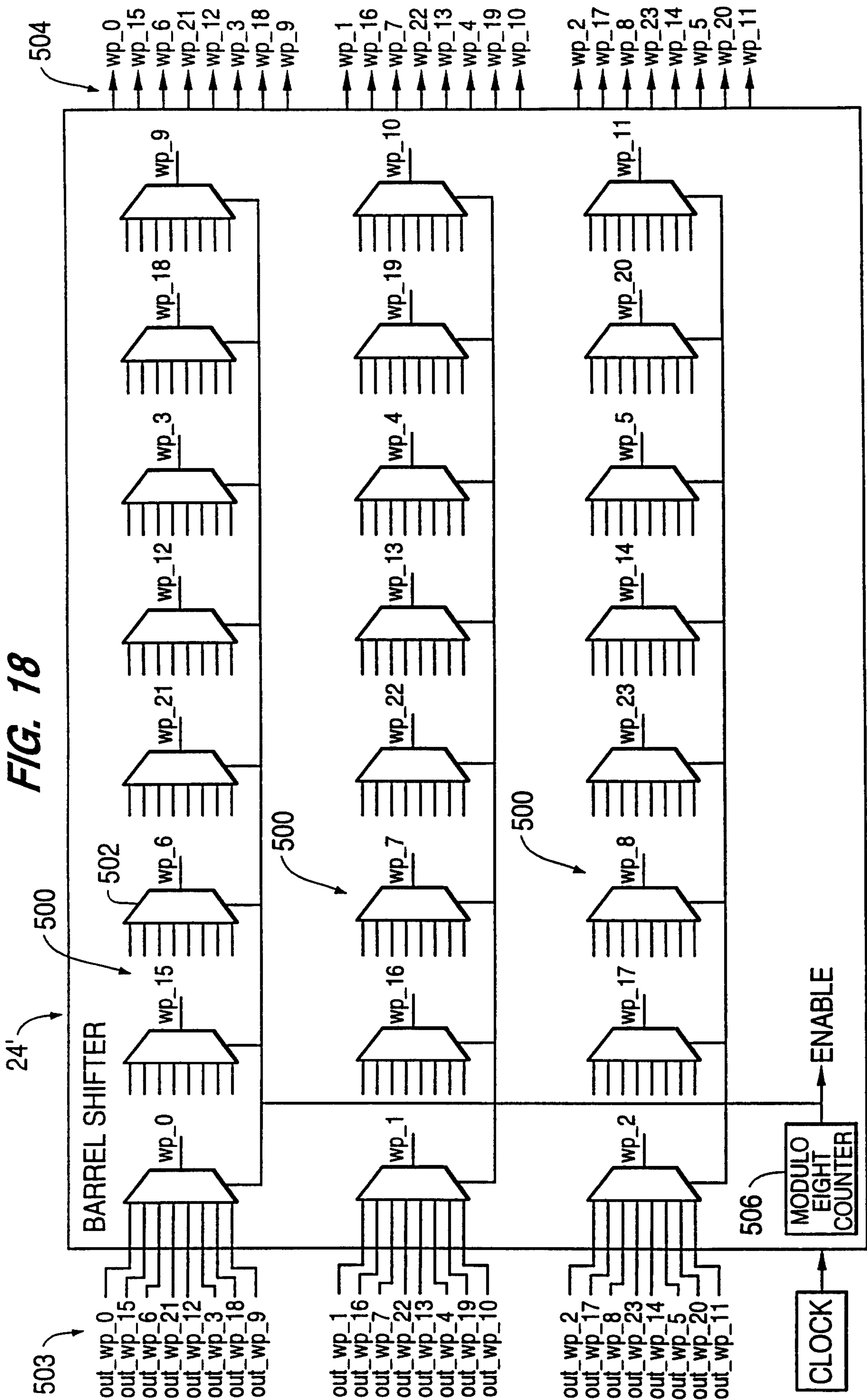
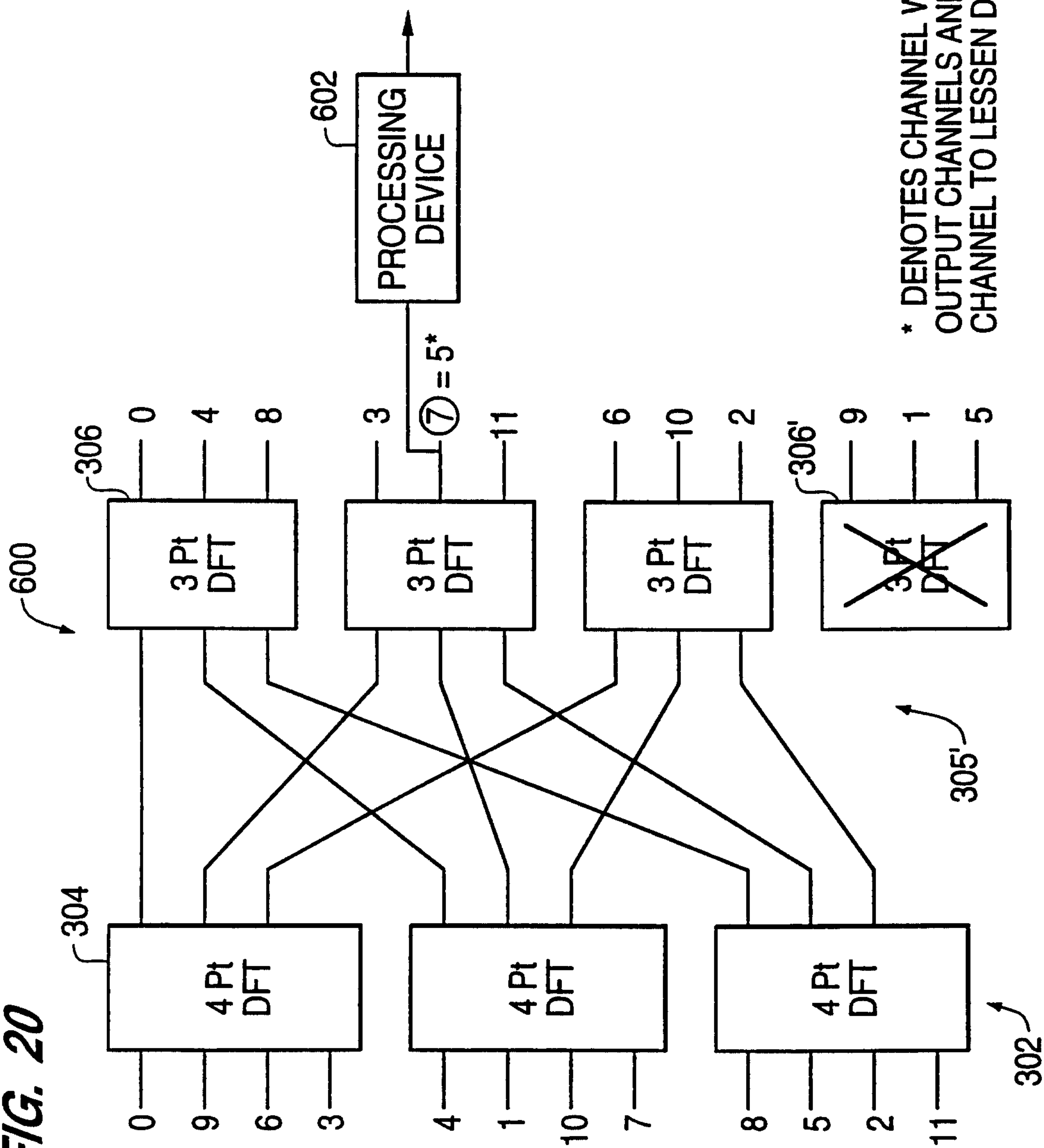


FIG. 19

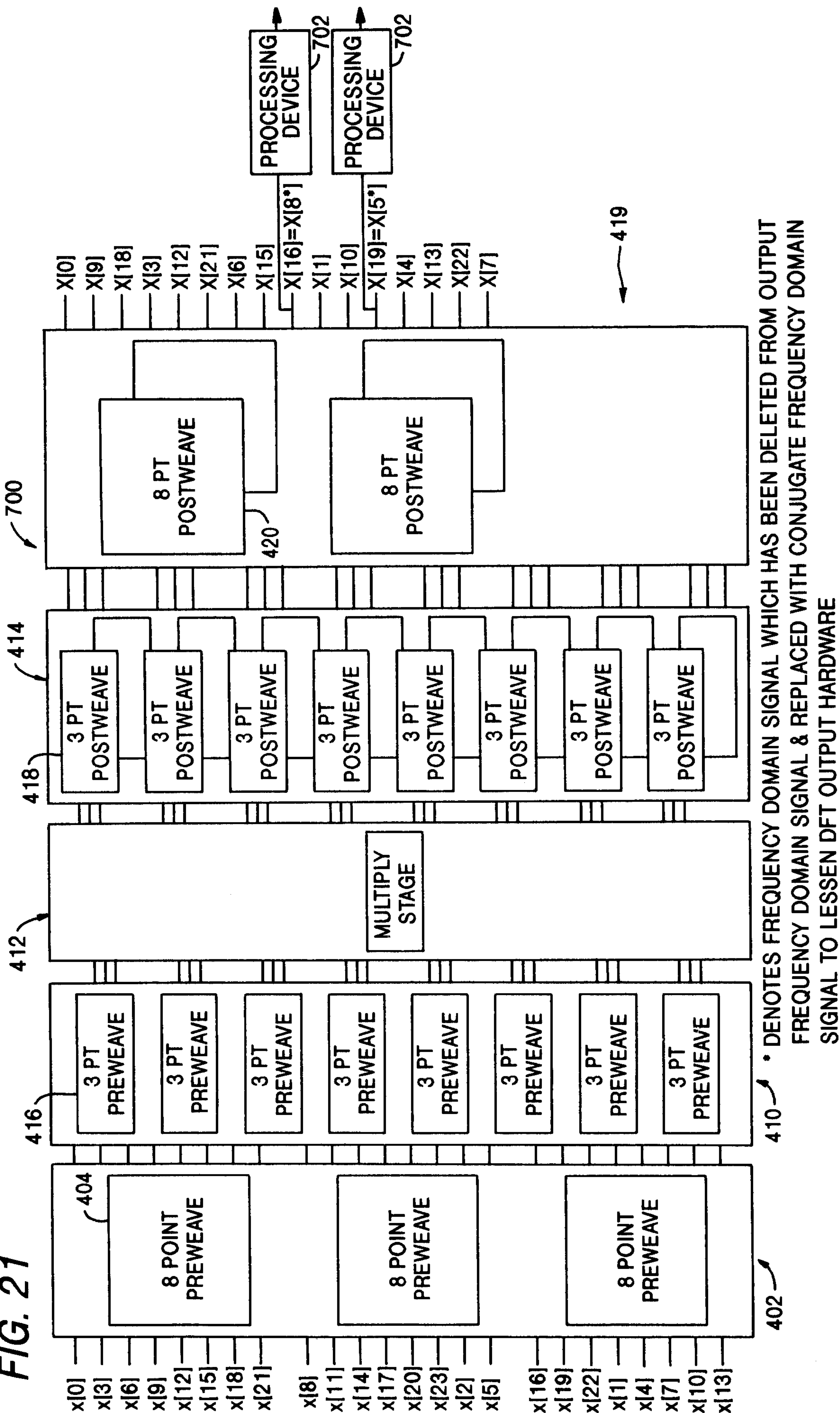
EXAMPLE:																							
M=15, N=24																							
GCD(15,24)=3																							
Q=N/GCD (15,24) = 8																							
Wp = 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23																							
CYCLE 1 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23																							
CYCLE 2 15 16 17 18 19 20 21 22 23 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14																							
CYCLE 3 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 0 1 2 3 4 5																							
CYCLE 4 21 22 23 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20																							
CYCLE 5 12 13 14 15 16 17 18 19 20 21 22 23 0 1 2 3 4 5 6 7 8 9 10 11																							
CYCLE 6 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 0 1 2																							
CYCLE 7 18 19 20 21 22 23 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17																							
CYCLE 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 0 1 2 3 4 5 6 7 8																							

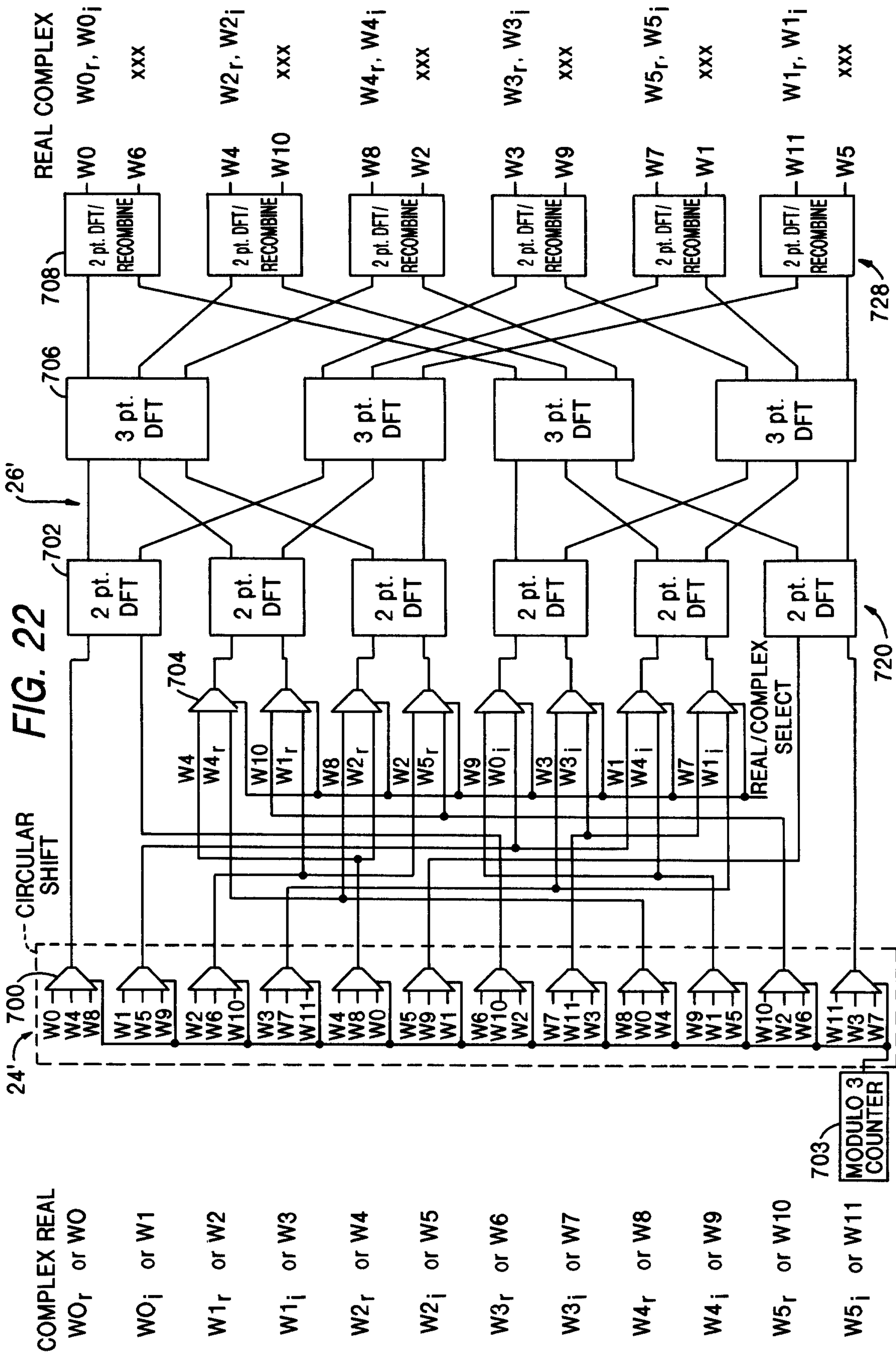
FIG. 20



* DENOTES CHANNEL WHICH HAS BEEN DELETED FROM
OUTPUT CHANNELS AND REPLACED WITH CONJUGATE
CHANNEL TO LESSEN DFT OUTPUT HARDWARE

FIG. 21





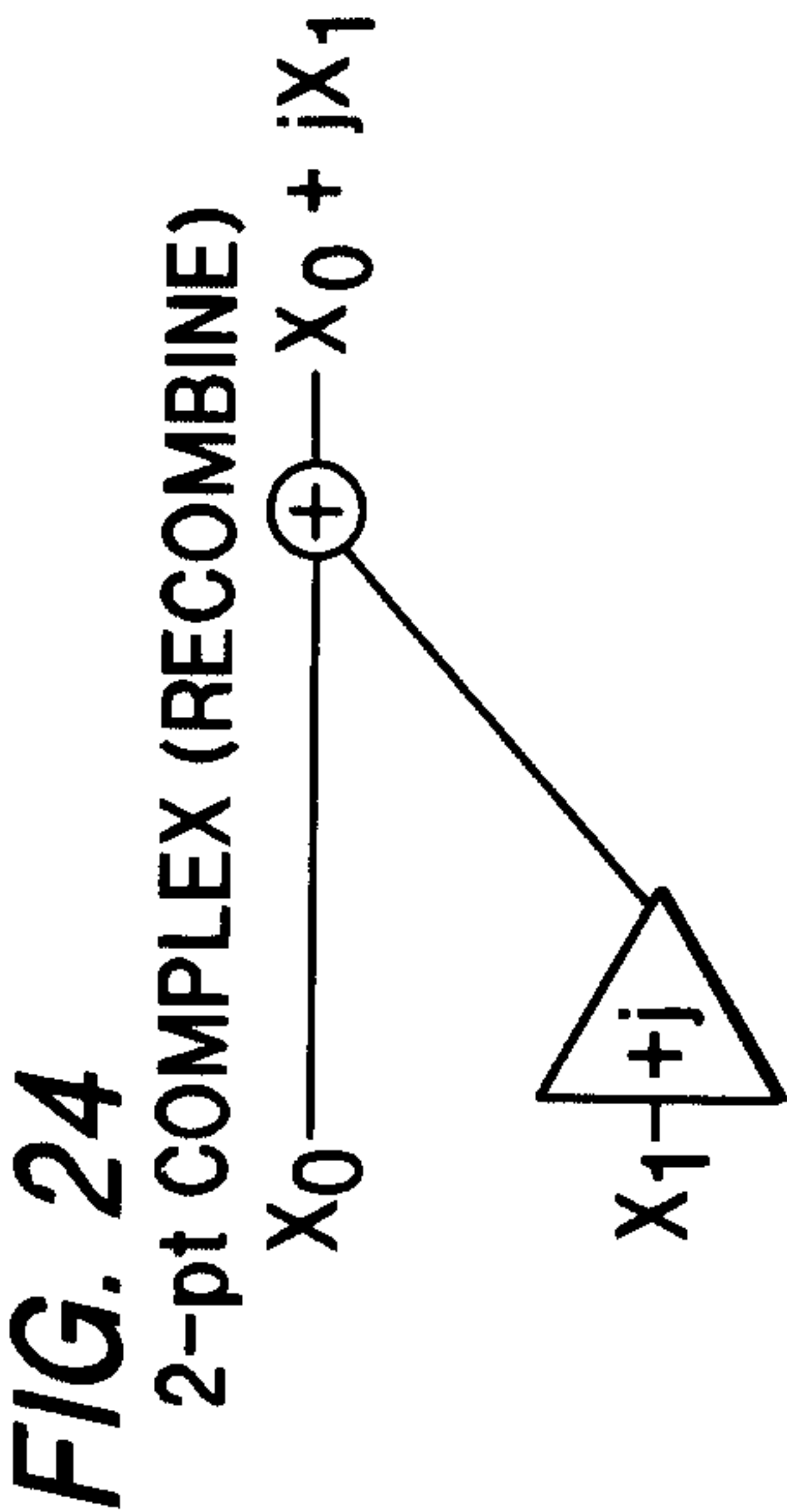
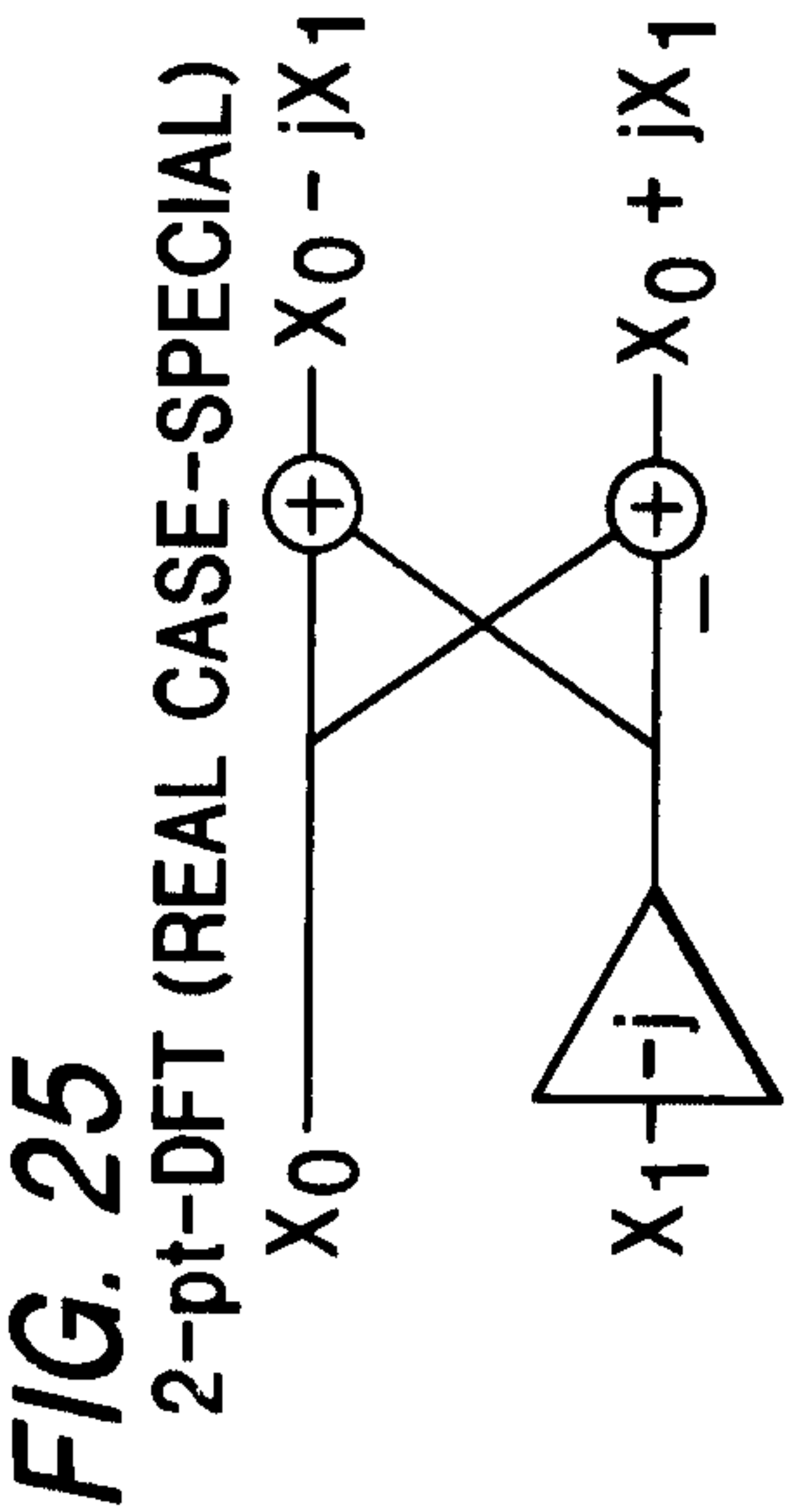
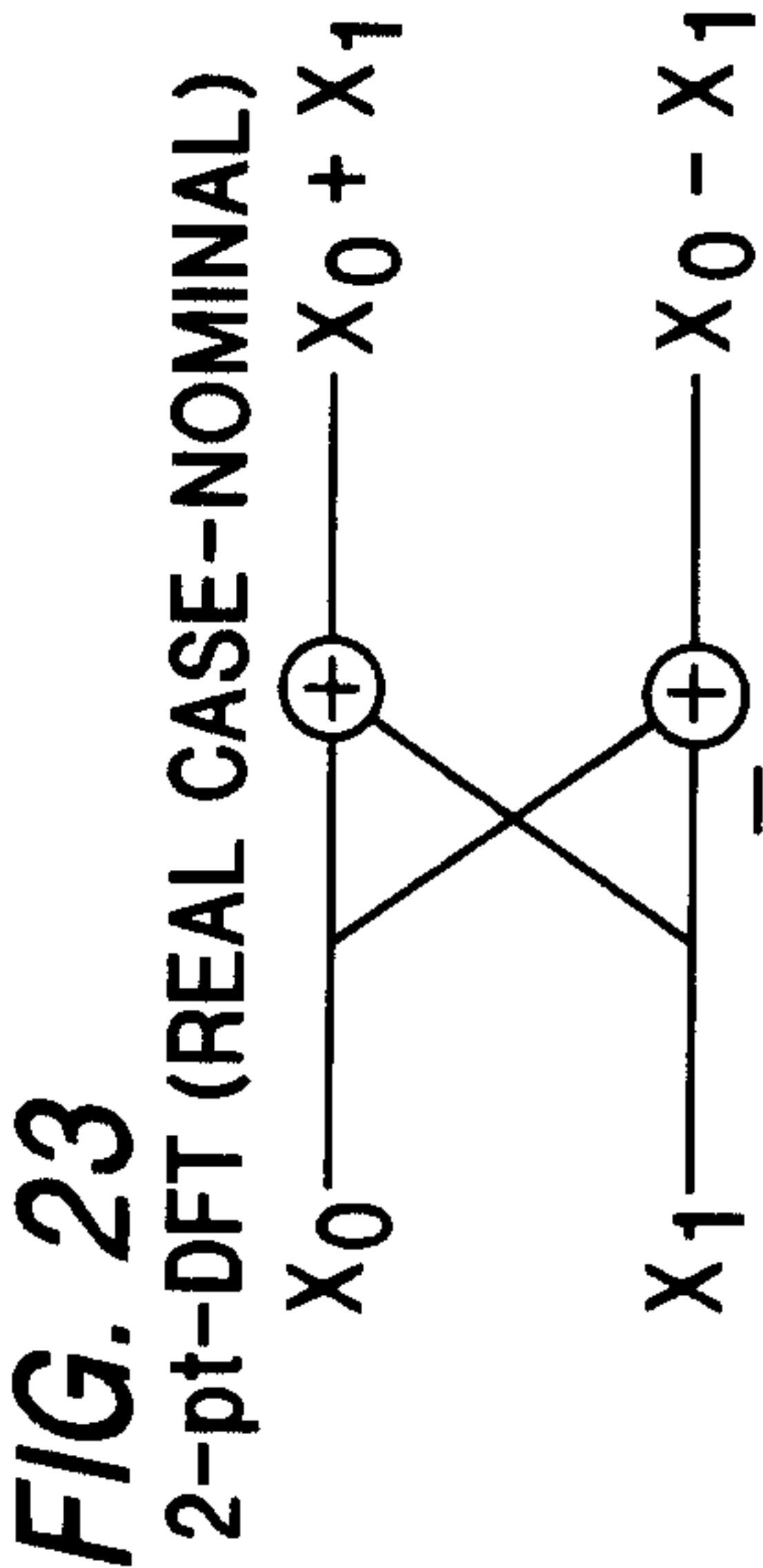


FIG. 26

