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We, Motorola, Inc., of 1303 East Algonquin Road, Schaumburg, Illincis 60196, United States of America, being the applicant in respect of Application No. 74789/94, state the following:

- 1. The nominated person has entitlement from the actual inventors by assignment from the inventors to Codex Corporation and by assignment from Codex Corporation to Motorola, Inc.
- 2. The nominated person is entitled to rely on the application listed in the declaration under Article 3 of the PCT.
- 3. The basic application referred to in the declaration under Article 8 of the PCT is the application first made in a Convention country in respect of the invention.

DATED: 2 October, 1995

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To: The Commissioner of Patents

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(57) Claim

- 1. A packet format adaptation method for ensuring portability of packet traffic from a first switching type packet network to a second switching type packet network, including the steps of:
- 1A) applying predetermined constraints to edge traffic type specific protocols to produce packets in a format compatible with the first switching type packet network and adaptable, one-to-one, to a format compatible with the second switching type packet network; and
- adapting, one-to-one, each packet from the first switching type packet network to a packet that is compatible with and hence portable across the second switching type packet network, wherein the first switching type packet network is a variable length packet network and the second switching type packet network is a fixed length packet network.

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- 11. A packet format adaptation method for ensuring portability of diverse traffic types (such as voice, video, data, etc.) from various sources of traffic from a first type packet switching network across a second type packet switching network, where the adaption at an edge of first packet network ensures that one packet is transported in said second packet network for every packet in the first network, including the steps of:
- 11A) limiting the payload size of the generated packets at an edge node of said first type packet switching network,
- 11B) receiving packets from the first type packet switching network at an inter-networking switching node connected to the second type switching network,
- 11C) translating the packets from the first type packet switching type network into packets of the second type packet switching network in accordance with a predetermined translation procedure, wherein the first type packet switching network is a variable length packet network and second type packet switching network is a fixed length packet network.



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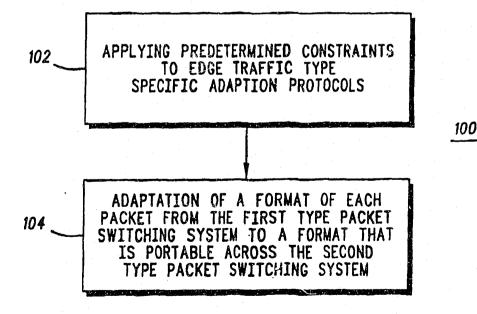
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(54) Title: FAST PACKET ADAPTATION METHOD FOR ENSURING PACKET PORTABILITY ACROSS DIVERSIFIED SWITCHING TYPE NETWORKS



### (57) Abstract

The present invention (100) provides for fast packet information transfer from a first type switching system to a diverse, i.e. different, type switching system by adaptation of a format of each fast packet from the first type switching system to a format that is portable across the second type switching system (104). For example, the present invention enables transport of fast packets from a fast packet switching system across a cell relay system and vice versa, from a fast packet switching system across an asynchronous transfer mode system and vice versa, and so on.

# FAST PACKET ADAPTATION METHOD FOR ENSURING PACKET PORTABILITY ACROSS DIVERSIFIED SWITCHING TYPE NETWORKS

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This invention is generally directed to fast packet transport and more particularly to diversified internetworking of fast packet and cell relay networks.

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Fast packet technology provides efficiency and resilience for transfer of information for networks containing a diversity of traffic types such as voice, constant bit rate (CBR), video, and multiple types of data traffic from diverse applications. Present communication systems rely on a particular network switching type to transport specific traffic types. For example, circuit switching is most appropriate for voice and other constant bit rate traffic sources. Frame switch technologies (X.25 and frame relay) are most appropriate for bursty data traffic (wide area LAN inter-connect and transaction based applications), but are not appropriate for video or voice traffic.

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The spectrum of the more well known switching technologies, listed progressively for fixed bit rate and simplicity to variable bit rate and complexity, are: circuit switching (CS), multirate circuit switching, fast circuit switching, cell relay (CR), fast packet switching (FP), and frame relay (FR) (frame switching).

FP and CR networking technologies are very similar in nature and solve many of the same networking problems. Both allow the construction of efficient and cost effective

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networks that carry diverse traffic types such as voice, video, CBR, and bursty data traffic. The major common element of FP and CR networks is that both convert (adapt) various types of traffic to a common format before transporting the traffic across the network. In both cases, the common form is a "small packet", where small is typically 64 or less octets. Each packet contains a network header containing destination or logical address, congestion level information, priority information, etc., and a payload portion containing the user's data. The main difference between a FP and CR network is that the size of the packet in a CR network is a fixed length, and the size of a packet in a FP network is variable in length.

The efficiency and diversity of CR and FP networks is obtained by the adaption of the traffic source, at the edge of the network, to a common form or packet. This allows network switches to handle the traffic in a common way inside each type of network independent of the source type. Note that at the edges of the network, different traffic types need to be adapted to the common form, or packets. These adaption procedures are traffic source type dependent. For example, CCITT AAL1 protocol is used for adaption of CBR traffic to ATM (asynchronous transfer mode) cells, and CCITT AAL5 protocol is used for adaption of HDLC framed traffic to ATM cells.

CR and FP networks both allow for efficient networks that carry diverse types of traffic. However, there has been no efficient means of inter-networking between CR and FP networks. The main hindrance is that the adaption procedures used by FP and CR networking technologies are similar, but different. Therefore, the only method of inter-networking has been to adapt the packets of one type packet to the original form, and then adapt the original form to the other packet type.

This level of processing at inter-networking points is costly and adds unacceptable transmission delay to the traffic.

As users and devices communicate with one another there develops a need to communicate also across information networking systems that use different types. Thus, there is a need for a fast packet adaptation method that ensures packet portability across diverse switching type networks.

According to one aspect of the present invention there is provided a packet format adaptation method for ensuring portability of packet traffic from a first switching type packet network to a second switching type packet network, including the steps of:

- A) applying predetermined constraints to edge traffic type specific protocols to produce packets in a format compatible with the first switching type packet network and adaptable, one-to-one, to a format compatible with the second switching type packet network; and
- B) adapting, one-to-one, each packet from the first switching type packet network to a packet that is compatible with and hence portable across the second switching type packet network, wherein the first switching type packet network is a variable length packet network and the second switching type packet network is a fixed length packet network.

According to a further aspect of the present invention there is provided a packet format adaptation method for ensuring portability of diverse traffic types (such as voice, video, data, etc.) from various sources of traffic from a first type packet switching network across a second type packet switching network, where the adaption at an edge of first packet network ensures that one packet is transported in said second packet network for every packet in the first network, including the steps of:

- A) limiting the payload size of the generated packets at an edge node of said first type packet switching network,
- B) receiving packets from the first type packet switching network at an inter-networking switching node connected to the second type switching network,
  - C) translating the packets from the first type packet switching type network into packets of the second type packet switching network in accordance

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with a predetermined translation procedure, wherein the first type packet switching network is a variable length packet network and second type packet switching network is a fixed length packet network.

According to a still further aspect of the present invention there is provided a connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:

- A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first networking edge node of the first switching type packet network,
- B) the first networking edge node, having a first edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network,
- C) the first switching type packet network, operably coupled to the first networking edge node, for transporting first type packets sequentially to at least a first inter-networking node,
- D) the first inter-networking node, operably coupled to the first switching type packet network, for operably coupling the first type packet network to at least the second type packet network and converting the first type packets to second type packets that are compatible with the second type packet network.
- E) the second type packet network, operably coupled to the first internetworking node, for transporting second type packets sequentially to at least a second inter-networking node,
- F) the second inter-networking node, operably coupled to the second type packet network, for operably coupling the second type packet network to at least a third type packet network and converting the second type packets to third type packets that are compatible with the third type packet network,
- G) the third type packet network, operably coupled to the second internetworking node, for transporting third type packets sequentially to at least a second edge node,

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- H) the second edge node, operably coupled to the third type packet network, having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and
- I) destination source(s), operably coupled to the second edge node, for receiving the traffic.

According to a still further aspect of the present invention there is provided a connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:

- A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network,
- B) the first edge node, having a first traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second switching type packets that are compatible with the second type packet network, then converting the first type packets to second type packets that are compatible with the second type packet network,
- C) the second type packet network, operably coupled to the first edge node, for transporting second type packets sequentially to at least a first internetworking node.
- D) the first inter-networking node, operably coupled to the second switching type packet network, for operably coupling the second type packet network to at least a third type packet network and converting the second type packets to third type packets that are compatible with the third switching type packet network,
- E) the third switching type packet network, operably coupled to the first inter-networking node, for transporting third type packets sequentially to at least a second edge node,



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- F) the second edge node, operably coupled to the third switching type packet network, having a second traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and
- G) destination source(s), operably coupled to the second edge node, for receiving the traffic.

According to a still further aspect of the present invention there is provided a connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:

- A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network,
- B) the first networking edge node, having a first edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network,
- C) the first switching type packet network, operably coupled to the first networking edge node, for transporting first type packets sequentially to at least a first inter-networking node,
- D) the first inter-networking node, operably coupled to the first switching type packet network, for operably coupling the first switching type packet network to at least the second switching type packet network and converting the first type packets to second type packets that are compatible with the second type packet network,
- E) the second switching type packet network, operably coupled to the first inter-networking node, for transporting second type packets sequentially to at least a second edge node,
- F) the second edge node, operably coupled to the second switching type packet network, converts the second type packets to third type packets that are compatible with a third switching type packet network, and having a second

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edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and

G) destination source(s), operably coupled to the second edge node, for receiving the traffic.

According to a still further aspect of the present invention there is provided a connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:

- A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network,
- B) the first edge node, having a first traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network, then converting the first type packets to second type packets that are compatible with the second switching type packet network,
- C) the second switching type packet network, operably coupled to the first edge node, for transporting second type packets sequentially to at least a second edge node,
- D) the second edge node, operably coupled to the second switching type packet network, converts the second type packets to third type packets that are compatible with a third switching type packet network, and having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and
- E) destination source(s), operably coupled to the second edge node, 30 for receiving the traffic.

A preferred embodiment of the present invention will now be described with reference to the accompanying drawings wherein:-



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This level of processing at inter-networking points is costly and adds unacceptable transmission delay to the traffic.

As users and devices communicate with one another there develops a need to sommunicate also across information networking systems that use different types. Thus, there is a need for a fast packet adaptation method that ensures packet portability across diverse switching type networks.

### Brief Description of the Drawings

- FIG. 1 is a flow chart showing steps taken in accordance with the method of the present invention in a connection-oriented packet communication network system having internetworking switching node(s).
- FIG. 2 is a diagram of a first embodiment of a connection-oriented communication network system in accordance with the present invention.
- FIG. 3, numeral 300, is a flow chart of the steps of a second embodiment in accordance with the method of the present invention.
- FIG. 4 shows selected further steps for encapsulating the fast packets into cells of FIG. 3.
  - FIG. 5 is a diagram of a second embodiment of a connection-oriented communication network system in accordance with the present invention.
  - FIG. 6 is a diagram of a third embodiment of a connection-oriented communication network system in accordance with the present invention.

- FIG. 7 illustrates exemplary fields for the fast packet format for the preferred embodiment, which is representative of fast packet formats in general.
- FIG. 8 illustrates a packet format for an ATM cell as specified by CCITT recommendation 1.361.
- 10 Fast packet and cell relay communication networks presently convey information (from various traffic types) from one system to a second system of a same switching type. This is done by performing traffic type specific adaption protocols which convert specific traffic types to packets, and vice versa, at edge networking nodes, and then transporting the packets across the network via packet switches of the same type.
- FIG. 1, numeral 100, is a flow chart showing steps taken 20 in accordance with the method of the present invention in a connection-oriented packet communication network system having inter-networking switching node(s). The method of the present invention allows packet information transfer from a first type packet switching system to a different second type 25 packet switching system by (a) applying predetermined constraints to the edge traffic type specific adaption protocols (102), and (b) adaptation of a format of each packet from the first type packet switching system to a format that is portable across the second type packet switching system (104). For example, the present invention allows variable 30 length fast packets from a fast packet switching system to be transported across a fixed length cell relay system and vice versa, from a fast packet switching system across an asynchronous transfer mode system and vice versa, from a cell

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relay switching system to be transported across a fast packet switch system, and so on.

FIG. 2, numeral 200, is a diagram of a first embodiment of a connection-oriented communication network system in accordance with the present invention. A traffic source (206) is a source of diverse traffic (voice, data, CBR (constant bit rate), video,...) that is operably coupled to a first networking edge node (208) of a first type packet network (210). The first edge node (208) includes a traffic adapter (224) for applying preselected constraints to edge traffic type specific adaption protocols. The first type packet network (210) is operably coupled by a first inter-networking node (218) to at least a first second type packet network (212). The second type packet network (212) is operably coupled to a third type packet network (214) by a second inter-networking node (220). The second edge node (222) includes a second traffic adapter (226) for reversing application of the preselected constraints to edge traffic type specific adaption protocols. The third type packet network (214) may be selected to be the same network as the first type packet network (210). The third type packet network (214) is operably coupled to destination device(s) (216) by a second edge node (222).

The traffic source (206) provides traffic to the first type packet network (210), and the first edge node (208) utilizes the first traffic adapter (224) to perform predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with first type packet network (210). Then, the first type packet network (210) transports the first type packets sequentially to the first inter-networking node (218) that operably couples the first type packet network (210) to the second type packet network (212). The first inter-networking node (218) adapts the first type packets to second type

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packets compatible with second type packet network (212). The second type network (212) transports the second type packets sequentially to the second inter-networking node (220) that operably couples the second type packet network (212) to the third type packet network (214), which may be selected to the the same type network as the first type packet network (210). The second inter-networking node (220) converts the second type packets from the second type network (212) to third type packets compatible with third type packet network (214). The third type packet network (214) transports the third type packets to the second edge node (222). Then, the second edge node (222) utilizes the second traffic adapter (226) to perform predetermined traffic type specific adaption protocols to convert the third type packets to the various traffic specific types and delivers traffic to destination device(s) (216). Not shown, is a possible reverse path from destination device(s) (216) to the traffic source (206), but which is a obvious extension of the above.

20 FIG. 3, numeral 300, is a flow chart of the steps of a second embodiment in accordance with the method of the present invention. The method provides for adaption of variable length fast packets such that portability of fast packet information is ensured across a cell relay network. The method assures that the adaption (typically at inter-25 networking nodes) is simple, i.e., is one-to-one, and traffic type independent. The method comprises the steps of: (A) determining a maximum allowable FP payload length (MAX\_FPPL) (302) in accordance with the function: MAX\_FPPL = CRPL-SALF-LCF, where CRPL represents cell relay payload 30 length in network, SALF represents storage area required for a length field to store the fast packet payload length, and LCF represents a length of any additional control fields, (B) enforcing a maximum fast packet length constraint (304) at a first node (for example, a fast packet traffic source edge node) 35

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where the traffic type specific adaption protocol is constrained to limit resulting fast packet payload sizes to MAX\_FPPL, (C) encapsulating the fast packets into cells at a first inter-networking node (306) (that may also be the first node), and transporting the cells across a network, and (D) mapping the cells back to fast packets at a second node (308) (for example, an edge node).

Step (A) above insures that fast packet to cell mapping is one-to-one and that no segmentation or re-assembly is required.

Typically, as shown in FIG. 4, numeral 400, encapsulating the fast packets into cells includes at least one of the following: mapping a logical address of the fast packet to a cell relay logical address (402), mapping a congestion level of the fast packet to a cell relay congestion level (404), mapping a priority level of the fast packet to a cell relay priority level (406), indicating a payload length of the fast packet in the cell (408), and copying the payload from the fast packet into the cell, and, where necessary, padding the cell to a predetermined length (410). Where the logical address of the fast packet is mapped to a cell relay logical address, a predetermined (where selected, multi-stage) table look-up procedure is typically used. Where a congestion level of the fast packet is mapped to a cell relay congestion level, and vice versa, and where a discard priority level of a fast packet is mapped discard priority level of a cell, and vice versa, the mapping may not be one to one and would typically require a boolean function to define suitable mappings between congestion and priority levels of the differing type networks.

Since variable length fast packet networks are optimized for networks with lower speed inter nodal links (56 Kbs to 2 Mbs), and cell relay networks are optimized for networks with

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high speed inter nodal links, it is not uncommon for fast packet networks to include more congestion and/or priority levels than cell relay networks. When this is the case, the multiple congestion and priority levels of the fast packet can be (also) mapped into a control field in the payload portion of the cell at an inter-networking node. In this case, the mapping of cells to fast packets at inter networking point would extract fast packet priority and congest information from the control field, combine through boolean function with cell congestion and priority information to derive new appropriate fast packet congestion and priority level information for fast packet network. Where the payload from the fast packet is copied into the cell, padding may be zero padding, or alternatively multiple code words, for example, the code word 01101010, binary, repeated. In addition, network control information may be mapped from the header of the fast packet into the payload of cell relay packet.

FIG. 5, numeral 500, is a diagram of a second embodiment of a connection-oriented communication network system in accordance with the present invention. The system ensures portability of packet traffic from a first switching type network to a second switching type network and includes: A) traffic source(s) (502), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network, B) the first edge node (504), having a first traffic adapter (506) for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with a first type packet network, where the packets are also convertible oneto-one to second type packets that are compatible with the second type packet network, C) the first edge node (504) converting the first type packets to second type packets that are compatible with the second type packet network, D) the

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second type packet network (508), operably coupled to the first edge node (504), for transporting second type packets sequentially to at least a first inter-networking node (510). E) the first inter-networking node (510), operably coupled to the second type packet network, for operably coupling the second type packet network (508) to at least a third type packet network (512) and converting the second type packets to third type packets that are compatible with the third type packet network (512). F) the third type packet network (512). operably coupled to the first inter-networking node (510), for transporting third type packets sequentially to at least a second edge node (514), G) the second edge node (514), operably coupled to the third type packet network (512), having a second traffic adapter (516) for performing predetermined traffic type specific adaption protocols to convert the third type packets to the diverse traffic types that are compatible with destination source(s), and H) destination source(s) (518), operably coupled to the second edge node, for receiving the traffic.

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Again, the third type packet network (512) may be selected to be a packet network of the first type such that, also, the third type of packets are selected to be packets of the first type.

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FIG. 6, numeral 600, is a diagram of a third embodiment of a connection-oriented communication network system in accordance with the present invention. The system ensures portability of packet traffic from a first switching type network to a second switching type network and includes: A) traffic source(s) (602), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network, B) the first edge node (604), having a first traffic adapter (606) for

performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with a first type packet network, where the packets are also convertible oneto-one to second type packets that are compatible with a 5 second type packet network. C) the first type packet network (608), operably coupled to the first edge node (604), for transporting first type packets sequentially to at least a first inter-networking node (610), D) the first inter-networking node (610), operably coupled to the first type packet network, 10 for operably coupling the first type packet network (608) to at least a second type packet network (612) and converting the first type packets to second type packets that are compatible with the second type packet network (612), E) the second type packet network (612), operably coupled to the first inter-15 networking node (610), for transporting second type packets sequentially to at least a second edge node (614), F) the second edge node (614), operably coupled to the second type packet network (612) converts the second type packets to third type 20 packets that are compatible with a third type packet network, G) the second edge node (614) having a second traffic adapter (616) for performing predetermined traffic type specific adaption protocols to convert the third type packets to the diverse traffic types that are compatible with destination source(s), and H) destination source(s) (618), operably coupled 25 to the second edge node, for receiving the traffic.

Again, the third type packet network may be selected to be a packet network of the first type (608) such that, also, the third type of packets are selected to be packets of the first type.

Alternatively, the connection-oriented communication network system for ensuring portability of packet traffic from

a first switching type network to a second switching type network, may be described as follows. The system comprises: A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network, B) the first 5 networking edge node, having a first edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with the first type packet network, where the packets are also convertible one-to-one to second type packets 10 that are compatible with the second type packet network, C) the first type packet network, operably coupled to the first networking edge node, for transporting first type packets sequentially to at least a first inter-networking node, D) the 15 first inter-networking node, operably coupled to the the first type packet network, for operably coupling the first type packet network to at least the second type packet network and converting the first type packets to second type packets that are compatible with the second type packet network, E) the 20 second type packet network, operably coupled to the first inter-networking node, for transporting second type packets sequentially to at least a second edge node, F) the second edge node, operably coupled to the the second type packet network, converts the second type packets to third type packets that are 25 compatible with a third type packet network, and having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and G) destination source(s), operably coupled to the second edge node, for 30 receiving the traffic. Clearly, the third type of packets may be selected to be packets of the first type where the third type packet network is a packet network of the first type.

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In addition, the system may be described as: a connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, comprising: A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network, B) the first edge node, having a first traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with the first type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second type packet network, then converting the first type packets to second type packets that are compatible with the second type packet network. C) the second type packet network, operably coupled to the first edge node, for transporting second type packets sequentially to at least a second edge node, D) the second edge node, operably coupled to the the second type packet network, converts the second type packets to third type packets that are compatible with a third type packet network, and having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and E) destination source(s), operably coupled to the second edge node, for receiving the traffic. Again, the third type of packets may be selected to be packets of the first type where the third type packet network is a packet network of the first type.

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FIG. 7, numeral 700, illustrates exemplary fields for the fast packet format (701) for the preferred embodiment, which is representative of fast packet formats in general. In the figure, several fields are shown. The first field is a four bit S field (702) and contains 2 bits to represent up to four levels of

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packet discard priority information, and 2 bits to represent up to four levels of network congestion information. The packet discard priority bits are set appropriately at the source edge node by the process that adapts the diverse traffic types to the common fast packet format (the so called fast packet adaption layer). The congestion bits are also initialized to the uncongested state at the source edge node, but can be modified appropriately, to indicate increasing levels of network congestion, by intermediate packet switches as the packet traverses the fast packet network. The LCN field (704) is the fast packet logical channel number and is used to differentiate various connections on connecting links between packet switches. The LCN in the preferred embodiment is a 12 bit field and allows 4096 connections per connecting link. The ECC field is an 8 bit error correction code (706) that protects the LCN field. The ECC can correct for single bit transmission errors in the LCN, and can detect multiple bit errors. The last field is the fast packet payload (708), and continue user's fast packet adapted traffic. Note that the S, LCN, and ISCC fields makeup the so called fast packet header, which is always 3 bytes in length, and the payload portion may be of variable size, from 1 to MAX FPPL bytes. Therefore the fast packet size may vary from 4 bytes to MAX\_FPPL+3 bytes.

FIG. 8 illustrates a packet format (800) for an ATM cell (801) as specified by CCITT recommendation I.361. The GFC (802) field is a 4 bit field that is presently for further study at the present time (encoded as 0000). The VPI field (804) is an 8 bit virtual path indentifier. The VCI field (806) is a 16 bit virtual channel indentifier. The VPI and VCI fields are used jointly to differentiate various connections on connecting links between packet (cell relay) switches (up to 16,777,216 connects per connecting link). The PTI field (808) is the payload type indicator field and is used to indicate up to 2 levels of network congestion. The PTI also contains other

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information not directly relevant to this invention. The PTI field is a 3 bit field, refer to CCITT recommendation 1.361 for further information. The single bit CLP (809) field contains up to 2 levels of cell discard priority information. The PTI and CLP bits are initialized by the source edge node by the process that adapts the diverse traffic types to the common ATM cell format (the so called ATM adaption layer). Both the PTI and CLP fields can be modified by intermediate cell relay switching nodes if the cell experiences congestion as it traverses the cell relay network. The HEC field (810) is the header error control field and is an 8 bit field which contains an error detection and correction code that protects the GFC, VPI, VCI, PTI, and CLP fields. The HEC field can be used to correct a single bit transmission error in these fields, or to detect multiple bit errors. The payload field (812) contains the user's cell adapted traffic. The payload size is fixed, and contains 48 bytes. Note that the GFC, VPI, VCI, PTI, CLP, and HEC fields makeup the so called cell header. The cell header is therefore always 5 bytes in length, the cell payload is always 48 bytes in length, and the total cell size is therefore always 53 bytes.

Although exemplary embodiments are described above, it will be obvious to those skilled in the art that many alterations and modifications may be made without departing from the invention. Accordingly, it is intended that all such alterations and modifications be included within the spirit and scope of the invention as defined in the appended claims.



### THE CLAIMS DEFINING THE INVENTION ARE AS FOLLOWS:

- 1. A packet format adaptation method for ensuring portability of packet traffic from a first switching type packet network to a second switching type packet network, including the steps of:
- 1A) applying predetermined constraints to edge traffic type specific protocols to produce packets in a format compatible with the first switching type packet network and adaptable, one-to-one, to a format compatible with the second switching type packet network; and
- 1B) adapting, one-to-one, each packet from the first switching type packet network to a packet that is compatible with and hence portable across the second switching type packet network, wherein the first switching type packet network is a variable length packet network and the second switching type packet network is a fixed length packet network.
  - 2. The packet format adaptation method of claim 1 wherein the steps of applying and adapting further comprise the steps of:
  - 2A) determining a maximum allowable variable length packet payload length (MAX\_FPPL) (applying constraints-Step 1A) in accordance with the function: MAX\_FPPL = CRPL-SALF-LCF, where CRPL represents the fixed length packet payload length in the fixed length packet network, SALF represents storage area required for a length field to store the variable length packet payload length, and LCF represents a length of any additional control fields,
  - 2B) enforcing a maximum variable length packet length constraint (applying constraints-Step 1A) at a variable length packet traffic source edge node where the traffic type specific adaption protocol is constrained to limit resulting variable length packet payload sizes to MAX\_FPPL, and
  - 2C) encapsulating the variable length packets into fixed length packets at a first inter-networking node (adapting-Step 1B).
  - 3. The packet format adaptation method of claim 2 wherein at leat one of:
  - 3A) a format of the variable length packets includes a variable length packet header and a payload field, and



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- 3B) a format of the fixed length packets includes a fixed length packet header, a control field, a length field, and a payload field.
- 4. The packet format adaptation method of comm 2 wherein encapsulating the variable length packets into fixed length packets further includes, for each packet, at least one of:
- 4A) mapping a logical address of the variable length packet to a fixed length packet logical address.
- 4B) mapping a congestion level of the variable length packet to a fixed length packet congestion level,
- 10 4C) mapping a priority level of the variable length packet to a fixed length packet priority level,
  - 4D) indicating a payload length of the variable length packet in the fixed length packet,
  - 4E) copying the payload from the variable length packet into the fixed length packet, and, where necessary, padding the fixed length packet to a predetermined length,
  - 4F) where step (4A) is utilized, further including using a predetermined multi-stage table look-up procedure for the mapping of step (4A),
  - 4G) where step (4B) is utilized, further including mapping a fixed length packet congestion field in fixed length packet header for 2 levels of congestion as a function of a variable length packet congestion field in the variable length packet header for 2 or more levels of congestion in the variable length packet network.
  - 4H) where step (4C) is utilized, further including mapping a fixed length packet loss priority field (CLP) in the fixed length packet header for 2 levels of fixed length packet discard priority as a function of a variable length packet discard priority field in the variable length packet header for 2 or more levels of discard priority in the variable length packet network, and
    - 4I) where step (4E) is utilized, padding utilizing multiple code words.
- 30 5. The packet format adaptation method of claim 2 wherein a first internetworking node performs translation of the variable length packets into fixed





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length packets, and a second inter-networking node subsequently performs translation of said fixed length packets back into variable length packets.

- 6. The packet format adaptation method of claim 5 wherein said subsequent translating the fixed length packets into variable length packets further includes, for each packet, at least one of:
- 6A) mapping a logical address of the fixed length packet to a variable length packet logical address,
- 6B) mapping a congestion level of the fixed length packet to a variable length packet congestion level,
- 6C) mapping a priority level of the fixed length packet to a variable length packet priority level,
  - 6D) reading a payload length indication field in the fixed length packet payload to determine the payload length of the variable length packet,
- 6E) copying the payload from the fixed length packet into the variable length packet,
  - 6F) where step (6A) is utilized, further including using a predetermined, possibly multi-stage, table look-up procedure for the mapping of step (6A),
  - 6G) where step (6B) is utilized, further including mapping a variable length packet congestion field in variable length packet header for 2 or more levels of congestion as a function of a fixed length packet congestion field in the fixed length packet header for 2 levels of congestion in the fixed length packet network, and
  - 6H) where step (6C) is utilized, further including mapping a variable length packet loss priority field in the variable length packet header for 2 or more levels of variable length packet discard priority as a function of a fixed length packet discard priority field in the fixed length packet header for 2 levels of discard priority in the fixed length packet network.
  - 7. The packet format adaptation method of claim 5 wherein:
- 7A) said first inter-networking node includes mapping a congestion and loss priority field from the header of the variable length packet into a control field in the payload of the fixed length packet, and





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- 7B) said second inter-networking node includes mapping the variable length packet congestion and loss priority field in the variable length packet header to be a Boolean function of said control information field in fixed length packet payload and fixed length packet congestion and loss priority information in fixed length packet header.
- 8. The packet format adaptation method of claim 4 wherein the padding is the code word 01101010, binary, repeated.
- 9. The packet format adaptation method of claim 2 wherein:
  - 9A) MAX FPPL is equal to 46 octets,
- 10 9B) CRPL = 48 octets,

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- 9C) SALF = 1 octet, and
- 9D) LCF = 1 octet.
- 10. The packet format adaptation method of claim 2 wherein:
  - 10A) MAX FPPL is equal to 47 octets,
- 15 10B) CRPL = 48 octets,
  - 10C) SALF = 1 octet, and
  - 10D) LCF = 0 octets.
  - 11. A packet format adaptation method for ensuring portability of diverse traffic types (such as voice, video, data, etc.) from various sources of traffic from a first type packet switching network across a second type packet switching network, where the adaption at an edge of first packet network ensures that one packet is transported in said second packet network for every packet in the first network, including the steps of:
  - 11A) limiting the payload size of the generated packets at an edge node of said first type packet switching network,
    - 11B) receiving packets from the first type packet switching network at an inter-networking switching node connected to the second type switching network,
  - 11C) translating the packets from the first type packet switching type network into packets of the second type packet switching network in accordance with a predetermined translation procedure, wherein the first type packet switching network is a variable length packet network and second type packet switching network is a fixed length packet network.

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- 12. The packet format adaptation method of claim 11 wherein the step of limiting comprises the steps of:
- 12A) determining a maximum allowable variable packet length payload length (MAX\_FPPL) in accordance with the function: MAX\_FPPL = CRPL-SALF-LCF, where CRPL represents fixed length packet payload length in network, SALF represents storage area required for a length field to store the variable length packet payload length, and LCF represents a length of any additional control fields.
- 12B) enforcing a maximum variable length packet length constraint at a variable length packet traffic source edge node where the traffic type specific adaption protocol is constrained to limit resulting variable length packet payload sizes to MAX FPPL, and

the step of translating comprises the step of:

- 12C) encapsulating the variable length packets into fixed length packets at a first inter-networking node.
  - 13. The packet format adaptation method of claim 12 wherein at least one of:
  - 13A) a format of the variable length packets (FPs) includes a variable length packet header and a payload field, and
- 13B) a format of the fixed length packets includes a fixed length packet header, a control field, a length field, and a payload field.
  - 14. The packet format adaptation method of claim 12 wherein translating the variable length packets into fixed length packets further includes, for each packet, at least one of:
- 14A) mapping a logical address of the variable length packet to a fixed length packet relay logical address,
  - 14B) mapping a congestion level of the variable length packet to a fixed length packet congestion level,
  - 14C) mapping a priority level of the variable length packet to a fixed length packet priority level,
- 30 14D) indicating a payload length of the variable length packet in the fixed length packet,



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- 14E) copying the payload from the variable length packet into the fixed length packet, and, where necessary, padding the fixed length packet to a predetermined length,
- 14F) where step (14A) is utilized, further including using a predetermined, where selected multi-stage, table look-up procedure for the mapping of step (14A),
- 14G) where step (14B) is utilized, further including mapping a fixed length packet congestion field in the fixed length packet header for 2 levels of congestion as a function of a variable length packet congestion field in the variable length packet header for 2 or more levels of congestion in the variable length packet network,
- 14H) where step (14C) is utilized, further including mapping a fixed length packet loss priority field (CLP) in the fixed length packet header for 2 levels of fixed length packet discard priority as a function of a variable length packet discard priority field in the variable length packet header for 2 or more levels of discard priority in the variable length packet network, and
  - 14I) where step (14E) is utilized, padding utilizing multiple code words.
- 15. The packet format adaptation method of claim 12 wherein a first internetworking node performs translation of the variable length packets into fixed length packets, and a second inter-networking node subsequently performs translation of said fixed length packets back into variable length packets.
- 16. The packet format adaptation method of claim 15 wherein said subsequent translating the fixed length packets into variable length packets further includes, for each packet, at least one of:
- 16A) mapping a logical address of the fixed length packet to a variable length packet logical address,
- 16B) mapping a congestion level of the fixed length packet to a variable length packet congestion level,
- 16C) mapping a priority level of the fixed length packet to a variable 30 length packet priority level,
  - 16D) reading a payload length indication field in the fixed length packet payload to determined the payload length of the variable length packet,

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- 16E) copying the payload from the fixed length packet into the variable length packet,
- 16F) where step (16A) is utilized, further including using a predetermined, possibly multi-stage, table look-up procedure for the mapping of step (16A),
- 16G) where step (16B) is utilized, further including mapping a variable length packet congestion field in variable length packet mader for 2 or more levels of congestion as a function of a fixed length packet congestion field in the fixed length packet header for 2 levels of congestion in the fixed length packet network, and
- 10 16H) where step (16C) is utilized, further including mapping a variable length packet loss priority field in the variable length packet header for 2 or more levels of variable length packet discard priority as a function of a fixed length packet discard priority field in the fixed length packet header for 2 levels of discard priority in the fixed length packet network.
- 15 17. The packet format adaption method of claim 15 wherein:
  - 17A) said first inter-networking node includes mapping a congestion and loss priority field from the header of the variable length packet into a control field in the payload of the fixed length packet, and
  - 17B) said second inter-networking node includes mapping variable length packet congestion and loss priority field in the variable length packet header to be a Boolean function of said control information field in fixed length packet payload and fixed length packet congestion and loss priority information in fixed length packet header.
  - 18. The packet format adaptation method of claim 14 wherein the padding is the code word 01101010, binary, repeated.
    - 19. The packet format adaptation method of claim 12 wherein:
      - 19A) MAX\_FPPL is equal to 46 octets,
      - 19B) CRPL = 48 octets,
      - 19C) SALF = 1 octet, and
- 30 19D) LCF = 1 octet.
  - 20. The packet format adaptation method of claim 12 wherein:
    - 20A) MAX\_FPPL is equal to 47 octets,



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- 20B) CRPL = 48 octets,
- 20C) SALF = 1 octet, and
- 20D) LCF = 0 octets.
- 21. A connection-oriented communication network system for ensuring
   5 portability of packet traffic from a first switching type network to a second switching type network, including:
  - 21A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first networking edge node of the first switching type packet network,
  - 21B) the first networking edge node, having a first edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network,
  - 21C) the first switching type packet network, operably coupled to the first networking edge node, for transporting first type packets sequentially to at least a first inter-networking node,
  - 21D) the first inter-networking node, operably coupled to the first switching type packet network, for operably coupling the first type packet network to at least the second type packet network and converting the first type packets to second type packets that are compatible with the second type packet network,
  - 21E) the second type packet network, operably coupled to the first internetworking node, for transporting second type packets sequentially to at least a second inter-networking node,
  - 21F) the second inter-networking node, operably coupled to the second type packet network, for operably coupling the second type packet network to at least a third type packet network and converting the second type packets to third type packets that are compatible with the third type packet network,
  - 21G) the third type packet network, operably coupled to the second internetworking node, for transporting third type packets sequentially to at least a second edge node,



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- 21H) the second edge node, operably coupled to the third type packet network, having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and
- 5 21I) destination source(s), operably coupled to the second edge node, for receiving the traffic.
  - 22. The connection-oriented communication network system of claim 21 wherein the third type of packets are packets of the first type, and the third switching type packet network is a packet network of the first type.
- 10 23. A connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:
  - 23A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network,
  - 23B) the first edge node, having a first traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second switching type packets that are compatible with the second type packet network, then converting the first type packets to second type packets that are compatible with the second type packet network,
  - 23C) the second type packet network, operably coupled to the first edge node, for transporting second type packets sequentially to at least a first internetworking node,
  - 23D) the first inter-networking node, operably coupled to the second switching type packet network, for operably coupling the second type packet network to at least a third type packet network and converting the second type packets to third type packets that are compatible with the third switching type packet network,



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- 23E) the third switching type packet network, operably coupled to the first inter-networking node, for transporting third type packets sequentially to at least a second edge node,
- 23F) the second edge node, operably coupled to the third switching type packet network, having a second traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and
- 23G) destination source(s), operably coupled to the second edge node, for receiving the traffic.
- 10 24. The connection-oriented communication network system of claim 23 wherein the third type of packets are packets of the first type, and the third type packet network is a packet network of the first type.
  - 25. A connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:
  - 25A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network.
  - 25B) the first networking edge node, having a first edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network,
  - 25C) the first switching type packet network, operably coupled to the first networking edge node, for transporting first type packets sequentially to at least a first inter-networking node,
  - 26D) the first inter-networking node, operably coupled to the first switching type packet network, for operably coupling the first switching type packet network to at least the second switching type packet network and converting the first type packets to second type packets that are compatible with the second type packet network,





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25E) the second switching type packet network, operably coupled to the first inter-networking node, for transporting second type packets sequentially to at least a second edge node,

25F) the second edge node, operably coupled to the second switching type packet network, converts the second type packets to third type packets that are compatible with a third switching type packet network, and having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and

10 25G) destination source(s), operably coupled to the second edge node, for receiving the traffic.

26. The connection-oriented communication network system of claim 25 wherein the third type of packets are packets of the first type, and the third switching type packet network is a packet network of the first type.

27. A connection-oriented communication network system for ensuring portability of packet traffic from a first switching type network to a second switching type network, including:

27A) traffic source(s), for providing diverse traffic (e.g., voice, data, CBR (constant bit rate), video,...) to a first edge node of the first switching type packet network,

27B) the first edge node, having a first traffic adapter for performing predetermined traffic type specific adaption protocols to convert various traffic types from the traffic sources to first type packets compatible with the first switching type packet network, where the packets are also convertible one-to-one to second type packets that are compatible with the second switching type packet network, then converting the first type packets to second type packets that are compatible with the second switching type packet network,

27C) the second switching type packet network, operably coupled to the first edge node, for transporting second type packets sequentially to at least a second edge node,

27D) the second edge node, operably coupled to the second switching type packet network, converts the second type packets to third type packets that

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are compatible with a third switching type packet network, and having a second edge traffic adapter for performing predetermined traffic type specific adaption protocols to convert the third type packets to the original diverse traffic types that are compatible with destination source(s), and

- 5 27E) destination source(s), operably coupled to the second edge node, for receiving the traffic.
  - 28. The connection-oriented communication network system of claim 27 wherein the third type of packets are packets of the first type, and the third switching type packet network is a packet network of the first type.
- 10 29. A packet format adaptation method substantially as herein described with reference to the accompanying drawings.
  - 30. A connection-oriented communication network system substantially as herein described with reference to the accompanying drawings.

15 DATED: 11 March, 1997

PHILLIPS ORMONDE & FITZPATRICK

Attorneys for:

MOTOROLA, INC.



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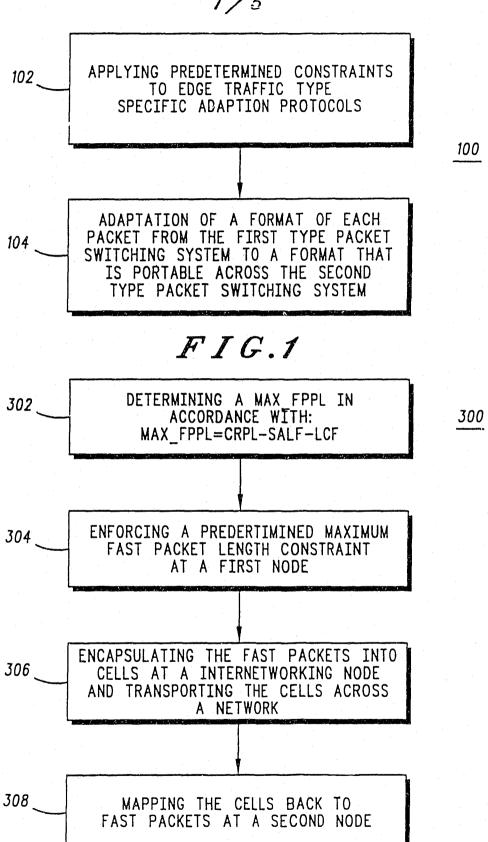


FIG.3

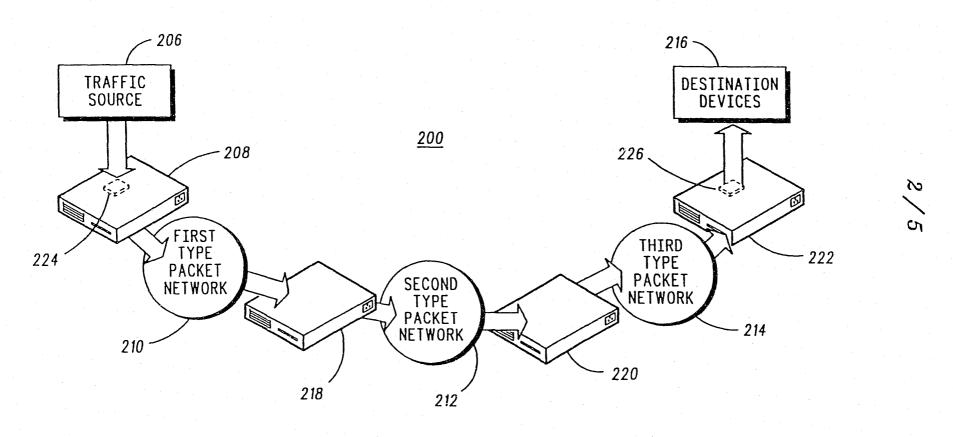
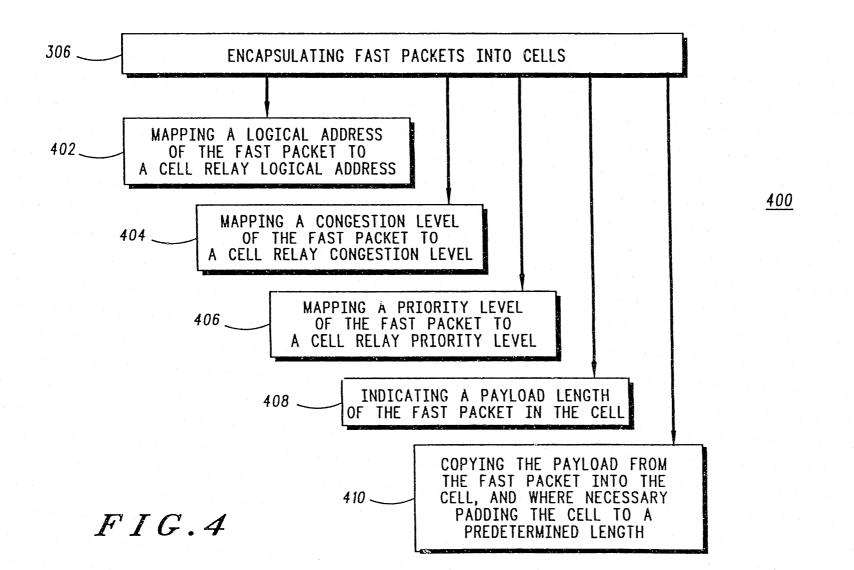


FIG.2



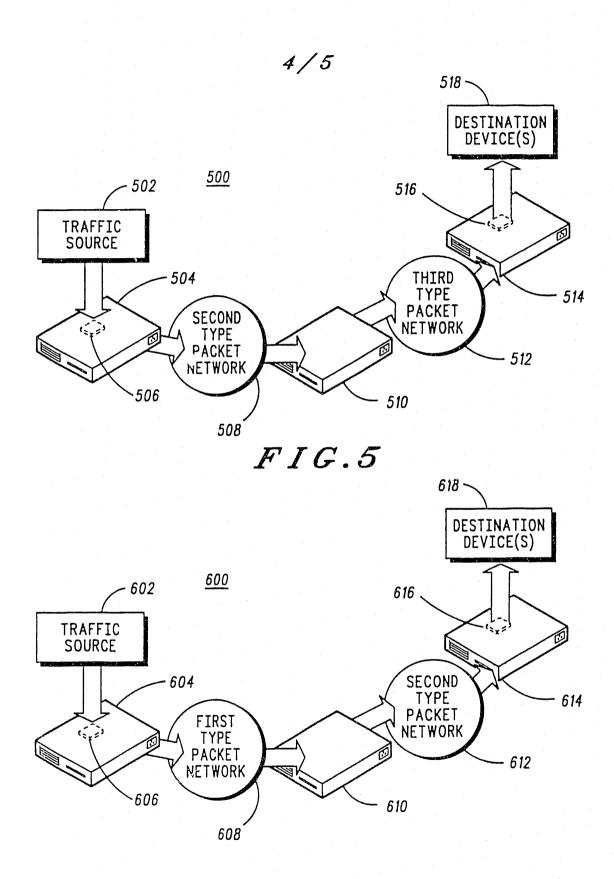


FIG.6

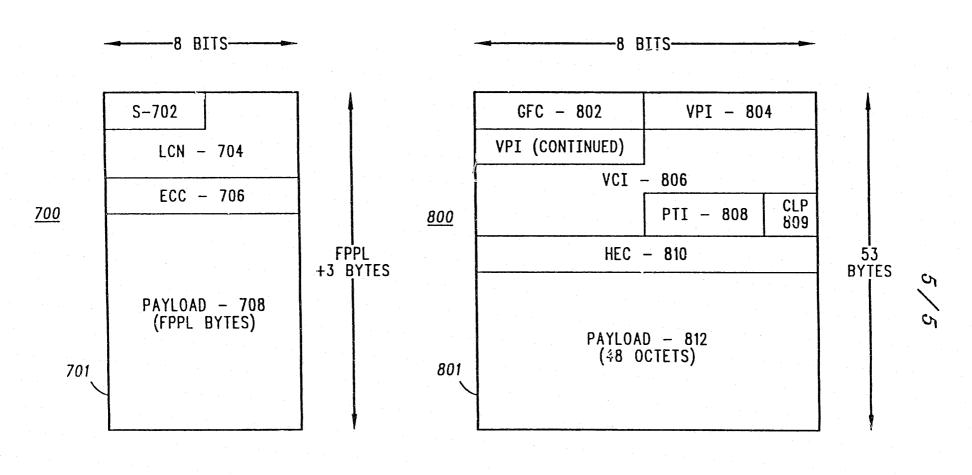


FIG.7

FIG.8

## INTERNATIONAL SEARCH REPORT

lnt ational application No.
PCT/US94/08703

A. CLASSIFICATION OF SUBJECT MATTER  IPC(5) :H04J 3/24, 3/00  US CL :370/94.1, 99  According to International Patent Classification (IPC) or to both national classification and IPC				
B. FIE	LDS SEARCHED			
Minimum o	documentation searched (classification system follow	ed by classification symbols)		
11 5. :	370/94.1, 99, 85.7, 85.13. 110.1, 54, 102, 85.6			
Documenta	tion searched other than minimum documentation to t	he extent that such documents are included	d in the fields searched	
Electronic o	data base consulted during the international search (	name of data base and, where practicable	, search terms used)	
APS search to	erms: packet, cell relay, format adaptation, pr	otocols, inter-network		
C. DOC	CUMENTS CONSIDERED TO BE RELEVANT			
Category*	Citation of document, with indication, where	appropriate, of the relevan passages	Relevant to claim No.	
X	US, A, 5,060,140 (BROWN et al.	) 22 October 1991 col 1	1	
	lines 12-13, 29-34, 48-51.	, 22 3010001 1001, 001. 1,		
Υ			2	
x	US, A, 5,291,485 (AFIFY et al.	) 01 March 1994 col 1	3	
,P	lines 47-52, col. 2, line 58 to col.	· · · · · · · · · · · · · · · · · · ·		
Υ ΄	37.		4	
Y,P	US, A, 5,224,099 (CORBALIS et lines 29-45.	al.) 29 June 1993, col. 1,	2, 4	
Y,P	US, A, 5,251,207 (ABENSOUR et 2, lines 25-42.	al.) 05 October 1993, col.	2, 4-10	
1				
X Furth	er documents are listed in the continuation of Box C	C. See patent family annex.		
Special categories of cited documents:  'A' document defining the general state of the art which is not considered		"T" later document published a fler the inte date and not in conflict with the applica principle or theory underlying the inve	tion but cited to understand the	
to be part of particular relevance  E* earlier document published on or after the international filing date  L* document which may throw doubts on priority claim(s) or which is		"X" document of particular relevance; the considered novel or cannot be consider when the document is taken alone		
cited	is reason (as specified)	"Y" document of particular relevance; the	claimed invention cannot be	
O* document referring to an oral disclosure, use, exhibition or other means		considered to involve an inventive combined with one or more other such being obvious to a person skilled in the	step when the document is documents, such combination	
	ument published prior to the international filing date but later than priority date claimed	*&* document member of the same patent family		
	ctual completion of the international search	Date of mailing of the international sear	rch report	
07 OCTOBER 1994		1 3 DEC 1994		
Name and mailing address of the ISA/US Authorized office				
Commissioner of Patents and Trademarks Box PCT		DOUGLAS OLMS		
Washington, D.C. 20231		Telephone No. (703) 305-4703		

Form PCT/ISA/210 (second sheet)(July 1992)\*

## INTERNATIONAL SEARCH REPORT

In. ational application No.
PCT/US94/08703

US, A, RE.33,426 (SUGIMOTO et al.) 06 November 1990, Fig. 5-10 1, col. 4, lines 11-58.	Category*	Citation of document, with indication, where appropriate, of the relevant passages		
US, A, 5,119,370 (TERRY) 02 June 1992,	? ?	US, A, RE.33,426 (SUGIMOTO et al.) 06 November 1990, Fig.		
		US, A, 5,119,370 (TERRY) 02 June 1992,	1-10	
	·			
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Form PCT/ISA/210 (continuation of second sheet)(July 1992)\*